Fifth Edition (February 2015)

This edition applies to Version 4 Release 5 of Enterprise PL/I for z/OS, 5655-W67, and to any subsequent releases until otherwise indicated in new editions or technical newsletters. Make sure you are using the correct edition for the level of the product.

Order publications through your IBM representative or the IBM branch office serving your locality. Publications are not stocked at the address below.

A form for readers’ comments is provided at the back of this publication. If the form has been removed, address your comments to:

IBM Corporation, Department H150/090
555 Bailey Ave
San Jose, CA, 95141-1099
United States of America

When you send information to IBM, you grant IBM a nonexclusive right to use or distribute the information in any way it believes appropriate without incurring any obligation to you.

© Copyright IBM Corporation 1999, 2015.

US Government Users Restricted Rights – Use, duplication or disclosure restricted by GSA ADP Schedule Contract with IBM Corp.
Contents

Tables ....................................... ix

Figures .................................... xi

Introduction ................................. xiii
About This Book .............................. xiii
Run-time environment for Enterprise PL/I for
z/OS ........................................... xiii
Using your documentation ................. xiii
PL/I information .............................. xiv
Notation conventions used in this book .. xiv
Conventions used ............................. xiv
How to read the syntax notation .......... xv
How to read the notational symbols ...... xvii
Summary of changes ........................ xviii
Enhancements in this release ............... xviii
Enhancements from V4R4 ..................... xix
Enhancements from V4R3 ............. xx
Enhancements from V4R2 ........... xxi
Enhancements from V4R1 .......... xxii
Enhancements from V3R9 .......... xiii
Enhancements from V3R8 .......... xxv
Enhancements from V3R7 .......... xxv
Enhancements from V3R6 .......... xxvi
Enhancements from V3R5 .......... xxviii
Enhancements from V3R4 .......... xxix
Enhancements from V3R3 .......... xxx
Enhancements from V3R2 .......... xxxi
Enhancements from V3R1 .......... xxxii
Enhancements from VisualAge PL/I .... xxiii
How to send your comments .............. xxxv
Accessibility .............................. xxxv
Interface information ..................... xxxv
Keyboard navigation ..................... xxxv
Accessibility of this information ...... xxxvi
IBM and accessibility ................... xxxvi

Part 1. Compiling your program ........................... 1

Chapter 1. Using compiler options and facilities ........................... 5
Compile-time option descriptions ........ 5
AGGREGATE ................................. 9
ARCH ..................................... 9
ATTRIBUTES .............................. 11
BACKREG ................................ 11
BIFPREC ................................. 11
BLANK ................................ 12
BLKOFF ................................ 13
CASERULES .................... 13
CEESTART .............................. 13
CHECK ................................ 14
CMAP ................................ 15
CODEPAGE .............................. 16
COMMON ................................. 16
COMPILE ................................ 17
COPYRIGHT ............................... 17
CSECT .................................. 18
CSECTCUT ............................... 18
CURRENCY ............................... 19
DBCS .................................. 19
DD .................................... 19
DDSQL ................................ 20
DECIMAL ............................... 20
DEFAULT ................................ 22
DEPREDATE ................................ 31
DEPRECATENEXT ..................... 32
DISPLAY ................................. 33
DLLINIT ................................. 33
EXIT .................................. 33
EXTRN ................................ 34
FILEREF ................................. 34
FLAG .................................. 34
FLOAT .................................. 35
FLOATINMATH ............................. 37
GOFF .................................. 37
GONUMBER ............................... 38
GRAPHIC ............................... 39
HGPR .................................. 39
IGNORE .................................. 40
INCAFTER ............................... 40
INCDIR ................................. 40
INCLUDE ................................ 41
INCPDS ................................. 41
INITAUTO ............................... 42
INITBASED ............................... 43
INITCTL ................................. 43
INITSINITIC ............................. 43
INSOURCE ......................... 43
INTERRUPT ......................... 44
JSON .................................. 45
LANGLEVEL ...................... 45
LIMITS .................................. 46
LINECOUNT ....................... 47
LINEDIR ............................... 47
LIST ................................ 47
LISTVIEW ............................... 48
MACRO ................................. 49
MAP .................................. 49
MARGINI ............................... 49
MARGINS ......................... 50
MAXBRANCH ......................... 51
MAXGEN ............................. 51
MAXMEM ............................. 52
MAXMSG ............................. 52
MAXNEST ......................... 53
MAXSTMT ......................... 53
MAXTEMP ......................... 53
MDELETE ......................... 54
MSGSUMMARY ..................... 54
© Copyright IBM Corp. 1999, 2015
| NAME .......... | 55 |
| NAMES .......... | 55 |
| NATLANG ........ | 55 |
| NEST .......... | 56 |
| NOT .......... | 56 |
| NUMBER .......... | 56 |
| OBJECT .......... | 57 |
| OFFSET .......... | 57 |
| ONSNAP .......... | 57 |
| OPTIMIZE ...... | 58 |
| OPTIONS .......... | 59 |
| OR .......... | 59 |
| PP .......... | 60 |
| PPCICS .......... | 61 |
| PPINCLUDE .... | 61 |
| PPLIST .......... | 62 |
| PPMACRO ........ | 62 |
| PPSQL .......... | 63 |
| PPTRACE .......... | 63 |
| PRECTYPE .......... | 64 |
| PROCESS .......... | 65 |
| QUOTE .......... | 65 |
| REDUCE .......... | 66 |
| RENT .......... | 66 |
| RESEXP .......... | 68 |
| RESPECT .......... | 68 |
| RTCHECK .......... | 68 |
| RULES .......... | 69 |
| SEMANTIC .......... | 77 |
| SERVICE .......... | 78 |
| SOURCE .......... | 78 |
| SPILL .......... | 78 |
| STATIC .......... | 78 |
| STDSYS .......... | 79 |
| STMT .......... | 79 |
| STORAGE .......... | 79 |
| STRINGOFGRAPHIC .......... | 80 |
| SYNTAX .......... | 80 |
| SYSPARM .......... | 81 |
| SYSTEM .......... | 81 |
| TERMINAL .......... | 82 |
| TEST .......... | 83 |
| UNROLL .......... | 85 |
| USAGE .......... | 86 |
| WIDECAR .......... | 87 |
| WINDOW .......... | 87 |
| WRITABLE .......... | 88 |
| XINFO .......... | 89 |
| XML .......... | 91 |
| XREF .......... | 91 |
| Blanks, comments and strings in options .......... | 92 |
| Changing the default options .......... | 92 |
| Specifying options in the %PROCESS or *PROCESS statements .......... | 93 |
| Using %PROCESS statements .......... | 93 |
| Using the %INCLUDE statement .......... | 94 |
| Using the compiler listing .......... | 94 |
| Heading information .......... | 96 |
| Options used for compilation .......... | 96 |
| Preprocessor input .......... | 97 |
| SOURCE program .......... | 97 |
| Statement nesting level .......... | 97 |
| ATTRIBUTE and cross-reference table .......... | 97 |
| Aggregate length table .......... | 98 |
| Statement offset addresses .......... | 99 |
| Storage offset listing .......... | 101 |
| File reference table .......... | 102 |
| Messages and return codes .......... | 102 |
| Example .......... | 104 |

## Chapter 2. PL/I preprocessors .......... 109

| Include preprocessor .......... | 109 |
| Macro preprocessor .......... | 110 |
| Macro preprocessor options .......... | 110 |
| Macro preprocessor example .......... | 113 |
| SQL preprocessor .......... | 114 |
| Programming and compilation considerations .......... | 115 |
| SQL preprocessor options .......... | 117 |
| Coding SQL statements in PL/I applications .......... | 120 |
| General information about LOBs .......... | 131 |
| Suppressing SQL preprocessor messages .......... | 135 |
| CICS Preprocessor .......... | 136 |
| Programming and compilation considerations .......... | 136 |
| CICS preprocessor options .......... | 137 |
| Coding CICS statements in PL/I applications .......... | 137 |
| Writing CICS transactions in PL/I .......... | 137 |
| Error-handling .......... | 138 |

## Chapter 3. Using PL/I cataloged procedures .......... 139

| IBM-supplied cataloged procedures .......... | 139 |
| Compile only (IBMZC) .......... | 140 |
| Compile and bind (IBMZCB) .......... | 141 |
| Compile, bind, and run (IBMZCBG) .......... | 143 |
| Compile, prelink, and link-edit (IBMZCPL) .......... | 145 |
| Compile, prelink, link-edit, and run (IBMZCPLG) .......... | 146 |
| Compile, prelink, load and run (IBMZCPLG) .......... | 148 |
| Invoking a cataloged procedure .......... | 150 |
| Specifying multiple invocations .......... | 150 |
| Modifying the PL/I cataloged procedures .......... | 151 |
| EXEC statement .......... | 152 |
| DD statement .......... | 152 |

## Chapter 4. Compiling your program .......... 155

| Invoking the compiler under z/OS UNIX .......... | 155 |
| Input files .......... | 155 |
| Specifying compile-time options under z/OS UNIX .......... | 156 |
| -qoption_keyword .......... | 156 |
| Single and multiletter flags .......... | 156 |
| Invoking the compiler under z/OS using JCL .......... | 157 |
| EXEC statement .......... | 157 |
| DD statements for the standard data sets .......... | 158 |
| Listing (SYSPRINT) .......... | 159 |
| Source Statement Library (SYSLIB) .......... | 160 |
| Specifying options .......... | 160 |
| Specifying options in the EXEC statement .......... | 161 |
Specifying options in the EXEC statement using an options file ........................................ 161

Chapter 5. Link-editing and running .................................................. 163

Link-edit considerations ......................................................... 163
Using the binder ................................................................ 163
Using the prelinker ................................................................ 163
Using the ENTRY card ............................................................. 164
Run-time considerations ......................................................... 164
Formatting conventions for PRINT files .................................. 164
Changing the format on PRINT files ....................................... 164
Automatic prompting ............................................................... 165
Punctuating long input lines ................................................... 166
Punctuating GET LIST and GET DATA statements ................. 166
ENDFILE ..................................................................... 167
SYSPRINT considerations ...................................................... 167
Using MSGFILE(SYSPRINT) .................................................. 169
Using FETCH in your routines ................................................ 169
Fetching Enterprise PL/I routines .......................................... 169
Fetching PL/I MAIN routines .............................................. 178
Fetching z/OS C routines ..................................................... 179
Fetching assembler routines ................................................ 179
Invoking MAIN under TSO/E ................................................ 179
Invoking MAIN under z/OS UNIX ........................................... 180

Part 2. Using I/O facilities .............................................................. 183

Chapter 6. Using data sets and files .............................................. 185
Allocating files ...................................................................... 185
Associating data sets with files under z/OS ......................... 186
  Associating several files with one data set ......................... 188
  Associating several data sets with one file ....................... 189
  Concatenating several data sets ....................................... 189
  Accessing HFS files under z/OS ........................................ 189
Associating data sets with files under z/OS UNIX ................. 191
  Using environment variables ........................................... 191
  Using the TITLE option of the OPEN statement ............... 192
  Attempting to use files not associated with data sets ......... 193
  How PL/I finds data sets ................................................ 193
  Specifying characteristics using DD_DDNAME environment variables ..................................... 194
Establishing data set characteristics .................................... 199
  Blocks and records .......................................................... 200
  Record formats ............................................................... 200
  Data set organization ...................................................... 202
  Labels ........................................................................ 203
  Data Definition (DD) statement ...................................... 203
  Using the TITLE option of the OPEN statement ............... 204
  Associating PL/I files with data sets ................................. 205
Specifying characteristics in the ENVIRONMENT attribute .... 207

Chapter 7. Using libraries .............................................................. 217
Types of libraries ................................................................. 217
  How to use a library ........................................................ 217
Creating a library ............................................................... 218
  SPACE parameter ........................................................ 218
Creating and updating a library member .............................. 219

Chapter 8. Defining and using consecutive data sets ................. 223
Using stream-oriented data transmission .............................. 223
  Defining files using stream I/O ........................................ 223
  Defining stream files using PL/I dynamic allocation .......... 224
Specifying ENVIRONMENT options ...................................... 224
Creating a data set with stream I/O ..................................... 226
Accessing a data set with stream I/O .................................... 230
Using PRINT files with stream I/O ...................................... 231
Using SYSIN and SYSPRINT files ........................................ 235
Controlling input from the terminal ..................................... 236
  Format of data ............................................................. 237
  Stream and record files .................................................. 238
  Defining QSAM files using PL/I dynamic allocation ......... 238
  Capital and lowercase letters ......................................... 239
End-of-file ..................................................................... 239
COPY option of GET statement .......................................... 239

Chapter 9. Controlling output to the terminal ............................ 241
Format of PRINT files .......................................................... 241
Stream and record files ...................................................... 241
Output from the PUT EDIT command ................................. 241

Chapter 10. Using record-oriented data transmission .................. 243
Specifying record format ..................................................... 244
Defining files using record I/O ............................................. 244
Specifying ENVIRONMENT options .................................... 244
  CONSECUTIVE ......................................................... 245
  ORGANIZATION(CONSECUTIVE) .................... 245
  CTLASA | CTL360 ................................................. 245
  LEAVE | REREAD .................................................. 246
Creating a data set with record I/O ...................................... 247
  Essential information .................................................... 248
Accessing and updating a data set with record I/O ............... 248
  Essential information .................................................... 249
  Example of consecutive data sets .................................... 249

Chapter 11. Defining and using regional data sets ..................... 253
Defining REGIONAL(1) data sets using PL/I dynamic allocation .............................. 255
Defining files for a regional data set .................................... 255
  Specifying ENVIRONMENT options ......................... 255
  Using keys with REGIONAL data sets ......................... 256
Using REGIONAL(1) data sets ........................................... 256
  Dummy Records ......................................................... 257
Creating a REGIONAL(1) data set ..................................... 257
Accessing and updating a REGIONAL(1) data set ............... 258
  Essential information for creating and accessing regional data sets ........................................ 261

Examples ........................................................................ 219
Extracting information from a library directory ..................... 221

Chapter 12. Using record-oriented data transmission ................. 243
Specifying record format ..................................................... 244
Defining files using record I/O ............................................. 244
Specifying ENVIRONMENT options .................................... 244
  CONSECUTIVE ......................................................... 245
  ORGANIZATION(CONSECUTIVE) .................... 245
  CTLASA | CTL360 ................................................. 245
  LEAVE | REREAD .................................................. 246
Creating a data set with record I/O ...................................... 247
  Essential information .................................................... 248
Accessing and updating a data set with record I/O ............... 248
  Essential information .................................................... 249
  Example of consecutive data sets .................................... 249

Chapter 13. Defining and using regional data sets ..................... 253
Defining REGIONAL(1) data sets using PL/I dynamic allocation .............................. 255
Defining files for a regional data set .................................... 255
  Specifying ENVIRONMENT options ......................... 255
  Using keys with REGIONAL data sets ......................... 256
Using REGIONAL(1) data sets ........................................... 256
  Dummy Records ......................................................... 257
Creating a REGIONAL(1) data set ..................................... 257
Accessing and updating a REGIONAL(1) data set ............... 258
  Essential information for creating and accessing regional data sets ........................................ 261
Chapter 12. Defining and using VSAM data sets ......... 265
Defining VSAM file using PL/I dynamic allocation 265
Using VSAM data sets .................................................. 265
How to run a program with VSAM data sets ............. 265
Pairing an Alternate Index Path with a File .................. 266
VSAM organization ......................................................... 266
Keys for VSAM data sets ............................................ 268
Choosing a data set type .............................................. 269
Defining files for VSAM data sets ................................. 271
Specifying ENVIRONMENT options ............................... 271
Performance options ..................................................... 274
Defining Files for Alternate Index Paths ......................... 274
Defining VSAM data sets ............................................. 275
Entry-sequenced data sets ........................................... 276
Loading an ESDS .......................................................... 276
Using a SEQUENTIAL file to access an ESDS ................ 277
Using Files Defined for non-VSAM Data Sets .................. 299
Part 3. Improving your program .............................. 301
Chapter 13. Improving performance ..................... 303
Selecting compiler options for optimal performance ............. 303
OPTIMIZE ................................................................. 303
GONUMBER .............................................................. 303
ARCH .................. .......................... .......................... 303
REDUCE ................................................................. 304
RULES ................................................................. 304
PREFIX ................................................................. 305
DEFAULT ................................................................. 306
Summary of compiler options that improve performance ........ 309
Coding for better performance ........................................... 309
DATA-directed input and output ................................... 309
Input-only parameters ................................................. 310
GOTO statements ....................................................... 310
String assignments ..................................................... 310
Loop control variables ............................................... 311
PACKAGEs versus nested PROCEDUREs ......................... 311
REPLACEABLE Functions ........................................... 312
DESCLOCATOR or DESCLIST .................................. 313
DEFINED versus UNION .......................................... 313
Named constants versus static variables ..................... 313
Avoiding calls to library routines .................................. 315
Preloading library routines ......................................... 315
Part 4. Using interfaces to other products ............ 317
Chapter 14. Using the Sort program .................... 319
Preparing to use Sort ................................................. 319
Choosing the type of Sort ........................................... 320
Specifying the sorting field ....................................... 323
Specifying the records to be sorted ......................... 324
Determining storage needed for Sort .......................... 325
Calling the Sort program ............................................ 326
Example 1 ............................................................... 327
Example 2 ............................................................... 328
Example 3 ............................................................... 328
Example 4 ............................................................... 328
Example 5 ............................................................... 328
Determining whether the Sort was successful .............. 328
Establishing data sets for Sort .................................. 329
Sort data input and output ...................................... 330
Data input and output handling routines .................... 330
E15—Input handling routine (Sort Exit E15) .................. 331
E35—Output handling routine (Sort Exit E35) ............... 334
Calling PLISRTA example ......................................... 335
Calling PLISRTH example .......................................... 336
Calling PLISRTE example ......................................... 337
Calling PLISRTD example ......................................... 338
Sorting variable-length records example ................... 339
Chapter 15. ILC with C ...................................... 341
Equivalent data types ................................................. 341
Simple type equivalence ........................................... 341
Struct type equivalence ............................................ 342
Enum type equivalence ............................................. 342
File type equivalence ............................................... 343
Using C functions ...................................................... 343
Matching simple parameter types .............................. 344
Matching string parameter types ................................. 347
Functions returning ENTRYs ....................................... 348
Linkages ................................................................. 349
Sharing output and input .......................................... 351
Sharing output .......................................................... 351
Sharing input ........................................................... 352
Using the ATTACH statement ................................... 352
Redirecting C standard streams ................................. 352
Summary ................................................................. 352
Chapter 16. Interfacing with Java ...................... 355
Java Native Interface (JNI) ........................................ 355
Calling PL/I program from Java .................................. 356
JNI Sample Program #1 - 'Hello World' ....................... 356
Writing Java Sample Program #1 ............................... 356
JNI Sample Program #2 - Passing a String ................... 359
Writing Java Sample Program #2 ............................... 359
JNI Sample Program #3 - Passing an Integer ................. 364
Writing Java Sample Program #3 ............................... 364
JNI Sample Program #4 - Java Invocation API ............... 368
Writing Java Sample Program #4 ............................... 368
Attaching programs to an existing Java VM .................. 372
Determining equivalent Java and PL/I data types .......... 373
Part 5. Specialized programming tasks ................ 375
Chapter 17. Using the PLISAXA and PLISAXB XML parsers .... 377
Overview ................................................................. 377
The PLISAXA built-in subroutine ................................ 378
The PLISAXB built-in subroutine ................................ 378
The SAX event structure ............................................. 379
start_of_document .................................................... 379
version_information .................................................. 380
encoding_declaration ............................................... 380
Contents

Chapter 18. Using the PLISAXC and PLISAXD XML parsers .......................... 407

Overview ........................................ 407
The PLISAXC built-in subroutine .................. 408
The PLISAXD built-in subroutine .................. 408
The SAX event structure ......................... 409

Parameters to the event functions .................. 409

Coded character sets for XML documents .......... 410
Supported EBCDIC code pages ....................... 410
Supported ASCII code pages ........................ 410
Specifying the code page ............................ 410

Exceptions ........................................ 410

Example .......................................... 410

Continuable exception codes ....................... 410

Terminating exception codes ....................... 410

Chapter 19. Using PLIDUMP ......................... 441

PLIDUMP usage notes............................. 442
Locating variables in the PLIDUMP output .......... 442
Locating AUTOMATIC variables ..................... 443
Locating STATIC variables ......................... 443
Locating CONTROLLED variables ................. 445
Saved Compilation Data ............................. 448
Copyright ......................................... 448
Timestamp ......................................... 449
Saved options string ............................... 449

Chapter 20. Interrupts and attention processing ............................... 455

Using ATTENTION ON-units ....................... 456
Interaction with a debugging tool .................. 456

Chapter 21. Using the Checkpoint/Restart facility ...................... 457

Requesting a checkpoint record .................... 457
Defining the checkpoint data set ................. 458
Requesting a restart ............................... 459
Automatic restart after a system failure ........... 459
Automatic restart within a program ............... 459
Getting a deferred restart ......................... 459
Modifying checkpoint/restart activity ............ 460

Chapter 22. Using user exits .......................... 461

Procedures performed by the compiler user exit .... 461
Activating the compiler user exit .................. 462
The IBM-supplied compiler exit, IBMUEXIT ........ 462
Customizing the compiler user exit ............... 462
Modifying SYSEXIT ................................ 463
Writing your own compiler exit .................... 463
Structure of global control blocks ................. 463
Writing the initialization procedure ............... 464
Writing the message filtering procedure .......... 465
Writing the termination procedure ................. 467
Example of suppressing SQL messages ............. 468

Chapter 23. PL/I descriptors ........................ 475

Passing an argument .............................. 475
Argument passing by descriptor list ............... 475
Argument passing by descriptor-locator ........... 476
CMPAT(V*) descriptors ........................... 476
String descriptors ............................... 476
Array descriptors .................................. 478
CMPAT(LE) descriptors ......................... 478
String descriptors ............................... 479
Array descriptors .................................. 479
# Tables

1. How to use Enterprise PL/I publications .......................... xiii
2. How to use z/OS Language Environment publications .... xiv
3. Compile-time options, abbreviations, and IBM-supplied defaults ... 5
4. SYSTEM option table .............................................. 82
5. Description of PL/I error codes and return codes .................. 103
6. Using the FLAG option to select the lowest message severity listed ... 103
7. SQL preprocessor options and IBM-supplied defaults ............ 117
8. SQL data types generated from PL/I declarations ................ 127
9. SQL data types generated from SQL TYPE declarations ........... 127
10. SQL data types mapped to PL/I declarations ...................... 128
11. SQL data types mapped to SQL TYPE declarations ............... 129
12. Compile-time option flags supported by Enterprise PL/I under z/OS UNIX ... 157
13. Compiler standard data sets ...................................... 158
14. Attributes of PL/I file declarations ................................ 207
15. A comparison of data set types available to PL/I record I/O .... 215
16. Information required when creating a library ...................... 218
17. Statements and options allowed for creating and accessing consecutive data sets ... 243
18. IBM machine code print control characters (CTL360) ............ 246
19. Effect of LEAVE and REREAD Options ............................ 247
20. Creating a consecutive data set with record I/O: essential parameters of the DD statement ... 247
21. Accessing a consecutive data set with record I/O: essential parameters of the DD statement ... 249
22. Statements and options allowed for creating and accessing regional data sets .......... 254
23. Creating a regional data set: essential parameters of the DD statement ........ 262
24. DCB subparameters for a regional data set ....................... 262
25. Accessing a regional data set: essential parameters of the DD statement ........ 263
26. Types and advantages of VSAM data sets ......................... 268
27. Processing Allowed on Alternate Index Paths ..................... 271
28. Statements and options allowed for loading and accessing VSAM entry-sequenced data sets .......... 276
29. Statements and options allowed for loading and accessing VSAM indexed data sets .......... 279
30. Statements and options allowed for loading and accessing VSAM relative-record data sets ........ 293
31. The entry points and arguments to PLISRTx (x = A, B, C, or D) .... 326
32. C and PL/I Type Equivalents .................................... 341
33. Java Primitive Types and PL/I Native Equivalents ............... 373
34. Continuable Exceptions ........................................ 398
35. Terminating Exceptions ........................................ 401

© Copyright IBM Corp. 1999, 2015
Figures

1. Including source statements from a library 95
2. Finding statement number (compiler listing example) 100
3. Finding statement number (runtime message example) 100
4. Compiler listing example 105
5. Using the macro preprocessor to produce a source deck 114
6. The PL/I declaration of SQLCA 121
7. The PL/I declaration of an SQL descriptor area 122
8. pliclob sample program 134
9. Invoking a cataloged procedure 140
10. Cataloged Procedure IBMZC 141
11. Cataloged Procedure IBMZCB 142
12. Cataloged Procedure IBMZCBG 144
13. Cataloged Procedure IBMZCPL 145
14. Cataloged Procedure IBMZCPLG 147
15. Cataloged Procedure IBMZCPG 149
16. Declaration of PLITABS 165
17. PAGELENGTH and PAGESIZE 165
18. Sample JCL to compile, link, and invoke the user exit 172
19. Sample program to display program arguments from the CPPL under TSO when using SYSTEM(STD) option 180
20. Sample program to display z/OS UNIX args and environment variables 181
21. Fixed-length records 201
22. How the operating system completes the DCB 206
23. Creating new libraries for compiled object modules 220
24. Placing a load module in an existing library 220
25. Creating a library member in a PL/I program 221
26. Updating a library member 221
27. Creating a data set with stream-oriented data transmission 228
28. Writing graphic data to a stream file 229
29. Accessing a data set with stream-oriented data transmission 231
30. Creating a print file via stream data transmission 234
31. PL/I structure PLITABS for modifying the preset tab settings 235
32. American National Standard print and card punch control characters (CTLASA) 246
33. Merge Sort—creating and accessing a consecutive data set 250
34. Printing record-oriented data transmission 252
35. Creating a REGIONAL(1) data set 258
36. Updating a REGIONAL(1) data set 260
37. VSAM data sets and allowed file attributes 270
38. Defining and loading an entry-sequenced data set (ESDS) 278
39. Updating an ESDS 279
40. Defining and loading a key-sequenced data set (KSDS) 282
41. Updating a KSDS 284
42. Creating a Unique Key Alternate Index Path for an ESDS 286
43. Creating a Nonunique Key Alternate Index Path for an ESDS 287
44. Creating a unique Key Alternate Index Path for a KSDS 288
45. Alternate Index Paths and Backward Reading with an ESDS 290
46. Using a Unique Alternate Index Path to Access a KSDS 292
47. Defining and loading a relative-record data set (RRDS) 296
48. Updating an RRDS 298
49. Flow of control for Sort program 322
50. Flowcharts for input and output handling subroutines 332
51. Skeletal code for an input procedure 333
52. Skeletal code for an output handling procedure 334
53. PLISRTA—sorting from input data set to output data set 335
54. PLISRTB—sorting from input handling routine to output data set 336
55. PLISRTC—sorting from input data set to output handling routine 337
56. PLISRTD—sorting from input handling routine to output handling routine 338
57. Sorting varying-length records using input and output handling routines 339
58. Simple type equivalence 341
59. Sample struct type equivalence 342
60. Sample enum type equivalence 343
61. Start of the C declaration for its FILE type 343
62. PL/I equivalent for a C file 343
63. Sample code to use fopen and fread to dump a file 344
64. Declarations for filedump program 344
65. C declaration of fread 345
66. First incorrect declaration of fread 345
67. Second incorrect declaration of fread 345
68. Third incorrect declaration of fread 345
69. Code generated for RETURNS BYADDR 346
70. Correct declaration of fread 346
71. Code generated for RETURNS BYVALUE 346
72. First incorrect declaration of fopen 347
73. Second incorrect declaration of fopen 347
74. Correct declaration of fopen 347
75. Optimal, correct declaration of fopen 347
76. Declaration of fclose 348
77. Commands to compile and run filedump 348
78. Output of running filedump 348
79. Sample compare routine for C qsort function 348
80. Sample code to use C qsort function 349

© Copyright IBM Corp. 1999, 2015
81. Incorrect declaration of qsort  . . . . . . . . . . . . . . . . . . . . 349
82. Correct declaration of qsort  . . . . . . . . . . . . . . . . . . . . . . 349
83. Code when parameters are BYADDR 350
84. Code when parameters are BYVALUE 351
85. Java Sample Program #2 - Passing a String 361
86. PL/I Sample Program #2 - Passing a String 363
87. Java Sample Program #3 - Passing an Integer 365
88. PL/I Sample Program #3 - Passing an Integer 367
89. Java Sample Program #4 - Receiving and printing a String 368
90. PL/I Sample Program #4 - Calling the Java Invocation API 371
91. Sample XML document  . . . . . . . . . . . . . . . . . . . . . . . . . . 379
92. PLISAXA coding example - type declarations 387
93. PLISAXA coding example - event structure 389
94. PLISAXA coding example - main routine 390
95. PLISAXA coding example - event routines 398
96. PLISAXA coding example - program output 398
97. Sample XML document  . . . . . . . . . . . . . . . . . . . . . . . . . . 409
98. PLISAXC coding example - type declarations 420
99. PLISAXC coding example - event structure 421
100. PLISAXC coding example - main routine 422
101. PLISAXC coding example - event routines 423
102. PLISAXC coding example - program output 429
103. PLISAXD coding example - event routines 430
104. Output from PLISAXD sample  . . . . . . . . . . . . . . . . . . . . . 440
105. Example PL/I routine calling PLIDUMP 441
106. Declare for the saved options string  . . . . . . . . . . . . . . . . . 450
107. Using an ATTENTION ON-unit  . . . . . . . . . . . . . . . . . . . . . 456
108. PL/I compiler user exit procedures  . . . . . . . . . . . . . . . . . . 462
109. Example of an user exit input file  . . . . . . . . . . . . . . . . . . . 463
110. Suppressing SQL messages  . . . . . . . . . . . . . . . . . . . . . . . 468
111. Record types encoded as an ordinal value 484
112. Declare for the header part of a record  . . . . . . . . . . . . . . . . 484
113. Declare for a summary record  . . . . . . . . . . . . . . . . . . . . . 485
114. Declare for a counter record  . . . . . . . . . . . . . . . . . . . . . . 485
115. Declare for a literal record  . . . . . . . . . . . . . . . . . . . . . . . 486
116. Declare for a file record  . . . . . . . . . . . . . . . . . . . . . . . . 486
117. Declare for a message record  . . . . . . . . . . . . . . . . . . . . . . 487
118. Declare for an ordinal type record  . . . . . . . . . . . . . . . . . . . 488
119. Declare for an ordinal element record 489
120. Symbol indices assigned to the elements of a structure  . . . . . . . . . 490
121. Data type of a variable  . . . . . . . . . . . . . . . . . . . . . . . . . . 491
122. Declare for a symbol record  . . . . . . . . . . . . . . . . . . . . . . . 492
123. Declare for xin_Bif.Kind  . . . . . . . . . . . . . . . . . . . . . . . . . 499
124. Declare for a source record  . . . . . . . . . . . . . . . . . . . . . . . 506
125. Declare for a token record  . . . . . . . . . . . . . . . . . . . . . . . . 506
126. Declare for the token record kind  . . . . . . . . . . . . . . . . . . . . 507
127. Node indices assigned to the blocks in a program  . . . . . . . . . . . . . 507
128. Declare for a syntax record  . . . . . . . . . . . . . . . . . . . . . . . 508
129. Declare for the syntax record kind  . . . . . . . . . . . . . . . . . . . 511
130. Node indices assigned to the syntax records in a program  . . . . . . . . . . . 512
131. Declare for the expression kind  . . . . . . . . . . . . . . . . . . . . . 513
132. Declare for the number kind  . . . . . . . . . . . . . . . . . . . . . . . 513
133. Declare for the lexeme kind  . . . . . . . . . . . . . . . . . . . . . . . 514
134. Declare for the voc kind  . . . . . . . . . . . . . . . . . . . . . . . . . 515
Introduction

About This Book

This book is for PL/I programmers and system programmers. It helps you understand how to use Enterprise PL/I for z/OS in order to compile PL/I programs. It also describes the operating system features that you might need to optimize program performance or handle errors.

**Important:** Enterprise PL/I for z/OS will be referred to as Enterprise PL/I throughout this book.

Run-time environment for Enterprise PL/I for z/OS

Enterprise PL/I uses Language Environment® as its run-time environment. It conforms to Language Environment architecture and can share the run-time environment with other Language Environment-conforming languages.

Language Environment provides a common set of run-time options and callable services. It also improves interlanguage communication (ILC) between high-level languages (HLL) and assembler by eliminating language-specific initialization and termination on each ILC invocation.

Using your documentation

The publications provided with Enterprise PL/I are designed to help you program with PL/I. The publications provided with Language Environment are designed to help you manage your run-time environment for applications generated with Enterprise PL/I. Each publication helps you perform a different task.

The following tables show you how to use the publications you receive with Enterprise PL/I and Language Environment. You’ll want to know information about both your compiler and run-time environment. For the complete titles and order numbers of these and other related publications, see “Bibliography” on page 533.

### PL/I information

**Table 1. How to use Enterprise PL/I publications**

<table>
<thead>
<tr>
<th>To...</th>
<th>Use...</th>
</tr>
</thead>
<tbody>
<tr>
<td>Evaluate Enterprise PL/I</td>
<td>Fact Sheet</td>
</tr>
<tr>
<td>Understand warranty information</td>
<td>Licensed Programming Specifications</td>
</tr>
<tr>
<td>Plan for and install Enterprise PL/I</td>
<td>Enterprise PL/I Program Directory</td>
</tr>
<tr>
<td>Understand compiler and run-time changes and adapt programs to Enterprise PL/I and Language Environment</td>
<td>Compiler and Run-Time Migration Guide</td>
</tr>
<tr>
<td>Prepare and test your programs and get details on compiler options</td>
<td>Programming Guide</td>
</tr>
<tr>
<td>Get details on PL/I syntax and specifications of language elements</td>
<td>Language Reference</td>
</tr>
</tbody>
</table>
Table 1. How to use Enterprise PL/I publications (continued)

<table>
<thead>
<tr>
<th>To...</th>
<th>Use...</th>
</tr>
</thead>
<tbody>
<tr>
<td>Diagnose compiler problems and report them to IBM</td>
<td>Diagnosis Guide</td>
</tr>
<tr>
<td>Get details on compile-time messages</td>
<td>Compile-Time Messages and Codes</td>
</tr>
</tbody>
</table>

Language Environment information

Table 2. How to use z/OS Language Environment publications

<table>
<thead>
<tr>
<th>To...</th>
<th>Use...</th>
</tr>
</thead>
<tbody>
<tr>
<td>Evaluate Language Environment</td>
<td>Concepts Guide</td>
</tr>
<tr>
<td>Install Language Environment on z/OS</td>
<td>z/OS Program Directory</td>
</tr>
<tr>
<td>Customize Language Environment on z/OS</td>
<td>Customization</td>
</tr>
<tr>
<td>Understand Language Environment program models and concepts</td>
<td>Concepts Guide Programming Guide</td>
</tr>
<tr>
<td>Find syntax for Language Environment run-time options and callable services</td>
<td>Programming Reference</td>
</tr>
<tr>
<td>Develop applications that run with Language Environment</td>
<td>Programming Guide and your language Programming Guide</td>
</tr>
<tr>
<td>Debug applications that run with Language Environment, get details on run-time messages, diagnose problems with Language Environment</td>
<td>Debugging Guide and Run-Time Messages</td>
</tr>
<tr>
<td>Develop interlanguage communication (ILC) applications</td>
<td>Writing Interlanguage Applications</td>
</tr>
<tr>
<td>Migrate applications to Language Environment</td>
<td>Run-Time Application Migration Guide and the migration guide for each Language Environment-enabled language</td>
</tr>
</tbody>
</table>

Notation conventions used in this book

This book uses the conventions, diagramming techniques, and notation described in "Conventions used" and "How to read the notational symbols" on page xvii to illustrate PL/I and non-PL/I programming syntax.

Conventions used

Some of the programming syntax in this book uses type fonts to denote different elements:

- Items shown in UPPERCASE letters indicate key elements that must be typed exactly as shown.
- Items shown in lowercase letters indicate user-supplied variables for which you must substitute appropriate names or values. The variables begin with a letter and can include hyphens, numbers, or the underscore character (_).
- The term digit indicates that a digit (0 through 9) should be substituted.
- The term do-group indicates that a do-group should be substituted.
- Underlined items indicate default options.
- Examples are shown in monocase type.
• Unless otherwise indicated, separate repeatable items from each other by one or more blanks.

**Note:** Any symbols shown that are not purely notational, as described in “How to read the notational symbols” on page xvii, are part of the programming syntax itself.

For an example of programming syntax that follows these conventions, see “Example of notation” on page xvii.

### How to read the syntax notation

The following rules apply to the syntax diagrams used in this book:

#### Arrow symbols
Read the syntax diagrams from left to right, from top to bottom, following the path of the line.

- ➤ Indicates the beginning of a statement.
- ➤ Indicates that the statement syntax is continued on the next line.
- ➤ Indicates that a statement is continued from the previous line.
- ➤ Indicates the end of a statement.

Diagrams of syntactical units other than complete statements start with the ➤ symbol and end with the ➤ symbol.

#### Conventions
- Keywords, their allowable synonyms, and reserved parameters, appear in uppercase for MVS and OS/2® platforms, and lowercase for UNIX® platforms. These items must be entered exactly as shown.
- Variables appear in lowercase italics (for example, column-name). They represent user-defined parameters or suboptions.
- When entering commands, separate parameters and keywords by at least one blank if there is no intervening punctuation.
- Enter punctuation marks (slashes, commas, periods, parentheses, quotation marks, equal signs) and numbers exactly as given.
- Footnotes are shown by a number in parentheses, for example, (1).
- A \_ symbol indicates one blank position.

#### Required items
Required items appear on the horizontal line (the main path).

- ➤ REQUIRED_ITEM ➤

#### Optional items
Optional items appear below the main path.

- ➤ REQUIRED_ITEM ➤ _optional_item_

If an optional item appears above the main path, that item has no effect on the execution of the statement and is used only for readability.
Multiple required or optional items
If you can choose from two or more items, they appear vertically in a stack. If you must choose one of the items, one item of the stack appears on the main path.

If choosing one of the items is optional, the entire stack appears below the main path.

Repeattable items
An arrow returning to the left above the main line indicates that an item can be repeated.

If the repeat arrow contains a comma, you must separate repeated items with a comma.

A repeat arrow above a stack indicates that you can specify more than one of the choices in the stack.

Default keywords
IBM-supplied default keywords appear above the main path, and the remaining choices are shown below the main path. In the parameter list following the syntax diagram, the default choices are underlined.

Fragments
Sometimes a diagram must be split into fragments. The fragments are represented by a letter or fragment name, set off like this: A. The fragment follows the end of the main diagram. The following example shows the use of a fragment.
Substitution-block

Sometimes a set of several parameters is represented by a substitution-block such as `<A>`. For example, in the imaginary /VERB command you could enter /VERB LINE 1, /VERB EITHER LINE 1, or /VERB OR LINE 1.

```
/VERB <A> LINE line#
```

where `<A>` is:

```
EITHER
OR
```

Parameter endings

Parameters with number values end with the symbol '#', parameters that are names end with 'name', and parameters that can be generic end with '*'.

```
/MSVERIFY MSNAME msname SYSID sysid#
```

The MSNAME keyword in the example supports a name value and the SYSID keyword supports a number value.

How to read the notational symbols

Some of the programming syntax in this book is presented using notational symbols. This is to maintain consistency with descriptions of the same syntax in other IBM publications, or to allow the syntax to be shown on single lines within a table or heading.

- **Braces**, `{ }`, indicate a choice of entry. Unless an item is underlined, indicating a default, or the items are enclosed in brackets, you must choose at least one of the entries.

- Items separated by a single **vertical bar**, `|`, are alternative items. You can select only one of the group of items separated by single vertical bars. (Double vertical bars, `||`, specify a concatenation operation, not alternative items. See the PL/I Language Reference for more information on double vertical bars.)

- Anything enclosed in **brackets**, `[ ]`, is optional. If the items are vertically stacked within the brackets, you can specify only one item.

- An **ellipsis**, `...`, indicates that multiple entries of the type immediately preceding the ellipsis are allowed.

Example of notation

The following example of PL/I syntax illustrates the notational symbols described in "How to read the notational symbols"
DCL file-reference FILE STREAM
   {INPUT | OUTPUT [PRINT]}
   ENVIRONMENT(option ...);

Interpret this example as follows:
- You must spell and enter the first line as shown, except for file-reference, for which you must substitute the name of the file you are referencing.
- In the second line, you can specify INPUT or OUTPUT, but not both. If you specify OUTPUT, you can optionally specify PRINT as well. If you do not specify either alternative, INPUT takes effect by default.
- You must enter and spell the last line as shown (including the parentheses and semicolon), except for option ..., for which you must substitute one or more options separated from each other by one or more blanks.

Summary of changes

Enhancements in this release
This release provides the following functional enhancements that are described in this and the other IBM® PL/I books.

Performance improvements
- The code that is generated for the EXEC CICS statement executes faster because one MVC is eliminated.
- Much faster code is generated for MOD and REM of FIXED DEC with precision greater than 15.
- The ARCH option now accepts 11 as its maximum value, and when ARCH(11) is specified, the compiler generates code that exploits the new hardware instructions on those z systems. This code especially improves the performance of some instances of the SEARCH, and VERIFY built-in functions.

SQL enhancements
- The validation of an EXEC SQL statement will not stop when the first invalid host variable is found, but will instead check all host variable references.
- The new SQL preprocessor option (NO)CODEPAGE determines how the compiler CODEPAGE option is honored by the SQL preprocessor.
- The new SQL preprocessor option (NO)WARNDECP allows you to reduce the amount of “noise” that is produced by the SQL preprocessor.
- A structure that is used as a host variable can now also have a structure as its indicator variable.
- A named constant, that is, a PL/I variable with the VALUE attribute can now be used as a host variable if SQL allows a constant in that setting.

Usability enhancements
- The new MAXBRANCH compiler option can help you find excessively complex code.
- The new FILEREF and NOFILEREF compiler options control whether the compiler produces a file reference table.
- The new JSON compiler option lets you choose the case of the names in the JSON text generated by the JSONPUT built-in functions and expected by the JSONGET built-in functions.
• The LIMITS compiler option now accepts STRING as a suboption with these values as the threshold for the length of a BIT, CHARACTER or WIDECHAR variable: 32K, 512K, 8M, and 128M.
• The XML compiler option now lets you choose whether XML attributes generated by the XMLCHAR built-in function are enclosed in apostrophes or quotation marks.
• The RULES(NOLAXRETURN) compiler option is enhanced so that ERROR will be raised if code falls through to the END statement in a PROCEDURE with the RETURNS attribute.
• The new ALL | SOURCE suboptions to the RULES(NOLAXNESTED) and RULES(NOPADDING) compiler options provide finer control over when the compiler flags questionable coding.
• The new FORCE(NOLAXQUAL) attribute and the new FORCE suboption of the RULES(NOLAXQUAL) option enable users to enforce the NOLAXQUAL rules in a structure-by-structure manner.
• The use of INITIAL on REFER objects will now be flagged with an E-level message.
• The MACRO preprocessor now supports DEPRECATE and DEPRECATENEXT options that allow you to remove selected macro procedures over a course of planned stages.
• If a FIXED or FLOAT variable has the VALUE attribute, the compiler lists its value as a character string in the ATTRIBUTES listing.
• The aggregate listing is enhanced to mark arrays of VARYING that contain padding between array members (for example, when a structure contains an array of CHAR(31) VARYING ALIGNED).
• When SUBSCRIPTRANGE is enabled and an array assignment includes arrays with non-constant bounds, the compiler now generates code that raises the SUBSCRIPTRANGE condition if the bounds do not match. If the bounds are all constant, the compiler continues to check at compile-time that they match.
• If an event in any of the PLISAX event structures is set to null (via the NULLENTRY built-in function or the UNSPEC pseudovariable), the XML parser will not call that event. This allows you to limit your XML parsing code to only the events in which you are interested and to improve the performance of the overall parse at the same time.
• The compiler now flags an attempt to FETCH or RELEASE the procedure that it is currently compiling.

Enhancements from V4R4
This release provides the following functional enhancements that are described in this and the other IBM PL/I books.

Performance improvements
• The generation of the pseudo assembler listing is faster than the previous releases.
• The compiler is now compiled with ARCH(7).
• The compiler now flags declarations for variables that do not have the STATIC attribute but have more than 100 INITIAL items.

Usability enhancements
• The SQL preprocessor now issues more messages when programs contain incorrect syntax. These messages make it easier to identify the source statement that is in error.
• The new NOINCLUDE option can be used to prohibit the use of the %INCLUDE and %XINCLUDE statements outside the MACRO preprocessor.
• The included files are now bracketed in the source listing when the %INCLUDE statement is processed by the final compiler pass.
• The NULLSTRPTR suboption of the DEFAULT option has the new STRICT suboption that flags the assignments and comparisons of "" to POINTERS as invalid.
• When the STMT option is specified, the source and message listings will include both the logical statement numbers and the source file numbers.

Enhancements from V4R3

This release provides the following functional enhancements described in this and the other IBM PL/I books.

Performance improvements
• The ARCH(10) option provides exploitation of the IBM zEnterprise® EC12 System instructions.
• The compiler now generates inline code for VERIFY and SEARCH when they have three arguments and the second argument is a single character.
• The compiler now generates inline code for more conversions from PICTURE to DFP.
• The compiler now generates inline code for more conversions from BIT to CHAR.
• The compiler now generates inline code for conversions of BIT to WIDECHAR.
• The code generated for TRIM of FIXED DEC is improved.

Usability enhancements
• The SQL preprocessor has the following improvements:
  – It supports the ONEPASS option.
  – It supports the use of some restricted expressions in the host variable declarations.
  – The listing of the EXEC SQL statement is displayed in a readable format that is similar to the original source.
  – It supports the use of host variables declared with the LIKE attribute.
  – It supports the new DEPRECATE option that causes the preprocessor to flag a list of statements that you want to deprecate.
• The ADATA file now records the use of the following attributes, built-in functions, and statement:
  – The INONLY, INOUT, and OUTONLY attributes
  – The XMLATTR and XMLOMIT attributes
  – The ALLCOMPARE, UTF8, UTF8TOCHAR, and UTF8TOWCHAR built-in functions
  – The ASSERT statement
• With the new CASERULES option, you can specify case rules for PL/I keywords. For example, you can specify the rule that all keywords must be in uppercase.
• The DEPRECATE option has a new STMT suboption that causes the compiler to flag a list of statements that you want to deprecate.
• The new DEPRECATENEXT option allows for staged deprecation of functions.
• The IGNORE option has a new ASSERT suboption that instructs the compiler to ignore all ASSERT statements.

• The new (NO)MSGSUMMARY option controls whether the compiler adds a summary of all messages that are issued during the compilation into the listing.

• The RTCHECK option has the new NULL370 suboption that checks whether a pointer that equals to the old NULL() value is dereferenced. Pointers that equal to the old NULL() value are those with the hexadecimal value of ‘FF000000’x.

• The RULES option now accepts (NO)CONTROLLED as a suboption that controls whether to flag the use of the CONTROLLED attribute.

• The RULES option now accepts (NO)LAXNESTED as a suboption that controls whether to flag programs where nested procedures exist between sections of executable code.

• The RULES option now accepts (NO)RECURSIVE as a suboption that determines whether to flag any use of the RECURSIVE attribute or any procedure that directly calls itself.

• The RULES(NOUNREF) option now accepts SOURCE | ALL as a suboption that determines whether to flag all unreferenced variables.

• RULES(NOLAXIF) now also causes the compiler to flag the assignments of the form \( x = y = z \).

• RULES(NOLAXSCALE) now also causes the compiler to flag \( \text{ROUND}(x, p) \) when \( p < 0 \).

Enhancements from V4R2

This release provides the following functional enhancements described in this and the other IBM PL/I books.

Performance improvements

• The ARCH(9) option provides further exploitation of the IBM zEnterprise 196 (z196) System instructions, including the high-word, floating-point extension, and population count facilities.

• The new UNROLL compiler option gives you control of loop unrolling.

• The compiler now generates inline code to resolve the ULENGTH and USUBSTR built-in functions for character strings.

• The compiler now generates inline code for MEMINDEX(\( p, n, x \)) where \( x \) is WCHAR(1).

• The compiler now generates inline code for STG(\( x \)) where \( x \) is BASED variable using REFER and meets both of the following conditions:
  – All NONVARYING BIT in \( x \) are specified with the ALIGNED attribute.
  – All other elements in \( x \) are specified with UNALIGNED.

Debugging improvements

• The compiler now supports typed structures in Debug Tool.

SQL support enhancements

• The SQL Preprocessor has the following significant changes:
  – It fully and correctly supports block scoping.
  – The preprocessor load module is more than eight times smaller.
  – The preprocessor runs faster.
  – It supports the use of the SQL TYPE attribute anywhere you can specify a PL/I data type.
– It now supports the following compiler options when processing declarations of host variables. It applies the defaults correctly and rejects unsuitable host variables as appropriate.

- DEFAULT(ANS | IBM)
- DEFAULT(ASCII | EBCDIC)
- DEFAULT((NO)EVENDEC)
- DEFAULT((NON)NATIVE)
- DEFAULT(SHORT(HEX | IEEE))
- RULES((NO)LAXCTL)

– It now correctly processes the PRECISION attribute.
– It now recognizes the UNSIGNED and COMPLEX attributes and rejects the use of these attributes in any host variable.
– It now ensures that DSNHMLTR is declared in the outermost procedure that contains codes that need DSNHMLTR.
– It now handles packages correctly.
– It now ensures that characters in the source code are printable when it generates code to set the SQLAVDAID.
– It no longer requires that indicator arrays have a lower bound of 1.
– It now generates an SQL parameter list structure that has fewer unions, fewer init clauses, and no additional declarations based on elements of the structure.

Usability enhancements
- The new PPLIST compiler option conditionally erases the part of the listing that is generated by any preprocessor phase that produces no messages.
- The compiler issues better messages when a comma is missing in a structure declaration.
- The compiler issues new messages when the source code contains invalid shift-in and shift-out bytes.
- Under RULES(NOLAXENTRY), the compiler does not flag names that start with DSN.
- Under RULES(NOUNREF), the compiler does not flag names that start with DSN or SQL.
- Under the new NOSELFASSIGN suboption of the RULES compiler option, the compiler flags assignments of variables to themselves.
- Under the new NOLAXRETURN suboption of the RULES compiler option, the compiler generates codes to raise the ERROR condition when a RETURN statement is used in some invalid ways.

Enhancements from V4R1
This release provides the following functional enhancements described in this and the other IBM PL/I books.

Debugging improvements
- Under TEST(SEPARATE), the compiler optionally places the statement number table in the debug file and thus reduces the size of the generated object deck.
- Under TEST(SEPARATE), the compiler includes information identifying the source lines for declarations, references, and assignments.
- Under TEST(SEPARATE), the compiler generates information to identify the implicit locator reference when a BASED variable is based on the ADDR of an array element or other complex references.
Performance improvements
- The ARCH(9) option provides exploitation of the IBM zEnterprise System instructions.
- If all the elements of structures using REFER are byte-aligned, the compiler inlines code to resolve references to these elements rather than through library calls.

Usability enhancements
- The new DEPRECATE compiler option flags variable names and included file names that you want to deprecate.
- The new SEPARATE suboption of the GONUMBER option is provided to place the generated statement number table in the separate debug file.
- The new NOGLOBALDO suboption of the RULES option flags all DO loop control variables that are declared in the parent block.
- The new NOPADDING suboption of the RULES option flags all structures that contain padding.
- The SQL preprocessor now supports the XREF option.
- The new PLISAXD built-in subroutine provides the ability to parse XML documents with validation against a schema by using the XML System Services parser.

Serviceability enhancements
- The use of either the GOSTMT option or the IMPRECISE option is now flagged as unsupported.
- The compiler now always lists all SQL preprocessor options that are in effect.

Enhancements from V3R9
This release provides the following functional enhancements described in this and the other IBM PL/I books.

Performance improvements
- UVALID will now be inlined for for strings of length 256 or less.
- Under ARCH(7) and higher, the CU12, CU14, CU21, CU24, CU41, and CU42 instructions will be used to provide for fast conversions between UTF-8, UFT-16 and UTF-32.
- Under ARCH(7) and higher, the TRTT, TROT, TRTO, and TROO instructions will be used to provide for fast translations between one- and two-byte buffers.
- Assignments of like arrays of scalars will now be handled as storage copy operations.
- All assignments of BIT VARYING to BIT VARYING will now be inlined.
- All assignments of byte-aligned BIT NONVARYING to BIT VARYING will now be inlined.
- The ROUND and ROUNDDEC built-in functions will be inlined when the argument to be rounded is DFP.
- To simplify the choosing of options for best performance,
  - the COMPACT option has been dropped
  - The default setting for DEFAULT(REORDER | ORDER) has been changed to DEFAULT(REORDER)
  - the TUNE option has been dropped
- Detection of the dereferencing of null pointers exploits the new compare-and-trap instruction under ARCH(8).
The compiler has been built with ARCH(6) to improve its performance.

**Usability enhancements**

- The CICS preprocessor now supports block-scoping and consequently will add the needed local CICS declares to all non-nested procedures.
- The SQL preprocessor now supports the PL/I rules for the scope of declarations when resolving host variable references through the new SCOPE option. NOSCOPE is the default for compatibility with previous releases.
- The MACRO preprocessor will now leave %include, %xinclude, %inscan, and %xinscan statements in the compiler listing as comments.
- The MACRO preprocessor now provides via the %DO SKIP; statement an easy and clear way to omit sections of code from a compile.
- The MACRO preprocessor now supports an option called NAMEPREFIX which allows the user to force macro procedures and variables to start with a specified character.
- The IGNORE compiler option provides the ability to suppress PUT FILE and/or DISPLAY statements (either of which may have been used for debug purposes but which should be compiled out of the production version).
- The NULLSTRPTR suboption of the DEFAULT compiler option provides user control of whether sysnull or null is assigned to a pointer when the source in the assignment is a null string.
- The new MAXGEN option specifies the maximum number of intermediate language statements that should be generated for any one user statement and will cause the compiler to flag any statement where this maximum is exceeded.
- The new ONSNAP option will allow the user to request the compiler to insert an ON STRINGRANGE SNAP; or an ON STRINGSIZE SNAP; statement into the prologue of a MAIN or FROMALIEN proc.
- The new SHORT suboption of the INITAUTO option will limit the INITAUTO option so that it does not duplicate all of the runtime STORAGE option, but does initialize variables that might be optimized to registers.
- The new RTCHECK option will generate code to test for the dereferencing of null pointers.
- The compiler will now flag various statements that may be risky:
  - code where the result of a FIXED operation has a scale factor less than zero
  - ENTRYs used as functions but declared without the RETURNS attribute
  - parameters declared as BYVALUE when doing so is ill-advised, e.g. declaring a FIXED DEC parameter BYVALUE
  - FIXED DECIMAL add and multiply operations that might raise FIXEDOVERFLOW
- The RULES option has been expanded to allow more control over and flagging of poor code:
  - NOPROCEDONLY will flag END statements for PROCs that don’t name the PROC they are closing
  - NOSTOP will flag the use of STOP and EXIT
  - NOLAXQUAL(STRICT) will flag variables not qualified with their level-1 name
  - NOLAXSCALE will flag declares of FIXED DEC(p,q) and FIXED BIN(p,q) where q < 0 or p < q
  - NOGOTO(LOOSE) will allow GOTOs only if in the same block
– More than one DELAY STATEMENTS can be concurrently executed in different procedures.

Serviceability enhancements
• When the compiler cannot open a file, the compiler will now, if possible, also include the related C runtime message in the message in the listing.
• If user code requires a DFP conversion at compile time but the compile is running on a machine without DFP hardware, this error will be trapped and a meaningful error issued.
• If the SQL preprocessor is invoked more than once without INCONONLY as its suboption, then the DBRM library created by the compiler will be empty, and now an E-level message will be issued to warn the user about this problem.

Enhancements from V3R8
This release provides the following functional enhancements described in this and the other IBM PL/I books.

Performance improvements
• The ARCH(8) and TUNE(8) options provide exploitation of the z/HE instructions.
• The HGPR option supports using 64-bit registers in 32-bit code.
• The GOFF option supports generation of GOFF objects.
• The PFPO instruction will be exploited in conversions between differing float formats.
• SRSTU will be used in code generated for UTF-16 INDEX.
• Calls to null internal procedures will now be completely removed.

Usability enhancements
• PLISAXC provides for access to the XML System Services parser via a SAX interface.
• The INCDIR option is supported under batch.
• The LISTVIEW option now provides the support previously provided by the AFTERMACRO etc suboptions of TEST.
• The NOLAXENTRY suboption of the RULES option allows for the flagging of unprototyped ENTRYs.
• The (NO)FOFLOMNULT suboption of the DECIMAL option allows for control of whether FOFL is raised in MULTIPLY of FIXED DECIMAL.
• The HEX and SUBSTR suboptions of the USAGE option provide more user control of the behavior of the corresponding built-in functions.
• The DDSQL compiler option provides the ability to specify an alternate DD name to be used for EXEC SQL INCLUDEs.
• The INCONONLY suboption provided by the MACRO and SQL preprocessors can request those preprocessors to perform only INCLUDEs.
• The integrated SQL preprocessor will now generate DB2 precompiler style declares for all *LOB_FILE, *LOCATOR, ROWID, BINARY and VARBINARY SQL types, in addition to the BLOB, CLOB and DBCLOB SQL types already supported, when the LOB(DB2) SQL preprocessor option is selected.

Enhancements from V3R7
This release provides the following functional enhancements described in this and the other IBM PL/I books.
Debugging improvements

- The TEST option has been enhanced so that users can choose to view the source in the listing and in the Debug Tool source window as that source would appear after a user-specified preprocessor had been run (or after all the preprocessors had been run).

Performance improvements

- The BASR instruction will now be used instead of the BALR instruction.
- The conversions of FIXED DEC with large precision to FLOAT will be inlined and speeded up by the use of FIXED BIN(63) as an intermediary.
- The CHAR built-in when applied to CHAR expressions will now always be inlined.
- The code generated for conversions of FIXED BIN(p,q) to unscaled FIXED DEC has been significantly improved.
- TRTR will be used, under ARCH(7), for SEARCHR and VERIFYR in the same situations where TRT would be used for SEARCH and VERIFY.
- UNPKU will be used to convert some PICTURE to WIDECHAR (rather than making a library call).

Usability enhancements

- IEEE Decimal Floating-Point (DFP) is supported.
- The new MEMCONVERT built-in function will allow the user to convert arbitrary lengths of data between arbitrary code pages.
- The new ONOFFSET built-in function will allow the user to have easy access to another piece of information formerly available only in the runtime error message or dump, namely the offset in the user procedure at which a condition was raised.
- The new STACKADDR built-in function will return the address of the current dynamic save area (register 13 on z/OS) and will make it easier for users to write their own diagnostic code.
- The length of the mnemonic field in the assembler listing will be increased to allow for better support of the new z/OS instructions that have long mnemonics.
- More of the right margin will be used in the attributes, cross-reference and message listings.
- The CODEPAGE option will now accept 1026 (the Turkish code page) and 1155 (the 1026 code page plus the Euro symbol).
- The new MAXNEST option allows the user to flag excessive nesting of BEGIN, DO, IF and PROC statements.
- Under the new (and non-default) suboption NOELSEIF of the RULES option, the compiler will flag any ELSE statement that is immediately followed by an IF statement and suggest that it be rewritten as a SELECT statement.
- Under the new (and non-default) suboption NOLAXSTG of the RULES option, the compiler will flag declares where a variable A is declared as BASED on ADDR(B) and STG(A) > STG(B) not only (as the compiler did before) when B is AUTOMATIC, BASED or STATIC with constant extents but now also when B is a parameter declared with constant extents.
- The new QUOTE option will allow the user to specify alternate code points for the quote (") symbol since this symbol is not code-page invariant.
- The new XML compiler option can be used to specify that the tags in the output of the XMLCHAR built-in function be either in all upper case or in the case in which they were declared.
• For compilations that produce no messages, the compiler will now include a line saying "no compiler messages" where the compiler messages would have been listed.

• The MACRO preprocessor will support a new suboption that will allow the user to specify whether it should process only `%INCLUDE statements or whether it should process all macro statements.

• The integrated SQL preprocessor will now generate DB2 precompiler style declares for all *LOB_FILE, *LOCATOR, ROWID, BINARY and VARBINARY SQL types, in addition to the BLOB, CLOB and DBCLOB SQL types already supported, when the LOB(DB2) SQL preprocessor option is selected.

Enhancements from V3R6

This release provides the following functional enhancements described in this and the other IBM PL/I books.

DB2 V9 support

• Support for STDSQL(YES/NO)
• Support for CREATE TRIGGER (aka multiple SQL statements)
• Support for FETCH CONTINUE
• Support for SQL style comments ('--') embedded in SQL statements
• Support for additional SQL TYPES including
  – SQL TYPE IS BLOB_FILE
  – SQL TYPE IS CLOB_FILE
  – SQL TYPE IS DBCLOB_FILE
  – SQL TYPE IS XML AS
  – SQL TYPE IS BIGINT
  – SQL TYPE IS BINARY
  – SQL TYPE IS VARBINARY
• The SQL preprocessor will also now list the DB2 coprocessor options

Debugging improvements

• Under TEST(NOSEPNAME), the name of the debug side file will not be saved in the object deck.

Performance improvements

• Support, under ARCH(7), of the z/OS extended-immediate facility
• Exploitation of the CLCLU, MVCLU, PKA, TP and UNPKA instructions under ARCH(6)
• Exploitation of the CVBG and CVDG instructions under ARCH(5)
• Expanded use of CLCLE
• Conversions involving DB2 date-time patterns are now inlined
• The ALLOCATION built-in function is now inlined
• Conversions with a drifting $ have been inlined
• Conditional code has been eliminated from assignments to PIC'(n)Z'
• Conversions from FIXED BIN to a PICTURE that specifies a scale factor have been inlined
• Assignments to BIT variables that have inherited dimensions but which have strides that are divisible by 8 will now be inlined
Usability enhancements

- The MAP output will now also include a list in order of storage offset (per block) of the AUTOMATIC storage used by the block
- Conformance checking has been extended to include structures
- Listings will now include 7 columns for the line number in a file
- The THREADID built-in function is now supported under z/OS
- The PICSPEC built-in function is now supported
- The new CEESTART option allows you to position the CEESTART csect at the start or end of the object deck
- The new PPCICS, PPMACRO and PPSQL options allow you to specify the default options to be used by the corresponding preprocessor
- The ENVIRONMENT option is now included in the ATTRIBUTES listing
- The DISPLAY option now supports a suboption to allow different DESC codes for DISPLAY with REPLY versus DISPLAY without REPLY
- The message flagging a semicolon in a comment will now include the line number of the line where the semicolon appears
- Flag assignments that can change a REFER object
- Flag use of KEYED DIRECT files without a KEY/KEYFROM clause
- Flag use of PICTURE as loop control variables

Enhancements from V3R5

This release provides the following functional enhancements described in this and the other IBM PL/I books.

Debugging improvements

- Under TEST(SEPARATE), the vast majority of debug information will be written to a separate debug file
- AUTOMONITOR will include the target in assignments
- The AT ENTRY hook will now be placed after AUTOMATIC has been initialized, thereby eliminating the need to step into the block before looking at any variables

Performance improvements

- Generation of branch-relative instructions so that the need for base registers and transfer vectors will be significantly eliminated
- Support, under ARCH(6), of the z/OS long displacement facility
- Simple structures using REFER will be mapped inlined rather than via a library call
- For structures using REFER still mapped via a library call, less code will be generated if the REFER specifies the bound for an array of substructures
- Faster processing of duplicate INCLUDEs
- Conversions to PICTURE variables with an I or R in the last position will now be inlined (such conversions had already been inlined when the last character was a T)
- Conversions to PICTURE variables ending with one or more B’s will now be inlined if the corresponding picture without the B’s would have been inlined
- Conversions to from CHARACTER to PICTURE variables consisting only of X’s will now be inlined
Usability enhancements

- All parts of the listing, etc will count the source file as file 0, the first include file as file 1, the second (unique) include file as file 2, etc
- Conformance checking extended to include arrays
- Listings will include the build dates for any preprocessors invoked
- One-byte FIXED BINARY arguments can be suppressed for easier ILC with COBOL
- Alternate DD names may be specified for SYSADATA, SYSXMLSD and SYSDEBUG
- RULES(NOLAXMARGINS) tolerates, under XNUMERIC, sequence numbers
- RULES(NOUNREF) flags unreferenced AUTOMATIC variables
- If an assignment to a variable is done via library call, the message flagging the library call will include the name of the target variable
- Flag one-time DO loops
- Flag labels used as arguments
- Flag ALLOCATE and FREE of non-PARAMETER CONTROLLED in FETCHABLE if PRV used
- Flag DEFINED and BASED larger than their base even if the base is declared later
- Flag implicit FIXED DEC to 8-byte integer conversions

Enhancements from V3R4

This release provides the following functional enhancements described in this and the other IBM PL/I books.

Migration enhancements

- Support sharing of CONTROLLED with old code
- Improve default initialization
- Ease decimal precision specification in ADD, DIVIDE and MULTIPLY
- Support old semantics for STRING of GRAPHIC
- Support old semantics for the DEFAULT statement
- Flag declares with storage overlay
- Lift restrictions on RETURN inside BEGIN
- Optionally flag semicolons in comments
- Support EXT STATIC initialized in assembler
- Flag invalid carriage control characters
- Flag more language misuse, especially with RETURN
- Support the REPLACEBY2 built-in function
- Optionally suppress FOFL on decimal assignments that would raise SIZE
- Flag more language handled differently than the old compiler

Performance improvements

- Improve code generated for INDEX and TRANSLATE
- Inline more assignments to pictures
- Improve the code generated for conversions of CHARACTER to PICTURE when the conversion would be done inline if the source were FIXED DEC
- Improve the code generated for packed decimal conversions
- Improve the code generated for some uses of REFER
• Inline compares of character strings of unknown length
• Reduce the amount of stack storage used for concatenates
• Inline more GET/PUT STRING EDIT statements
• Short-circuit more LE condition handling
• Inline more Or and And of BIN FIXED
• Inline SIGNED FIXED BIN(8) to ALIGNED BIT(8)
• Flag statements where the compiler generates a call to a library routine to map a structure at run time
• Lessen the amount of I/O used to produce the listing

Usability enhancements
• Optionally provide offsets in the AGGREGATE listing in hex
• Support DEC(31) only when needed via the LIMITS(FIXEDDEC(15,31)) option
• Allow comments in options
• Optionally flag FIXED DEC declares with even precision
• Optionally flag DEC to DEC assignments that could raise SIZE
• Flag DEC/PIC to PIC assignments that could raise SIZE
• Support LIKE without INIT via the NOINIT attribute
• Ease includes from PDS’s under z/OS UNIX
• Support the LOWERCASE, MACNAME, TRIM and LOWERCASE built-in functions in the MACRO preprocessor
• Ease introduction of options via PTF
• Optionally disallow use of *PROCESS
• Optionally keep *PROCESS in MDECK
• Support one-time INCLUDE
• Support macro-determined INCLUDE name
• Support runtime string parameter checking
• Flag more possibly uninitialized variables
• Flag unusual compares that are likely to be coding errors
• The output of the STORAGE option is now formatted more nicely, and the output of the LIST option will now include the hex offset for each block from the start of the compilation unit.

Debugging improvements
• Better support for overlay hooks
• Easier resolution of CONTROLLED variables in LE dumps
• Always include user specified options in the listing

Enhancements from V3R3
This release also provides all of the functional enhancements offered in Enterprise PL/I V3R3, including the following:

More XML support
The XMLCHAR built-in function will write XML with the names and values of the elements of a referenced structure to a buffer and return the number of bytes written. This XML can then be passed to other applications, including code using the PL/I SAX parser, which want to consume it.
Improved performance

• The compilation time under OPT(2) will be significantly less than under Enterprise PL/I V3R2, especially for large programs.
• The compiler now uses the ED and EDMK instructions for inlined numeric conversions to PICTURE and CHARACTER. This results in faster, shorter code sequences and also in faster compilations.
• The compiler now generates better code for string comparisons. This also results in faster, shorter code sequences.
• The compiler now generates shorter, faster code for conversion from FIXED DECIMAL to PICTURE with trailing overpunch characters.
• The ARCH and TUNE compiler options now accept 5 as a valid sub-option. Under ARCH(5), the compiler will generate, when appropriate, some new z/Architecture instructions such as NILL, NILH, OILL, OILH, LLILL, and LLILH.

Easier migration

• The new BIFPREC compiler option controls the precision of the FIXED BIN result returned by various built-in functions and thus provides for better compatibility with the OS PL/I compiler.
• The new BACKREG compiler option controls which register the compiler uses as the backchain register and thus allows for easier mixing of old and new object code.
• The new RESEXP compiler option controls the evaluation of restricted expressions in code, and thus provides for better compatibility with the OS PL/I compiler.
• The new BLKOFF compiler option provides for controlling the way offsets in the compiler's pseudo-assembler listing are calculated.
• The STORAGE compiler option causes the compiler to produce, as part of the listing, a summary, similar to that produced by the OS PL/I compiler, of the storage used by each procedure and begin-block.

Improved usability

• The new LAXDEF suboption of the RULES compiler option allows the use of so-called illegal defining without having the compiler generate E-level messages.
• The new FLOATINMATH compiler option offers easier control of the precision with which math functions are evaluated.
• The new MEMINDEX, MEMSEARCH(R) and MEMVERIFY(R) built-in functions provide the ability to search strings larger than 32K.
• The new ROUTCDE and DESC suboptions of the DISPLAY(WTO) compiler option offers control of the corresponding elements of the WTO.
• The compiler now stores in each object a short string that will be in storage even when the associated code runs and that records all the options used to produce that object. This allows various tools to produce better diagnostics.
• The compiler now issues messages identifying more of the places where statements have been merged or deleted.
• The PLIDUMP output now includes a hex dump of user static.
• The PLIDUMP output now includes the options used to compile each program in the Language Environment traceback.
• The PLIDUMP output now includes more information on PL/I files.
**Improved debug support**

- BASED structures using REFER are now supported in DebugTool and in data-directed I/O statements (with the same restrictions as on all other BASED variables).
- BASED structures that are BASED on scalar members of other structures (which, in turn, may be BASED, etc) are now supported in DebugTool and in data-directed I/O statements (with the same restriction as on all other BASED variables).

**Enhancements from V3R2**

This release also provides all of the functional enhancements offered in Enterprise PL/I V3R2, including the following:

**Improved performance**

- The compiler now handles even more conversions by generating inline code which means these conversions will be done much faster than previously. Also, all conversions done by library call are now flagged by the compiler.
- The compiler-generated code now uses, in various situations, less stack storage.
- The compiler now generates much better code for references to the TRANSLATE built-in function.
- The compiler-generated code for SUBSCRIPTRANGE checking is now, for arrays with known bounds, twice as fast as before.
- The ARCH and TUNE options now support 4 as a suboption, thereby allowing exploitation of instructions new to the zSeries machines.
- ARCH(2), FLOAT(AFP) and TUNE(3) are now the default.

**Easier migration**

- Compiler defaults have been changed for easier migration and compatibility. The changed defaults are:
  - CSECT
  - CMPAT(V2)
  - LIMITS(EXTNAME(7))
  - NORENT
- The compiler now honors the NOMAP, NOMAPIN and NOMAP attributes for PROCs and ENTRYs with OPTIONS(COBOL).
- The compiler now supports PROCs with ENTRY statements that have differing RETURNS attribute in the same manner as did the old host compiler.
- The compiler will now assume OPTIONS(RETCODE) for PROCs and ENTRYs with OPTIONS(COBOL).
- The SIZE condition is no longer promoted to ERROR if unhandled.
- Various changes have been made to reduce compile time and storage requirements.
- The OFFSET option will now produce a statement offset table much like the ones it produced under the older PL/I compilers.
- The FLAG option now has exactly the same meaning as it had under the old compilers, while the new MAXMSG option lets you decide if the compiler should terminate after a specified number of messages of a given severity. For example, with FLAG(I) MAXMSG(E,10), you can now ask to see all I-level messages while terminating the compilation after 10 E-level messages.
- The AGGREGATE listing now includes structures with adjustable extents.
• The STMT option is now supported for some sections of the listing.
• The maximum value allowed for LINESIZE has been changed to 32759 for F-format files and to 32751 for V-format files.

**Improved usability**
• The defaults for compiler options may now be changed at installation.
• The integrated SQL preprocessor now supports DB2 Unicode.
• The compiler now generates information that allows Debug Tool to support Auto Monitor, whereby immediately before each statement is executed, all the values of all the variables used in the statement are displayed.
• The new NOWRITABLE compiler option lets you specify that even under NORENT and at the expense of optimal performance, the compiler should use no writable static when generating code to handle FILEs and CONTROLLED.
• The new USAGE compiler option gives you full control over the IBM or ANS behavior of the ROUND and UNSPEC built-in function without the other effects of the RULES(IBM | ANS) option.
• The new STDSYS compiler option lets you specify that the compiler should cause the SYSPRINT file to be equated to the C stdout file.
• The new COMPACT compiler option lets you direct the compiler to favor those optimizations which tend to limit the growth of the code.
• The LRECL for SYSPRINT has been changed to 137 to match that of the C/C++ compiler.
• POINTERs are now allowed in PUT LIST and PUT EDIT statements: the 8-byte hex value will be output.
• If specified on a STATIC variable, the ABNORMAL attribute will cause that variable to be retained even if unused.

**Enhancements from V3R1**
This release also provides all of the functional enhancements offered in Enterprise PL/I V3R1, including the following:
• Support for multithreading on z/OS
• Support for IEEE floating-point on z/OS
• Support for the ANSWER statement in the macro preprocessor
• SAX-style XML parsing via the PLISAXA and PLISAXB built-in subroutines
• Additional built-in functions:
  – CS
  – CDS
  – ISMAIN
  – LOWERCASE
  – UPPERCASE

**Enhancements from VisualAge PL/I**
This release also provides all of the functional enhancements offered in VisualAge PL/I V2R2, including the following:
• Initial UTF-16 support via the WIDECHAR attribute
  There is currently no support yet for
  – WIDECHAR characters in source files
  – W string constants
  – use of WIDECHAR expressions in stream I/O
  – implicit conversion to/from WIDECHAR in record I/O
– implicit endianness flags in record I/O

If you create a WIDECHAR file, you should write the endianness flag ('fe_ff'wx) as the first two bytes of the file.

• DESCRIPTORS and VALUE options supported in DEFAULT statements

• PUT DATA enhancements
  – POINTER, OFFSET and other non-computational variables supported
  – Type-3 DO specifications allowed
  – Subscripts allowed

• DEFINE statement enhancements
  – Unspecified structure definitions
  – CAST and RESPEC type functions

• Additional built-in functions:
  – CHARVAL
  – ISIGNED
  – IUNSIGNED
  – ONWCHAR
  – ONWSOURCE
  – WCHAR
  – WCHARVAL
  – WHIGH
  – WIDECHAR
  – WLOW

• Preprocessor enhancements
  – Support for arrays in preprocessor procedures
  – WHILE, UNTIL and LOOP keywords supported in %DO statements
  – %ITERATE statement supported
  – %LEAVE statement supported
  – %REPLACE statement supported
  – %SELECT statement supported
  – Additional built-in functions:
    – COLLATE
    – COMMENT
    – COMPILEDATE
    – COMPILETIME
    – COPY
    – COUNTER
    – DIMENSION
    – HBOUND
    – INDEX
    – LBOUND
    – LENGTH
    – MACCOL
    – MACLMAR
    – MACRMAR
    – MAX
    – MIN
    – PARMSET
    – QUOTE
    – REPEAT
    – SUBSTR
    – SYSPARM
    – SYSTEM
    – SYSVERSION
How to send your comments

Your feedback is important in helping us to provide accurate, high-quality information. If you have comments about this document or any other PL/I documentation, contact us in one of these ways:

- Use the Online Readers’ Comment Form at www.ibm.com/software/awdtools/rcf/

or send an e-mail to comments@us.ibm.com

Be sure to include the name of the document, the publication number of the document, the version of PL/I, and, if applicable, the specific location (for example, page number) of the text that you are commenting on.

- Fill out the Readers’ Comment Form at the back of this document, and return it by mail or give it to an IBM representative. If the form has been removed, address your comments to:

  International Business Machines Corporation
  Reader Comments
  H150/090
  555 Bailey Avenue
  San Jose, CA 95141-1003
  USA

- Fax your comments to this U.S. number: (800)426-7773.

When you send information to IBM, you grant IBM a nonexclusive right to use or distribute the information in any way it believes appropriate without incurring any obligation to you.

Accessibility

Accessibility features help users who have a disability, such as restricted mobility or limited vision, to use information technology products successfully. The accessibility features in z/OS provide accessibility for Enterprise PL/I.

The major accessibility features in z/OS are:

- Interfaces that are commonly used by screen readers and screen-magnifier software
- Keyboard-only navigation
- Ability to customize display attributes such as color, contrast, and font size

Interface information

Assistive technology products work with the user interfaces that are found in z/OS. For specific guidance information, see the documentation for the assistive technology product that you use to access z/OS interfaces.

Keyboard navigation

Users can access z/OS user interfaces by using TSO/E or ISPF. For information about accessing TSO/E or ISPF interfaces, see the following publications:
Accessibility

- z/OS TSO/E Primer
- z/OS TSO/E Primer
- z/OS TSO/E Primer

These guides describe how to use TSO/E and ISPF, including the use of keyboard shortcuts or function keys (PF keys). Each guide includes the default settings for the PF keys and explains how to modify their functions.

Accessibility of this information

The English-language XHTML format of this information that will be provided in the IBM System z® Enterprise Development Tools & Compilers Information Center at publib.boulder.ibm.com/infocenter/pdthelp/index.jsp is accessible to visually impaired individuals who use a screen reader.

To enable your screen reader to accurately read syntax diagrams, source code examples, and text that contains the period or comma PICTURE symbols, you must set the screen reader to speak all punctuation.

IBM and accessibility

See the IBM Human Ability and Accessibility Center at www.ibm.com/able for more information about the commitment that IBM has to accessibility.
# Part 1. Compiling your program

## Chapter 1. Using compiler options and facilities

<table>
<thead>
<tr>
<th>Option</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>AGGREGATE</td>
<td>5</td>
</tr>
<tr>
<td>ARCH</td>
<td>9</td>
</tr>
<tr>
<td>ATTRIBUTES</td>
<td>11</td>
</tr>
<tr>
<td>BACKREG</td>
<td>11</td>
</tr>
<tr>
<td>BIFPREC</td>
<td>11</td>
</tr>
<tr>
<td>BLANK</td>
<td>12</td>
</tr>
<tr>
<td>BLKOFF</td>
<td>13</td>
</tr>
<tr>
<td>CASERULES</td>
<td>13</td>
</tr>
<tr>
<td>CEESTART</td>
<td>13</td>
</tr>
<tr>
<td>CMPAT</td>
<td>14</td>
</tr>
<tr>
<td>CODEPAGE</td>
<td>16</td>
</tr>
<tr>
<td>COMMON</td>
<td>16</td>
</tr>
<tr>
<td>COMPILE</td>
<td>17</td>
</tr>
<tr>
<td>COPYRIGHT</td>
<td>17</td>
</tr>
<tr>
<td>CSECT</td>
<td>18</td>
</tr>
<tr>
<td>CSECTCUT</td>
<td>18</td>
</tr>
<tr>
<td>CURRENCY</td>
<td>19</td>
</tr>
<tr>
<td>DBCS</td>
<td>19</td>
</tr>
<tr>
<td>DD</td>
<td>19</td>
</tr>
<tr>
<td>DDCS</td>
<td>20</td>
</tr>
<tr>
<td>DECIMAL</td>
<td>20</td>
</tr>
<tr>
<td>DEFAULT</td>
<td>22</td>
</tr>
<tr>
<td>DEPRECATE</td>
<td>31</td>
</tr>
<tr>
<td>DEPRECATENEXT</td>
<td>32</td>
</tr>
<tr>
<td>DISPLAY</td>
<td>33</td>
</tr>
<tr>
<td>DLLINIT</td>
<td>33</td>
</tr>
<tr>
<td>EXIT</td>
<td>33</td>
</tr>
<tr>
<td>EXTRN</td>
<td>34</td>
</tr>
<tr>
<td>FILEREF</td>
<td>34</td>
</tr>
<tr>
<td>FLAG</td>
<td>34</td>
</tr>
<tr>
<td>FLOAT</td>
<td>35</td>
</tr>
<tr>
<td>FLOATINMATH</td>
<td>37</td>
</tr>
<tr>
<td>GOF</td>
<td>37</td>
</tr>
<tr>
<td>GONUMBER</td>
<td>38</td>
</tr>
<tr>
<td>GRAPHIC</td>
<td>39</td>
</tr>
<tr>
<td>HGPR</td>
<td>39</td>
</tr>
<tr>
<td>IGNORE</td>
<td>40</td>
</tr>
<tr>
<td>INAFTER</td>
<td>40</td>
</tr>
<tr>
<td>INCDIR</td>
<td>40</td>
</tr>
<tr>
<td>INCLUDE</td>
<td>41</td>
</tr>
<tr>
<td>INCDDS</td>
<td>41</td>
</tr>
<tr>
<td>INITAUTO</td>
<td>42</td>
</tr>
<tr>
<td>INITBASED</td>
<td>43</td>
</tr>
<tr>
<td>INITCTL</td>
<td>43</td>
</tr>
<tr>
<td>INITSTATIC</td>
<td>43</td>
</tr>
<tr>
<td>INSOURCE</td>
<td>43</td>
</tr>
<tr>
<td>INTERRUPT</td>
<td>44</td>
</tr>
<tr>
<td>JSON</td>
<td>45</td>
</tr>
<tr>
<td>LANGLEVEL</td>
<td>45</td>
</tr>
<tr>
<td>LIMITS</td>
<td>46</td>
</tr>
<tr>
<td>LINECOUNT</td>
<td>47</td>
</tr>
<tr>
<td>LINEDIR</td>
<td>47</td>
</tr>
<tr>
<td>LIST</td>
<td>47</td>
</tr>
<tr>
<td>LISTVIEW</td>
<td>48</td>
</tr>
<tr>
<td>MACRO</td>
<td>49</td>
</tr>
<tr>
<td>MAP</td>
<td>49</td>
</tr>
<tr>
<td>MARGINI</td>
<td>49</td>
</tr>
<tr>
<td>MARGINS</td>
<td>50</td>
</tr>
<tr>
<td>MAXBRANCH</td>
<td>51</td>
</tr>
<tr>
<td>MAXGEN</td>
<td>51</td>
</tr>
<tr>
<td>MAXMEM</td>
<td>52</td>
</tr>
<tr>
<td>MAXMSG</td>
<td>52</td>
</tr>
<tr>
<td>MAXNEST</td>
<td>53</td>
</tr>
<tr>
<td>MAXSTMNT</td>
<td>53</td>
</tr>
<tr>
<td>MAXTEMP</td>
<td>53</td>
</tr>
<tr>
<td>MDECK</td>
<td>54</td>
</tr>
<tr>
<td>MSGSUMMARY</td>
<td>54</td>
</tr>
<tr>
<td>NAME</td>
<td>55</td>
</tr>
<tr>
<td>NAMES</td>
<td>55</td>
</tr>
<tr>
<td>NATLANG</td>
<td>55</td>
</tr>
<tr>
<td>NEST</td>
<td>56</td>
</tr>
<tr>
<td>NOT</td>
<td>56</td>
</tr>
<tr>
<td>NUMBER</td>
<td>56</td>
</tr>
<tr>
<td>OBJECT</td>
<td>57</td>
</tr>
<tr>
<td>OFFSET</td>
<td>57</td>
</tr>
<tr>
<td>ONSNAP</td>
<td>57</td>
</tr>
<tr>
<td>OPTIMIZE</td>
<td>58</td>
</tr>
<tr>
<td>OPTIONS</td>
<td>59</td>
</tr>
<tr>
<td>OR</td>
<td>59</td>
</tr>
<tr>
<td>PP</td>
<td>60</td>
</tr>
<tr>
<td>PPCCIC</td>
<td>61</td>
</tr>
<tr>
<td>PPINCLUDE</td>
<td>61</td>
</tr>
<tr>
<td>PPLIST</td>
<td>62</td>
</tr>
<tr>
<td>PPMACRO</td>
<td>62</td>
</tr>
<tr>
<td>PPSQL</td>
<td>63</td>
</tr>
<tr>
<td>PPTTRACE</td>
<td>63</td>
</tr>
<tr>
<td>PRECTYPE</td>
<td>63</td>
</tr>
<tr>
<td>PREFIX</td>
<td>64</td>
</tr>
<tr>
<td>PROCEED</td>
<td>64</td>
</tr>
<tr>
<td>PROCESS</td>
<td>65</td>
</tr>
<tr>
<td>QUOTE</td>
<td>65</td>
</tr>
<tr>
<td>REDUCE</td>
<td>66</td>
</tr>
<tr>
<td>RENT</td>
<td>66</td>
</tr>
<tr>
<td>RESEX</td>
<td>68</td>
</tr>
<tr>
<td>RESPECT</td>
<td>68</td>
</tr>
<tr>
<td>RTCHECK</td>
<td>68</td>
</tr>
<tr>
<td>RULES</td>
<td>69</td>
</tr>
<tr>
<td>SEMANTIC</td>
<td>77</td>
</tr>
<tr>
<td>SERVICE</td>
<td>78</td>
</tr>
<tr>
<td>SOURCE</td>
<td>78</td>
</tr>
<tr>
<td>SPILL</td>
<td>78</td>
</tr>
<tr>
<td>STATIC</td>
<td>78</td>
</tr>
<tr>
<td>STDENV</td>
<td>79</td>
</tr>
<tr>
<td>STMT</td>
<td>79</td>
</tr>
<tr>
<td>STORAGE</td>
<td>79</td>
</tr>
<tr>
<td>STRINGOFGRAPHIC</td>
<td>80</td>
</tr>
<tr>
<td>SYNTAX</td>
<td>80</td>
</tr>
<tr>
<td>SYS瘠</td>
<td>81</td>
</tr>
<tr>
<td>SYSTEM</td>
<td>81</td>
</tr>
</tbody>
</table>

© Copyright IBM Corp. 1999, 2015
Punctuating GET LIST and GET DATA statements ........................................ 166
Automatic padding for GET EDIT ......................................................... 167
ENDFILE ......................................................................................... 167
SYSPRINT considerations ................................................................. 167
Using MSGFILE(SYS PRINT) ............................................................. 169
Using FETCH in your routines ............................................................ 169
Fetching Enterprise PL/I routines ....................................................... 169
Fetching PL/I MAIN routines ............................................................ 178
Fetching z/OS C routines ................................................................. 179
Fetching assembler routines ............................................................. 179
Invoking MAIN under TSO/E .............................................................. 179
Invoking MAIN under z/OS UNIX ....................................................... 180
Chapter 1. Using compiler options and facilities

This chapter describes the options that you can use for the compiler, along with their abbreviations and IBM-supplied defaults. It’s important to remember that PL/I requires access to Language Environment run time when you compile your applications. You can override most defaults when you compile your PL/I program. You can also override the defaults when you install the compiler.

Compile-time option descriptions

There are three types of compiler options; however, most compiler options have a positive and negative form. The negative form is the positive with 'NO' added at the beginning (as in TEST and NOTEST). Some options have only a positive form (as in SYSTEM). The three types of compiler options are:

1. Simple pairs of keywords: a positive form that requests a facility, and an alternative negative form that inhibits that facility (for example, NEST and NONEST).
2. Keywords that allow you to provide a value list that qualifies the option (for example, FLAG(W)).
3. A combination of 1 and 2 above (for example, NOCOMPILE(E)).

Table 3 lists all the compiler options with their abbreviations (if any) and their IBM-supplied default values. If an option has any suboptions which may be abbreviated, those abbreviations are described in the full description of the option.

For the sake of brevity, some of the options are described loosely in the table (for example, only one suboption of LANGLVL is mandatory, and similarly, if you specify one suboption of TEST, you do not have to specify the other). The full and completely accurate syntax is described in the pages that follow.

The paragraphs following Table 3 describe the options in alphabetical order. For those options specifying that the compiler is to list information, only a brief description is included; the generated listing is described under “Using the compiler listing” on page 96.

Table 3. Compile-time options, abbreviations, and IBM-supplied defaults

<table>
<thead>
<tr>
<th>Compile-Time Option</th>
<th>Abbreviated Name</th>
<th>z/OS Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>AGGREGATE[(DEC</td>
<td>HEX)]</td>
<td>NOAGGREGATE</td>
</tr>
<tr>
<td>ARCH(n) - ARCH(7)</td>
<td>-</td>
<td>ARCH(7)</td>
</tr>
<tr>
<td>ATTRIBUTES[(FULL</td>
<td>SHORT)]</td>
<td>NOATTRIBUTES</td>
</tr>
<tr>
<td>BACKREG(5</td>
<td>11) - BACKREG(5)</td>
<td>-</td>
</tr>
<tr>
<td>BIFPREC(15</td>
<td>31) - BIFPREC(15)</td>
<td>-</td>
</tr>
<tr>
<td>BLANK('c') - BLANK('t')</td>
<td>-</td>
<td>BLANK('t')</td>
</tr>
<tr>
<td>BLKOFF</td>
<td>NOBLKOFF</td>
<td>-</td>
</tr>
<tr>
<td>CASERULES(KEYWORD(LOWER</td>
<td>MIXED</td>
<td>START</td>
</tr>
<tr>
<td>CEESTART(FIRST</td>
<td>LAST) - CEESTART(FIRST)</td>
<td>-</td>
</tr>
</tbody>
</table>

© Copyright IBM Corp. 1999, 2015
<table>
<thead>
<tr>
<th>Compile-Time Option</th>
<th>Abbreviated Name</th>
<th>z/OS Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>CHECK(STORAGE</td>
<td>NOSTORAGE, CONFORMANCE</td>
<td>NOCONFORMANCE)</td>
</tr>
<tr>
<td>CMPAT(LE</td>
<td>V1</td>
<td>V2</td>
</tr>
<tr>
<td>CODEPAGE(n)</td>
<td>CP</td>
<td>CODEPAGE(1140)</td>
</tr>
<tr>
<td>COMMON</td>
<td>NOCOMMON</td>
<td>-</td>
</tr>
<tr>
<td>COMPIL</td>
<td>NOCOMPIL[(W</td>
<td>E</td>
</tr>
<tr>
<td>COPYRIGHT(‘string’)</td>
<td>NOCOPYRIGHT</td>
<td>-</td>
</tr>
<tr>
<td>CSECT</td>
<td>NOCSECT</td>
<td>CSE</td>
</tr>
<tr>
<td>CSECTCUT(n)</td>
<td>-</td>
<td>CSECTCUT(4)</td>
</tr>
<tr>
<td>CURRENCY(‘c’)</td>
<td>NOCOPYRIGHT</td>
<td>-</td>
</tr>
<tr>
<td>DBCS</td>
<td>NODBCS</td>
<td>-</td>
</tr>
<tr>
<td>DD(ddname-list)</td>
<td>-</td>
<td>DD(SYSPRINT, SYSIN, SYSLIB, SYSDEBUG)</td>
</tr>
<tr>
<td>DECIMAL(FOFLONASGN</td>
<td>NOFOFLONASGN, FLOMUL</td>
<td>NOFLOMULMUL, FORCEDSIGN, NOFORCEDSIGN)</td>
</tr>
<tr>
<td>DEFAULT(attribute</td>
<td>option)</td>
<td>DFT</td>
</tr>
<tr>
<td>DEPRECATE(BUILTIN(built-in-name)</td>
<td>ENTRY(entry-name)</td>
<td>INCLUDE(filename)</td>
</tr>
<tr>
<td>DEPRECATENEXT(BUILTIN(built-in-name)</td>
<td>ENTRY(entry-name)</td>
<td>INCLUDE(filename)</td>
</tr>
<tr>
<td>DISPLAY</td>
<td>STD</td>
<td>WTO(ROUTCDE(x) DESC(y) REPLY(z))</td>
</tr>
<tr>
<td>DLLINIT</td>
<td>NODLLINIT</td>
<td>-</td>
</tr>
<tr>
<td>EXIT</td>
<td>NOEXIT</td>
<td>-</td>
</tr>
<tr>
<td>EXTRN(FULL</td>
<td>SHORT)</td>
<td>-</td>
</tr>
<tr>
<td>FILEREF</td>
<td>NOFILEREF</td>
<td>-</td>
</tr>
<tr>
<td>FLAG[(I</td>
<td>W</td>
<td>E</td>
</tr>
<tr>
<td>FLOAT(AFP(NOVOLATILE</td>
<td>VOLATILE)</td>
<td>NOAFP, DFP</td>
</tr>
<tr>
<td>FLOATINMATH(ASIS</td>
<td>LONG</td>
<td>EXTENDED)</td>
</tr>
<tr>
<td>GOFF</td>
<td>NOGOFF</td>
<td>-</td>
</tr>
<tr>
<td>GONUMBER(SEPARATE</td>
<td>NOSEPARATE)</td>
<td>NOGONUMBER</td>
</tr>
<tr>
<td>GRAPHIC</td>
<td>NOGRAPHIC</td>
<td>GR</td>
</tr>
<tr>
<td>HGPR[(PRESERVE</td>
<td>NOPRESEVE)]</td>
<td>NOHGPR</td>
</tr>
<tr>
<td>IGNORE(ASSERT</td>
<td>DISPLAY</td>
<td>PUT)</td>
</tr>
<tr>
<td>INCAFTER([PROCESS(filename)])</td>
<td>-</td>
<td>INCAFTER()</td>
</tr>
</tbody>
</table>
Table 3. Compile-time options, abbreviations, and IBM-supplied defaults (continued)

<table>
<thead>
<tr>
<th>Compile-Time Option</th>
<th>Abbreviated Name</th>
<th>z/OS Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>INCDIR('directory name')</td>
<td>NOINCDIR</td>
<td>-</td>
</tr>
<tr>
<td>INCLUDE</td>
<td>NOINCLUDE</td>
<td>-</td>
</tr>
<tr>
<td>INCPDS('PDS name')</td>
<td>NOINCPDS</td>
<td>-</td>
</tr>
<tr>
<td>INITAUTO([ SHORT</td>
<td>FULL ])</td>
<td>NOINITAUTO</td>
</tr>
<tr>
<td>INITBASED</td>
<td>NOINITBASED</td>
<td>-</td>
</tr>
<tr>
<td>INITCTLC</td>
<td>NOINITCTLC</td>
<td>-</td>
</tr>
<tr>
<td>INITSTATIC</td>
<td>NOINITSTATIC</td>
<td>-</td>
</tr>
<tr>
<td>INSOURCE([FULL</td>
<td>SHORT])</td>
<td>NOINSOURCE</td>
</tr>
<tr>
<td>INTERRUPT</td>
<td>NOINTERRUPT</td>
<td>INT</td>
</tr>
<tr>
<td>JSON(CASE( UPPER</td>
<td>ASIS))</td>
<td>-</td>
</tr>
<tr>
<td>LANGLVL(NOEXT</td>
<td>OS)</td>
<td>-</td>
</tr>
<tr>
<td>LIMITS(options)</td>
<td>-</td>
<td>See “LIMITS” on page 46</td>
</tr>
<tr>
<td>LINECOUNT(n)</td>
<td>LC</td>
<td>LINECOUNT(60)</td>
</tr>
<tr>
<td>LINEDIR</td>
<td>NOLINEDIR</td>
<td>-</td>
</tr>
<tr>
<td>LIST</td>
<td>NOLIST</td>
<td>-</td>
</tr>
<tr>
<td>LISTVIEW(SOURCE</td>
<td>AFTERMACRO</td>
<td>AFTERCICS</td>
</tr>
<tr>
<td>MACRO</td>
<td>NOMACRO</td>
<td>M</td>
</tr>
<tr>
<td>MAP</td>
<td>NOMAP</td>
<td>-</td>
</tr>
<tr>
<td>MARGINI(‘c’)</td>
<td>NOMARGINI</td>
<td>MI</td>
</tr>
<tr>
<td>MARGINS(m,n[,c])</td>
<td>NOMARGINS</td>
<td>MAR(m,n)</td>
</tr>
<tr>
<td>MAXBRANCH(max)</td>
<td>-</td>
<td>MAXBRANCH(2000)</td>
</tr>
<tr>
<td>MAXGEN(n)</td>
<td>-</td>
<td>MAXGEN(100000)</td>
</tr>
<tr>
<td>MAXMEM(n)</td>
<td>MAXM</td>
<td>MAXMEM(1048576)</td>
</tr>
<tr>
<td>MAXMSG(I</td>
<td>W</td>
<td>E</td>
</tr>
<tr>
<td>MAXNEST(BLOCK(x) DO(y) IF(z))</td>
<td>-</td>
<td>MAXNEST(BLOCK(17) DO(17) IF(17))</td>
</tr>
<tr>
<td>MAXSTMT(n)</td>
<td>-</td>
<td>MAXSTMT(4096)</td>
</tr>
<tr>
<td>MAXTEMP(n)</td>
<td>-</td>
<td>MAXTEMP(50000)</td>
</tr>
<tr>
<td>MDECK</td>
<td>NOMDECK</td>
<td>MD</td>
</tr>
<tr>
<td>MSGSUMMARY([XREF</td>
<td>NOXREF])</td>
<td>NOSMSGSUMMARY</td>
</tr>
<tr>
<td>NAME(‘external name’)</td>
<td>NONAME</td>
<td>N</td>
</tr>
<tr>
<td>NAMES(‘lower’</td>
<td>‘upper’)</td>
<td>-</td>
</tr>
<tr>
<td>NATLANG(ENU</td>
<td>UEN)</td>
<td>-</td>
</tr>
<tr>
<td>NEST</td>
<td>NONEST</td>
<td>-</td>
</tr>
<tr>
<td>NOT</td>
<td>-</td>
<td>NOT(“~”)</td>
</tr>
<tr>
<td>NUMBER</td>
<td>NONUMBER</td>
<td>NUM</td>
</tr>
<tr>
<td>OBJECT</td>
<td>NOOBJECT</td>
<td>OBJ</td>
</tr>
<tr>
<td>OFFSET</td>
<td>NOOFFSET</td>
<td>OF</td>
</tr>
</tbody>
</table>
Table 3. Compile-time options, abbreviations, and IBM-supplied defaults (continued)

<table>
<thead>
<tr>
<th>Compile-Time Option</th>
<th>Abbreviated Name</th>
<th>z/OS Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>ONSNAP(STRINGRANGE, STRINGSIZE)</td>
<td>NOONSNAP</td>
<td>NOONSNAP</td>
</tr>
<tr>
<td>OPTIMIZE(0</td>
<td>2</td>
<td>3)</td>
</tr>
<tr>
<td>OPTIONS[(ALL</td>
<td>DOC)]</td>
<td>NOOPTIONS</td>
</tr>
<tr>
<td>OR('c')</td>
<td>-</td>
<td>OR('c')</td>
</tr>
<tr>
<td>PP(pp-name)</td>
<td>NOPP</td>
<td>NOPP</td>
</tr>
<tr>
<td>PPCICS('string ')</td>
<td>NOPPCICS</td>
<td>NOPPCICS</td>
</tr>
<tr>
<td>PPINCLUDE('string ')</td>
<td>NOPPINCLUDE</td>
<td>NOPPINCLUDE</td>
</tr>
<tr>
<td>PPLIST(KEEP</td>
<td>ERASE)</td>
<td>PPLIST(KEEP)</td>
</tr>
<tr>
<td>PP MACRO('string')</td>
<td>NOPPMACRO</td>
<td>NOPPMACRO</td>
</tr>
<tr>
<td>PP SQL('string')</td>
<td>NOPPSQL</td>
<td>NOPPSQL</td>
</tr>
<tr>
<td>PPTRACE</td>
<td>NOPPT RACE</td>
<td>NOPPT RACE</td>
</tr>
<tr>
<td>PREFIX(condition)</td>
<td>-</td>
<td>PREFIX(condition)</td>
</tr>
<tr>
<td>PRECTYPE(ANS</td>
<td>DEC DIGIT</td>
<td>DEC RESULT)</td>
</tr>
<tr>
<td>PROCEED</td>
<td>NOPROCEED[(W</td>
<td>E</td>
</tr>
<tr>
<td>PROCESS[(KEEP</td>
<td>DELETE)]</td>
<td>NOPROCESS</td>
</tr>
<tr>
<td>QUOTE('')</td>
<td>-</td>
<td>QUOTE('')</td>
</tr>
<tr>
<td>REDUCE</td>
<td>NOREDUCE</td>
<td>REDUCE</td>
</tr>
<tr>
<td>RENT</td>
<td>NOREN T</td>
<td>REDUCE</td>
</tr>
<tr>
<td>REEXP</td>
<td>NOREEXP</td>
<td>REEXP</td>
</tr>
<tr>
<td>RESPECT([DATE])</td>
<td>-</td>
<td>RESPECT()</td>
</tr>
<tr>
<td>RTCHECK(NULLPTR</td>
<td>NONULLPTR</td>
<td>NULL370)</td>
</tr>
<tr>
<td>RULES(options)</td>
<td>-</td>
<td>RULES(options)</td>
</tr>
<tr>
<td>SEMANTIC</td>
<td>NOSEMANTIC[(W</td>
<td>E</td>
</tr>
<tr>
<td>SERVICE('service string')</td>
<td>NOSERVICE</td>
<td>SERV</td>
</tr>
<tr>
<td>SOURCE</td>
<td>NOSOURCE</td>
<td>S</td>
</tr>
<tr>
<td>SPILL(n)</td>
<td>SP</td>
<td>SPILL(512)</td>
</tr>
<tr>
<td>STATIC(FULL</td>
<td>SHORT)</td>
<td>-</td>
</tr>
<tr>
<td>STDSYS</td>
<td>NOSTDSYS</td>
<td>NOSTDSYS</td>
</tr>
<tr>
<td>STMT</td>
<td>NOSTMT</td>
<td>-</td>
</tr>
<tr>
<td>STORAGE</td>
<td>NOSTORAGE</td>
<td>STG</td>
</tr>
<tr>
<td>STRINGOFGRAPHIC(CHAR</td>
<td>GRAPHIC)</td>
<td>-</td>
</tr>
<tr>
<td>SYNTAX</td>
<td>NOSYNTAX[(W</td>
<td>E</td>
</tr>
<tr>
<td>SYSPARM('string')</td>
<td>SYSPARM('')</td>
<td>SYSPARM('')</td>
</tr>
<tr>
<td>SYSTEM(MVS</td>
<td>CICS</td>
<td>IMS</td>
</tr>
<tr>
<td>TERMINAL</td>
<td>NOTERMINAL</td>
<td>TERM</td>
</tr>
<tr>
<td>TEST(options)</td>
<td>NOTEST</td>
<td>TEST(options)</td>
</tr>
<tr>
<td>UNROLL(AUTO</td>
<td>NO)</td>
<td>UNROLL(AUTO)</td>
</tr>
<tr>
<td>USAGE(options)</td>
<td>-</td>
<td>USAGE(options)</td>
</tr>
<tr>
<td>WIDECHAR(BIGENDIAN</td>
<td>LITTLEENDIAN)</td>
<td>WCHAR</td>
</tr>
</tbody>
</table>
Table 3. Compile-time options, abbreviations, and IBM-supplied defaults (continued)

<table>
<thead>
<tr>
<th>Compile-Time Option</th>
<th>Abbreviated Name</th>
<th>z/OS Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>WINDOW(w)</td>
<td>-</td>
<td>WINDOW(1950)</td>
</tr>
<tr>
<td>WRITABLE</td>
<td>NOWRITABLE[(FWS</td>
<td>PRV)]</td>
</tr>
<tr>
<td>XINFO(options)</td>
<td>-</td>
<td>XINFO(NODEF,NOMSG, NOSYM, NOSYN, NOXMI, NOXML)</td>
</tr>
<tr>
<td>XML(CASE( UPPER</td>
<td>ASIS))</td>
<td>-</td>
</tr>
<tr>
<td>XREF[(FULL</td>
<td>SHORT)]</td>
<td>NOXREF</td>
</tr>
</tbody>
</table>

Notes:
1. FULL is the default suboption if the suboption is omitted with ATTRIBUTES or XREF.
2. The default value for the BLANK character is the tab character with value '05'x.
3. (ALL,SYM) is the default suboption if the suboption is omitted with TEST.

**AGGREGATE**

The AGGREGATE option creates an Aggregate Length Table that gives the lengths of arrays and major structures in the source program in the compiler listing.

ABBREVIATIONS: AG, NAG

The suboptions of the AGGREGATE option determine how the offsets of subelements are displayed in the Aggregate Length Table:

**DECIMAL**

All offsets are displayed in decimal.

**HEXADEC**

All offsets are displayed in hexadecimal.

In the Aggregate Length Table, the length of an undimensioned major or minor structure is always expressed in bytes, but the length might not be accurate if the major or minor structure contains unaligned bit elements.

The Aggregate Length Table includes structures but not arrays that have non-constant extents. However, the sizes and offsets of elements within structures with non-constant extents may be inaccurate or specified as *.

**ARCH**

The ARCH option specifies the architecture for which the instructions of executable programs are to be generated. It allows the optimizer to take advantage of specific hardware instruction sets.

The current values that can be specified for the ARCH level are:
7 Produces code that uses instructions available on the 2096-xxx (IBM System z9 BC) and 2094-xxx (IBM System z9 EC) models in z/Architecture mode. Specifically, these ARCH(7) machines and their follow-ons add instructions that are supported by facilities such as extended-immediate and extended translation facilities, which can be used by the compiler. For further information about these facilities, see the z/Architecture Principles of Operation.

ARCH(7) machines also support the DFP instructions.

8 Produces code that uses instructions available on the 2097-xxx models (IBM System z10 EC) in z/Architecture mode.

ARCH(8) produces code that uses instructions available on the 2098-xxx models (IBM System z10<sup>®</sup> BC) and 2097-xxx models (IBM System z10 EC) in z/Architecture mode. Specifically, these ARCH(8) machines and subsequent generations instructions supported by the general instruction extensions facility, which may be exploited by the compiler.

9 Produces code that uses instructions available on the 2817-xxx (IBM zEnterprise 196 (z196)) and 2818-xxx models (IBM zEnterprise 114 (z114)) models in z/Architecture mode.

Specifically, these ARCH(9) machines and their follow-ons add instructions supported by the following facilities, which may be exploited by the compiler:

- The high-word facility
- The population count facility
- The distinct-operands facility
- The floating-point extension facility
- The load/store-on-condition facility

For further information about these facilities, see the z/Architecture Principles of Operation.

10 Produces code that uses instructions available on the 2827-xxx (IBM zEnterprise EC12) models in z/Architecture mode.

Specifically, these ARCH(10) machines and their follow-ons add instructions supported by the following facilities, which may be exploited by the compiler:

- The execution-hint facility
- The load-and-trap facility
- The miscellaneous-instructions-extension facility
- The transactional-execution facility

For further information about these facilities, see the z/Architecture Principles of Operation.

11 Produces code that uses instructions available on the 2964-xxxx (IBM z13) models in z/Architecture mode.

Specifically, these ARCH(11) machines and their follow-ons add instructions supported by the following facilities:

- Vector facility

Notes:
1. If you specify an ARCH value less than 7, the compiler resets it to 7.
2. The x in the model numbers above (such as 2084-xxx) is a wildcard and stands for any alphanumeric machine of that type.
3. Code that is compiled at ARCH(n) runs on machines in the ARCH(m) group if and only if m \geq n.
4. You can mix code that is compiled with different ARCH levels without any restrictions.

**ATTRIBUTES**

The ATTRIBUTES option specifies that the compiler includes a table of source-program identifiers and their attributes in the compiler listing.

```
NOATTRIBUTES
ATTRIBUTES
    (FULL)
    (SHORT)
```

**ABBREVIATIONS:** A, NA, F, S

**FULL**

All identifiers and attributes are included in the compiler listing. FULL is the default.

**SHORT**

Unreferenced identifiers are omitted, making the listing more manageable.

If you include both ATTRIBUTES and XREF (creates a cross-reference table), the two tables are combined. However, if the SHORT and FULL suboptions are in conflict, the last option specified is used. For example, if you specify ATTRIBUTES(SHORT) XREF(FULL), FULL applies to the combined listing.

**BACKREG**

The BACKREG option controls the backchain register, which is the register used to pass the address of the automatic storage of a parent routine when a nested routine is invoked.

```
BACKREG (5)
BACKREG (11)
```

For best compatibility with PL/I for MVS & VM, OS PL/I V2R3 and earlier compilers, BACKREG(5) should be used.

All routines that share an ENTRY VARIABLE must be compiled with the same BACKREG option, and it is strongly recommended that all code in application be compiled with the same BACKREG option.

Note that code compiled with VisualAge PL/I for OS/390 effectively used the BACKREG(11) option. Code compiled with Enterprise PL/I V3R1 or V3R2 also used the BACKREG(11) option by default.

**BIFPREC**

The BIFPREC option controls the precision of the FIXED BIN result returned by various built-in functions.
For best compatibility with PL/I for MVS & VM, OS PL/I V2R3 and earlier compilers, BIFPREC(15) should be used.

BIFPREC affects the following built-in functions:
- COUNT
- INDEX
- LENGTH
- LINENO
- ONCOUNT
- PAGENO
- SEARCH
- SEARCHR
- SIGN
- VERIFY
- VERIFYR

The effect of the BIFPREC compiler option is most visible when the result of one of the above built-in functions is passed to an external function that has been declared without a parameter list. For example, consider the following code fragment:

```pli
dcl parm char(40) var;
dcl funky ext entry( pointer, fixed bin(15) );
dcl beans ext entry;
call beans( addr(parm), verify(parm), ' ' );
```

If the function `beans` actually declares its parameters as POINTER and FIXED BIN(15), then if the code above were compiled with the option BIFPREC(31) and if it were run on a big-endian system such as z/OS, the compiler would pass a four-byte integer as the second argument and the second parameter would appear to be zero.

Note that the function `funky` would work on all systems with either option.

The BIFPREC option does not affect the built-in functions DIM, HBOUND and LBOUND. The CMPAT option determines the precision of the FIXED BIN result returned these three functions: under CMPAT(V1), these array-handling functions return a FIXED BIN(15) result, while under CMPAT(V2) and CMPAT(LE), they return a FIXED BIN(31) result. Under CMPAT(V3), they return a FIXED BIN(63) result.

**BLANK**

The BLANK option specifies up to ten alternate symbols for the blank character.

```plaintext
BLANK( 'char' )
```

**Note:** Do not code any blanks between the quotes.
The IBM-supplied default code point for the BLANK symbol is '05X.

\textbf{char}

A single SBCS character.

You cannot specify any of the alphabetic characters, digits, and special characters defined in the \textit{PL/I Language Reference}.

If you specify the BLANK option, the standard blank symbol is still recognized as a blank.

\textbf{BLKOFF}

The BLKOFF option controls whether the offsets shown in the pseudo-assembler listing (produced by the LIST option) and the statement offset listing (produced by the OFFSET option) are from the start of the current module or from the start of the current procedure.

The pseudo-assembler listing also includes the offset of each block from the start of the current module (so that the offsets shown for each statement can be translated to either block or module offsets).

\textbf{CASERULES}

The CASERULES option controls the enforcement of case rules for keywords.

\textbf{LOWER}

Instructs the compiler to flag any keyword that is not in lowercase.

\textbf{MIXED}

Instructs the compiler to accept all keywords as they are coded. MIXED is the default.

\textbf{START}

Instructs the compiler to flag any keyword whose first letter is not in uppercase or whose remaining letters are not in lowercase.

\textbf{UPPER}

Instructs the compiler to flag any keyword that is not in uppercase.

\textbf{Notes:}

1. The CASERULES option does not apply to elements of the OPTIONS and ENVIRONMENT attributes.
2. The CASERULES option does not apply to any of the preprocessors.

\textbf{CEESTART}

The CEESTART option specifies whether the compiler should place the CEESTART csect before or after all the other generated object code.
Under the CEESTART(FIRST) option, the compiler places the CEESTART csect before all the other generated object code; however, under the CEESTART(LAST) option, the compiler places it after all the other generated object code.

Using CEESTART(FIRST) will cause the binder to choose CEESTART as the entry point for a module if no ENTRY card is specified during the bind step.

Those users who want to use linker CHANGE cards must use the CEESTART(LAST) option.

However, MAIN routines should be linked with an ENTRY CEESTART linkage editor card. But, if you use the CEESTART(LAST) option, then you must include an ENTRY CEESTART card when linking your MAIN routine.

CHECK

The CHECK option specifies whether the compiler should generate special code to detect various programming errors.

ABBREVIATIONS: STG, NSTG

Specifying CHECK(CONFORMANCE) causes the compiler to generate, under the following circumstances, code that checks at run time if the attributes of the arguments passed to a procedure match those of the declared parameters:

- If a parameter is a string (or an array of strings) declared with a constant length, then the STRINGSIZE condition will be raised if the argument passed does not have matching length.
- If a parameter is a string (or an array of strings), then the STRINGSIZE condition will be raised if the argument does not have the same length type (VARYING, NONVARYING or VARYINGZ).
- If a parameter is an array (of scalars or structures), then the SUBSCRIPTRANGE condition will be raised if any constant bounds do not match those of the passed argument. The SUBSCRIPTRANGE condition will also be raised if all the extents are constant and the size and spacing of the array elements in the argument do not match those in the parameter. Arrays inside a structure are not checked.
- If a parameter is a structure or union with constant extents, then the SUBSCRIPTRANGE condition will be raised if the offset of the last element does not match that of the passed argument.
- If the procedure has the RETURNS BYADDR attribute and that attribute specifies a string type, then the STRINGSIZE condition will be raised if the string passed for the RETURNS value does not have matching length.
This extra code will not be generated if the NODESCRIPTOR option applies to the procedure or if the block contains ENTRY statements or if the CMPAT(LE) option is in effect.

When you specify CHECK(STORAGE), the compiler calls slightly different library routines for ALLOCATE and FREE statements (except when these statements occur within an AREA). The following built-in functions, described in the PL/I Language Reference, can be used only when CHECK(STORAGE) has been specified:

- ALLOCSIZE
- CHECKSTG
- UNALLOCATED

AMODE(24) is not recommended for Enterprise PL/I applications. For code compiled with the CHECK(STORAGE) option, if you have to use AMODE(24), then you must also specify the HEAP(,,BELOW) runtime option.

**CMPAT**

The CMPAT option specifies whether object compatibility with OS PL/I Version 1, OS PL/I Version 2, PL/I for MVS and VM, VisualAge PL/I for OS/390 or Enterprise PL/I for z/OS is to be maintained for programs sharing strings, AREAs, arrays and/or structures.

/SM590000/SM590000

CMPAT

\[
\begin{array}{c}
V2 \\
LE \\
V1 \\
V3 \\
\end{array}
\]

**ABBREVIATIONS:** CMP

**LE** Under CMPAT(LE), your program can share strings, AREAs, arrays and/or structures only with programs compiled with VisualAge PL/I for OS/390 or Enterprise PL/I for z/OS and only as long as the CMPAT(V1) and CMPAT(V2) options were not used when they were compiled.

**V1** Under CMPAT(V1), you can share strings, AREAs, arrays and/or structures with programs compiled with the OS PL/I compiler and with programs compiled with later PL/I compilers as long as the CMPAT(V1) option was used. The JSON built-in functions are not supported under CMPAT(V1).

**V2** Under CMPAT(V2), you can share strings, AREAs, arrays and/or structures with programs compiled with the OS PL/I compiler and with programs compiled with later PL/I compilers as long as the CMPAT(V2) option was used.

**V3** Under CMPAT(V3), you can share strings with programs compiled with the OS PL/I compiler and with programs compiled with later PL/I compilers as long as one of the CMPAT(V*) options were used. However, you can not share AREAs, arrays or structures with any code not compiled with CMPAT(V3).

DB2 stored procedures must not be compiled with CMPAT(LE).

All the modules in an application must be compiled with the same CMPAT option.

Mixing old and new code still has some restrictions. For information about these restrictions see the *Enterprise PL/I for z/OS Compiler and Run-Time Migration Guide*.
The DFT(DESCLIST) option conflicts with the CMPAT(V*) options, and if it is specified with any of them, a message will be issued and the DFT(DESCLOCATOR) option assumed.

Under CMPAT(V3), arrays may be declared with any value that an 8-byte integer could assume. However, the total size of an array currently still has the same limit as an array declared under CMPAT(V2).

Under CMPAT(V3), the following built-in functions will always return a FIXED BIN(63) result:
- CURRENTSIZE/CSTG
- DIMENSION
- HBOUND
- LBOUND
- LOCATION
- SIZE/STG

Because these functions will be returning 8-byte integer values, under CMPAT(V3), the second option in the FIXEDBIN suboption of the LIMITS option must be 63.

However, even under CMPAT(V3), statement and format label constants must be specified using 4-byte integers.

**CODEPAGE**

The CODEPAGE option specifies the code page used for:
- conversions between CHARACTER and WIDECHAR
- the default code page used by the PLISAX built-in subroutines

```
<table>
<thead>
<tr>
<th>ccсид</th>
</tr>
</thead>
<tbody>
<tr>
<td>01047</td>
</tr>
<tr>
<td>01140</td>
</tr>
<tr>
<td>01141</td>
</tr>
<tr>
<td>01142</td>
</tr>
<tr>
<td>01143</td>
</tr>
<tr>
<td>01144</td>
</tr>
<tr>
<td>01025</td>
</tr>
</tbody>
</table>
```

The default CCSID 1140 is an equivalent of CCSID 37 (EBCDIC Latin-1, USA) but includes the Euro symbol.

**COMMON**

The COMMON option directs the compiler to generate CM linkage records for EXTERNAL STATIC variables.
Under the COMMON option, if the NORENT option applies, then CM linkage records will be generated for EXTERNAL STATIC variables that are not RESERVED and which contain no INITIAL values. This matches what the OS PL/I compiler did.

Under the NOCOMMON option, SD records will be written as was true in earlier releases of Enterprise PL/I.

The COMMON option must not be used with the RENT option or with LIMITS(EXTNAME(n)) if n > 7.

**COMPILE**

The COMPILE option causes the compiler to stop compiling after all semantic checking of the source program if it produces a message of a specified severity during preprocessing or semantic checking. Whether the compiler continues or not depends on the severity of the error detected, as specified by the NOCOMPILE option in the list below. The NOCOMPILE option specifies that processing stops unconditionally after semantic checking.

**ABBREVIATIONS: C, NC**

**COMPILE**

Generates code unless a severe error or unrecoverable error is detected. Equivalent to NOCOMPILE(S).

**NOCOMPILE**

Compilation stops after semantic checking.

**NOCOMPILE(W)**

No code generation if a warning, error, severe error, or unrecoverable error is detected.

**NOCOMPILE(E)**

No code generation if an error, severe error, or unrecoverable error is detected.

**NOCOMPILE(S)**

No code generation if a severe error or unrecoverable error is detected.

If the compilation is terminated by the NOCOMPILE option, the cross-reference listing and attribute listing can be produced; the other listings that follow the source program will not be produced.

**COPYRIGHT**

The COPYRIGHT option places a string in the object module, if generated. This string is loaded into memory with any load module into which this object is linked.
The string is limited to 1000 characters in length. However, if the string is longer than 100 characters, it will not be shown in the options listing.

To ensure that the string remains readable across locales, only characters from the invariant character set should be used.

**CSECT**

The CSECT option ensures that the object module, if generated, contains named CSECTs. Use this option if you will be using SMP/E to service your product or to aid in debugging your program.

```
CSECT

```

**ABBREVIATIONS: CSE, NOCSE**

Under the NOCSECT option, the code and static sections of your object module are given default names.

Under the CSECT option, the code and static sections of your object module are given names that depend on the "package name" which is defined as follows:

- if the package statement was used, the "package name" is the leftmost label on the package statement
- otherwise, the "package name" is the leftmost label on the first procedure statement.

A "modified package name" of length 7 is then formed as follows:

- when the package name is less than 7 characters long, "*"s are prefixed to it to make a modified package name that is 7 characters long
- when the package name is more than 7 characters long, the first \( n \) and last \( 7 - n \) characters are used to make the modified package name, where the value \( n \) is set by the CSECTCUT option
- otherwise the package name is copied to the modified package name

The code csect name is built by taking the modified package name and appending a '1' to it.

The static csect name is built by taking the modified package name and appending a '2' to it.

So, for a package named "SAMPLE", the code csect name would be "SAMPLE1", and the static csect name would be "SAMPLE2".

**CSECTCUT**

The CSECTCUT option controls how the compiler, when processing the CSECT option, handles long names.

```
CSECTCUT ( n )
```
The CSECTCUT option has no effect unless you specify the CSECT option. It also has no effect if the "package name" used by the CSECT option has 7 or fewer characters.

The value \( n \) in the CSECTCUT option must be between 0 and 7.

If the "package name" used by the CSECT option has more than 7 characters, the compiler will collapse the name to 7 characters by taking the first \( n \) and last \( 7 - n \) characters.

For example, for a compilation consisting of one procedure with the name BEISPIEL,
- under CSECTCUT(3), the compiler would collapse the name to BEIPIEL
- under CSECTCUT(4), the compiler would collapse the name to BEISIEL

**CURRENCY**

The CURRENCY option allows you to specify an alternate character to be used in picture strings instead of the dollar sign.

ABBREVIATIONS: CURR

\( x \) Character that you want the compiler and run time to recognize and accept as the dollar sign in picture strings.

**DBCS**

The DBCS option ensures that the listing, if generated, is sensitive to the possible presence of DBCS even though the GRAPHIC option has not been specified.

Under z/OS UNIX System Services, the NODBCS option causes the listing, if generated, to show all DBCS shift-codes as ".".

Under both batch and z/OS UNIX System Services, the NODBCS option ensures that the header text in the listing page does not contain unmatched shift codes.

The NODBCS option should not be specified if the GRAPHIC option is also specified.

**DD**

The DD option allows you to specify alternate DD names for the various datasets used by the compiler.
Up to eight DD names may be specified. In order, they specify alternate DD names for:
- SYSPRINT
- SYSIN
- SYSLIB
- SYSPUNCH
- SYSLIN
- SYSADATA
- SYSXMLSD
- SYSDEBUG

If you wanted to use ALTIN as the DD name for the primary compiler source file, you would have to specify DD(SYSPRINT,ALTIN). If you specified DD(ALTIN), SYSIN would be used as the DDNAME for the primary compiler source file and ALTIN would be used as the DD name for the compiler listing.

You can also use * to indicate that the default DD name should be used. Thus DD(*,ALTIN) is equivalent to DD(SYSPRINT,ALTIN).

**DDSQ**

The DDSQ option allows you to specify an alternate DD name for the dataset to be used by the SQL preprocessor when resolving EXEC SQL INCLUDE statements.

```
>> DDSQ( ddname )
```

Under the DDSQ("") option, the DD name that is used to resolve EXEC SQL INCLUDE statements is the DD name for SYSLIB from the DD compiler option.

This option may be useful when moving from the SQL precompiler to the integrated SQL preprocessor.

**DECIMAL**

The DECIMAL option specifies how the compiler should handle certain FIXED DECIMAL operations and assignments.

```
>> DECIMAL( )
```

- **FOFLONASGN**
  - The FOFLONASGN option requires the compiler to generate code that will raise the FIXEDOVERFLOW condition, if it is enabled and the SIZE condition is disabled, whenever a FIXED DECIMAL expression is assigned to a FIXED DECIMAL target and significant digits are lost.
Conversely, under the NOFOFLONASGN option, the compiler will generate code that will not raise the FIXEDOVERFLOW condition when significant digits are lost in such an assignment.

So, for example, given a variable A declared as FIXED DEC(5), the assignment A = A + 1 might raise FOFL under the FOFLONASGN option, but would never raise it under the NOFOFLONASGN option.

Note, however, that under the NOFOFLONASGN option, the FIXEDOVERFLOW condition can still be raised by operations that produce a result with more digits than allowed by the FIXEDDEC suboption of the LIMITS option. For example, given a variable B declared as FIXED DEC(15) with the value 999_999_999_999_999 and given that the FIXEDDEC suboption of the LIMITS specifies the maximum precision as 15, then the assignment B = B + 1 will raise the FIXEDOVERFLOW condition (if FOFL is enabled, of course) since the addition B + 1 will raise the condition.

**FOFLONMULT**

The FOFLONMULT option requires the compiler to generate code that will raise the FIXEDOVERFLOW condition for any use of the MULTIPLY built-in function that would produce a FIXED DEC result that is too large for the precision specified in the built-in function.

Conversely, under the NOFOFLONMULT option, the compiler will generate code that will produce a truncated result for any such use of the MULTIPLY built-in function.

Note that the use of the FOFLONMULT option changes the default language semantics (which would be to truncate a too large result of the MULTIPLY built-in function applied to FIXED DEC - unless the SIZE condition were enabled).

**FORCEDSIGN**

The FORCEDSIGN option will force the compiler to generate extra code to insure that whenever a FIXED DECIMAL result with the value zero is generated, the sign nibble of the result will have the value 'C'X. This option can cause the compiler to generate code that will perform much, much worse than the code generated under the NOFORCEDSIGN suboption.

Also, when this option is in effect, more data exceptions may occur when you run your code. For example, if you assign one FIXED DEC(5) variable to another FIXED DEC(5) variable, the compiler would normally generate a MVC instruction to perform the move. However if this option is in effect, in order to insure that the result has the preferred sign, the compiler will generate a ZAP instruction to perform the move. If the source contains invalid packed decimal data, the ZAP instruction, but not the MVC instruction, will raise a decimal data exception.

Under this option, data exceptions may also be raised when one PICTURE variable is assigned to another PICTURE variable since that conversion usually involves an implicit conversion to FIXED DEC which, under this option, will generate a ZAP instruction that will raise a data exception if the source contains invalid data.

Under DECIMAL(NOFORCEDSIGN), a "negative zero" may be produced by certain calculations. However, a programmer should rely on getting a negative zero only when assigning a negative literal to a FIXED DEC that cannot hold a value of that small a magnitude (such as assigning -.001 to a FIXED DEC(5,2) )
DEFAULT

The DEFAULT option specifies defaults for attributes and options. These defaults are applied only when the attributes or options are not specified or implied in the source code.
ABBREVIATIONS: DFT, ASGN, NONASGN, NONCONN, CONN, INL, NOINL

**ALIGNED | UNALIGNED**

This suboption allows you to force byte-alignment on all of your variables.

If you specify ALIGNED, all variables other than character, bit, graphic, and picture are given the ALIGNED attribute unless the UNALIGNED attribute is explicitly specified (possibly on a parent structure) or implied by a DEFAULT statement.

If you specify UNALIGNED, all variables are given the UNALIGNED attribute unless the ALIGNED attribute is explicitly specified (possibly on a parent structure) or implied by a DEFAULT statement.

ALIGNED is the default.

**IBM | ANS**

Use IBM or ANS SYSTEM defaults. The arithmetic defaults for IBM and ANS are the following:

<table>
<thead>
<tr>
<th>Attributes</th>
<th>DEFAULT(IBM)</th>
<th>DEFAULT(ANS)</th>
</tr>
</thead>
<tbody>
<tr>
<td>FIXED DECIMAL</td>
<td>(5,0)</td>
<td>(10,0)</td>
</tr>
<tr>
<td>FIXED BINARY</td>
<td>(15,0)</td>
<td>(31,0)</td>
</tr>
<tr>
<td>FLOAT DECIMAL</td>
<td>(6)</td>
<td>(6)</td>
</tr>
<tr>
<td>FLOAT BINARY</td>
<td>(21)</td>
<td>(21)</td>
</tr>
</tbody>
</table>

Under the IBM suboption, variables with names beginning from I to N default to FIXED BINARY and any other variables default to FLOAT DECIMAL. If you select the ANS suboption, the default for all variables is FIXED BINARY.

IBM is the default.

**ASCII | EBCDIC**

Use this option to set the default for the character set used for the internal representation of character problem program data.

Specify ASCII only when compiling programs that depend on the ASCII character set collating sequence. Such a dependency exists, for example, if your program relies on the sorting sequence of digits or on lowercase and uppercase alphabatics. This dependency also exists in programs that create an uppercase alphabetic character by changing the state of the high-order bit.

**Note:** The compiler supports A and E as suffixes on character strings. The A suffix indicates that the string is meant to represent ASCII data, even if the EBCDIC compiler option is in effect. Alternately, the E suffix indicates that the string is EBCDIC, even when you select DEFAULT(ASCII).

'123'A is the same as '313233'X
'123'E is the same as 'F1F2F3'X

EBCDIC is the default.

**ASSIGNABLE | NONASSIGNABLE**

This option causes the compiler to apply the specified attribute to all static variables that are not declared with the ASSIGNABLE or NONASSIGNABLE attribute. The compiler flags statements in which NONASSIGNABLE variables are the targets of assignments.

ASSIGNABLE is the default.
INONLY
Specifying NONASSIGNABLE(INONLY) indicates that parameters declared with the INONLY attribute are given the NONASSIGNABLE attribute.

STATIC
Specifying NONASSIGNABLE(STATIC) indicates that STATIC variables are given the NONASSIGNABLE attribute.

The INONLY and STATIC suboptions have no effect on either variables with the ASSIGNABLE or NONASSIGNABLE attribute or structure members that inherit the ASSIGNABLE or NONASSIGNABLE attribute from a parent.

BYVALUE parameters are given the INONLY attribute after the resolution of the (NON)ASSIGNABLE attribute, and hence the NONASSIGNABLE(INONLY) suboption has no effect on BYVALUE parameters.

To specify the NONASSIGNABLE attribute to both STATIC and INONLY variables, you must specify the suboption NONASSIGNABLE(STATIC INONLY).

The NONASSIGNABLE attribute can be specified without any suboptions, in which case it means NONASSIGNABLE(STATIC).

BIN1ARG | NOBIN1ARG
This suboption controls how the compiler handles 1-byte REAL FIXED BIN arguments passed to an unprototyped function.

Under BIN1ARG, the compiler passes a FIXED BIN argument as is to an unprototyped function.

But under NOBIN1ARG, the compiler assigns any 1-byte REAL FIXED BIN argument passed to an unprototyped function to a 2-byte FIXED BIN temporary and pass that temporary instead.

Consider the following example:
```
 dcl f1 ext entry;
 dcl f2 ext entry( fixed bin(15) );

call f1( 1b );
call f2( 1b );
```

If you specified DEFAULT(BIN1ARG), the compiler would pass the address of a 1-byte FIXED BIN(1) argument to the routine f1 and the address of a 2-byte FIXED BIN(15) argument to the routine f2. However, if you specified DEFAULT(NOBIN1ARG), the compiler would pass the address of a 2-byte FIXED BIN(15) argument to both routines.

Note that if the routine f1 was a COBOL routine, passing a 1-byte integer argument to it would cause problems since COBOL has no support for 1-byte integers. In this case, using DEFAULT(NOBIN1ARG) might be helpful; but it would be better to specify the argument attributes in the entry declare.

BIN1ARG is the default.

BYADDR | BYVALUE
Set the default for whether arguments or parameters are passed by address or by value. BYVALUE applies only to certain arguments and parameters. See the PL/I Language Reference for more information.

BYADDR is the default.
CONNECTED | NONCONNECTED
Set the default for whether parameters are connected or nonconnected.
CONNECTED allows the parameter to be used as a target or source in record-oriented I/O or as a base in string overlay defining.
NONCONNECTED is the default.

DESCLIST | DESCLOCATOR
When you specify DEFAULT(DESCLIST), the compiler passes all descriptors in a list as a 'hidden' last parameter.
If you specify DEFAULT(DESCLOCATOR), parameters requiring descriptors are passed using a locator or descriptor in the same way as previous releases of PL/I. This allows old code to continue to work even if it passed a structure from one routine to a routine that was expecting to receive a pointer.
The DFT(DESCLIST) option conflicts with the CMPAT(V*) options, and if it is specified with any of them, a message will be issued and the DFT(DESCLOCATOR) option assumed.
DESCLOCATOR is the default.

DESCRIPTOR | NODESCRIPTOR
Using DESCRIPTOR with a PROCEDURE indicates that a descriptor list was passed, while DESCRIPTOR with ENTRY indicates that a descriptor list should be passed. NODESCRIPTOR results in more efficient code, but has the following restrictions:
• For PROCEDURE statements, NODESCRIPTOR is invalid if any of the parameters have:
  – An asterisk (*) specified for the bound of an array, the length of a string, or the size of an area except if it is a VARYING or VARYINGZ string with the NONASSIGNABLE attribute
  – The NONCONNECTED attribute
  – The UNALIGNED BIT attribute
• For ENTRY declarations, NODESCRIPTOR is invalid if an asterisk (*) is specified for the bound of an array, the length of a string, or the size of an area in the ENTRY description list.
DESCRIPTOR is the default.

DUMMY(ALIGNED | UNALIGNED)
This suboption reduces the number of situations in which dummy arguments get created.
DUMMY(ALIGNED) indicates that a dummy argument should be created even if an argument differs from a parameter only in its alignment.
DUMMY(UNALIGNED) indicates that no dummy argument should be created for a scalar (except a nonvarying bit) or an array of such scalars if it differs from a parameter only in its alignment.
Consider the following example:

dcl
  1 a1 unaligned,
     2 b1 fixed bin(31),
     2 b2 fixed bin(15),
     2 b3 fixed bin(31),
     2 b4 fixed bin(15);

dcl x entry( fixed bin(31) );
call x( b3 );
If you specified DEFAULT(DUMMY(ALIGNED)), a dummy argument would be created, while if you specified DEFAULT(DUMMY(UNALIGNED)), no dummy argument would be created.

DUMMY(ALIGNED) is the default.

**E (HEXADEC | IEEE)**
The E suboption determines how many digits will be used for the exponent in E-format items.

If you specify E(IEEE), 4 digits will be used for the exponent in E-format items.

If you specify E(HEXADEC), 2 digits will be used for the exponent in E-format items.

If DFT(E(HEXADEC)) is specified, an attempt to use an expression whose exponent has an absolute value greater than 99 will cause the SIZE condition to be raised.

If the compiler option DFT(IEEE) is in effect, you should normally also use the option DFT(E(IEEE)). However, under this option, some E format items that would be valid under DFT(E(HEXADEC)) would not be valid. For instance, under DFT(E(IEEE)), the statement "put skip edit(x) (e(15,8));" would be flagged because the E format item is invalid.

E(HEXADEC) is the default.

**EVENDEC | NOEVENDEC**
This suboption controls the compiler’s tolerance of fixed decimal variables declared with an even precision.

Under NOEVENDEC, the precision for any fixed decimal variable is rounded up to the next highest odd number.

If you specify EVENDEC and then assign 123 to a FIXED DEC(2) variable, the SIZE condition is raised. If you specify NOEVENDEC, the SIZE condition is not raised.

EVENDEC is the default.

**HEXADEC | IEEE**
This suboption allows you to specify the default representation that is used to hold all FLOAT variables and all floating-point intermediate results. This suboption also determines whether the compiler evaluates floating-point expressions using the hexadecimal or IEEE float instructions and math routines.

Programs that communicate with JAVA should probably use the IEEE option, and programs that pass data to or receive data from platforms where IEEE is the default representation for floating-point data might also want to use the IEEE option.

HEXADEC is the default.

**INITFILL | NONINITFILL**
This suboption controls the default initialization of automatic variables.

If you specify INITFILL with a hex value (nn), that value is used to initialize the storage that is used by all automatic variables in a block each time that block is entered. If you do not enter a hex value, the default is ‘00’.

Note that the hex value can be specified without or without quotation marks, but if it is specified with quotation marks, the string should not have an X suffix.
Under NOINITFILL, the storage that is used by an automatic variable can hold arbitrary bit patterns unless the variable is explicitly initialized.

INITFILL can cause programs to run slower and should not be specified in production programs. However, the INITFILL option produces code that runs faster than the LE STORAGE option. Also, during program development, this option is useful for detecting uninitialized automatic variables: a program that runs correctly with DFT(INITFILL('00')) and with DFT(INITFILL('ff')) probably has no uninitialized automatic variables.

NOINITFILL is the default.

**INLINE | NOINLINE**

This option sets the default for the inline procedure option.

Specifying INLINE allows your code to run faster but, in some cases, also creates a larger executable file. For more information on how inlining can improve the performance of your application, see Chapter 13, “Improving performance,” on page 303.

NOINLINE is the default.

**LINKAGE**

The linkage convention for procedure invocations is:

**OPTLINK**

The default linkage convention for Enterprise PL/I. This linkage provides the best performance.

**SYSTEM**

The standard linking convention for system APIs.

LINKAGE(OPTLINK) should be used for all routines called by or calling to JAVA, and it should also be used for all routines called by or calling to C (unless the C code has been compiled with a non-default linkage).

LINKAGE(SYSTEM) should be used for all non-PL/I routines that expect the high-order bit to be on in the address of the last (and only the last) parameter.

Note that specifying OPTIONS(ASSEMBLER) for a PROCEDURE or an ENTRY forces LINKAGE(SYSTEM) regardless of the setting of this option.

LINKAGE(OPTLINK) is the default.

**LOWERINC | UPPERINC**

If you specify LOWERINC, the compiler requires that the actual file names of INCLUDE files are in lowercase. If you specify UPPERINC, the compiler requires that the names are in uppercase.

Note: This suboption applies only to compilations under z/OS UNIX.

Under z/OS UNIX, the include name is built using the extension `.inc`. So, for example, under the DFT(LOWERINC) option, the statement `%INCLUDE STANDARD;` will cause the compiler to try to include `standard.inc`. But, under the DFT(UPPERINC) option, the statement `%INCLUDE STANDARD;` will cause the compiler to try to include `STANDARD.INC`.

LOWERINC is the default.

**NATIVE | NONNATIVE**

This option affects only the internal representation of fixed binary, ordinal,
offset, area, and varying string data. When the NONNATIVE suboption is in effect, the NONNATIVE attribute is applied to all such variables not declared with the NATIVE attribute.

You should specify NONNATIVE only to compile programs that depend on the nonnative format for holding these kind of variables.

If your program bases fixed binary variables on pointer or offset variables (or conversely, pointer or offset variables on fixed binary variables), specify either:

- Both the NATIVE and NATIVEADDR suboptions
- Both the NONNATIVE and NONNATIVEADDR suboptions.

Other combinations produce unpredictable results.

NATIVE is the default.

NATIVEADDR | NONNATIVEADDR
This option affects only the internal representation of pointers. When the NONNATIVEADDR suboption is in effect, the NONNATIVE attribute is applied to all pointer variables not declared with the NATIVE attribute.

If your program bases fixed binary variables on pointer or offset variables (or conversely, pointer or offset variables on fixed binary variables), specify either:

- Both the NATIVE and NATIVEADDR suboptions
- Both the NONNATIVE and NONNATIVEADDR suboptions.

Other combinations produce unpredictable results.

NATIVEADDR is the default.

NULLSYS | NULL370
This suboption determines which value is returned by the NULL built-in function. If you specify NULLSYS, binvalue(null()) is equal to 0. If you want binvalue(null()) to equal 'ff_00_00_00'xn as is true with previous releases of PL/I, specify NULL370.

NULL370 is the default.

NULLSTRADDR | NONULLSTRADDR
This suboption controls how the compiler handles null strings when passed as arguments.

Under NULLSTRADDR, when a null string is specified as an argument in an entry invocation, the compiler will pass the address of an initialized piece of automatic storage. This is compatible with what the OS PL/I and PL/I for MVS compilers did.

But under NONULLSTRADDR, when a null string is specified as an argument in an entry invocation, the compiler passes a null pointer as the address of the argument. This is compatible with what early releases of the Enterprise PL/I compiler did.

NULLSTRADDR is the default.

NULLSTRPTR
This suboption controls how the compiler handles null strings when assigned to POINTERs.

Under NULLSTRPTR(SYSNULL), the result of assigning " to a POINTER is the same as assigning SYSNULL() to the pointer.

Under NULLSTRPTR(NULL), the result of assigning " to a POINTER is the same as assigning NULL() to the pointer.
Under NULLSTRPTR(STRICT), assignments and comparisons of '' to
POINTERs are flagged as invalid.
NULLSTRPTR(NULL) is the default.

ORDER | REORDER
This suboption affects the optimization of the object code. Specifying
REORDER allows more optimization of your code. For detailed information,
see Chapter 13, “Improving performance,” on page 303.
REORDER is the default.

ORDINAL(MIN | MAX)
If you specify ORDINAL(MAX), all ordinals whose definition does not include
a PRECISION attribute is given the attribute PREC(31). Otherwise, they are
given the smallest precision that covers their range of values.
ORDINAL(MIN) is the default.

OVERLAP | NOOVERLAP
If you specify OVERLAP, the compiler presumes the source and target in an
assignment can overlap and generates, as needed, extra code in order to ensure
that the result of the assignment is okay.
The OVERLAP suboption applies only to string variables. It has no effect on
assignments of FIXED DECIMAL or other variable types. In those assignments,
the source and target must not overlap.
NOOVERLAP produces code that performs better; however, if you use
NOOVERLAP, you must ensure that the source and target never overlap.
NOOVERLAP is the default.

PSEUDODUMMY | NOPSEUDODUMMY
This suboption determines if dummy arguments are created when a SUBSTR
reference is specified as an argument to an unprototyped function.
If you specify PSEUDODUMMY, dummy arguments are created when a
SUBSTR reference is specified as an argument to an unprototyped function.
If you specify NOPSEUDODUMMY, dummy arguments are not created when
a SUBSTR reference is specified as an argument to an unprototyped function.
PSEUDODUMMY is the default.

RECURSIVE | NONRECURSIVE
When you specify DEFAULT(RECURSIVE), the compiler applies the
RECURSIVE attribute to all procedures. If you specify
DEFAULT(NONRECURSIVE), all procedures are nonrecursive except
procedures with the RECURSIVE attribute.
NONRECURSIVE is the default.

RETCODE | NORETCODE
If you specify RETCODE, for any external procedure that does not have the
RETURNS attribute, the compiler will generate extra code so that the
procedure returns the integer value obtained by invoking the PLIRETV built-in
function just before returning from that procedure.
If you specify NORETCODE, no special code is generated for procedures that
do not have the RETURNS attribute.
NORETCODE is the default.
RETURNS (BYVALUE | BYADDR)
Sets the default for how values are returned by functions. See the PL/I Language Reference for more information.

You should specify RETURNS(BYADDR) if your application contains ENTRY statements and the ENTRY statements or the containing procedure statement have the RETURNS option. You must also specify RETURNS(BYADDR) on the entry declarations for such entries.

RETURNS(BYADDR) is the default.

SHORT (HEXADEC | IEEE)
This suboption improves compatibility with other non-IBM UNIX compilers. SHORT (HEXADEC) maps FLOAT BIN (p) to a short (4-byte) floating point number for p <= 21. SHORT (IEEE) maps FLOAT BIN (p) to a short (4-byte) floating point number for p <= 24.

SHORT (HEXADEC) is the default.

Default: DEFAULT(ALIGNED IBM EBCDIC ASSIGNABLE BIN1ARG BYADDR NONCONNECTED DESCLOCATOR DESCRIPTOR DUMMY(ALIGNED) E(HEXADEC) EVENDEC HEXADEC NOINITFILL NOINLINE LINKAGE(OPTLINK) LOWERINC NATIVE NATIVEADDR NULL370 NULLSTRPTR(NULL) NULLSTRADDR REORDER ORDINAL(MIN) NOOVERLAP PSEUDODUMMY NONRECURSIVE NORETCODE RETURNS(BYADDR) SHORT(HEXADEC))

DEPRECATE
This option flags variable names, included file names, and statement names that you want to deprecate with error messages.

BUILTIN
Flags any declaration of built-in-name with the BUILTIN attribute.

built-in-name
Name of the BUILTIN variable

ENTRY
Flags any declaration of entry-name with the ENTRY attribute.

entry-name
Level-1 name

INCLUDE
Flags any %INCLUDE statement that includes filename.

Chapter 1. Using compiler options and facilities  31
**filename**
Name of the file

**STMT**
Flags all statements with the name as `statement-name`.

**statement-name**
Name of the statement. The names are identified by the initial keywords of PL/I statements. The STMT option accepts the following keywords:

allocate assert attach begin call close delay delete
detach display exit fetch flush free get goto
iterate leave locate on open put read release
resignal revert rewrite signal stop wait write

**VARIABLE**
Flags any declaration of `variable-name` that does not have the BUILTIN or ENTRY attributes.

**variable-name**
Level-1 name

To specify the DEPRECATE option, you must specify at least one of these suboptions with a possible empty suboption list. For example, both of the following two specifications are invalid:

- DEPRECATE
- DEPRECATE(BUILTIN)

Specifying one of the suboptions does not change the setting of any of the other suboptions that are specified previously.

Specifying a suboption a second time replaces the previous specifications.

In all cases, there is no checking of the suboption lists.

Examples:
- The following specifications are equivalent:
  
  DEPRECATE(ENTRY(old)) DEPRECATE(BUILTIN(cos))
  DEPRECATE(ENTRY(old) BUILTIN(cos))

- In the following example, x in the first specification is replaced by y:
  
  DEPRECATE(BUILTIN(x)) DEPRECATE(BUILTIN(y))
  DEPRECATE(BUILTIN(y))

**Default:** DEPRECATE(BUILTIN() ENTRY() INCLUDE() STMT() VARIABLE())

**DEPRECATENEXT**

The purpose of specifying the DEPRECATENEXT option and usage of this option is the same with the DEPRECATE option, except that the compiler issues warning messages rather than error messages for items that you will deprecate in a future development phase.

**Default:** DEPRECATENEXT(BUILTIN() ENTRY() INCLUDE() STMT() VARIABLE())
The DISPLAY option determines how the DISPLAY statement performs I/O.

**STD**

All DISPLAY statements are completed by writing the text to stdout and reading any REPLY text from stdin.

**WTO**

All DISPLAY statements without REPLY are completed via WTOs, and all DISPLAY statements with REPLY are completed via WTORs. This is the default.

The following suboptions are supported

**ROUTCDE**

Specifies one or more values to be used as the ROUTCDE in the WTO. The default ROUTCDE is 2.

**DESC**

Specifies one or more values to be used as the DESC in the WTO. The default DESC is 3.

**REPLY**

Specifies one or more values to be used as the DESC in the WTOR. If omitted, the value from the DESC option (or default) is used.

All values specified for the ROUTCDE, DESC and REPLY must between 1 and 16.

**DLLINIT**

The DLLINIT option applies OPTIONS(FETCHABLE) to all external procedures that are not MAIN. It should be used only on compilation units containing one external procedure, and then that procedure should be linked as a DLL.

**NOEXIT**

EXIT

The EXIT option enables the compiler user exit to be invoked.

NODLLINIT has no effect on your programs.

For detailed information, see “DEPRECATE” on page 31.
**inparm_string**

A string that is passed to the compiler user exit routine during initialization. The string can be up to 31 characters long.

See Chapter 22, “Using user exits,” on page 461 for more information on how to exploit this option.

**EXTRN**

The EXTRN option controls when EXTRNs are emitted for external entry constants.

- **FULL**
  EXTRNs are emitted for all declared external entry constants. This is the default.
- **SHORT**
  EXTRNs are emitted only for those constants that are referenced.

**FILEREF**

The (NO)FILEREF option controls whether the compiler produces a file reference table. The NOFILEREF option eliminates the file reference table from the listing if the compiler issues no messages for the specified FLAG setting.

The default is FILEREF.

**FLAG**

The FLAG option specifies the minimum severity of error that requires a message listed in the compiler listing.

- **F** List all messages.
- **W** List all except information messages.
- **E** List all except warning and information messages.
- **S** List only severe error and unrecoverable error messages.

If messages are below the specified severity or are filtered out by a compiler exit routine, they are not listed.
**FLOAT**

The FLOAT option controls the use of the additional floating-point registers and whether Decimal Floating Point is supported.

![Diagram of FLOAT options](image)

**NOAFP**

Compiler-generated code uses the traditional 4 floating-point registers.

**AFP**

Compiler-generated code may use all 16 floating-point registers.

**VOLATILE**

The compiler will not expect that the registers FPR8-FPR15 will be preserved by any called program. Consequently, the compiler will generate extra code to protect these registers.

Extra code will be generated even if no floating-point registers are used, and this will hurt performance. If your application makes little or no use of floating-point, it would be much better to compile with FLOAT(NOAFP) than with FLOAT(AFP(VOLATILE)).

**NOVOLATILE**

The compiler will expect that the registers FPR8-FPR15 will be preserved by any called program (as z/OS says they should be). This option will produce the most optimal code and is highly recommended if you have no CICS code. However, it must be used with care in any code included in a CICS application. For a CICS application,

- either: all code containing EXEC CICS statements must be compiled with FLOAT(AFP(VOLATILE))
- or: all code using floating-point must be compiled with FLOAT(NOAFP) or with FLOAT(AFP(VOLATILE))

It is safe and probably simplest not to use the FLOAT(AFP(NOVOLATILE)) option in any code that runs as part of a CICS application.

However, if you are using CICS TS V4.1 or later and are using either z/OS 1.11 or have the PTF for APAR PK71900 applied to z/OS 1.9 or 1.10, then these restrictions no longer apply to CICS programs.

**DFP**

The DFP facility is exploited: all DECIMAL FLOAT data will be held in the DFP format described in the z/OS Principles of Operations manual and operations using DECIMAL FLOAT will be carried using the DFP hardware instructions described therein.

The ARCH option must be 7 (or greater) or this option will be rejected.

**NODFP**

The DFP facility is not exploited.

Under FLOAT(DFP),
• the maximum precision for extended DECIMAL FLOAT will be 34 (not 33 as it is for hex float)
• the maximum precision for short DECIMAL FLOAT will be 7 (not 6 as it is for hex float)
• the values for DECIMAL FLOAT for the following built-ins will all have the appropriate changes
  – EPSILON
  – HUGE
  – MAXEXP
  – MINEXP
  – PLACES
  – RADIX
  – TINY
• the following built-in functions will all return the appropriate values for DECIMAL FLOAT (and values that will be much easier for a human to understand; for example SUCC(1D0) will be 1.000_000_000_000_001, and the ROUND function will round at a decimal place of course)
  – EXPONENT
  – PRED
  – ROUND
  – SCALE
  – SUCC
• decimal floating-point literals, when converted to "right-units-view", i.e. when the exponent has been adjusted, if needed, so that no non-zero digits follow the decimal point (for example, as would be done when viewing 3.1415E0 as 31415E-4), must have an exponent within the range of the normal numbers for the precision given by the literal. These bounds are given by the value of MINEXP-1 and MAXEXP-1. In particular, the following must hold
  – For short float, -95 <= exponent <= 90
  – For long float, -383 <= exponent <= 369
  – For extended float, -6143 <= exponent <= 6111
• when a DECIMAL FLOAT is converted to CHARACTER, the string will hold 4 digits for the exponent (as opposed to the 2 digits used for hexadecimal float)
• the IEEE and HEXADEC attributes will be accepted only if applied to FLOAT BIN, and the DEFAULT(IEEE/HEXADEC) option will apply only to FLOAT BIN.
• the mathematical built-in functions (ACOS, COS, SQRT, etc) will accept DECIMAL FLOAT arguments and use corresponding Language Environment functions to evaluate them. DECIMAL FLOAT exponentiation will be handled in a similar way.
• users of DFP need to be wary of the conversions that will arise in operations where one operand is FLOAT DECIMAL and the other is binary (i.e. FIXED BINARY, FLOAT BINARY or BIT). In such operations, the PL/I language rules dictate that the FLOAT DECIMAL operand be converted to FLOAT BINARY, and that conversion will require a library call. So, for example, for an assignment of the form \( A = A + B \); where \( A \) is FLOAT DECIMAL and \( B \) is FIXED BINARY, 3 conversions will be done and 2 of those will be library calls:
  1. \( A \) will be converted via library call from FLOAT DECIMAL to FLOAT BINARY
  2. \( B \) will be converted via inline code from FIXED BINARY to FLOAT BINARY
3. the sum $A + B$ will be converted via library call from FLOAT BINARY to FLOAT DECIMAL. 

The use of the DECIMAL built-in function would help here: if the statement were changed to $A = A + \text{DEC}(B)$, the library calls would be eliminated. The library calls could be eliminated by previously assign $B$ to a FLOAT DECIMAL temporary variable and then adding that to $A$.

- the built-in function SQRTF will not be supported for DECIMAL FLOAT arguments (as there is no hardware instruction to which it can be mapped)
- DFP is not supported by the CAST type function.

**FLOATINMATH**

The FLOATINMATH option specifies that the precision that the compiler should use when invoking the mathematical built-in functions.

```
FLOATINMATH ( { ASIS } LONG | EXTENDED )
```

**ASIS**
Arguments to the mathematical built-in functions will not be forced to have long or extended floating-point precision.

**LONG**
Any argument to a mathematical built-in function with short floating-point precision will be converted to the maximum long floating-point precision to yield a result with the same maximum long floating-point precision.

**EXTENDED**
Any argument to a mathematical built-in function with short or long floating-point precision will be converted to the maximum extended floating-point precision to yield a result with the same maximum extended floating-point precision.

A FLOAT DEC expression with precision $p$ has short floating-point precision if $p \leq 6$, long floating-point precision if $6 < p \leq 16$ and extended floating-point precision if $p > 16$.

A FLOAT BIN expression with precision $p$ has short floating-point precision if $p \leq 21$, long floating-point precision if $21 < p \leq 53$ and extended floating-point precision if $p > 53$.

The maximum extended floating-point precision depends on the platform.

**GOFF**

The GOFF option instructs the compiler to produce an object file in the Generalized Object File Format (GOFF).

```
GOFF
```

When the GOFF and OBJECT options are in effect, the compiler produces an object file in GOFF format.
When the NOGOF and OBJECT options are in effect, the compiler produces an object file in XOBJ format.

The GOFF format supersedes the S/370 Object Module and Extended Object Module formats. It removes various limitations of the previous format (for example, 16 MB section size) and provides a number of useful extensions, including native z/OS support for long names and attributes. GOFF incorporates some aspects of industry standards such as XCOFF and ELF.

When you specify the GOFF option, you must use the binder to bind the output object. You cannot use the prelinker to process GOFF objects.

The following options are not supported with the GOFF option:

- COMMON
- NOWRITABLE(PRIV)

Note: When using GOFF and source files with duplicate file names, the linker may emit an error and discard one of the code sections. In this case, turn off the CSECT option by specifying NOCSECT.

**GONUMBER**

The GONUMBER option specifies that the compiler produces additional information that allows line numbers from the source program to be included in runtime messages.

ABBREVIATIONS: GN, NGN

**SEPARATE**

Places the generated statement number table in the separate debug file if the TEST(SEPARATE) option is specified. When you use GONUMBER(SEPARATE), the statement numbers are not available for inclusion in runtime messages.

**NOSEPARATE**

Places the generated statement number table in the object deck.

Alternatively, the line numbers can be derived by using the offset address, which is always included in runtime messages, and either the table produced by the OFFSET option or the assembler listing produced by the LIST option.

GONUMBER is forced by the ALL and STMT suboptions of the TEST option.

Note that the GOSTMT option does not exist. The only option that produces information at run time identifying where an error has occurred is the GONUMBER option. When the GONUMBER option is used, the term statement in the runtime error messages refers to the line numbers as used by the NUMBER compiler option, even if the STMT option is in effect.

If the GONUMBER(SEPARATE) option is specified without TEST(SEPARATE), it is changed to GONUMBER(NOSEPARATE).
If both TEST and NOGONUMBER are specified, the NOGONUMBER option is changed to GONUMBER(NOSEPARATE).

The default is NOGONUMBER.

For compatibility, when the GONUMBER option is specified, the default suboption is SEPARATE.

**GRAPHIC**

The GRAPHIC option specifies that the source program can contain double-byte characters. The hexadecimal codes '0E' and '0F' are treated as the shift-out and shift-in control codes, respectively, wherever they appear in the source program, including occurrences in comments and string constants.

**ABBREVIATIONS: GR, NGR**

The GRAPHIC option will also cause the GRAPHIC ENVIRONMENT option to be applied to any STREAM file used in the compilation.

The GRAPHIC option must be specified if the source program uses any of the following:
- DBCS identifiers
- Graphic string constants
- Mixed-string constants
- Shift codes anywhere else in the source

**HGPR**

The HGPR option specifies that the compiler is permitted to exploit 64-bit General Purpose Registers (GPRs) in 32-bit programs targeting z/Architecture hardware.

**PRESERVE**

PRESERVE instructs the compiler to preserve the high halves of the 64-bit GPRs that a function is using, by saving them in the prolog for the function and restoring them in the epilog. The PRESERVE suboption is necessary only if the caller is not known to be Enterprise PL/I and/or the z/OS XL C/C++ compiler-generated code.

**NOPRESERVE**

NOPRESERVE lets the compiler omit preserving the high halves of the 64-bit GPRs that a function is using. Because of performance considerations, the default suboption for HGPR is NOPRESERVE.

HGPR means "High half of 64-bit GPR", which refers to the use of native 64-bit instructions. In particular, if the application use 8-byte integers, it should benefit from the native 64-bit instructions.
**IGNORE**

The IGNORE option controls whether ASSERT, DISPLAY, and PUT statements are ignored. When a statement is ignored, it is as if the statement is replaced by a semicolon.

*NOIGNORE*

*IGNORE ( ASSERT ) DISPLAY PUT*

**ASSERT**

The compiler ignores all ASSERT statements, including any function references contained in those statements.

**DISPLAY**

The compiler ignores all DISPLAY statements, including any function references contained in those statements.

**PUT**

The compiler ignores all PUT FILE statements.

**INCAFTER**

The INCAFTER option specifies a file to be included after a particular statement in your source program.

*INCAFTER ( PROCESS ( filename ) )*

*filename*

Name of the file to be included after the last PROCESS statement.

Currently, PROCESS is the only suboption and specifies the name of a file to be included after the last PROCESS statement.

Consider the following example:

`INCAFTER(PROCESS(DFTS))`

This example is equivalent to having the statement `%INCLUDE DFTS;` after the last PROCESS statement in your source.

**INCDIR**

The INCDIR compiler option specifies a directory to be added to the search path used to locate of include files.

*NOINCDIR INCDIR ( directory name )*

*directory name*

Name of the directory that should be searched for include files. You can specify the INCDIR option more than once and the directories are searched in order.
Except under batch, the compiler looks for INCLUDE files in the following order:
1. Current directory
2. Directories specified with the –I flag or with the INCDIR compiler option
3. /usr/include directory
4. PDS specified with the INCPDS compiler option

Under batch, this option is probably best used with the DFT(LOWERINC) option, and it affects only includes of the form "%include x;" For these includes, an hfs file with the name x.inc will be sought first in the directories specified in this option. If not found, x must be a member of the pds(e) specified in the syslib dd. For includes of the form "%include dd(x);" no hfs file will ever be included: the member x must always be a member of the pds named by the specified dd.

**INCLUDE**

The INCLUDE compiler option controls whether the final pass of the compiler handles %INCLUDE and %XINCLUDE statements.

```
INCLUDE NOINCLUDE
```

**INCLUDE**

Both the MACRO preprocessor and the final pass of the compiler handle %INCLUDE and %XINCLUDE statements.

**NOINCLUDE**

Only the MACRO preprocessor handles %INCLUDE and %XINCLUDE statements.

INCLUDE is the default.

**INCPDS**

The INCPDS option specifies a PDS from which the compiler will include files when compiling a program under z/OS UNIX.

**Note:** This option applies only to compilations under z/OS UNIX.

```
INCPDS('PDS name')
```

**PDS name**

Name of a PDS from which files will be included.

For example, if you wanted to compile the program TEST from a PDS named SOURCE.PLI and wanted to use the INCLUDE files from the PDS SOURCE.INC, then you could specify the following command

```
pli -c -qincpds="SOURCE.INC" "//'SOURCE.PLI(TEST)"
```

The compiler looks for INCLUDE files in the following order:
1. Current directory
2. Directories specified with the –I flag or with the INCDIR compiler option
3. /usr/include directory
4. PDS specified with the INCPDS compiler option
The INITAUTO option directs the compiler to add an INITIAL attribute to any AUTOMATIC variable declared without an INITIAL attribute.

Under INITAUTO(FULL), the compiler adds an INITIAL attribute to any AUTOMATIC variable that does not have an INITIAL attribute according to its data attributes:

- INIT( (*) 0 ) if it is FIXED or FLOAT
- INIT( (*) " ) if it is PICTURE, CHAR, BIT, GRAPHIC or WIDECHAR
- INIT( (*) SYSNULL() ) if it is POINTER or OFFSET
- INIT( (*) NULLENTRY() ) if it is ENTRY

The compiler will not add an INITIAL attribute to variables with other attributes.

INITAUTO will cause more code to be generated in the prologue for each block containing any AUTOMATIC variables that are not fully initialized (but unlike the DFT(INITFILL) option, those variables will now have meaningful initial values) and will have a negative impact on performance.

The INITAUTO option does not apply an INITIAL attribute to any variable declared with the NOINIT attribute.

Under INITAUTO(SHORT), the compiler adds an INITIAL attribute to an AUTOMATIC variable that does not have an INITIAL attribute only if it is also a scalar and has one of the following attributes:

- POINTER
- OFFSET
- FIXED BIN
- FLOAT
- NONVARYING BIT
- NONVARYING CHAR(1)
- NONVARYING WCHAR(1)

As under INITAUTO(FULL), the added INITIAL attribute will be appropriate to the data type.

If you are using the runtime STORAGE option to zero-out all storage, the optimizer may still generate undesired code for uninitialized AUTOMATIC variables and especially for those that are optimized to registers. Scalar variables with the data types listed above are those most likely to be optimized to registers, and hence using INITAUTO(SHORT) might have less of a performance impact than using DFT(INITFILL); however, the latter would leave the optimizer with no ambiguous code. Moreover, your code is correct only if all variables are explicitly initialized before use.
**INITBASED**

The INITBASED option directs the compiler to add an INITIAL attribute to any BASED variable declared without an INITIAL attribute.

This option performs the same function as INITAUTO except for BASED variables.

The INITBASED option will cause more code to be generated for any ALLOCATE of a BASED variable that is not fully initialized and will have a negative impact on performance.

The INITBASED option does not apply an INITIAL attribute to any variable declared with the NOINIT attribute.

**INITCTL**

The INITCTL option directs the compiler to add an INITIAL attribute to any CONTROLLED variable declared without an INITIAL attribute.

This option performs the same function as INITAUTO except for CONTROLLED variables.

The INITCTL option will cause more code to be generated for any ALLOCATE of a CONTROLLED variable that is not fully initialized and will have a negative impact on performance.

The INITCTL option does not apply an INITIAL attribute to any variable declared with the NOINIT attribute.

**INITSTATIC**

The INITSTATIC option directs the compiler to add an INITIAL attribute to any STATIC variable declared without an INITIAL attribute.

This option performs the same function as INITAUTO except for STATIC variables.

The INITSTATIC option could make some objects larger and some compilations longer, but should otherwise have no impact on performance.

The INITSTATIC option does not apply an INITIAL attribute to any variable declared with the NOINIT attribute.

**INSOURCE**

The INSOURCE option specifies that the compiler should include a listing of the source program before the PL/I macro, CICS, or SQL preprocessors translate it.
ABBREVIATION: IS, NIS

FULL
The INSOURCE listing will ignore %NOPRINT statements and will contain all the source before the preprocessor translates it.

FULL is the default.

SHORT
The INSOURCE listing will heed %PRINT and %NOPRINT statements.

If the INSOURCE option is in effect, and the compiler invokes more than one of the macro, CICS or SQL preprocessors, the listing contains the source for each of these preprocessors.

Under the INSOURCE option, text is included in the listing not according to the logic of the program, but as each file is read. So, for example, consider the following simple program which has a %INCLUDE statement between its PROC and END statements.

```
insource: proc options(main);
  %include member;
end;
```

The INSOURCE listing will contain all of the main program before any of the included text from the file "member" (and it would contain all of that file before any text included by it - and so on).

Under the INSOURCE(SHORT) option, text included by a %INCLUDE statement inherits the print/noprint status that was in effect when the %INCLUDE statement was executed, but that print/noprint status is restored at the end of the included text (however, in the SOURCE listing, the print/noprint status is not restored at the end of the included text).

INTERRUPT

The INTERRUPT option causes the compiled program to respond to attention requests (interrupts).

ABBREVIATION: INT, NINT

This option determines the effect of attention interrupts when the compiled PL/I program runs under an interactive system. This option will have an effect only on programs running under TSO. If you have written a program that relies on raising the ATTENTION condition, you must compile it with the INTERRUPT option. This option allows attention interrupts to become an integral part of programming. This gives you considerable interactive control of the program.
If you specify the INTERRUPT option, an established ATTENTION ON-unit gets control when an attention interrupt occurs. When the execution of an ATTENTION ON-unit is complete, control returns to the point of interrupt unless directed elsewhere by means of a GOTO statement. If you do not establish an ATTENTION ON-unit, the attention interrupt is ignored.

If you specify NOINTERRUPT, an attention interrupt during a program run does not give control to any ATTENTION ON-units.

If you require the attention interrupt capability only for testing purposes, use the TEST option instead of the INTERRUPT option. For more information see “TEST” on page 83.

See Chapter 20, “Interrupts and attention processing,” on page 455 for more information about using interrupts in your programs.

**JSON**

The JSON option lets you choose the case of the names in the JSON text generated by the JSONPUT built-in functions and expected by the JSONGET built-in functions.

```
CASE(UPPER | ASIS)
```

Under the CASE(UPPER) suboption, the names in the JSON text generated by the JSONPUT built-in functions and expected by the JSONGET built-in functions will all be in upper case.

Under the CASE(ASIS) suboption, the names in the JSON text generated by the JSONPUT built-in functions and expected by the JSONGET built-in functions will be in the case used in their declares. Note that if you use the MACRO preprocessor without using the macro preprocessor option CASE(ASIS), then the source seen by the compiler will have all the names in upper case - and that would make specifying the JSON(CASE(ASIS)) option useless.

**LANGLVL**

The LANGLVL option specifies the level of PL/I language definition that you want the compiler to accept.

```
LANGLVL(OS | NOEXT)
```

The only ENVIRONMENT options accepted are:

```
Bkwd       Genkey       Keyloc       Relative
Consecutive Graphic Organization Scalarvarying
Ctlasa     Indexed      Recsize       Vsam
Deblock    Keylength    Regional
```

0S  All ENVIRONMENT options are allowed. For a complete list of the ENVIRONMENT options, see Table 14 on page 207.
LIMITS

The LIMITS option specifies various implementation limits.

EXTNAME
Specifies the maximum length for EXTERNAL name. The maximum value for n is 100; the minimum value is 7.

FIXEDDEC
Specifies the maximum precision for FIXED DECIMAL to be either 15 or 31. The default is 15.

If FIXEDDEC(15,31) is specified, then you may declare FIXED DECIMAL variables with precision greater than 15, but unless an expression contains an operand with precision greater than 15, all arithmetic will done using 15 as the maximum precision.

FIXEDDEC(15,31) will provide much better performance than FIXEDDEC(31).

FIXEDDEC(15) and FIXEDDEC(15,15) are equivalent; similarly, FIXEDDEC(31) and FIXEDDEC(31,31) are equivalent.

FIXEDDEC(31,15) is not allowed.

FIXEDBIN
Specifies the maximum precision for SIGNED FIXED BINARY to be either 31 or 63. The default is (31,63).

If FIXEDBIN(31,63) is specified, then you may declare 8-byte integers, but unless an expression contains an 8-byte integer, all integer arithmetic will done using 4-byte integers.

Note however, that specifying the FIXEDBIN(31,63) or FIXEDBIN(63) option may cause the compiler to use 8-byte integer arithmetic for expressions mixing data types. For example, if a FIXED BIN(31) value is added to a FIXED DEC(13) value, then the compiler will produce a FIXED BIN result and under LIMITS(FIXEDBIN(31,63)) that result would have a precision greater than 31 (because the FIXED DEC precision is greater than 9). When this occurs, the compiler will issue informational message IBM2809.

FIXEDBIN(31,63) will provide much better performance than FIXEDBIN(63).
FIXEDBIN(31) and FIXEDBIN(31,31) are equivalent; similarly, FIXEDBIN(63) and FIXEDBIN(63,63) are equivalent.

FIXEDBIN(63,31) is not allowed.

The maximum precision for UNSIGNED FIXED BINARY is one greater, that is, 32 and 64.

**NAME**
Specifies the maximum length of variable names in your program. The maximum value for \( n \) is 100; the minimum value is 31.

**STRING**
Accepts these values as the threshold for the length of a BIT, CHARACTER or WIDECHAR variable: 32K, 512K, 8M, and 128M. This means that the length must be less than the threshold, and so the corresponding limits are 32767, 524287, 8388607, and 134217727.

- 32K is the default value. A larger value is accepted only if the CMPAT(V3) and BIFPREC(31) options are also specified.
- This limit applies to NONVARYING, VARYINGZ, and VARYING4, but not to VARYING. The maximum value allowed for VARYING is 32K.

**LINECOUNT**
The LINECOUNT option specifies the number of lines per page for compiler listings, including blank and heading lines.

\[
\text{LINECOUNT}(n)
\]

**ABBREVIATION:** LC

- \( n \) The number of lines in a page in the listing. The value can be from 10 to 32,767.

**LINEDIR**
This option specifies that the compiler should accept \%LINE directives.

\[
\text{LINEDIR}
\]

If the LINEDIR option is specified, the compiler will reject all \%INCLUDE statements. If the LINEDIR option is specified, the compiler will also reject the use of the SEPARATE suboption of the TEST option.

**LIST**
The LIST option specifies that the compiler should produce a pseudo-assembler listing.

\[
\text{LIST}
\]
Specifying the LIST option will increase the time and region required for a compilation. The OFFSET and MAP options may provide the information you need at much less cost.

The pseudo-assembler listing will also include at the end of the listing for each block the offset of the first instruction in that block from the start of the whole compilation unit.

**LISTVIEW**

The LISTVIEW option specifies whether the compiler should show the source in the source listing or whether it should show the source after it has been processed by one or more of the preprocessors.

```
LISTVIEW ( SOURCE - AFTERALL - AFTERCICS - AFTERMACRO - AFTERSQL )
```

The LISTVIEW option is ignored if the NOSOURCE option is in effect.

**SOURCE**

Causes the source listing to show the unadulterated source and, more importantly perhaps, it will cause Debug Tool to bring up this as the source view.

**AFTERALL**

Causes the source listing to show the source as if it came from the MDECK from the last invocation, if any, of the last preprocessor, and, more importantly perhaps, it will cause Debug Tool to bring up this as the source view if the SEPARATE suboption of the TEST compiler option is also specified.

AALL may be used as an abbreviation for AFTERALL.

**AFTERCICS**

Causes the source listing to show the source as if it came from the MDECK from the last invocation, if any, of the CICS preprocessor, and, more importantly perhaps, it will cause Debug Tool to bring up this as the source view if the SEPARATE suboption of the TEST compiler option is also specified.

ACICS may be used as an abbreviation for AFTERCICS.

**AFTERMACRO**

Causes the source listing to show the source as if it came from the MDECK from the last invocation, if any, of the MACRO preprocessor, and, more importantly perhaps, it will cause Debug Tool to bring up this as the source view if the SEPARATE suboption of the TEST compiler option is also specified.

AMACRO may be used as an abbreviation for AFTERMACRO.

**AFTERSQL**

Causes the source listing to show the source as if it came from the MDECK from the last invocation, if any, of the SQL preprocessor, and, more importantly perhaps, it will cause Debug Tool to bring up this as the source view if the SEPARATE suboption of the TEST compiler option is also specified.

ASQL may be used as an abbreviation for AFTERSQL.
If the TEST option is specified and a suboption other than SOURCE is specified for LISTVIEW then the SEPARATE suboption must also be specified for the TEST option.

As an example of the differing effects of the AFTERMACRO, AFTERSQL and AFTERALL suboptions suppose the PP option was PP( MACRO('INCONLY'), SQL, MACRO). Then:

- Under LISTVIEW(AFTERMACRO), the "source" in the listing and in the Debug Tool source window if TEST(SEP) were specified would appear as if it came from the MDECK that the second invocation of the MACRO preprocessor would have produced.
- Under LISTVIEW(AFTERSQL), the "source" in the listing and in the Debug Tool source window if TEST(SEP) were specified would appear as if it came from the MDECK that the invocation of the SQL preprocessor would have produced (and hence %DCL and other macro statements would still be visible).
- Under LISTVIEW(AFTERALL), the "source" would be as under the LISTVIEW(AFTERMACRO) option since the MACRO preprocessor is the last in the PP option.

MACRO

The MACRO option invokes the MACRO preprocessor.

ABBREVIATIONS: M, NM

You may also invoke the MACRO preprocessor via the PP(MACRO) option. However, the use of both the MACRO option and the PP(MACRO) option in the same compilation is not recommended.

For more discussion of the PP option, see "PP" on page 60.

For more discussion of the MACRO preprocessor, see "Macro preprocessor" on page 110.

MAP

The MAP option specifies that the compiler produces additional information that can be used to locate static and automatic variables in dumps.

MARGINI

The MARGINI option specifies a character that the compiler will place in the column preceding the left-hand margin, and also in the column following the right-hand margin, of the listings produced by the INSOURCE and SOURCE options.
ABBREVIATIONS: MI, NMI

c The character to be printed as the margin indicator.

Note: NOMARGINI is equivalent to MARGINI(‘ ’).

MARGINS

The MARGINS option specifies which part of each compiler input record contains PL/I statements, and the position of the ANS control character that formats the listing, if the SOURCE and/or INSOURCE options apply. The compiler does not process data that is outside these limits, but it does include it in the source listings.

The PL/I source is extracted from the source input records so that the first data byte of a record immediately follows the last data byte of the previous record. For variable records, you must ensure that when you need a blank you explicitly insert it between margins of the records.

![MARGINS Diagram]

ABBREVIATION: MAR

m The column number of the leftmost character (first data byte) that is processed by the compiler. It must not exceed 100.

n The column number of the rightmost character (last data byte) that is processed by the compiler. It should be greater than \( m \), but must not exceed 200, except under MVS batch where it must not exceed 100.

Variable-length records are effectively padded with blanks to give them the maximum record length.

c The column number of the ANS printer control character. It must not exceed 200, except under MVS batch where it must not exceed 100, and it should be outside the values specified for \( m \) and \( n \). A value of 0 for \( c \) indicates that no ANS control character is present. Only the following control characters can be used:

- (blank) Skip one line before printing
- 0 Skip two lines before printing
- - Skip three lines before printing
- + No skip before printing
- 1 Start new page

Any other character is an error and is replaced by a blank.

Do not use a value of \( c \) that is greater than the maximum length of a source record, because this causes the format of the listing to be unpredictable. To avoid this problem, put the carriage control characters to the left of the source margins for variable-length records.

Specifying MARGINS(\( m,c \)) is an alternative to using \%PAGE and \%SKIP statements (described in PL/I Language Reference).
The IBM-supplied default for fixed-length records is MARGINS(2,72). For variable-length and undefined-length records, the IBM-supplied default is MARGINS(10,100). This specifies that there is no printer control character.

Use the MARGINS option to override the default for the primary input in a program. The secondary input must have the same margins as the primary input.

The NOMARGINS option will suppress any previously encountered instance of the MARGINS option. The purpose of this option is allow your installation to have a default set of compile-time options which use a MARGINS option tailored for their fixed format source preferences while retaining the ability to use variable source format files.

You would usually specify the NOMARGINS option, if you use it at all, as part of the parameter-string passed to the compiler. The compiler will ignore NOMARGINS if it finds the option in a %PROCESS statement.

**MAXBRANCH**

The MAXBRANCH option flags blocks that have too many branches. Branches include all conditional jumps and each WHEN in a SELECT statement that can be turned into a branch table.

```
max
```

The limit that measures the cyclomatic or conditional complexity of the block. The default is 2000.

A statement of the form "if a then ...; else ..." adds 1 to the total number of branches in its containing block, and a statement of the form "if a = 0 | b = 0 then ..." adds 2.

**MAXGEN**

The MAXGEN option specifies the maximum number of intermediate language statements that should be generated for any user statement and will cause the compiler to flag any statement where this maximum is exceeded.

The number of intermediate language statements generated for any user statement may vary depending on the compiler release, the compiler maintenance level, and the compiler options in effect. This option is intended to be used only to help find statements for which excessive amounts of code are generated and which thus may indicate that they are perhaps are poorly coded.

However, it should be noted that using a preprocessor may cause the number of intermediate language statements generated for some statements to be very large. In such a situation, it may be better either to set the MAXGEN threshold to be larger or to use the LISTVIEW(AFTERALL) option.

The default is MAXGEN(100000).
MAXMEM

When compiling with OPTIMIZE, the MAXMEM option limits the amount of memory used for local tables of specific, memory intensive optimizations to the specified number of kilobytes. The minimum number of kilobytes that may be specified is 1. The maximum number of kilobytes that may be specified is 2097152, and the default is 1048576.

If you specify the maximum value of 2097152, the compiler will assume that unlimited memory is available. If you specify any smaller value for MAXMEM, the compiler, especially when the OPT(2) option is in effect, may issue a message saying that optimization is inhibited and that you should try using a larger value for MAXMEM.

Use the MAXMEM option if you know that less (or more) memory is available than implied by the default value.

If the memory specified by the MAXMEM option is insufficient for a particular optimization, the compilation is completed in such a way that the quality of the optimization is reduced, and a warning message is issued.

```
MAXMEM (size)
```

ABBREVIATIONS: MAXM

When a large size is specified for MAXMEM, compilation may be aborted because of insufficient virtual storage, depending on the source file being compiled, the size of the subprogram in the source, and the virtual storage available for the compilation.

The advantage of using the MAXMEM option is that, for large and complex applications, the compiler produces a slightly less-optimized object module and generates a warning message, instead of terminating the compilation with an error message of "insufficient virtual storage”.

MAXMSG

The MAXMSG option specifies the maximum number of messages with a given severity (or higher) that the compilation should produce.

```
MAXMSG
```

I  Count all messages.
W  Count all except information messages.
E  Count all except warning and information messages.
S  Count only severe error and unrecoverable error messages.
n Terminate the compilation if the number of messages exceeds this value. If messages are below the specified severity or are filtered out by a compiler exit routine, they are not counted in the number. The value of n can range from 0 to 32767. If you specify 0, the compilation terminates when the first error of the specified severity is encountered.

**MAXNEST**

The MAXNEST option specifies the maximum nesting of various kinds of statements that should be allowed before the compiler flags your program as too complex.

```
MAXNEST    ( BLOCK(17) DO(17) IF(17) )
```

**BLOCK**

Specifies the maximum nesting of BEGIN and PROCEDURE statements.

**DO**

Specifies the maximum nesting of DO statements.

**IF**

Specifies the maximum nesting of IF statements.

The value of any nesting limit must be between 1 and 50 (inclusive).

The default is MAXNEST( BLOCK(17) DO(17) IF(17) ).

**MAXSTMT**

Under the MAXSTMT option, optimization will be turned off for any block that has more than the specified number of statements. Use the MAXSTMT option - with a reasonable limit to the number of statements - if you want the compiler to optimize the code generated for a program and are willing for the compiler to optimize only the reasonably sized blocks in that program.

```
MAXSTMT    (size)
```

When a large size is specified for MAXSTMT and some blocks have a large number of statements, compilation may be aborted if there is not enough virtual storage available.

The default for MAXSTMT is 4096.

**MAXTEMP**

The MAXTEMP option determines when the compiler flags statements using an excessive amount of storage for compiler-generated temporaries.

```
MAXTEMP    (max)
```

The limit for the number of bytes that can be used for compiler-generated
temporaries. The compiler flags any statement that uses more bytes than those specified by $max$. The default for $max$ is 50000.

You should examine statements that are flagged under this option - if you code them differently, you may be able to reduce the amount of stack storage required by your code.

**MDECK**

The MDECK option specifies that the preprocessor produces a copy of its output either on the file defined by the SYSPUNCH DD statement under z/OS, or on the .dek file under z/OS UNIX.

**ABBREVIATIONS: MD, NMD**

The MDECK option allows you to retain the output from the preprocessor as a file of 80-column records. This option is applicable only when the MACRO option is in effect.

**MSGSUMMARY**

The MSGSUMMARY option determines whether the compiler adds a summary of all messages that are issued during the compilation into the listing.

**MSGSUMMARY(NOXREF)**

The compiler adds a message summary to the listing. The summary is after the file reference table in the listing. It is sorted by compiler component and within each component by severity and then by message number.

The summary includes the following information:
- One instance of each message that is produced in the compilation
- The number of times that each message is produced

**MSGSUMMARY(XREF)**

The compiler adds a message summary to the listing. The summary is the same with the one added when MSGSUMMARY(XREF) is specified with one difference: after each message the summary lists all the line or statement numbers where the message is issued.

**NOMSGSUMMARY**

No message summary is produced.

NOMSGSUMMARY is the default. When MSGSUMMARY is specified, MSGSUMMARY(NOXREF) is the default.

For the compiler listing example with a message summary generated by using MSGSUMMARY, see [Figure 4 on page 105](#)
NAME

The NAME option specifies that the TEXT file created by the compiler will contain a NAME record.

ABBREVIATIONS: N

If no 'name' is specified as a suboption of the NAME option, then the 'name' used is determined as follows

- if there is a PACKAGE statement, the leftmost name on it is used
- otherwise, the leftmost name on the first PROCEDURE statement is used

The length of the 'name' must not be greater than 8 characters if the LIMITS(EXTNAME(n)) option is used with \( n \leq 8 \).

NAMES

The NAMES option specifies the extralingual characters that are allowed in identifiers. Extralingual characters are those characters other than the 26 alphabetic, 10 digit, and special characters defined in PL/I Language Reference.

extralingual_char

An extralingual character

upp_extraling_char

The extralingual character that you want interpreted as the uppercase version of the corresponding character in the first suboption.

If you omit the second suboption, PL/I uses the character specified in the first suboption as both the lowercase and the uppercase values. If you specify the second suboption, you must specify the same number of characters as you specify in the first suboption.

The default is NAMES('#@$' '#@$').

NATLANG

The NATLANG option specifies the "language" for compiler messages, headers, etc.

Chapter 1. Using compiler options and facilities  55
ENU
All compiler messages, headers etc will be in mixedcase English.

UEN
All compiler messages, headers etc will be in uppercase English.

NEST
The NEST option specifies that the listing resulting from the SOURCE option indicates the block level and the do-group level for each statement.

NOT
The NOT option specifies up to seven alternate symbols that can be used as the logical NOT operator.

char
A single SBCS character.

You cannot specify any of the alphabetic characters, digits, and special characters defined in PL/I Language Reference, except for the standard logical NOT symbol (~). You must specify at least one valid character.

When you specify the NOT option, the standard NOT symbol is no longer recognized unless you specify it as one of the characters in the character string.

For example, NOT(~) means that the tilde character, 'A1'X, will be recognized as the logical NOT operator, and the standard NOT symbol, '~', '5F'X, will not be recognized. Similarly, NOT('~') means that either the tilde or the standard NOT symbol will be recognized as the logical NOT operator.

The IBM-supplied default code point for the NOT symbol is '5F'X. The logical NOT sign might appear as a logical NOT symbol (~) or a caret symbol (^) on your keyboard.

NUMBER
The number option specifies that statements in the source program are to be identified by the line and file number of the file from which they derived and that this pair of numbers is used to identify statements in the compiler listings resulting from the AGGREGATE, ATTRIBUTES, LIST, MAP, OFFSET, SOURCE and XREF options. The File Reference Table at the end of the listing shows the number assigned to each of the input files read during the compilation.

If a preprocessor has been used, more than one line in the source listing may be identified by the same line and file numbers. For example, almost every EXEC
CICS statement generates several lines of code in the source listing, but these would all be identified by one line and file number.

In the pseudo-assembler listing produced by the LIST option, the file number is left blank for the first file.

NUMBER and STMT are mutually exclusive and specifying one will negate the other.

The default is NUMBER.

**OBJECT**

The OBJECT option specifies that the compiler creates an object module. Under batch z/OS, the compiler stores the object in the data set defined by the SYSLIN DD, while under z/OS UNIX, the compiler creates a .o file.

---

**ABBREVIATIONS: OBJ, NOBJ**

Under the NOOBJECT option, the compiler does not create an object module. However, under the NOOBJECT option, the compiler will complete all of its syntactic and semantic analysis phases as well as all of its detection of uninitialized variables, and hence it could produce more messages than under the NOCOMPILE, the NOSEMANTIC or the NOSYNTAX options.

Under the NOOBJECT option, the LIST, MAP, OFFSET and STORAGE options will be ignored.

**OFFSET**

The OFFSET option specifies that the compiler is to print a table of line numbers for each procedure and BEGIN block with their offset addresses relative to the primary entry point of the procedure. This table can be used to identify a statement from a run-time error message if the GONUMBER option is not used.

---

**ONSNAP**

For a PROCEDURE with either the OPTIONS(MAIN) or the OPTIONS(FROMALIEN) attribute, the ONSNAP option specifies that the compiler should insert an ON STRINGRANGE SNAP; or ON STRINGSIZE SNAP; statement (or both) into the prologue code for that PROCEDURE. This can make it easier to
determine the calling chain if the corresponding conditions are raised in other
routines called from such a PROCEDURE.

The ONSNAP option has no affect on a PROCEDURE without either of these
attributes.

OPTIMIZE

The OPTIMIZE option specifies the type of optimization required:

ABBREVIATIONS: OPT, NOPT

OPTIMIZE(0)
Specifies fast compilation speed, but inhibits optimization.

OPTIMIZE(2)
Optimizes the machine instructions generated to produce a more efficient
object program. This type of optimization can also reduce the amount of main
storage required for the object module.

OPTIMIZE(3)
Performs all the optimizations done under OPTIMIZE(2) plus some additional
optimizations. Under OPTIMIZE(3), the compiler will generally, but especially
for programs with large blocks and many variables, generate smaller and more
efficient object code. However, it may also take considerably more time and
region to complete compiles under OPTIMIZE(3) than under OPTIMIZE(2).

It is strongly recommended that the DFT(REORDER) option be used with the
OPTIMIZE option. In fact, the effect of OPTIMIZE is severely limited for any
PROCEDURE or BEGIN-block for which all of the following are true:
• The ORDER option applies to the block.
• The block contains ON-units for hardware-detected conditions (such as
  ZERODIVIDE).
• The block has labels that are the (potential) target of branches out of those
  ON-units.

The use of OPTIMIZE(2) could result in a substantial increase in compile time over
NOOPTIMIZE and a substantial increase in the space required. For example,
compiling a large program at OPTIMIZE(2) might take several minutes and could
require a region of 100M or more.

The use of OPTIMIZE(3) will increase the time and region needed for a compile
over what is needed under OPTIMIZE(2). For large programs, the time to compile
a program under OPTIMIZE(3) can be more than twice as long as under
OPTIMIZE(2).

During optimization the compiler can move code to increase runtime efficiency. As
a result, statement numbers in the program listing might not correspond to the
statement numbers used in runtime messages.
NOOPTIMIZE is the equivalent of OPTIMIZE(0).

OPTIMIZE(TIME) is the equivalent of OPTIMIZE(2).

Note that the use of OPTIMIZE(2) or OPTIMIZE(3) severely limits the functionality of the TEST option. For example:

- if the HOOK suboption of TEST is in effect, then only block hooks will be generated
- if the NOHOOK suboption of TEST is in effect, attempts to list or change a variable may fail (because the variable may have been optimized into a register) and attempts to stop at a particular statement may cause the debugger to stop several times (because the statement may have split up into several parts)

The use of the PREFIX option with one or more of the checkout conditions (SIZE, STRINGRANGE, STRINGSIZE and SUBSCRIPTRANGE) can significantly increase the time and space needed for a compile.

For more information on choosing the best options to improve the performance of your code, see Chapter 13, “Improving performance,” on page 303.

OPTIONS

The OPTIONS option specifies that the compiler includes a list showing the compiler options to be used during this compilation in the compiler listing.

ABBREVIATIONS: OP, NOP

This list includes all options applied by default, those specified in the PARM parameter of an EXEC statement or in the invoking command (pli), those specified in a %PROCESS statement, those specified in the IBM_OPTIONS environment variable under z/OS, and all those incorporated from any options file.

Under OPTIONS(DOC), the OPTIONS listing will include only those options (and suboptions) documented in this document at the time of the compiler’s release.

Under OPTIONS(ALL), the OPTIONS listing will also include any option added by PTF after the compiler’s release.

OR

The OR option specifies up to seven alternate symbols as the logical OR operator. These symbols are also used as the concatenation operator, which is defined as two consecutive logical OR symbols.

Note: Do not code any blanks between the quotes.
The IBM-supplied default code point for the OR symbol (|) is '4F'X.

char
A single SBCS character.

You cannot specify any of the alphabetic characters, digits, and special characters defined in the PL/I Language Reference, except for the standard logical OR symbol (|). You must specify at least one valid character.

If you specify the OR option, the standard OR symbol is no longer recognized unless you specify it as one of the characters in the character string.

For example, OR(\') means that the backslash character, 'E0'X, will be recognized as the logical OR operator, and two consecutive backslashes will be recognized as the concatenation operator. The standard OR symbol, '4F'X, will not be recognized as either operator. Similarly, OR(\|) means that either the backslash or the standard OR symbol will be recognized as the logical OR operator, and either symbol or both symbols can be used to form the concatenation operator.

PP

The PP option specifies which (and in what order) preprocessors are invoked before compilation.

\[PP\] \[pp-name\] \[\text{-}pp-string\]

pp-name
The name given to a particular preprocessor. CICS, INCLUDE, MACRO and SQL are the only preprocessors currently supported. Using an undefined name causes a diagnostic error.

pp-string
A string, delimited by quotation marks, of up to 100 characters representing the options for the corresponding preprocessor. For example, PP(MACRO('CASE(ASIS)')) invokes the MACRO preprocessor with the option CASE(ASIS).

Preprocessor options are processed from left to right, and if two options conflict, the last (rightmost) option is used. For example, if you invoke the MACRO preprocessor with the option string 'CASE(ASIS) CASE(UPPER)', then the option CASE(UPPER) is used.

You can specify a maximum of 31 preprocessor steps, and you can specify the same preprocessor more than once with the exception of the CICS and SQL preprocessors. The CICS preprocessor must be invoked at most once, and the SQL preprocessor must be invoked no more than twice. The SQL preprocessors can be invoked twice only if the first specification specifies INONLY as its option.

The MACRO option and the PP(MACRO) option both cause the macro facility to be invoked before compilation. If both MACRO and PP(MACRO) are specified, the macro facility is invoked twice. However, the use of both the MACRO option and the PP(MACRO) option in the same compilation is not recommended.
If you specify the PP option more than once, the compiler effectively concatenates them. So specifying PP(SQL) PP(CICS) is the same as specifying PP(SQL CICS). This also means that if you specified PP(MACRO SQL('CCSID0')) and PP(MACRO SQL('CCSID0 DATE(ISO)')), then the resulting PP option would be PP(MACRO SQL('CCSID0') MACRO SQL('CCSID0 DATE(ISO)')), both the MACRO and SQL preprocessor would be invoked twice, and the second invocation of the SQL preprocessor would be in error. If you were doing this in an attempt to override the earlier SQL options, it might be better not to specify the preprocessor options in the PP option, but rather to specify them through the PPSQL option, that is, specify PP(MACRO SQL) PPSQL('CCSID0 DATE(ISO)').

For more discussion of the preprocessors, see Chapter 2, “PL/I preprocessors,” on page 109.

**PPCICS**

The PPCICS option specifies options to be passed to the CICS preprocessor if it is invoked.

```plaintext
-NOPPCICS
PPCICS-("options string")
```

So, specifying PPCICS('EDF') PP(CICS) is the same as specifying PP(CICS('EDF')).

This option has no effect unless the PP(CICS) option is specified. However, if you want to specify a set of CICS preprocessor options that should be used if and when the CICS preprocessor is invoked, you could specify this option in the installation options exit. Then whenever you specified PP(CICS), the set of options specified in the PPCICS option would be used.

Also, any options specified when the preprocessor is invoked overrule those specified in the PPCICS option. So specifying PPCICS('EDF') PP(CICS('NOEDF')) is the same as specifying PP(CICS('EDF NOEDF')) or the even simpler PP(CICS('NOEDF')).

The options string is limited to 1000 characters in length. However, if the string is longer than 100 characters, it will not be shown in the options listing.

**PPINCLUDE**

The PPINCLUDE option specifies options to be passed to the INCLUDE preprocessor if it is invoked.

```plaintext
-NOPPINCLUDE
PPINCLUDE-("options string")
```

So, specifying PPINCLUDE('ID(-inc)') PP(INCLUDE) is the same as specifying PP(INCLUDE('ID(-inc)')).

This option has no effect unless the PP(INCLUDE) option is specified. However, if you want to specify a set of INCLUDE preprocessor options that should be used if and when the INCLUDE preprocessor is invoked, you could specify this option in the installation options exit. Then whenever you specified PP(INCLUDE), the set of options specified in the PPINCLUDE option would be used.
Also, any options specified when the preprocessor is invoked overrule those specified in the PPINCLUDE option. So specifying PPINCLUDE('ID(-inc)') PP(INCLUDE('ID(+include)')) is the same as specifying PP(INCLUDE('ID(-inc) ID(+include)')) or the even simpler PP(INCLUDE('ID(+include)')).

The options string is limited to 1000 characters in length. However, if the string is longer than 100 characters, it will not be shown in the options listing.

**PPLIST**

The PPLIST option controls whether the compiler keeps or erases the part of the listing that is generated by each preprocessor phase.

When you specify PPLIST(KEEP), the compiler keeps the part of the listing that is generated by each preprocessor phase.

When you specify PPLIST(ERASE), the compiler erases the part of the listing that is generated by any preprocessor phase that produces no messages.

The compiler does not count messages that are suppressed by the EXIT and FLAG options. Therefore, specifying both FLAG(W) and PPLIST(ERASE) causes the compiler to suppress all output from any preprocessor that produces no warning, error or severe messages.

PPLIST(KEEP) is the default.

**PPMACRO**

The PPMACRO option specifies options to be passed to the MACRO preprocessor if it is invoked.

So, specifying PPMACRO('CASE(ASIS)') PP(MACRO) is the same as specifying PP(MACRO('CASE(ASIS)')).

This option has no effect unless the PP(MACRO) option is specified. However, if you want to specify a set of MACRO preprocessor options that should be used if and when the MACRO preprocessor is invoked, you could specify this option in the installation options exit. Then whenever you specified the MACRO or PP(MACRO) options, the set of options specified in the PPMACRO option would be used.

Also, any options specified when the preprocessor is invoked overrule those specified in the PPMACRO option. So specifying PPMACRO('CASE(ASIS)') PP(MACRO('CASE(UPPER)')) is the same as specifying PP(MACRO('CASE(ASIS) CASE(UPPER)')) or the even simpler PP(MACRO('CASE(UPPER)')).

The options string is limited to 1000 characters in length. However, if the string is longer than 100 characters, it will not be shown in the options listing.
PPSQL
The PPSQL option specifies options to be passed to the SQL preprocessor.

PPSQL('options string')

Specifying PPSQL('APOSTSQL') PP(SQL) is the same as specifying PP(SQL('APOSTSQL'))

This option has no effect unless the PP(SQL) option is specified. However, if you want to specify a set of options that should be used if the SQL preprocessor is invoked, you can specify this option in the installation options exit. Then, whenever you specify PP(SQL), the set of options in the PPSQL option are used.

Also, any options that are specified when the preprocessor is invoked overrule those that are specified in the PPSQL option. Therefore, specifying PPSQL('APOSTSQL') PP(SQL('QUOTESQL')) is the same as specifying PP(SQL('APOSTSQL QUOTESQL')) or the even simpler PP(SQL('QUOTESQL')).

The options string is limited to 1000 characters in length. However, if the string is longer than 100 characters, it will not be shown in the options listing.

PPTRACE
The PPTRACE option specifies that, when a deck file is written for a preprocessor, every nonblank line in that file is preceded by a line containing a %LINE directive. The directive indicates the original source file and line to which the nonblank line should be attributed.

PRECTYPE
The PRECTYPE option determines how the compiler derives the attributes for the MULTIPLY, DIVIDE, ADD and SUBTRACT built-in functions when the operands are FIXED and at least one is FIXED BIN.

ANS
under PRECTYPE(ANS), the value p in BIF(x, y, p) and in BIF( x, y, p, 0) is interpreted as specifying a binary number of digits, the operation is performed as a binary operation and the result has the attributes FIXED BIN(p,0).

However, for BIF( x, y, p, q) if q is not-zero, then the operation will be performed as a decimal operation and the result will have the attributes FIXED DEC(t,u) where t and u are the decimal equivalents of p and q, namely t = 1 + ceil(p / 3.32) and u = ceil(q / 3.32). In this case, x, y, p and q are effectively all converted to decimal (in contrast to the next suboption which converts only x and y to decimal and does so even if q is zero). The compiler will issue the Informational message 1BM1053 in this situation.
DECDIGIT
under PRECTYPE(DECDIGIT), the value p in BIF(x, y, p) and BIF(x, y, p, 0) is interpreted as specifying a decimal number of digits, the operation is performed as a binary operation, and the result has the attributes FIXED BIN(s) where s is the corresponding binary equivalent to p (namely s = ceil(3.32*p)). For an instance of BIF(x, y, p, q) where q is not-zero, the results under PRECTYPE(DECDIGIT) are the same as described below under PRECTYPE(DECRESULT).

DECRESULT
under PRECTYPE(DECRESULT), the value p in BIF(x, y, p) and the values p and q in BIF(x, y, p, q) are interpreted as specifying a decimal number of digits, the operation is performed as a decimal operation and the result has the attributes FIXED DEC(p,0) or FIXED DEC(p,q) respectively. The result is the same as would be produced if the DECIMAL built-in were applied to x and y.

PREFIX
The PREFIX option enables or disables the specified PL/I conditions in the compilation unit being compiled without your having to change the source program. The specified condition prefixes are logically prefixed to the beginning of the first PACKAGE or PROCEDURE statement.

condition
Any condition that can be enabled/disabled in a PL/I program, as explained in PL/I Language Reference.

The use of the PREFIX option with one or more of the checkout conditions (SIZE, STRINGRANGE, STRINGSIZE and SUBSCRIPTRANGE) can significantly increase the time and space needed for a compile.

Default PREFIX(CONVERSION FIXEDOVERFLOW INVALIDOP OVERFLOW NOSIZE NOSTRINGRANGE NOSTRINGSIZE NOSUBSCRIPTRANGE UNDERFLOW ZERODIVIDE)

PROCEED
The PROCEED option stops the compiler after processing by a preprocessor is completed depending on the severity of messages issued by previous preprocessors.

ABBREVIATIONS: PRO, NPRO

PROCEED
Is equivalent to NOPROCEED(S).
**NOPROCEED**

Ends the processing after the preprocessor has finished compiling.

**NOPROCEED(S)**

The invocation of preprocessors and the compiler does not continue if a severe or unrecoverable error is detected in this stage of preprocessing.

**NOPROCEED(E)**

The invocation of preprocessors and the compiler does not continue if an error, severe error, or unrecoverable error is detected in this stage of preprocessing.

**NOPROCEED(W)**

The invocation of preprocessors and the compiler does not continue if a warning, error, severe error, or unrecoverable error is detected in this stage of preprocessing.

**PROCESS**

The PROCESS option determines if *PROCESS statements are allowed and, if they are allowed, if they are written to the MDECK file.

Under the NOPROCESS option, the compiler will flag any *PROCESS statement with an E-level message.

Under the PROCESS(KEEP) option, the compiler will not flag *PROCESS statements, and the compiler will retain any *PROCESS statements in the MDECK output.

Under the PROCESS(DELETE) option, the compiler will not flag *PROCESS statements, but the compiler will not retain any *PROCESS statements in the MDECK output.

**QUOTE**

The QUOTE option specifies an alternate symbol that can be used as the quote character.

Note: Do not code any blanks between the quotes.

The IBM-supplied default code point for the QUOTE symbol is "".

char

A single SBCS character.

You cannot specify any of the alphabetic characters, digits, and special characters defined in the PL/I Language Reference, except for the standard QUOTE symbol (").

You must specify a valid character.
The QUOTE option is ignored if the GRAPHIC option is also specified.

**REDUCE**

The REDUCE option specifies that the compiler is permitted to reduce an assignment of a null string to a structure into simpler operations - even if that means padding bytes might be overwritten.

The REDUCE option also allows the compiler to reduce the assignment of matching structures into a simple aggregate move - even if the structures contain POINTER fields.

The NOREDUCE option specifies that the compiler must decompose an assignment of a null string to a structure into a series of assignments of the null string to the base members of the structure.

Under the NOREDUCE option, BY NAME assignments that can be reduced to aggregate moves are not reduced if the elements that would be moved together had the AREA or VARYING(Z) attributes.

The REDUCE option causes fewer lines of code to be generated for an assignment of a null string to a structure, and that usually means your compilation is quicker and your code runs much faster. However, padding bytes may be zeroed out.

For instance, in the following structure, there is one byte of padding between field12 and field13.

dcl sample ext,
  5 field10 bin fixed(31),
  5 field11 bin fixed(15),
  5 field12 bit(8),
  5 field13 bin fixed(31);

Now consider the assignment sample = ";"

Under the NOREDUCE option, it will cause four assignments to be generated, and the padding byte will be unchanged.

However, under REDUCE, the assignment would be reduced to one operation, but the padding byte will be zeroed out.

The NOREDUCE option makes the compiler act more like the OS PL/1 and the PL/1 for MVS compilers. These compilers would reduce an assignment of matching structures into a simple aggregate move unless the structures contain POINTER fields. The NOREDUCE option will make this compiler act the same way.

**RENT**
Your code is “naturally reentrant” if it does not alter any of its static variables.

The RENT option specifies that the compiler is to take code that is not naturally reentrant and make it reentrant. For a detailed description of reentrancy, see the z/OS Language Environment Programming Guide. If you use the RENT option, the Linkage Editor cannot directly process the object module that is produced: you must use either the prelinker or PDSE’s.

The NORENT option specifies that the compiler is not to specifically generate reentrant code from non-reentrant code. Any naturally reentrant code remains reentrant.

If you link a module (either MAIN or FETCHABLE) containing one or more programs compiled with the RENT option, you must specify DYNAM=DLL and REUS=RENT on the link step.

If you specify the options NORENT and LIMITS(EXITNAME(n)) (with \( n \leq 7 \)), then the text decks generated by the compiler will have the same format as those generated by the older PL/I compilers. This means that the prelinker would not be needed to create a PDS-style load module. If you use any other options, you must use either the prelinker or PDSE’s.

The code generated under the NORENT option may not be reentrant unless the NOWRITABLE option is also specified.

The use of the NORENT does preclude the use of some features of the compiler. In particular:
- DLLs cannot be built
- reentrant, writeable static is not supported
- a STATIC ENTRY VARIABLE cannot have an INITIAL value

You may mix RENT and NORENT code subject to the following restrictions:
- code compiled with RENT cannot be mixed with code compiled with NORENT if they share any EXTERNAL STATIC variables
- code compiled with RENT cannot call an ENTRY VARIABLE set in code compiled with NORENT
- code compiled with RENT cannot call an ENTRY CONSTANT that was fetched in code compiled with NORENT
- code compiled with RENT can fetch a module containing code compiled with NORENT if one of the following is true
  - all the code in the fetched module was compiled with NORENT
  - the code containing the entry point to the module was compiled with RENT
- code compiled with NORENT code cannot fetch a module containing any code compiled with RENT
- code compiled with NORENT WRITABLE cannot be mixed with code compiled with NORENT NOWRITABLE if they share any external CONTROLLED variables or any external FILEs

Given the above restrictions, the following is still valid:
- a NORENT routine, called say mnorent, statically links and calls a RENT routine, called say mrent
- the RENT routine mrent then fetches and CALLs a separately-linked module with an entry point compiled with RENT
RESEXP

The RESEXP option specifies that the compiler is permitted to evaluate all restricted expressions at compile time even if this would cause a condition to be raised and the compilation to end with S-level messages.

Under the NORESEXP compiler option, the compiler will still evaluate all restricted expression occurring in declarations, including those in INITIAL value clauses.

For example, under the NORESEXP option, the compiler would not flag the following statement (and the ZERODIVIDE exception would be raised at run time)

```pli
if preconditions_not_met then
  x = 1 / 0;
```

RESPECT

The RESPECT option causes the compiler to honor any specification of the DATE attribute and to apply the DATE attribute to the result of the DATE built-in function.

Using the default, RESPECT(), causes the compiler to ignore any specification of the DATE attribute and ensures that the compiler does not apply the DATE attribute to the result of the DATE built-in function.

RTCHECK

The RTCHECK option specifies that extra code is generated to force the ERROR condition to be raised if a null pointer is dereferenced, that is, if the pointer is used to change or obtain the value of a variable.

NULLPTR

Extra code is generated to force the ERROR condition to be raised if a null pointer, that is, a pointer equal to SYSNULL(), is dereferenced. ARCH(8) or greater must be specified to use NULLPTR.

NULL370

Extra code is generated to force the ERROR condition to be raised if a pointer equal to the old NULL() value is dereferenced. The old NULL() value is the hexadecimal value of 'FF000000'x. ARCH(8) or greater must be specified to use NULL370.

NONULLPTR

No extra code is generated to force the ERROR condition to be raised if a null pointer is dereferenced.
RTCHECK(NONULLPTR) is the default.

**Note:** When a null pointer is dereferenced, a compare-and-trap data exception occurs.

**RULES**

The RULES option allows or disallows certain language capabilities and lets you choose semantics when alternatives are available. It can help you diagnose common programming errors.
IBM | ANS

Under the IBM suboption:

- For operations requiring string data, data with the BINARY attribute is converted to BIT.
- Conversions in arithmetic operations or comparisons occur as described in the PL/I Language Reference.
- Conversions for the ADD, DIVIDE, MULTIPLY, and SUBTRACT built-in functions occur as described in the PL/I Language Reference except that operations specified as scaled fixed binary are evaluated as scaled fixed decimal.
- Nonzero scale factors are permitted in FIXED BIN declares.
- If the result of any precision-handling built-in function (ADD, BINARY, and so on) has FIXED BIN attributes, the specified or implied scale factor can be nonzero.
- Even if all arguments to the MAX or MIN built-in functions are UNSIGNED FIXED BIN, all the arguments are converted to SIGNED, the function is evaluated with these converted arguments, and the result is always SIGNED.
- Even when you add, multiply, or divide two UNSIGNED FIXED BIN operands, all the operands are converted to SIGNED, the operation is evaluated with these converted arguments, and the result has the SIGNED attribute.
- Even when you apply the MOD or REM built-in functions to two UNSIGNED FIXED BIN operands, all the arguments are converted to SIGNED, the function is evaluated with these converted arguments, and the result has the SIGNED attribute.
- Declaring a variable with the OPTIONS attribute implies the ENTRY attribute.

Under the ANS suboption:

- For operations requiring string data, data with the BINARY attribute is converted to CHARACTER.
- Conversions in arithmetic operations or comparisons occur as described in the PL/I Language Reference.
- Conversions for the ADD, DIVIDE, MULTIPLY, and SUBTRACT built-in functions occur as described in the PL/I Language Reference.
- Nonzero scale factors are not permitted in FIXED BIN declares.
- If the result of any precision-handling built-in function (ADD, BINARY, and so on) has FIXED BIN attributes, the specified or implied scale factor must be zero.
- If all arguments to the MAX or MIN built-in functions are UNSIGNED FIXED BIN, the result is UNSIGNED.
- When you add, multiply or divide two UNSIGNED FIXED BIN operands, the result has the UNSIGNED attribute.
- When you apply the MOD or REM built-in functions to two UNSIGNED FIXED BIN operands, the result has the UNSIGNED attribute.
- Declaring a variable with the OPTIONS attribute does not imply the ENTRY attribute.

Also, under RULES(ANS), the following errors, which the old compilers ignored, produce E-level messages:

- Specifying a string constant as the argument to the STRING built-in
- Giving too many asterisks as subscripts in an array reference
- Qualifying a CONTROLLED variable with a POINTER reference (as if the CONTROLLED variable were BASED)

**BYNAME | NOBYNAME**
Specifying NOBYNAME causes the compiler to flag all BYNAME assignments with an E-level message.

**CONTROLLED | NOCONTROLLED**
Specifying NOCONTROLLED causes the compiler to flag any use of the CONTROLLED attribute.
Specifying CONTROLLED causes the compiler not to flag the use of the CONTROLLED attribute.

**DECSIZE | NODECSIZE**
Specifying DECSIZE causes the compiler to flag any assignment of a FIXED DECIMAL expression to a FIXED DECIMAL variable when the SIZE condition is disabled if the SIZE condition could be raised by the assignment.
Specifying RULES(DECSIZE) might cause the compiler to produce many messages since if SIZE is disabled, then any statement of the form \( X = X + 1 \) will be flagged if \( X \) is FIXED DECIMAL.

**ELSEIF | NOELSEIF**
Specifying NOELSEIF causes the compiler to flag any ELSE statement that is immediately followed by an IF statement and suggest that it be rewritten as a SELECT statement.
This option can be useful in enforcing that SELECT statements be used rather than a series of nested IF-THEN-ELSE statements.

**EVENDEC | NOEVENDEC**
Specifying NOEVENDEC causes the compiler to flag any FIXED DECIMAL declaration that specifies an even precision.

**GLOBALDO | NOGLOBALDO**
Specifying NOGLOBALDO instructs the compiler to flag all DO loops with control variables that are declared in a parent block.

**GOTO | NOGOTO**
Specifying NOGOTO(LOOSE) causes the compiler to flag any GOTO statement to a label constant unless the GOTO is exiting an ON-unit or unless the target label constant is in the same block as the GOTO statement.
Specifying NOGOTO(STRICT) causes the compiler to flag any GOTO statement to a label constant unless the GOTO is exiting an ON-unit.

**LAXBIF | NOLAXBIF**
Specifying LAXBIF causes the compiler to build a contextual declaration for built-in functions, such as NULL, even when used without an empty parameter list.

**LAXCTL | NOLAXCTL**
Specifying LAXCTL allows a CONTROLLED variable to be declared with a constant extent and yet to be allocated with a differing extent. NOLAXCTL requires that if a CONTROLLED variable is to be allocated with a varying extent, then that extent must be specified as an asterisk or as a non-constant expression.
The following code is illegal under NOLAXCTL:
dcl a bit(8) ctl;
alloc a;
alloc a bit(16);

But this code would still be valid under NOLAXCTL:
dcl b bit(n) ctl;
dcl n fixed bin(31) init(8);
alloc b;
alloc b bit(16);

**LAXDCL | NOLAXDCL**
Specifying LAXDCL allows implicit declarations. NOLAXDCL disallows all implicit and contextual declarations except for BUILTINs and for files SYSIN and SYSPRINT.

**LAXDEF | NOLAXDEF**
Specifying LAXDEF allows so-called illegal defining to be accepted without any compiler messages (rather than the E-level messages that the compiler would usually produce).

**LAXENTRY | NOLAXENTRY**
Specifying LAXENTRY allows unprototyped entry declarations. Specifying NOLAXENTRY causes the compiler to flag all unprototyped entry declarations, that is, all ENTRY declares that do not specify a parameter list. Note that this would mean that if an ENTRY should have no parameters, it should be declared as ENTRY() rather than simply as ENTRY.

**STRICT**
Specifying RULES(NOLAXENTRY(STRICT)) causes the compiler to flag unprototyped entry declarations that have the OPTIONS(ASM) attribute.

**LOOSE**
Specifying RULES(NOLAXENTRY(LOOSE)) causes the compiler not to flag unprototyped entry declarations that have the OPTIONS(ASM) attribute.

RULES(LAXENTRY) is the default. When you specify NOLAXENTRY, STRICT is the default.

**LAXIF | NOLAXIF**
Specifying RULES(NOLAXIF) causes the compiler to flag any IF, WHILE, UNTIL, and WHEN clauses that do not have the attributes BIT(1) NONVARYING. It also causes the compiler to flag the assignments of the form \(x = y = z\), but it does not flag assignments of the \(x = (y = z)\) form.

The following would all be flagged under NOLAXIF:
dcl i fixed bin;
dcl b bit(8);
...
if i then ...
if b then ...

**LAXINOUT | NOLAXINOUT**
Specifying NOLAXINOUT causes the compiler to assume that all ASSIGNABLE BYADDR parameters are input (and possibly output) parameters and hence to issue a warning if it thinks such a parameter has not been initialized.

**LAXLINK | NOLAXLINK**
Specifying NOLAXLINK causes the compiler to flag any assign or compare of two ENTRY variables or constants if any of the following do not match:
• the parameter description lists
  For instance if A1 is declared as ENTRY(CHAR(8)) and A2 as ENTRY(POINTER) VARIABLE, then under RULES(NOLAXLINK), the compiler would flag an attempt to assign A1 to A2.
• the RETURNS attribute
  For instance if A3 is declared as ENTRY RETURNS(FIXED BIN(31)) and A4 as an ENTRY VARIABLE without the RETURNS attribute, then under RULES(NOLAXLINK), the compiler would flag an attempt to assign A3 to A4.
• the LINKAGE and other OPTIONS suboptions
  For instance if A5 is declared as ENTRY OPTIONS(ASM) and A6 as an ENTRY VARIABLE without the OPTIONS attribute, then under RULES(NOLAXLINK), the compiler would flag an attempt to assign A5 to A6 (since the OPTIONS(ASM) in the declare of A5 implies that A5 has LINKAGE(SYSTEM)) while since A6 has no OPTIONS attribute, it will have LINKAGE(OPTLINK) by default).

**LAXMARGINS | NOLAXMARGINS**
Specifying NOLAXMARGINS causes the compiler to flag, depending on the setting of the STRICT and XNUMERIC suboption, lines containing non-blank characters after the right margin. This can be useful in detecting code, such as a closing comment, that has accidentally been pushed out into the right margin.

If the NOLAXMARGINS and STMT options are used together with one of the preprocessors, then any statements that would be flagged because of the NOLAXMARGINS option will be reported as statement zero (since statement numbering occurs only after all the preprocessors are finished, but the detection of text outside the margins occurs as soon as the source is read).

**STRICT**
Under STRICT, the compiler flags any line that contains non-blank characters after the right margin.

**XNUMERIC**
Under XNUMERIC, the compiler flags any line that contains non-blank characters after the right margin except if the right margin is column 72 and columns 73 through 80 all contain numeric digits.

**LAXNESTED | NOLAXNESTED**
Specifying RULES(NOLAXNESTED) causes the compiler to flag any executable code in a procedure that follows any subprocedures.

Specifying RULES(LAXNESTED) causes the compiler not to flag the executable code in a procedure that follows any subprocedures.

**ALL**
Under ALL, all violations of RULES(NOLAXNESTED) are flagged. ALL is the default.

**SOURCE**
Under SOURCE, only those violations that occur in the primary source file are flagged.

**LAXPUNC | NOLAXPUNC**
Specifying NOLAXPUNC causes the compiler to flag with an E-level message any place where it assumes punctuation that is missing.

For instance, given the statement "I = (1 * (2);", the compiler assumes that a closing right parenthesis was meant before the semicolon. Under
RULES(NOLAXPUNC), this statement would be flagged with an E-level message; otherwise it would be flagged with a W-level message.

**LAXQUAL | NOLAXQUAL**

Specifying NOLAXQUAL(LOOSE) causes the compiler to flag any reference to structure members that are not level 1 and are not dot qualified. Consider the following example:

```
dcl
  1 a,
    2 b fixed bin,
    3 c fixed bin;

c = 11;  /* would be flagged */
b.c = 13; /* would not be flagged */
a.c = 17; /* would not be flagged */
```

Specifying NOLAXQUAL(STRICT) causes the compiler to flag any reference to structure members that do not include the level-1 name. Consider the following example:

```
dcl
  1 a,
    2 b fixed bin,
    3 c fixed bin;

c = 11;  /* would be flagged */
b.c = 13; /* would be flagged */
a.c = 17; /* would not be flagged */
```

**ALL**

Under ALL, all violations of RULES(NOLAXQUAL) are flagged. ALL is the default.

**FORCE**

Under FORCE, only those violations that occur in structures with the FORCE(NOLAXQUAL) attribute are flagged.

**LAXRETURN | NOLAXRETURN**

Specifying NOLAXRETURN causes the compiler to generate code to raise the ERROR condition when a RETURN statement is used in either of the following ways.

- With an expression in a procedure that is coded without the RETURNS option
- Without an expression in a procedure that is coded with the RETURNS option

ERROR will also be raised if the code falls through to the END statement in a PROCEDURE with the RETURNS attribute. RULES(LAXRETURN) is the default.

**LAXSCALE | NOLAXSCALE**

Specifying NOLAXSCALE causes the compiler to flag any FIXED BIN(p, q) or FIXED DEC(p, q) declaration where q < 0 or p < q.

It also causes the compiler to flag ROUND(x, p) when p < 0.

The message issued when flagging ROUND(x, p) is different from that issued when flagging the FIXED BIN(p, q) or FIXED DEC(p, q) declaration. Therefore, you can use the EXIT option to suppress the message issued when ROUND(x, p) is flagged and keep the message for other questionable declarations.
LAXSEMI | NOLAXSEMI
Specifying NOLAXSEMI causes the compiler to flag any semicolons appearing inside comments.

LAXSTG | NOLAXSTG
Specifying NOLAXSTG causes the compiler to flag declares where a variable A is declared as BASED on ADDR(B) and STG(A) > STG(B) even (and this is the key part) if B is a parameter.

Note that even with NOLAXSTG specified, if B has subscripts, no IBM2402I E-level message will be produced.

The compiler would already flag this kind of problem if B were in AUTOMATIC or STATIC storage, but it does not, by default, flag this when B is a parameter (since some customers declare B with placeholder attributes that do not describe the actual argument). For those customers whose parameter and argument declares match (or should match), specifying RULES(NOLAXSTG) can help detect more storage overlay problems.

LAXSTRZ | NOLAXSTRZ
Specifying LAXSTRZ causes the compiler not to flag any bit or character variable that is initialized to or assigned a constant value that is too long if the excess bits are all zeros (or if the excess characters are all blank).

MULTICLOSE | NOMULTICLOSE
Specifying NOMULTICLOSE causes the compiler to flag all statements that force the closure of multiple groups of statement with an E-level message.

PADDING | NOPADDING
Specifying NOPADDING causes the compiler to flag all structures that contain padding.

ALL
Specifying NOPADDING(ALL) causes the compiler to flag all violations of RULES(NOPADDING). ALL is the default.

SOURCE
Specifying NOPADDING(SOURCE) causes the compiler to flag only those violations that occur in the primary source files.

PROCENDONLY | NOPROCENDONLY
Specifying NOPROCENDONLY causes any END statement closing a PROCEDURE to be flagged if the END statement does not name the PROCEDURE, that is, if the END keyword is immediately followed by a semicolon.

RECURSIVE | NORECURSIVE
Specifying NORECURSIVE causes the compiler not to flag the use of the RECURSIVE attribute or any procedure that directly calls itself.

Specifying RECURSIVE causes the compiler to flag any use of the RECURSIVE attribute or any procedure that directly calls itself.

Note: Do not use RULES(NORECURSIVE) and DFT(RECURSIVE) together.

SELFASSIGN | NOSELFASSIGN
Specifying NOSELFASSIGN causes the compiler to flag all assignments where the source and the target are the same.

RULES(SELFASSIGN) is the default.

UNREF | NOUNREF
Specifying NOUNREF causes the compiler to flag any level-1 AUTOMATIC
variable, which is not referenced and that, if it is a structure or union, contains no subelement that is referenced. NOUNREF ignores variables with names that start with any of the following prefixes: DSN, DFH, EYU, and SQL.

**ALL**

Specifying RULES(NOUNREF(ALL)) causes the compiler to flag all unreferenced variables. When NOUNREF is specified, ALL is the default.

**SOURCE**

Specifying RULES(NOUNREF(SOURCE)) causes the compiler to flag unreferenced variables that are not declared in an INCLUDE file.

*Default:* RULES (IBM BYNAME CONTROLLED NODECSIZE EVENDEC ELSEIF GLOBALDO GOTO NOLAXBIF NOLAXCTL LAXDCL NOLAXDEF LAXENTRY LAXIF LAXINOUT LAXLINK LAXNESTED LAXPUNC LAXMARGINS(STRICT) LAXQUAL LAXRETURN LAXSCALE LAXSEMI LAXSTG NOLAXSTRZ MULTICLOSE PADDING PROCENDONLY RECURSIVE SELFASSIGN UNREF)

**SEMANTIC**

The SEMANTIC option specifies that the execution of the compiler’s semantic checking stage depends on the severity of messages issued prior to this stage of processing.

ABBREVIATIONS: SEM, NSEM

**SEMANTIC**

Equivalent to NOSEMANTIC(S).

**NOSEMANTIC**

Processing stops after syntax checking. No semantic checking is performed.

**NOSEMANTIC (S)**

No semantic checking is performed if a severe error or an unrecoverable error has been encountered.

**NOSEMANTIC (E)**

No semantic checking is performed if an error, a severe error, or an unrecoverable error has been encountered.

**NOSEMANTIC (W)**

No semantic checking is performed if a warning, an error, a severe error, or an unrecoverable error has been encountered.

Semantic checking is not performed if certain kinds of severe errors are found. If the compiler cannot validate that all references resolve correctly (for example, if built-in function or entry references are found with too few arguments) the suitability of any arguments in any built-in function or entry reference is not checked.
**SERVICE**

The SERVICE option places a string in the object module, if generated. This string is loaded into memory with any load module into which this object is linked, and if the LE dump includes a traceback, this string will be included in that traceback.

```
NOSERVICE
```

```
SERVICE('service string')
```

**ABBREVIATIONS:** SERV, NOSERV

The string is limited to 64 characters in length.

To ensure that the string remains readable across locales, only characters from the invariant character set should be used.

**SOURCE**

The SOURCE option specifies that the compiler includes a listing of the source program in the compiler listing. The source program listed is either the original source input or, if any preprocessors were used, the output from the last preprocessor.

```
NOSOURCE
```

```
SOURCE
```

**ABBREVIATIONS:** S, NS

**SPILL**

The SPILL option specifies the size of the spill area to be used for the compilation. When too many registers are in use at once, the compiler dumps some of the registers into temporary storage that is called the spill area.

```
SPILL(size)
```

**ABBREVIATIONS:** SP

If you have to expand the spill area, you will receive a compiler message telling you the size to which you should increase it. Once you know the spill area that your source program requires, you can specify the required size (in bytes) as shown in the syntax diagram above. The maximum spill area size is 3900. Typically, you will need to specify this option only when compiling very large programs with OPTIMIZE.

**STATIC**

The STATIC option controls whether INTERNAL STATIC variables are retained in the object module even if unreferenced.

```
SHORT
```

```
FULL
```

```
FULL
```

```
NOSERVICE
```

```
SERVICE('service string')
```

```
NOSOURCE
```

```
SOURCE
```

```
SPILL(size)
```

```
SHORT
```

```
FULL
```

```
FULL
```
SHORT
INTERNAL STATIC will be retained in the object module only if used.

FULL
All INTERNAL STATIC with INITIAL will be retained in the object module.

If INTERNAL STATIC variables are used as "eyecatchers", you should specify the STATIC(FULL) option to insure that they will be in the generated object module.

STDSYS
The STDSYS option specifies that the compiler should cause the SYSPRINT file to be equated to the C stdout file and the SYSIN file to be equated to the C stdin file.

Using the STDSYS option may make it easier to develop and debug a mixed PL/I and C application.

When SYSPRINT is equated to stdout, its LINESIZE cannot be greater than 132 (the largest value allowed by C).

STMT
The STMT option specifies that statements in the source program are to be counted and that this "statement number" is used to identify statements in the compiler listings resulting from the AGGREGATE, ATTRIBUTES, SOURCE and XREF options.

The default is NOSTMT.

When the STMT option is specified, the source and message listings will include both the logical statement numbers and the source file numbers.

Note that the GOSTMT option does not exist. The only option that produces information at run time identifying where an error has occurred is the GONUMBER option. When the GONUMBER option is used, the term statement in the runtime error messages refer to line numbers as used by the NUMBER compiler option even if the STMT option is in effect.

NUMBER and STMT are mutually exclusive and specifying one will negate the other.

STORAGE
The STORAGE option directs the compiler to produce as part of the listing a summary of the storage used by each procedure and begin-block.
ABBREVIATIONS: STG, NSTG

The STORAGE output also includes the amount of storage used for the internal static for the compilation.

The compiler now generates inline code for STG(x) where x has the BASED attribute with the REFER option and meets both of the following conditions:
- Elements in x with the NONVARYING and BIT attributes have the ALIGNED attribute.
- All other elements in x have the UNALIGNED attribute.

STRINGOFGRAPHIC

The STRINGOFGRAPHIC option determines whether the result of the STRING built-in function when applied to a GRAPHIC aggregate has the attribute CHARACTER or GRAPHIC.

SYNTAX

The SYNTAX option specifies that the compiler continues into syntax checking after preprocessing when you specify the MACRO option, unless an unrecoverable error has occurred. Whether the compiler continues with the compilation depends on the severity of the error, as specified by the NOSYNTAX option.

SYNTAX

Continues syntax checking after preprocessing unless a severe error or an unrecoverable error has occurred. SYNTAX is equivalent to NOSYNTAX(S).

ABBREVIATIONS: SYN, NSYN

SYNTAX

Continues syntax checking after preprocessing unless a severe error or an unrecoverable error has occurred. SYNTAX is equivalent to NOSYNTAX(S).

NOSYNTAX

Processing stops unconditionally after preprocessing.

NOSYNTAX(N)

No syntax checking if a warning, error, severe error, or unrecoverable error is detected.
NOSYNTAX(E)
No syntax checking if the compiler detects an error, severe error, or unrecoverable error.

NOSYNTAX(S)
No syntax checking if the compiler detects a severe error or unrecoverable error.

If the NOSYNTAX option terminates the compilation, no cross-reference listing, attribute listing, or other listings that follow the source program is produced.

You can use this option to prevent wasted runs when debugging a PL/I program that uses the preprocessor.

If the NOSYNTAX option is in effect, any specification of the CICS preprocessor via the CICS, XOPT or XOPTS options will be ignored. This allows the MACRO preprocessor to be invoked before invoking the CICS translator.

SYSPARM
The SYSPARM option allows you to specify the value of the string that is returned by the macro facility built-in function SYSPARM.

```
SYSPARM("string")
```

string
Can be up to 64 characters long. A null string is the default.

For more information about the macro facility, see PL/I Language Reference.

SYSTEM
The SYSTEM option specifies the format used to pass parameters to the MAIN PL/I procedure, and generally indicates the host system under which the program runs.

```
SYSTEM(MVS)(CICS)(IMS)(OS)(TSO)
```

Table 4 on page 82 shows the type of parameter list you can expect, and how the program runs under the specified host system. It also shows the implied settings of NOEXECOPS. Your MAIN procedure must receive only those types of parameter lists that are indicated as valid in this table. Additional runtime information for the SYSTEM option is provided in z/OS Language Environment Programming Guide.
### Table 4. SYSTEM option table

<table>
<thead>
<tr>
<th>SYSTEM option</th>
<th>Type of parameter list</th>
<th>Program runs as</th>
<th>NOEXECOPS implied</th>
</tr>
</thead>
<tbody>
<tr>
<td>SYSTEM(MVS)</td>
<td>Single CHARACTER string or no parameters.</td>
<td>z/OS application program</td>
<td>NO</td>
</tr>
<tr>
<td></td>
<td>Otherwise, arbitrary parameter list.</td>
<td></td>
<td>YES</td>
</tr>
<tr>
<td>SYSTEM(CICS)</td>
<td>Pointer(s)</td>
<td>CICS® transaction</td>
<td>YES</td>
</tr>
<tr>
<td>SYSTEM(IMS)</td>
<td>Pointer(s)</td>
<td>IMS™ application program</td>
<td>YES</td>
</tr>
<tr>
<td>SYSTEM(OS)</td>
<td>z/OS UNIX parameter list</td>
<td>z/OS UNIX application program</td>
<td>YES</td>
</tr>
<tr>
<td>SYSTEM(TSO)</td>
<td>Pointer to CPPL processor</td>
<td>TSO command processor</td>
<td>YES</td>
</tr>
</tbody>
</table>

**Note:**
1. See ["Invoking MAIN under TSO/E" on page 179](#) for more details about how to invoke a MAIN under TSO.

Under SYSTEM(IMS), all pointers are presumed to be passed BYVALUE, but under SYSTEM(MVS) they are presumed to be passed BYADDR.

MAIN procedures run under CICS must be compiled with SYSTEM(CICS) or SYSTEM(MVS).

It is highly recommended that NOEXECOPS be specified in the MAIN procedure OPTIONS option for code, such as a DB2 stored procedure, compiled with SYSTEM(MVS) but run where runtime options would not be passed.

The compiler will flag any MAIN program compiled with SYSTEM(MVS) if it has either more than one parameter or a single parameter that is not CHARACTER VARYING. It is probably better to compile such MAIN programs with SYSTEM(OS) because then the library simply passes on to MAIN the parameter list without any scanning for options or other massaging of the parameter list.

### TERMINAL

The TERMINAL option determines whether or not diagnostic and information messages produced during compilation are displayed on the terminal.

**Note:** This option applies only to compilations under z/OS UNIX.

**ABBREVIATIONS:** TERM, NTERM

**TERMINAL**

Messages are displayed on the terminal.

**NOTERMINAL**

No information or diagnostic compiler messages are displayed on the terminal.
**TEST**

The TEST option specifies the level of testing capability that the compiler generates as part of the object code. It allows you to control the location of test hooks and to control whether or not the symbol table will be generated.

ABBREVIATIONS: AALL, ACICS, AMACRO, ASQL

**STMT**

Causes the compiler to generate a statement table, and if the HOOK suboption is in effect, inserts hooks at statement boundaries and block boundaries.

**PATH**

Causes the compiler to generate a statement table, and if the HOOK suboption is in effect, tells the compiler to insert hooks:

- Before the first statement enclosed by an iterative DO statement
- Before the first statement of the true part of an IF statement
- Before the first statement of the false part of an IF statement
- Before the first statement of a true WHEN or OTHERWISE statement of a SELECT group
- Before the statement following a user label, excluding labeled FORMAT statements. If a statement has multiple labels, only one hook is inserted.
- At CALLs or function references - both before and after control is passed to the routine
- At block boundaries

**BLOCK**

Causes the compiler to generate a statement table, and if the HOOK suboption is in effect, tells the compiler to insert hooks at block boundaries (block entry and block exit).

**ALL**

Causes the compiler to generate a statement table, and if the HOOK suboption is in effect, inserts hooks at all possible locations and generates a statement table.

**NONE**

No hooks are put into the program.

*Note:* Under opt(2) and opt(3), hooks are set only at block boundaries.
SYM
Creates a symbol table that allows you to examine variables by name.

NOSYM
No symbol table is generated.

NOTEST
Suppresses the generation of all testing information.

HOOK
Causes the compiler to insert hooks into the generated code if any of the TEST suboptions ALL, STMT, BLOCK or PATH are in effect.

NOHOOK
Causes the compiler not to insert hooks into the generated code.

For Debug Tool to generate overlay hooks, one of the suboptions ALL, PATH, STMT or BLOCK must be specified, but HOOK need not be specified, and NOHOOK would in fact be recommended.

If NOHOOK is specified, ENTRY and EXIT breakpoints are the only PATH breakpoints at which Debug Tool will stop.

SEPARATE
Causes the compiler to place most of the debug information it generates into a separate debug file. Using this option will substantially reduce the size of the object deck created by the compiler when the TEST option is in effect.

If your program contains GET or PUT DATA statements, the separate debug file will contain less debug information since those statements require that symbol table information be placed into the object deck.

The generated debug information always includes a compressed version of the source that is passed to the compiler. This means that source may be specified by using SYSIN DD * or that the source may be a temporary data set that is created by an earlier job step (for example, the source may be the output of the old SQL or CICS precompilers). The suboptions that you specify for the LISTVIEW option control the content of the source.

If SEPARATE is used in a batch compile, then the JCL for that compile must include a DD card for SYSDEBUG which must name a data set with RECFM=FB and with 80 <= LRECL <= 1024.

This suboption may not be used with the LINEDIR compiler option.

NOSEPARATE
Causes the compiler to place all of the debug information it generates into the object deck.

Under this option, the generated debug information will not include a compressed version of the source passed to the compiler. This means that source must in a data set that can be found by DebugTool when you try to debug the program.

SEPNAME
Causes the compiler to place the name of the separate debug file into the object deck.

This option is ignored if the SEPARATE option is not in effect.

NOSEPNAME
Causes the compiler not to place the name of the separate debug file into the object deck.
This option is ignored if the SEPARATE option is not in effect.

Notes:

- Under opt(2) or opt(3), hooks are set only at block boundaries. This means that debugging of optimized code is effectively limited to tracing entry and exit to PROCEDUREs and BEGIN blocks.
- You must use Debug Tool Version 6 (or later) to debug code compiled with the SEPARATE compiler option.
- There is no support for an input file that spans concatenated data sets.

Specifying TEST(NONE,NOSYM) causes the compiler to set the option to NOTEST.

Use of TEST(NONE,SYM) is strongly discouraged, and it is unclear what is intended when you specify these settings. You would probably be much better off if you specified TEST(ALL,SYM,NOHOOK) or TEST(STMT,SYM,NOHOOK).

Any TEST option other than NOTEST and TEST(NONE,NOSYM) will automatically provide the attention interrupt capability for program testing.

If the program has an ATTENTION ON-unit that you want invoked, you must compile the program with either of the following:
- The INTERRUPT option
- A TEST option other than NOTEST or TEST(NONE,NOSYM)

Note: ATTENTION is supported only under TSO.

The TEST option will imply GONUMBER.

Because the TEST option can increase the size of the object code and can affect performance, you might want to limit the number and placement of hooks.

If the TEST option is specified, no inlining will occur.

Structures with REFER are supported in the symbol table.

If TEST(SYM) is in effect, the compiler will generate tables to enable the automonitor feature of DebugTool. These tables may substantially increase the size of the object module unless the TEST(SEPARATE) option is in effect. When the automonitor feature of DebugTool is activated, these tables will be used to display the values of the variables used in a statement before the statement executes - as long as the variable has computational type or has the attribute POINTER, OFFSET or HANDLE. If the statement is an assignment statement, the value of the target will also be displayed; however, if the target has not been initialized or assigned previously, its value will be meaningless.

Any variable declared with an * for its name will not be visible when using DebugTool. Additionally, if an * is used as the name of a parent structure or substructure, then all of its children will also be invisible. For these purposes, it might be better to use a single underscore for the name of any structure elements that you want to leave “unnamed”.

**UNROLL**

The UNROLL option controls loop unrolling under optimization. Loop unrolling is an optimization that replicates a loop body multiple times and adjusts the loop control code accordingly.
AUTO
  Indicates that the compiler is permitted to unroll loops that it determines are
  appropriate for unrolling.

Specifying the UNROLL option can increase the size of the object code that is
generated.

NO  Indicates that the compiler is not permitted to unroll loops.

The UNROLL option is ignored when the NOOPTIMIZE option is in effect.

Loop unrolling improves the performance of a program by exposing instruction
level parallelism for instruction scheduling and software pipelining. It also creases
code in the new loop body, which might increase pressure on register allocation,
cause register spilling, and thus cause a loss in performance.

Therefore, before you unroll a loop, take the following steps to check if the
UNROLL option improves the performance of a particular application:
1. Compile the program with the usual options.
2. Run the program with a representative workload.
3. Recompile the program with the UNROLL option.
4. Rerun the program under the same conditions.

UNROLL(AUTO) is the default.

USAGE
  The USAGE option lets you choose different semantics for selected built-in
  functions.

HEX( SIZE | CURRENTSIZE )
  Under the HEX(SIZE) suboption, when HEX is applied to a VARYING or
  VARYINGZ string, it will return a hex string that represents the maximum
  amount of storage used by the string.

  Under the HEX(CURRENTSIZE) suboption, when HEX is applied to a
  VARYING or VARYINGZ string, it will return a hex string that represents the
  current amount of storage used by the string.

ROUND( IBM | ANS )
  Under the ROUND(IBM) suboption, the second argument to the ROUND
  built-in function is ignored if the first argument has the FLOAT attribute.
Under the ROUND(ANS) suboption, the ROUND built-in function is implemented as described in the PL/I Language Reference.

**SUBSTR( STRICT | LOOSE )**

Under the SUBSTR(STRICT) suboption, if x has CHARACTER type, a SUBSTR(x,y,z) built-in function reference will return a string whose length is equal to MIN( z, MAXLENGTH(x) ).

Under the SUBSTR(LOOSE) suboption, the same reference would return a string whose length is z.

The SUBSTR(LOOSE) suboption may be useful for those who have SUBSTR(x,y,z) references where x is a CHAR(1) BASED variable.

**UNSPEC( IBM | ANS )**

Under the UNSPEC(IBM) suboption, UNSPEC cannot be applied to a structure and, if applied to an array, returns an array of bit strings.

Under the UNSPEC(ANS) suboption, UNSPEC can be applied to structures and, when applied to a structure or an array, UNSPEC returns a single bit string.

*Default:* USAGE( HEX(SIZE) ROUND(IBM) SUBSTR(STRICT) UNSPEC(IBM) )

**WIDECHAR**

The WIDECHAR option specifies the format in which WIDECHAR data will be stored.

-WIDECHAR-(BIGENDIAN|LITTLEENDIAN)-

**BIGENDIAN**

Indicates that WIDECHAR data will be stored in big-endian format. For instance, the WIDECHAR value for the UTF-16 character '1' will be stored as '0031'x.

**LITTLEENDIAN**

Indicates that WIDECHAR data will be stored in little-endian format. For instance, the WIDECHAR value for the UTF-16 character '1' will be stored as '3100'x.

WX constants should always be specified in big-endian format. Thus the value '1' should always be specified as '0031'wx, even if under the WIDECHAR(LITTLEENDIAN) option, it is stored as '3100'x.

**WINDOW**

The WINDOW option sets the value for the \( w \) window argument used in various date-related built-in functions.

-WINDOW-(1950)-

\( w \) Either an unsigned integer that represents the start of a fixed window or a negative integer that specifies a “sliding” window. For example, \( \text{WINDOW}(-20) \) indicates a window that starts 20 years prior to the year when the program runs.
**WRITABLE**

The WRITABLE option specifies that the compiler may treat static storage as writable (and if it does, this would make the resultant code non-reentrant).

This option has no effect on programs compiled with the RENT option.

The NORENT WRITABLE options allow the compiler to write on static storage to implement the following constants or variables:

- CONTROLLED variables
- FETCHABLE ENTRY constants
- FILE constants

Under the NORENT WRITABLE options, a module using CONTROLLED variables, performing I/O, or using FETCH is not reentrant.

The NORENT NOWRITABLE options require the compiler not to write on static storage for the following constants or variables:

- CONTROLLED variables
- FETCHABLE ENTRY constants
- FILE constants

Under the NORENT NOWRITABLE options, a module using CONTROLLED variables, performing I/O, or using FETCH is reentrant.

The FWS and PRV suboptions determine how the compiler handles CONTROLLED variables:

**FWS**

Upon entry to an EXTERNAL procedure, the compiler makes a library call to find storage it can use to address the CONTROLLED variables in that procedure (and any subprocedures).

**PRV**

The compiler will use the same pseudoregister variable mechanism used by the old OS PL/I compiler to address CONTROLLED variables.

Hence, under the NORENT NOWRITABLE(PRV) options, old and new code may share CONTROLLED variables.

However, this also means that under the NORENT NOWRITABLE(PRV) options, the use of CONTROLLED variables is subject to all the same restrictions as under the old compiler.

Under the NORENT NOWRITABLE(FWS) options, the following application might not perform as well as if they were compiled with the RENT or WRITABLE options:

- Applications that use CONTROLLED variables
- Applications that assign FILE CONSTANTs to FILE VARIABLEs
The performance of an application under NORENT NOWRITABLE(FWS) may be especially bad if it uses many CONTROLLED variables in many PROCEDUREs.

Under the NOWRITABLE option, the following may not be declared in a PACKAGE outside a PROCEDURE:

- CONTROLLED variables
- FETCHABLE ENTRY constants
- FILE constants

Code compiled with NORENT WRITABLE cannot be mixed with code compiled with NORENT NOWRITABLE if they share any external CONTROLLED variables. In general, you should avoid mixing code compiled with WRITABLE with code compiled with NOWRITABLE.

**XINFO**

The XINFO option specifies that the compiler should generate additional files with extra information about the current compilation unit.

A definition side-deck file is created. This file lists, for the compilation unit, all:

- defined EXTERNAL procedures
- defined EXTERNAL variables
- statically referenced EXTERNAL routines and variables
- dynamically called fetched modules

Under batch, this file is written to the file specified by the SYSDEFSD DD statement. Under z/OS UNIX Systems Services, this file is written to the same directory as the object deck and has the extension "def".

For instance, given the program:

```assembly
defs: proc;
   dcl (b,c) ext entry;
   dcl x ext fixed bin(31) init(1729);
   dcl y ext fixed bin(31) reserved;
   call b(y);
   fetch c;
   call c;
end;
```

The following def file would be produced:

```assembly
EXPORTS CODE
DEFS
EXPORTS DATA
```

Chapter 1. Using compiler options and facilities  89
The def file can be used to build a dependency graph or cross-reference analysis of your application.

**NODEF**
No definition side-deck file is created.

**MSG**
Message information is generated to the ADATA file. See the appendix for more details on the format of the ADATA file.

Under batch, the ADATA file is generated to the file specified by the SYSADATA DD statement. Under z/OS UNIX, the ADATA is generated in the same directory as the object file and has an extension of adt.

**NOMSG**
No message information is generated to the ADATA file. If neither MSG nor SYM is specified, no ADATA file is generated.

**SYM**
Symbol information is generated to the ADATA file. See the appendix for more details on the format of the ADATA file.

Under batch, the ADATA file is generated to the file specified by the SYSADATA DD statement. Under z/OS UNIX, the ADATA file is generated in the same directory as the object file and has an extension of adt.

**NOSYM**
No symbol information is generated to the ADATA file.

**SYN**
Syntax information is generated to the ADATA file. See the appendix for more details on the format of the ADATA file. Specifying the XINFO(SYN) option can greatly increase the amount of storage, both in memory and for the file produced, required by the compiler.

Under batch, the ADATA file is generated to the file specified by the SYSADATA DD statement. Under z/OS UNIX, the ADATA file is generated in the same directory as the object file and has an extension of adt.

**NOSYN**
No syntax information is generated to the ADATA file.

**XML**
An XML side-file is created. This XML file includes:
- the file reference table for the compilation
- the block structure of the program compiled
- the messages produced during the compilation

Under batch, this file is written to the file specified by the SYSXMLSD DD statement. Under z/OS UNIX Systems Services, this file is written to the same directory as the object deck and has the extension "xml".

The DTD file for the XML produced is:
```xml
<?xml encoding="UTF-8"?>
<!ELEMENT PACKAGE ((PROCEDURE)*,(MESSAGE)*,FILEREFERNCETABLE)>
<!ELEMENT PROCEDURE (BLOCKFILE,BLOCKLINE,(PROCEDURE)*,(BEGINBLOCK)*)>
```
NOXML
No XML side-file is created.

XML
The XML option lets you choose the case of the names and the character enclosing attributes in the XML generated by the XMLCHAR built-in function.

CASE( UPPER | ASIS )
Under the CASE(UPPER) suboption, the names in the XML generated by the XMLCHAR built-in function will all be in upper case.

Under the CASE(ASIS) suboption, the names in the XML generated by the XMLCHAR built-in function will be in the case used in their declares. Note that if you use the MACRO preprocessor without using the macro preprocessor option CASE(ASIS), then the source seen by the compiler will have all the names in upper case - and that would make specifying the XML(CASE(ASIS)) option useless.

XMLATTR( APOSTROPHE | QUOTE )
Under the XMLATTR(APOSTROPHE) suboption, all XML attributes generated by the XMLCHAR built-in function will be enclosed in apostrophes.

Under the XMLATTR(QUOTE) suboption, all XML attributes generated by the XMLCHAR built-in function will be enclosed in quotation marks. Under CODEPAGE(1026) and COEPAGE(1155), the hex value for the quote character is 'FC'x, and under all other supported EBCDIC code pages, it is '7F'x.

XREF
The XREF option provides a cross-reference table of names used in the program together with the numbers of the statements in which they are declared or referenced in the compiler listing.
ABBREVIATIONS: X, NX

FULL
  Includes all identifiers and attributes in the compiler listing.

SHORT
  Omits unreferenced identifiers from the compiler listing.

The only names not included in the cross reference listing created when using the XREF option are label references on END statements. For example, assume that statement number 20 in the procedure PROC1 is END PROC1; In this situation, statement number 20 does not appear in the cross reference listing for PROC1.)

If you specify both the XREF and ATTRIBUTES options, the two listings are combined. If there is a conflict between SHORT and FULL, the usage is determined by the last option specified. For example, ATTRIBUTES(SHORT) XREF(FULL) results in the FULL option for the combined listing.

For a description of the format and content of the cross-reference table, see “Cross-reference table” on page 98.

Blanks, comments and strings in options

Wherever you may use a blank when specifying an option, you may also specify as many blanks or comments as you wish.

However, if a comment is specified in %PROCESS line or in a line in an options file, then the comment must end on the same line as which it begins.

Similarly, if a comment starts in the command line or the PARM= specification, then it must also end there.

The same rule applies to strings: if a string is specified in %PROCESS line or in a line in an options file, then the string must end on the same line as which it begins. Similarly, if a string starts in the command line or the PARM= specification, then it must also end there.

Changing the default options

If you want to change the supplied default compiler options, during installation of the compiler you should edit and submit sample job IBMZWIOP.

This job will let you specify options that will be applied before any other options, thus effectively changing the default options. This job will also let you specify options that will be applied after all other options, thus effectively changing the default options and preventing them from being overridden.

If you want to change the defaults for the macro preprocessor options, you can also do this at installation time by specifying the appropriate PPMACRO option as part of this job. The PPCICS and PPSQL options let you make the corresponding changes for the CICS and SQL preprocessors respectively.
Consult the instructions in the sample job for more information.

**Specifying options in the %PROCESS or *PROCESS statements**

You can use either %PROCESS or *PROCESS in your program; they are equally acceptable. For consistency and readability in this book, we will always refer to %PROCESS.

The %PROCESS statement identifies the start of each external procedure and allows compiler options to be specified for each compilation. The options you specify in adjacent %PROCESS statements apply to the compilation of the source statements to the end of input, or the next %PROCESS statement.

To specify options in the %PROCESS statement, code as follows:

```plaintext
%PROCESS options;
```

where `options` is a list of compiler options. You must end the list of options with a semicolon, and the options list should not extend beyond the default right-hand source margin. The percent sign (%) or asterisk (*) must appear in the first column of the record. The keyword PROCESS can follow in the next byte (column) or after any number of blanks. You must separate option keywords by a comma or at least one blank.

The number of characters is limited only by the length of the record. If you do not wish to specify any options, code:

```plaintext
%PROCESS;
```

If you find it necessary to continue the %PROCESS statement onto the next record, terminate the first part of the list after any delimiter, and continue on the next record. You cannot split keywords or keyword arguments across records. You can continue a %PROCESS statement on several lines, or start a new %PROCESS statement. An example of multiple adjacent %PROCESS statements is as follows:

```plaintext
%PROCESS INT F(I) AG A(F) OP STG NEST X(F) SOURCE ;
%PROCESS LIST TEST ;
```

Compile-time options, their abbreviated syntax, and their IBM-supplied defaults are shown in [Table 3 on page 5](#).

How the compiler determines if there are any %PROCESS statements depends on the format of the initial source file:

**F or FB format**

if the first character in the record is a "*" or a "%", then the compiler will check if the next non-blank characters are "PROCESS".

**V or VB format**

if the first character in the record is a numeric, then the compiler will assume that the first 8 characters are sequence numbers, and if the ninth character is a "*" or a "%", then it will check if the next non-blank characters are "PROCESS". However, if the first character is not a numeric, then if the first character is a "*" or "%", the compiler will check if the next non-blank characters are "PROCESS".

**U format**

if the first character in the record is a "*" or a "%", then the compiler will check if the next non-blank characters are "PROCESS".
Using % statements

Statements that direct the operation of the compiler begin with a percent (%) symbol. % statements allow you to control the source program listing and to include external strings in the source program. % statements must not have label or condition prefixes and cannot be a unit of a compound statement. You should place each % statement on a line by itself.

The usage of each % control statement is listed below. For a complete description of these statements, see PL/I Language Reference.

%INCLUDE
Directs the compiler to incorporate external text into the source program.

%XINCLUDE
Directs the compiler to incorporate external text into the source program if it has not been previously included.

%PRINT
Directs the compiler to resume printing the source and insource listings.

%NOPRINT
Directs the compiler to suspend printing the source and insource listings until a %PRINT statement is encountered.

%PAGE
Directs the compiler to print the statement immediately after a %PAGE statement in the program listing on the first line of the next page.

%POP
Directs the compiler to restore the status of the %PRINT and %NOPRINT saved by the most recent %PUSH.

%PUSH
Saves the current status of the %PRINT and %NOPRINT in a push down stack on a last-in, first-out basis.

%SKIP
Specifies the number of lines to be skipped.

Using the %INCLUDE statement

%INCLUDE statements are used to include additional PL/I files at specified points in a compilation unit. The PL/I Language Reference describes how to use the %INCLUDE statement to incorporate source text from a library into a PL/I program.

For a batch compile

A library is an z/OS partitioned data set that can be used to store other data sets called members. Source text that you might want to insert into a PL/I program using a %INCLUDE statement must exist as a member within a library. “Source Statement Library (SYSLIB)” on page 160 further describes the process of defining a source statement library to the compiler.

The statement:

%INCLUDE DD1 (INVERT);

specifies that the source statements in member INVERT of the library defined by the DD statement with the name DD1 are to be inserted consecutively into the source program. The compilation job step must include appropriate DD statements.
If you omit the ddname, the ddname SYSLIB is assumed. In such a case, you must include a DD statement with the name SYSLIB. (The IBM-supplied cataloged procedures do not include a DD statement with this name in the compilation procedure step.)

For a z/OS UNIX compile

The name of the actual include file must be lowercase, unless you specify UPPERINC. For example, if you used the include statement %include sample, the compiler would find the file sample.inc, but would not find the file SAMPLE.inc. Even if you used the include statement %include SAMPLE, the compiler would still look for sample.inc.

The compiler looks for INCLUDE files in the following order:
1. Current directory
2. Directories specified with the –I flag or with the INCDIR compiler option
3. /usr/include directory
4. PDS specified with the INCPDS compiler option

The first file found by the compiler is included into your source.

A %PROCESS statement in source text included by a %INCLUDE statement results in an error in the compilation.

Figure 1 shows the use of a %INCLUDE statement to include the source statements for FUN in the procedure TEST. The library HPU8.NEWLIB is defined in the DD statement with the qualified name PLI.SYSLIB, which is added to the statements of the cataloged procedure for this job. Since the source statement library is defined by a DD statement with the name SYSLIB, the %INCLUDE statement need not include a ddname.

It is not necessary to invoke the preprocessor if your source program, and any text to be included, does not contain any macro statements.

```pli
//OPT4#9 JOB
//STEP3 EXEC IBMZCBG,PARM.PLI='INC,S,A,X,NEST'
//PLI.SYSLIB DD DSN=HPU8.NEWLIB,DISP=OLD
//PLI.SYSIN DD *
TEST: PROC OPTIONS(MAIN) REORDER;
   DCL ZIP PIC '99999';        /* ZIP CODE */
   DCL EOF BIT INIT('0'B);
   ON ENDFILE(SYSIN) EOF = '1'B;
   GET EDIT(ZIP) (COL(1), P'99999');
   DO WHILE(¬EOF);
      PUT SKIP EDIT(ZIP, CITYFUN(ZIP)) (P'99999', A(16));
      GET EDIT(ZIP) (COL(1), P'99999');
   END;
%PAGE;
%INCLUDE FUN;
END;            /* TEST */
//GO.SYSIN DD *
95141
95030
94101
/*
Figure 1. Including source statements from a library
```
Using the compiler listing

During compilation, the compiler generates a listing, most of which is optional, that contains information about the source program, the compilation, and the object module. The following description of the listing refers to its appearance on a printed page.

Note:

Although the compiler listing is for your use, it is not a programming interface and is subject to change.

Of course, if compilation terminates before reaching a particular stage of processing, the corresponding listings do not appear.

Heading information

The first page of the listing is identified by the product number, the compiler version number, a string specifying when the compiler was built, and the date and the time compilation commenced. This page and subsequent pages are numbered.

The listing will then show any options that have been specified for this compilation. These options will be shown even if the NOOPTIONS option has been specified and will include, in order, the following:

- the initial install options (those are the options set at install time and which are applied before any other options)
- under z/OS UNIX, any options specified in the IBM_OPTIONS environment variable
- any options specified in the parameter string passed to the compiler (i.e. in the command line under z/OS UNIX or in the PARM= under batch)
- any options specified in options files named in the compiler parameter string
  This will include the name of each options file and its contents in the same form as read by the compiler.
- any options specified in *PROCESS or %PROCESS lines in the source
- the final install options (those are the options set at install time and which are applied after any other options)

Near the end of the listing you will find either a statement that no errors or warning conditions were detected during the compilation, or a message that one or more errors were detected. The format of the messages is described under "Messages and return codes" on page 102. The second to the last line of the listing shows the time taken for the compilation. The last line of the listing is END OF COMPILATION OF xxxx, where xxxx is the external procedure name. If you specify the NOSYNTAX compiler option, or the compiler aborts early in the compilation, the external procedure name xxxx is not included and the line truncates to END OF COMPILATION.

The following sections describe the optional parts of the listing in the order in which they appear.

Options used for compilation

If you specify the OPTIONS option, a complete list of the options specified for the compilation, including the default options, appears on the following pages. It will show the settings of all the options finally in effect during the compilation. If the
setting of an option differs from the default setting after the initial install options were applied, then that line will be marked with a +

**Preprocessor input**

If you specify both the MACRO and INSOURCE options, the compiler lists input to the preprocessor, one record per line, each line numbered sequentially at the left.

If the preprocessor detects an error, or the possibility of an error, it prints a message on the page or pages following the input listing. The format of these messages is the same as the format for the compiler messages described under “Messages and return codes” on page 102.

**SOURCE program**

If you specify the SOURCE option, the compiler lists one record per line. These records will always include the source line and file numbers. However, if a file contains 999,999 or more lines, then the compiler will flag the file as too large and will list only the last 6 digits in the source line numbers for that file.

If the input records contain printer control characters, or %SKIP or %PAGE statements, the lines are spaced accordingly. Use %NOPRINT and %PRINT statements to stop and restart the printing of the listing.

If you specify the MACRO option, the source listing shows the included text in place of the %INCLUDE statements in the primary input data set.

**Statement nesting level**

If you specify the NEST option, the block level and the DO-level are printed to the right of the statement or line number under the headings LEV and NT respectively, as in the following example:

<table>
<thead>
<tr>
<th>Line</th>
<th>File</th>
<th>LV</th>
<th>NT</th>
</tr>
</thead>
<tbody>
<tr>
<td>1.0</td>
<td>A:</td>
<td></td>
<td></td>
</tr>
<tr>
<td>2.0</td>
<td></td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>3.0</td>
<td></td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>4.0</td>
<td></td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>5.0</td>
<td></td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>6.0</td>
<td></td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>7.0</td>
<td></td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>8.0</td>
<td></td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>9.0</td>
<td></td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>10.0</td>
<td></td>
<td>3</td>
<td></td>
</tr>
<tr>
<td>11.0</td>
<td></td>
<td>3</td>
<td></td>
</tr>
<tr>
<td>12.0</td>
<td></td>
<td>3</td>
<td></td>
</tr>
<tr>
<td>13.0</td>
<td></td>
<td>1</td>
<td></td>
</tr>
</tbody>
</table>

**ATTRIBUTE and cross-reference table**

If you specify the ATTRIBUTES option, the compiler prints an attribute table containing a list of the identifiers in the source program together with their declared and default attributes.

If you specify the XREF option, the compiler prints a cross-reference table containing a list of the identifiers in the source program together with the file and line numbers of the statements in which they appear.

If you specify both ATTRIBUTES and XREF, the two tables are combined. In these tables, if you explicitly declare an identifier, the compiler will list file number and
line number of its DECLARE. Contextually declared variables are marked by ++++, and other implicitly declared variables are marked by *****.

Attribute table
The compiler never includes the attributes INTERNAL and REAL. You can assume them unless the respective conflicting attributes, EXTERNAL and COMPLEX, appear.

For a file identifier, the attribute FILE always appears, and the attribute EXTERNAL appears if it applies; otherwise, the compiler lists only explicitly declared attributes.

The OPTIONS attribute will not appear unless the ENTRY attribute applies, and then only the following options would appear (as appropriate)
- ASSEMBLER
- COBOL
- FETCHABLE
- FORTRAN
- NODESCRIPTOR
- RETCODE

The compiler prints the dimension attribute for an array first. It prints the bounds as in the array declaration, but expressions are replaced by asterisks unless they have been reduced by the compiler to a constant, in which case the value of the constant is shown.

For a character string, a bit string, a graphic string, or an area variable, the compiler prints the length, as in the declaration, but expressions are replaced by asterisks unless they have been reduced by the compiler to a constant, in which case the value of the constant is shown.

Cross-reference table
If you combine the cross-reference table with the attribute table, the list of attributes for a name is identified by file number and line number. An identifier appears in the Sets: part of the cross-reference table if it is:
- The target of an assignment statement
- Used as a loop control variable in DO loops
- Used in the SET option of an ALLOCATE or LOCATE statement
- Used in the REPLY option of a DISPLAY statement

If you specify ATTRIBUTES and XREF, the two tables are combined.

If there are unreferenced identifiers, they are displayed in a separate table.

Aggregate length table
An aggregate length table is obtained by using the AGGREGATE option. The table includes structures but not arrays that have non-constant extents, but the sizes and offsets of elements within structures with non-constant extents may be inaccurate or specified as *. For the aggregates listed, the table contains the following information:
- Where the aggregate is declared.
- The name of the aggregate and each element within the aggregate.
- The byte offset of each element from the beginning of the aggregate.
- The length of each element.
• The total length of each aggregate, structure, and substructure.
• The total number of dimensions for each element.

Be careful when interpreting the data offsets indicated in the data length table. An odd offset does not necessarily represent a data element without halfword, fullword, or even double word alignment. If you specify or infer the aligned attribute for a structure or its elements, the proper alignment requirements are consistent with respect to other elements in the structure, even though the table does not indicate the proper alignment relative to the beginning of the table.

If there is padding between two structure elements, a /*PADDING*/ comment appears, with appropriate diagnostic information.

**Statement offset addresses**

If the LIST compile option is used, the compiler includes a pseudo-assembler listing in the compiler listing. This listing includes, for each instruction, an offset whose meaning depends on the setting of the BLKOFF compiler option:
• Under the BLKOFF option, this offset is the offset of the instruction from the primary entry point for the function or subroutine to which it belongs. Thus under this option, the offsets are reset with each new block.
• Under the NOBLKOFF option, this offset is the offset of the instruction from the start of the compilation unit. Thus under this option, the offsets are cumulative.

The pseudo-assembler listing also includes, at the end of the code for each block, the offset of the block from the start of the current module (so that the offsets shown for each statement can be translated to either block or module offsets).

These offsets can be used with the offset given in a runtime error message to determine the statement to which the message applies.

The OFFSET option produces a table that gives for each statement, the offset of the first instruction belonging to that statement.

In the example shown in Figure 2 on page 100, the message indicates that the condition was raised at offset +98 from the SUB1 entry. The compiler listing excerpt shows this offset associated with line number 8. The runtime output from this erroneous statement is shown if Figure 3 on page 100.
Compiler Source

Line File

1.0 TheMain: proc options ( main );
2.0 call sub1();
3.0 Sub1: proc;
4.0 dcl (i, j) fixed bin(31);
5.0 i = 0; j = 0;
6.0 j = j / i;
7.0 put skip data( j );
8.0 end Sub1;
9.0 end TheMain;
10.0...

OFFSET OBJECT CODE LINE# FILE# PSEUDO ASSEMBLY LISTING
000000 000002 | THEMAIN DS 0D
000004 01C3C5C5 36(,r15)
000008 000000B0 CEE eyecatcher
00000C 000001F8 DSA size
000010 47F0 F001 000002 | B 1(,r15)

OFFSET OBJECT CODE LINE# FILE# PSEUDO ASSEMBLY LISTING
000000 000002 | THEMAIN DS 0D
000004 01C3C5C5 36(,r15)
000008 000000B0 CEE eyecatcher
00000C 000001F8 DSA size
00000C 00000190 =A(PPA1-THEMAIN)
000010 47F0 F001 000002 | B 1(,r15)

000008 5020 D0B8 000007 | ST r2,I(,r13,184)
00000A 1842 000007 | SRDA r4,32
00000C 5040 D0BC 000007 | ST r4,J(,r13,188)
000010 5800 D0B8 000008 | L r0,I(,r13,184)
000014 0E40 0020 000008 | DR r4,r0
000016 1805 000008 | LR r0,r5
000018 5000 D0BC 000008 | ST r0,J(,r13,188)
00001A 4100 D0C0 000009 | LA r0_.temp1(,r13,192)
00001C 5000 D130 000009 | ST r0_.temp2(,r13,304)
00001E A708 5AB8 000009 | LHI r0,H’23176’
000022 4000 D0EC 000009 | STH r0_.temp1(,r13,236)
000024 5800 6004 000009 | L r0SYSPRINT(,r6,4)
000028 5000 D12C 000009 | ST r0_.temp2(,r13,300)
00002A 4100 0001 000009 | LA r0.1
00002C 5000 D0C0 000009 | ST r0_.temp1(,r13,192)
00002E 4100 D128 000009 | LA r0_.temp2(,r13,296)

Figure 2. Finding statement number (compiler listing example)

Message:

IBM0301S ONCODE=320 The ZERO DIVIDE condition was raised. From entry point SUB1 at compile unit offset +00000098 at entry offset +00000098 at address 0EB00938.

Figure 3. Finding statement number (runtime message example)
Entry offsets given in dump and ON-unit SNAP error messages can be compared with this table and the erroneous statement discovered. The statement is identified by finding the section of the table that relates to the block named in the message and then finding the largest offset less than or equal to the offset in the message. The statement number associated with this offset is the one needed.

Storage offset listing

If the MAP compile option is used, the compiler includes a storage offset listing in the compiler listing. This listing gives the location in storage of the following level-1 variables if they are used in the program:

- AUTOMATIC
- CONTROLLED except for PARAMETERs
- STATIC except for ENTRY CONSTANTs that are not FETCHABLE

The listing may also include some compiler generated temporaries.

For an AUTOMATIC variable with adjustable extents, there will be two entries in this table:

- an entry with '_addr' prefixing the variable name - this entry gives the location of the address of the variable
- an entry with '_desc' prefixing the variable name - this entry gives the location of the address of the variable’s descriptor

For STATIC and CONTROLLED variables, the storage location will depend on the RENT/NORENT compiler option, and if the NORENT option is in effect, the location of CONTROLLED variables will also depend on the WRITABLE/NOWRITABLE compiler option.

The first column in the Storage Offset Listing is labeled IDENTIFIER and holds the name of the variable whose location appears in the fourth column.

The second column in the Storage Offset Listing is labeled DEFINITION and holds a string in the format "B-F:N" where

- B is the number of the block where the variable is declared
- F is the number of the source file where the variable is declared
- N is the number of the source line where the variable is declared in that source file

You can find the name of the block corresponding to this block number in the Block Name List which will proceed the Storage Offset Listing (and the Pseudo Assembly Listing, if any)

You can find the name of the file corresponding to this file number in the File Reference Table which will appear very near the end of the entire compilation listing.

The third column in the Storage Offset Listing is labeled ATTRIBUTES and indicates the storage class of the variable.

The fourth column in the Storage Offset Listing is unlabeled and tells how to find the location of the variable.

This storage offset listing is sorted by block and by variable name, and it also includes only user variables. However, specifying the MAP option also causes the compiler to produce
• a "static map" which lists all STATIC variables but sorted by hex offset
• an "automatic map" which lists, for each block, all AUTOMATIC variables but sorted by hex offset

The PL/L language’s mapping rules may require that a structure be offset by up to 8 bytes from where it would seem to start. For example, consider the AUTOMATIC structure A declared as

```plaintext
dcl
  1 A,
  2 B char(2),
  2 C fixed bin(31);
```

Since C must be aligned on a 4-byte boundary, 2 bytes of padding will be needed for this structure. However, PL/I places those 2 bytes not after B, but before B. These 2 bytes of “padding” before a structure starts are referred to as the “hang bytes” for the structure.

These hang bytes will also be reflected in the "automatic map" generated by the compiler. While the "storage offset listing" would show the offset and length for A without including its hang bytes

```
A Class = automatic, Location = 186 : 0xBA(r13), Length = 6
```

the "automatic map" would show the offset and length for A with its hang bytes included

```
OFFSET (HEX) LENGTH (HEX) NAME
98  8  #MX_TEMP1
A0 18  _Sfi
B8  8  A
```

**File reference table**

The file reference table consists of three columns which list the following information about the files read during the compile:

• the number assigned by the compiler to the file
• the included-from data for the file
• the name of the file

The first entry in the included-from column is blank because the first file listed is the source file. Subsequent entries in this column show the line number of the include statement followed by a period and the file number of the source file containing the include.

If the file is a member of a PDS or PDSE, the file name lists the fully qualified dataset name and the member name.

If the file is included via a subsystem (such as Librarian), then the file name will have the form `DD:ddname(member)`, where

• `ddname` is the ddname specified on the `%INCLUDE` statement (or SYSLIB if no ddname was specified)
• `member` is the member name specified on the `%INCLUDE` statement.

**Messages and return codes**

If the preprocessor or the compiler detects an error, or the possibility of an error, messages are generated. Messages generated by the preprocessor appear in the
listing immediately after the listing of the statements processed by the 
preprocessor. You can generate your own messages in the preprocessing stage by 
use of the %NOTE statement. Such messages might be used to show how many 
times a particular replacement had been made. Messages generated by the 
compiler appear at the end of the listing.

For compilations that produce no messages, the compiler will include a line saying 
"no compiler messages" where the compiler messages would have been listed.

Messages are displayed in the following format:

PPPnnnnI X

where PPP is the prefix identifying the origin of the message (for example, IBM 
indicates the PL/I compiler), nnnn is the 4-digit message number, and X identifies 
the severity code. All messages are graded according to their severity, and the 
severity codes are I, W, E, S, and U.

For every compilation job or job step, the compiler generates a return code that 
indicates to the operating system the degree of success or failure it achieved. For 
z/OS, this code appears in the end-of-step message that follows the listing of the job 
control statements and job scheduler messages for each step.

Table 5 provides an explanation of the severity codes and the comparable return 
code for each:

<table>
<thead>
<tr>
<th>Severity Code</th>
<th>Return Code</th>
<th>Message Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>I</td>
<td>0000</td>
<td>Informational</td>
<td>The compiled program should run correctly. The compiler might inform you of a possible inefficiency in your code or some other condition of interest.</td>
</tr>
<tr>
<td>W</td>
<td>0004</td>
<td>Warning</td>
<td>A statement might be in error (warning) even though it is syntactically valid. The compiled program should run correctly, but it might produce different results than expected or be significantly inefficient.</td>
</tr>
<tr>
<td>E</td>
<td>0008</td>
<td>Error</td>
<td>A simple error fixed by the compiler. The compiled program should run correctly, but it might produce different results than expected.</td>
</tr>
<tr>
<td>S</td>
<td>0012</td>
<td>Severe</td>
<td>An error not fixed by the compiler. If the program is compiled and an object module is produced, it should not be used.</td>
</tr>
<tr>
<td>U</td>
<td>0016</td>
<td>Unrecoverable</td>
<td>An error that forces termination of the compilation. An object module is not successfully created.</td>
</tr>
</tbody>
</table>

Note: Compiler messages are printed in groups according to these severity levels.

The compiler lists only messages that have a severity equal to or greater than that 
specified by the FLAG option, as shown in Table 6.

Table 6. Using the FLAG option to select the lowest message severity listed

<table>
<thead>
<tr>
<th>Type of Message</th>
<th>Option</th>
</tr>
</thead>
<tbody>
<tr>
<td>Information</td>
<td>FLAG(I)</td>
</tr>
<tr>
<td>Warning</td>
<td>FLAG(W)</td>
</tr>
</tbody>
</table>
Table 6. Using the FLAG option to select the lowest message severity listed (continued)

<table>
<thead>
<tr>
<th>Type of Message</th>
<th>Option</th>
</tr>
</thead>
<tbody>
<tr>
<td>Error</td>
<td>FLAG(E)</td>
</tr>
<tr>
<td>Severe Error</td>
<td>FLAG(S)</td>
</tr>
<tr>
<td>Unrecoverable Error</td>
<td>Always listed</td>
</tr>
</tbody>
</table>

The text of each message, an explanation, and any recommended programmer response, are given in *Enterprise PL/I Messages and Codes*.

**Example**

This example of the compiler listing is generated when the compiler compiles the following `msgsumm` program with these options: PP(SQL,MACRO,CICS), SOURCE, FLAG(I), INSOURCE, MSGSUMMARY(XREF).

```pli
msgsumm: proc;

    exec sql include sqlca;

    exec cics what now;

    exec cics not this;

    %dcl z0 fixed bin;
    %dcl z1 fixed dec;
end;
```

**Note:** The program is intentionally incorrect. Because the MSGSUMMARY option is specified, the compiler includes the Summary of Messages section at the end of the listing. This section also includes the line numbers associated with each of the messages in the summary because the XREF suboption to the MSGSUMMARY option is specified.
Chapter 1. Using compiler options and facilities  105
MacRO (Built: 20141120) Preprocessor Source

Line.File
1.0 
2.0 msgsumm: proc;
3.0
4.0 /*$*$*$*/
4.0 exec sql include sqlca
4.0 $*$*$*/
4.0 DO;
4.0 DCL
4.0 1 SQLCA ,
4.0 2 SQLCAID CHAR(8),
4.0 2 SQLCBC FIXED BIN(31),
4.0 2 SQLCODE FIXED BIN(31),
4.0 2 SQLERRM CHAR(70) VAR,
4.0 2 SQLERRP CHAR(8),
4.0 2 SQLERRD(6) FIXED BIN(31),
4.0 2 SQLWARN,
4.0 3 SQLWARN0 CHAR(1),
4.0 3 SQLWARN1 CHAR(1),
4.0 3 SQLWARN2 CHAR(1),
4.0 3 SQLWARN3 CHAR(1),
4.0 3 SQLWARN4 CHAR(1),
4.0 3 SQLWARN5 CHAR(1),
4.0 3 SQLWARN6 CHAR(1),
4.0 3 SQLWARN7 CHAR(1),
4.0 2 SQLEXT,
4.0 3 SQLWARN8 CHAR(1),
4.0 3 SQLWARN9 CHAR(1),
4.0 3 SQLWARNA CHAR(1),
4.0 3 SQLSTATE CHAR(5);
4.0 END;
5.0
6.0 exec cics what now;
7.0 exec cics not this;
8.0
9.0 %dcl z0 fixed bin;
10.0 %dcl z1 fixed dec;
11.0 end;

Compiler listing example (continued)
MACRO Messages

<table>
<thead>
<tr>
<th>Message</th>
<th>Line</th>
<th>File</th>
<th>Message Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>IBM3552I E 9.0</td>
<td>The statement element BIN is invalid. The statement will be ignored.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>IBM3552I E 10.0</td>
<td>The statement element DEC is invalid. The statement will be ignored.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>IBM3258I W 9.0</td>
<td>Missing ; assumed before BIN.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>IBM3258I W 10.0</td>
<td>Missing ; assumed before DEC.</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

CICS (Built:20140829) Preprocessor Source

```plaintext
1.0 MSGSUMM: PROC;
2.0 /
2.0 /*$*$*$
2.0 exec sql include sqlca
2.0 $*$*$$/
2.0 DQ;
2.0 DCL
4.0 1 SQLCA,
4.0 2 SQLCAID CHAR(8),
4.0 2 SQLABC FIXED BIN(31),
4.0 2 SQLCODE FIXED BIN(31),
4.0 2 SQLERRM CHAR(70) VAR,
4.0 2 SQLERRP CHAR(8),
4.0 2 SQLERRD(6) FIXED BIN(31),
4.0 2 SQLWARN,
4.0 3 SQLWARN0 CHAR(1),
4.0 3 SQLWARN1 CHAR(1),
4.0 3 SQLWARN2 CHAR(1),
4.0 3 SQLWARN3 CHAR(1),
4.0 3 SQLWARN4 CHAR(1),
4.0 3 SQLWARN5 CHAR(1),
4.0 3 SQLWARN6 CHAR(1),
4.0 3 SQLWARN7 CHAR(1),
4.0 2 SQLEXT,
4.0 3 SQLWARN8 CHAR(1),
4.0 3 SQLWARN9 CHAR(1),
4.0 3 SQLWARNA CHAR(1),
4.0 3 SQLSTATE CHAR(5);
4.0 END;
4.0 END;
```

Compiler listing example (continued)
### CICS Messages

<table>
<thead>
<tr>
<th>Message</th>
<th>Line.File</th>
<th>Message Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>DFH7059I S 6.0</td>
<td>WHAT COMMAND IS NOT VALID AND IS NOT TRANSLATED.</td>
<td></td>
</tr>
<tr>
<td>DFH7059I S 7.0</td>
<td>WHAT COMMAND IS NOT VALID AND IS NOT TRANSLATED.</td>
<td></td>
</tr>
</tbody>
</table>

### No Compiler Messages

### File Reference Table

<table>
<thead>
<tr>
<th>File Included From Name</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>DD:SYSIN</td>
<td>0</td>
</tr>
</tbody>
</table>

### Summary of Messages

<table>
<thead>
<tr>
<th>Component</th>
<th>Message</th>
<th>Total</th>
<th>Default Message Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SQL</td>
<td>IBM3250I W</td>
<td>1</td>
<td>DSNH053I DSNHPSRV NO SQL STATEMENTS WERE FOUND</td>
</tr>
<tr>
<td>MACRO</td>
<td>IBM3552I W</td>
<td>2</td>
<td>The statement element %1 is invalid. The statement will be ignored.</td>
</tr>
<tr>
<td>MACRO</td>
<td>IBM3258I W</td>
<td>2</td>
<td>Missing %1 assumed before %2.</td>
</tr>
<tr>
<td>CICS</td>
<td>IBM3750I S</td>
<td>2</td>
<td>DFH7059I S WHAT COMMAND IS NOT VALID AND IS NOT TRANSLATED.</td>
</tr>
</tbody>
</table>

### Component Return Code Messages (Total/Suppressed) Time

<table>
<thead>
<tr>
<th>Component</th>
<th>Return Code</th>
<th>Messages (Total/Suppressed)</th>
<th>Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>SQL</td>
<td>4</td>
<td>1 / 0</td>
<td>0 secs</td>
</tr>
<tr>
<td>MACRO</td>
<td>8</td>
<td>4 / 0</td>
<td>0 secs</td>
</tr>
<tr>
<td>CICS</td>
<td>12</td>
<td>2 / 0</td>
<td>0 secs</td>
</tr>
<tr>
<td>Compiler</td>
<td>0</td>
<td>0 / 0</td>
<td>0 secs</td>
</tr>
</tbody>
</table>

End of compilation

---

*Compiler listing example (continued)*
Chapter 2. PL/I preprocessors

The PL/I compiler allows you to select one or more of the integrated preprocessors as required for use in your program. You can select the include preprocessor, the macro preprocessor, the SQL preprocessor, or the CICS preprocessor and the order in which you would like them to be called.

- The include preprocessor processes special include directives and incorporates external source files.
- The macro preprocessor, based on %statements and macros, modifies your source program.
- The SQL preprocessor modifies your source program and translates EXEC SQL statements into PL/I statements.
- The CICS preprocessor modifies your source program and translates EXEC CICS statements into PL/I statements.

Each preprocessor supports a number of options to allow you to tailor the processing to your needs.

The three compile-time options MDECK, INSOURCE, and SYNTAX are meaningful only when you also specify the PP option. For more information about these options, see "MDECK" on page 54, "INSOURCE" on page 43, and "SYNTAX" on page 80.

Include preprocessor

The include preprocessor allows you to incorporate external source files into your programs by using include directives other than the PL/I directive %INCLUDE.

The following syntax diagram illustrates the options supported by the INCLUDE preprocessor:

```
PP ( INCLUDE ( 'ID(<directive>)' ) )
```

**ID** Specifies the name of the include directive. Any line that starts with this directive as the first set of nonblank characters is treated as an include directive.

The specified directive must be followed by one or more blanks, an include member name, and finally an optional semicolon. Syntax for ddname(membername) is not supported.

In the following example, the first include directive is valid and the second one is not:

```
++include payroll
++include syslib(payroll)
```

This first example causes all lines that start with -INC (and possibly preceding blanks) to be treated as include directives:

```
pp( include( 'id(-inc)'))
```

This second example causes all lines that start with ++INCLUDE (and possibly preceding blanks) to be treated as include directives:
Macro preprocessor

Macros allow you to write commonly used PL/I code in a way that hides implementation details and the data that is manipulated, and exposes only the operations. In contrast with a generalized subroutine, macros allow generation of only the code that is needed for each individual use.

The macro preprocessing facilities of the compiler are described in the *PL/I Language Reference*.

You can invoke the macro preprocessor by specifying either the MACRO option or the PP(MACRO) option. You can specify PP(MACRO) without any options or with options from the list below.

The defaults for all these options cause the macro preprocessor to behave the same as the OS PL/I V2R3 macro preprocessor.

If options are specified, the list must be enclosed in quotes (single or double, as long as they match); for example, to specify the FIXED(BINARY) option, you must specify PP(MACRO('FIXED(BINARY)')).

If you want to specify more than one option, you must separate them with a comma and/or one or more blanks. For example, to specify the CASE(ASIS) and RESCAN(UPPER) options, you can specify PP(MACRO('CASE(ASIS) RESCAN(UPPER)')) or PP(MACRO("CASE(ASIS),RESCAN(UPPER)")). You may specify the options in any order.

**Macro preprocessor options**

This section describes the options that the macro preprocessor supports.

**CASE**

This option specifies if the preprocessor should convert the input text to uppercase.

- **ASIS** the input text is left "as is".
- **UPPER** the input text is to be converted to upper case.

**DBCS**

This option specifies if the preprocessor should normalize DBCS during text replacement.

- **INEXACT** the input text is left "as is", and the preprocessor will treat <kk.B> and <kk>B as different names.
**INEXACT**

the input text is "normalized", and the preprocessor will treat <kk.B> and<br> <kk>B as two versions of the same name.

**DEPRECATE**

This option flags the usage of macro procedures that you want to deprecate with<br> error messages.

**ENTRY**

Flags any usage of a macro procedure with name *entry-name*.

**DEPRECATENEXT**

This option flags the usage of macro procedures that you want to deprecate with<br> warning messages.

**ENTRY**

Flags any usage of a macro procedure with name *entry-name*.

**FIXED**

This option specifies the default base for FIXED variables.

**DECIMAL**

FIXED variables will have the attributes REAL FIXED DEC(5).

**BINARY**

FIXED variables will have the attributes REAL SIGNED FIXED BIN(31).

**INCONLY**

The INCONLY option specifies that the preprocessor should process only<br> %INCLUDE and %XINCLUDE statements.

The NOINCONLY option specifies that the preprocessor should process all<br> preprocessor statements and not only %INCLUDE and %XINCLUDE statements.
When the INCONLY option is in effect, you may use neither INCLUDE nor XINCLUDE as a macro:

- procedure name
- statement label
- variable name

The INCONLY option and the NOINCONLY option are mutually exclusive.

For compatibility, the default is NOINCONLY.

**NAMEPREFIX**

The NAMEPREFIX option specifies that the names of preprocessor procedures and variables must start with the specified character.

The NONAMEPREFIX option specifies that the names of preprocessor procedures and variables are not required to start with one particular character.

\[ \text{NAMEPREFIX} \text{character} \]

The character should be specified "as is" and should not be enclosed in quotes.

The default is NONAMEPREFIX.

**RESCAN**

This option specifies how the preprocessor should handle the case of identifiers when rescanning text.

\[ \text{RESCAN (UPPER)} \]

**UPPER**

- rescans will not be case-sensitive.

**ASIS**

- rescans will be case-sensitive.

To see the effect of this option, consider the following code fragment:

```pli
%dcl eins char ext;
%dcl text char ext;
%eins = 'zwei';
%text = 'EINS';
display( text );
%text = 'eins';
display( text );
```

When compiled with PP(MACRO('RESCAN(ASIS)')), in the second display statement, the value of `text` is replaced by `eins`, but no further replacement occurs since under RESCAN(ASIS), `eins` does not match the macro variable `eins` since the former is left as is while the latter is uppercased. Hence the following text would be generated...
But when compiled with PP(MACRO('RESCAN(UPPER)')), in the second display statement, the value of text is replaced by eins, but further replacement does occur since under RESCAN(UPPER), eins does match the macro variable eins since both are uppercased. Hence the following text would be generated

DISPLAY( zwei );
DISPLAY( eins );

In short: RESCAN(UPPER) ignores case while RESCAN(ASIS) does not.

**Macro preprocessor example**

A simple example of the use of the preprocessor to produce a source deck is shown in Figure 5 on page 114. According to the value assigned to the preprocessor variable USE, the source statements will represent either a subroutine (CITYSUB) or a function (CITYFUN).

The DSNAMES used for SYSPUNCH specifies a source program library on which the preprocessor output will be placed. Normally compilation would continue and the preprocessor output would be compiled.
In general, the coding for your PL/I program will be the same whether or not you want it to access a DB2 database. However, to retrieve, update, insert, and delete DB2 data and use other DB2 services, you must use SQL statements. You can use dynamic and static EXEC SQL statements in PL/I applications.

To communicate with DB2, you need to do the following:

- Code any SQL statements you need, delimiting them with **EXEC SQL**.
- Use the DB2 precompiler or, if using DB2 for z/OS Version 9 Release 1 or later, compile with the PL/I PP(SQL()) compiler option.

Before you can take advantage of EXEC SQL support, you must have authority to access a DB2 system. Contact your local DB2 Database Administrator for your authorization.
Programming and compilation considerations

When you use the PL/I SQL Preprocessor, the PL/I compiler handles your source program containing embedded SQL statements at compile time, without your having to use a separate precompile step. Although the use of a separate precompile step continues to be supported, use of the PL/I SQL Preprocessor is recommended.

Interactive debugging with Debug Tool is enhanced when you use the preprocessor because you see only the SQL statements while debugging (and not the generated PL/I source). However, you must have DB2 for z/OS Version 9 Release 1 or later to use the SQL preprocessor.

Using the preprocessor lifts some of the DB2 precompiler’s restrictions on SQL programs. When you process SQL statements with the preprocessor, you can:

- Use nested SQL INCLUDE statements.
- Use fully-qualified names for structured host variables.
- Include SQL statements at any level of a nested PL/I program, instead of in only the top-level source file.
- Use the SQL TYPE attribute anywhere you can specify a PL/I data type. All such attributes can be factored and can be used in structure elements, in arrays, and with any storage class like BASED.
- Use variables declared with the LIKE attribute as host variables.

All PL/I statements must be syntactically correct, because the SQL preprocessor scans the source looking for EXEC SQL statements, DECLARE statements, and statements that delimit blocks of declarations. If a statement is coded incorrectly, this might mislead the preprocessor when looking for the END statement to a BEGIN, DO, PACKAGE, PROCEDURE, or SELECT statement, and that can cause the preprocessor to be unable to resolve some host variable references correctly. To help identify such incorrect code, the SQL preprocessor flags errors as follows:

- Left parentheses without matching right parentheses
- SELECT statements that do not end with a semicolon
- IF statements that do not contain a THEN keyword
- Statements that start with an invalid symbol

If the source uses %INCLUDE or other macro statements, invoke the MACRO preprocessor before the SQL preprocessor.

The SQL preprocessor supports DBCS in the same manner as the PL/I compiler does. When the GRAPHIC PL/I compiler option is in effect, some source language elements can be written using DBCS and SBCS characters. In particular, you can use DBCS characters in the source program in following places:

- Inside comments
- As part of statement labels and identifiers
- In G or M literals

The following restrictions apply to the use of PL/I built-in functions, compiler options, and statements when you program and compile SQL statements:

- When EXEC SQL statements are translated to PL/I, the following built-in functions might be included in the generated code. If you use any of the following built-in functions as the names of elements in structures, you must also explicitly declare them as BUILTIN:
### Programming and compilation considerations

- ADDR
- LENGTH
- MAXLENGTH
- PTRVALUE
- SYSNULL

- You must not use PL/I type functions, such as BIND(: t, p ;), in EXEC SQL statements.
- When compiling with the preprocessor, you must not use these compiler options:
  - DFT(ASCII)
  - DFT(IEEE)
- Do not use DECLARE STATEMENT statements in SQL queries, because the PL/I preprocessor always ignores these statements.

Compiling with the SQL preprocessor option generates a DB2 database request module (DBRM) along with the usual PL/I compiler outputs such as object module and listing. As input to the DB2 bind process, the DBRM data set contains information about the SQL statements and host variables in the program. Not all of the information in the DBRM is important in terms of the bind or runtime processing, however. For example, if the HOST value in the DBRM specifies a language other than PL/I, there is no reason to be concerned. All this means that the other language is selected as the installation default for the HOST value, which does not effect the bind or runtime processing of your program.

If the EMPTYDBRM option is in effect and if the source meets one of the following conditions, the preprocessor issues a message indicating that no statements required translation and does not create a DBRM:

- It contains no EXEC SQL statements.
- It contains only EXEC SQL INCLUDE statements other than EXEC SQL INCLUDE SQLCA and EXEC SQL INCLUDE SQLDA.

The PL/I compiler listing includes the error diagnostics (such as syntax errors in the SQL statements) that the preprocessor generates. The listing of the EXEC SQL statement is displayed in a readable format that is similar to the original source.

To use the preprocessor, you need to do the following things:

- Specify the following option when you compile your program:
  ```pp(SQL('options'))
  ```
  This compiler option indicates that you want the compiler to invoke the integrated SQL preprocessor. Specify a list of SQL processing options in the parenthesis after the SQL keyword. The options can be separated by a comma or by a space and the list of options must be enclosed in quotation marks (single or double, as long as they match).

  For example, `PP(SQL('DATE(USA),TIME(USA)'))` tells the preprocessor to use the USA format for both DATE and TIME data types.

  In addition, for LOB support you must specify the following option:
  ```LIMITS( FIXEDBIN(31,63) FIXEDDEC(15,31) )
  ```

  An alternative way to specify SQL preprocessor options is to use the PPSQL compiler option. For information about how to use it, see "PPSQL" on page 63.

- Include DD statements for the following data sets in the JCL for your compile step:
Programming and compilation considerations

- DB2 load library (prefix.SDSNNLOAD)
  The SQL preprocessor calls DB2 modules to do the SQL statement processing. You therefore need to include the name of the DB2 load library data set in the STEPLIB concatenation for the compile step.

- Library for SQL INCLUDE statements
  If your program contains SQL INCLUDE member-name statements that specify secondary input to the source program, you need to include the name of the data set that contains member-name in the SYSLIB concatenation for the compile step.

- DBRM library
  The compilation of the PL/I program generates a DB2 database request module (DBRM), and the DBRMLIB DD statement is required to designate the data set to which the DBRM is written.
  The DBRMLIB DD statement must name a data set that can be opened and closed more than once during the compilation.
  For example, you might have the following lines in your JCL:
    
    //STEPLIB DD DSN=DSNA10.SDSNLOAD,DISP=SHR
    //SYSLIB DD DSN=PAYROLL.MONTHLY.INCLUDE,DISP=SHR
    //DBRMLIB DD DSN=PAYROLL.MONTHLY.DBRMLIB.DATA(MASTER),DISP=SHR

SQL preprocessor options

When you specify SQL preprocessor options, the list of options must be enclosed in a pair of quotation marks. For example, to specify the CCSID0 option, you must specify PP(SQL('CCSID0')).

Two groups of options can be passed to the SQL preprocessor: those handled by the PL/I SQL preprocessor and those handled by the DB2 coprocessor. You must specify these options in the options string of the PP(SQL('option-list')) option, and the options can be intermingled in the options string.

Table 7 lists the PL/I SQL preprocessor options with their abbreviations (if any) and their IBM-supplied default values. This table uses a vertical bar (|) to separate mutually exclusive options, and brackets ([ ]) to indicate that you can sometimes omit the enclosed option.

For more information about DB2 coprocessor options, see DB2 for z/OS Application Programming and SQL Guide.

Table 7. SQL preprocessor options and IBM-supplied defaults

<table>
<thead>
<tr>
<th>SQL Preprocessor Option</th>
<th>Abbreviated Name</th>
<th>z/OS Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>CCSID0</td>
<td>NOCCSID0</td>
<td>-</td>
</tr>
<tr>
<td>CODEPAGE</td>
<td>NOCODEPAGE</td>
<td>-</td>
</tr>
<tr>
<td>DEPRECATE(STMT([EXPLAIN</td>
<td>GRANT</td>
<td>REVOKE</td>
</tr>
<tr>
<td>EMPTYDBRM</td>
<td>NOENTRYDBRM</td>
<td>-</td>
</tr>
<tr>
<td>INCONLY</td>
<td>NOINCONLY</td>
<td>-</td>
</tr>
<tr>
<td>WARNDECP</td>
<td>NOWARNDECP</td>
<td>-</td>
</tr>
</tbody>
</table>
SQL preprocessor options

CCSID0
The CCSID0 option specifies that no host variable other than WIDECHAR is to be assigned a CCSID value by the PL/I SQL preprocessor.

The NOCCSID0 option allows host variables to be assigned a CCSID value by the PL/I SQL preprocessor.

If your program updates FOR BIT DATA columns with a data type that is not BIT data, you will want to choose CCSID0. CCSID0 tells DB2 that the host variable is not associated with a CCSID, allowing the assignment to be made. Otherwise, if a host variable that is associated with a CCSID that is not BIT data is assigned to a FOR BIT DATA column, a DB2 error occurs.

WIDECHAR is always assigned a CCSID value of 1200.

For compatibility with older PL/I programs which used the DB2 precompiler, CCSID0 should be chosen.

CCSID0 and NOCCSID0 are mutually exclusive options.

The default is CCSID0.

CODEPAGE
When the CODEPAGE option is in effect, the compiler CODEPAGE option is always used as the CCSID for SQL host variables of character type.

When the NOCODEPAGE option is in effect, the compiler CODEPAGE option is used as the CCSID for SQL host variables of character type only if the SQL preprocessor option NOCCSID0 is also in effect.

The default is NOCODEPAGE.

DEPRECATE
This option indicates that the preprocessor flags the specified statements as deprecated.
SQL preprocessor options

STMT
Specifies a list of statements that the preprocessor should flag as deprecated. The list can be empty.

EXPLAIN
The EXPLAIN SQL statement.

GRANT
The GRANT SQL statement.

REVOKE
The REVOKE SQL statement.

SET_CURRENT_SQLID
The SET CURRENT SQLID SQL statement.

The default is DEPRECATE(STMT()).

EMPTYDBRM
The EMPTYDBRM option specifies that the SQL preprocessor always creates a DBRM even if the code contains no EXEC SQL statements. However, the SQL preprocessor does not create a DBRM when the INCONLY option is invoked.

The NOENTRYDBRM option specifies that the SQL preprocessor does not create a DBRM.

INCONLY
The INCONLY option specifies that the SQL preprocessor should process only EXEC SQL INCLUDE statements except for includes of the SQLCA and SQLDA. No code is generated by the SQL preprocessor when this option is in effect.

The NOINCONLY option specifies that the SQL preprocessor should process all statements and not only EXEC SQL INCLUDE statements.

When you specify the INCONLY option, the compiler does not produce the SQL options listing, because all other options are ignored under INCONLY.

The INCONLY option and the NOINCONLY option are mutually exclusive.

For compatibility, the default is NOINCONLY.

WARNDECP
When the WARNDECP option is in effect, the preprocessor issues a warning message if the compilation uses the DB2-supplied DSNHDECP module.

When the NOWARNDECP option is in effect, no warning message is issued.
The default is NOWARNDECP.

**PL/I-specific notes for SQL processor options**

When you specify the following SQL processor options for the PL/I compiler, the following rules apply:

**FLOAT**

An error message is issued if the FLOAT option is different from the PL/I compiler's DEFAULT(HEXADEC | IEEE) option.

**ONEPASS | TWOPASS**

Under the option ONEPASS, you must declare all host variables before they are used by any SQL statement.

**STDSQL**

Under the option STDSQL(YES), you must declare all host variables between SQL BEGIN DECLARE SECTION and SQL END DECLARE SECTION statements.

**Coding SQL statements in PL/I applications**

You can code SQL statements in your PL/I applications using the language defined in *DB2 UDB for z/OS SQL Reference*. Specific requirements for your SQL code are described in the sections that follow.

**Defining the SQL communications area**

A PL/I program that contains SQL statements must include either an SQLCODE variable (if the STDSQL(86) preprocessor option is used) or an SQL communications area (SQLCA). As shown in Figure 6 on page 121, part of an SQLCA consists of an SQLCODE variable and an SQLSTATE variable.

- The SQLCODE value is set by the Database Services after each SQL statement is executed. An application can check the SQLCODE value to determine whether the last SQL statement was successful.
- The SQLSTATE variable can be used as an alternative to the SQLCODE variable when analyzing the result of an SQL statement. Like the SQLCODE variable, the SQLSTATE variable is set by the Database Services after each SQL statement is executed.

The SQLCA declaration should be included by using the EXEC SQL INCLUDE statement:

```sql
exec sql include sqlca;
```

The SQLCA structure must not be defined within an SQL declare section. The scope of the SQLCODE and SQLSTATE declarations must include the scope of all SQL statements in the program.
Defining SQL descriptor areas

The following statements require an SQLDA:

- `PREPARE statement-name INTO descriptor-name FROM host-variable`
- `EXECUTE...USING DESCRIPTOR descriptor-name`
- `FETCH...USING DESCRIPTOR descriptor-name`
- `OPEN...USING DESCRIPTOR descriptor-name`
- `DESCRIBE statement-name INTO descriptor-name`

Unlike the SQLCA, there can be more than one SQLDA in a program, and an SQLDA can have any valid name. An SQLDA should be included by using the `EXEC SQL INCLUDE` statement:

```
exec sql include sqlda;
```

The SQLDA must not be defined within an SQL declare section.

---

Figure 6. The PL/I declaration of SQLCA

```
Dcl
1 Sqlca,
2   sqlcaid char(8), /* Eyecatcher = 'SQLCA ' */
2   sqlcabc fixed binary(31), /* SQLCA size in bytes = 136 */
2   sqlcode fixed binary(31), /* SQL return code */
2   sqlerrmc char(70) var, /* Error message tokens */
2   sqlerrp char(8), /* Diagnostic information */
2   sqlerrd(0:5) fixed binary(31), /* Diagnostic information */
2   sqlwarn, /* Warning flags */
   3   sqlwarn0 char(1),
   3   sqlwarn1 char(1),
   3   sqlwarn2 char(1),
   3   sqlwarn3 char(1),
   3   sqlwarn4 char(1),
   3   sqlwarn5 char(1),
   3   sqlwarn6 char(1),
   3   sqlwarn7 char(1),
2   sqlext,
   3   sqlwarn8 char(1),
   3   sqlwarn9 char(1),
   3   sqlwarna char(1),
   3   sqlstate char(5); /* State corresponding to SQLCODE */
```
Embedding SQL statements

The first statement of your program must be a PROCEDURE or a PACKAGE statement. You can add any SQL statement to your program wherever executable statements can appear. You can also add the following SQL statements in a PACKAGE and outside of any procedure:

- EXEC SQL BEGIN DECLARE SECTION
- EXEC SQL INCLUDE
- EXEC SQL DECLARE if no executable code needs to be generated for the statement
- EXEC SQL END DECLARE SECTION

Each SQL statement must begin with EXEC (or EXECUTE) SQL and end with a semicolon (;).

For example, an UPDATE statement might be coded as follows:

```
exec sql update DSN8A10.DEPT
set Mgrno = :Mgr_Num
where Deptno = :Int_Dept;
```

**Comments:** In addition to SQL statements, comments can be included in embedded SQL statements wherever a blank is allowed.

---

**Figure 7. The PL/I declaration of an SQL descriptor area**
If a comment appears inside a SQL statement, then the forward slash (/) that closes the comment is shown as a greater than sign (>) in the listing. For example, given the following statement:

```sql
exec sql insert into table /* some text */ values(:data);
```

The compiler will show it in the source listing as follows:

```sql
/*$*$*$
 exec sql insert into table /* some text *> values(:data)
$*$*$*/
```

SQL style comments ('--') are supported when embedded in SQL statements.

**Continuation for SQL statements:** The line continuation rules for SQL statements are the same as those for other PL/I statements.

**Including code:** SQL statements or PL/I host variable declaration statements can be included by placing the following SQL statement in the source code. Place it at the point where the statements are to be embedded.

```sql
exec sql include member;
```

**Margins:** SQL statements must be coded in columns $m$ through $n$ where $m$ and $n$ are specified in the MARGINS($m$, $n$) compiler option.

**Names:** Any valid PL/I variable name can be used for a host variable. The length of a host variable name must not exceed the value $n$ specified in the LIMITS(NAME($n$)) compiler option.

**Statement labels:** With the exception of the END DECLARE SECTION statement, and the INCLUDE text-file-name statement, executable SQL statements, like PL/I statements, can have a label prefix.

**WHENEVER statement:** The target for the GOTO clause in an SQL WHENEVER statement must be a label in the PL/I source code and must be within the scope of any SQL statements affected by the WHENEVER statement.

**Using host variables**

All host variables used in SQL statements must be explicitly declared, and all host variables within an SQL statement must be preceded by a colon (:).

Subscripts must not be used in host variable references.

The following topics describe the details of using host variables:

- “Using arrays as host variables”
- “Declaring host variables” on page 124
- “Determining equivalent SQL and PL/I data types” on page 126
- “Determining compatibility of SQL and PL/I data types” on page 129

**Using arrays as host variables:** You can use an array as a host variable only in the following two ways:

- As an array of indicator variables for a host structure
- As an array of host variables when used in any of the following statements:
  - A FETCH statement for a multiple row fetch
  - An INSERT statement with a multiple row insert
  - A multiple row MERGE statement
All such arrays must be one-dimensional, have the CONNECTED attribute, and have constant bounds.

All other use of arrays as host variables is invalid.

Declaring host variables: Host variable declarations can be made at the same place as regular PL/I variable declarations.

Only a subset of valid PL/I declarations are recognized as valid host variable declarations. The preprocessor does not use the data attribute defaults specified in the PL/I DEFAULT statement. The preprocessor ignores the DEFINE ALIAS and DEFINE STRUCT statements, so variables with attributes that are dependent on them may not be used as host variables. If the declaration for a variable is not recognized, any statement that references the variable might result in the message: 'The host variable token ID is not valid'

To use a structure that is declared with LIKE or an element of this structure as a host variable, the declaration for the LIKE object must be visible to the SQL preprocessor. For example, the LIKE object must not be in a %INCLUDE file that has not been included.

You can use restricted expressions in host variable declarations to define the bounds of an array or the length of a string as long as the expression has one of the following forms:

- A prefix operator applied to an expression where the expression can be collapsed to an integer
- An add or sub operator applied to two expressions where both expressions can be collapsed to integers
- A multiply operator applied to two expressions where both expressions can be collapsed to integers
- A reference to a named constant where the reference can be collapsed to an integer
- One of these built-in functions: INDICATORS, HBOUND, LENGTH, and MAXLENGTH
- A number that is an integer

Although you can use a named constant to define the bounds and lengths of a host variable, you cannot use a named constant itself as a host variable except if both of the following conditions apply:

- DB2 allows a simple, unnamed constant at that place in the EXEC SQL statement.
- The named constant has any of these attributes:
  - CHARACTER, in which case the VALUE attribute of the named constant must specify a character string.
  - FIXED, in which case the VALUE attribute of the named constant must specify a decimal number or an expression that can be reduced, with the same restrictions as above, to an integer constant.

Only the names and data attributes of the variables are used by the preprocessor; the alignment, scope, and storage attributes are ignored.

Declaring scalar host variables: You must declare a scalar host variable with one of the following data attributes:
CHARACTER, GRAPHIC, or WIDECHAR
Host variables that are declared with the CHARACTER, GRAPHIC, or WIDECHAR attributes are called string host variables. The following restrictions apply to the string host variable:

- It must have either the NONVARYING or VARYING attribute.
- If it has the VARYING attribute, it must have the NATIVE attribute.

FIXED BINARY, FIXED DECIMAL, or FLOAT
Host variables that are declared with the FIXED and BINARY, FIXED and DECIMAL, or FLOAT attributes are called numeric host variables. The following restrictions apply to the numeric host variable:

- It must have the REAL attribute.
- If it has the FIXED and BINARY attributes, it must have the SIGNED and NATIVE attributes, a zero scale factor, and a precision greater than 7.
- If it has the FIXED and DECIMAL attributes, it must have a non-negative scale factor that is smaller than its precision.
- If it has the FLOAT and DECIMAL attributes, it must have a precision that is less than 17 unless the FLOAT(DFP) option is in effect.
- If it has the FLOAT and BINARY attributes, it must have a precision that is less than 54.

SQL TYPE
Host variables that are declared with the SQL TYPE attribute are called SQL TYPE host variables. The attribute specification must conform to one of the following syntax diagrams:

BINARY

\[
\text{SQL TYPE IS BINARY (length)}
\]

VARBINARY

\[
\text{SQL TYPE IS VARBINARY (length)}
\]

Result set locator

\[
\text{SQL TYPE IS RESULT_SET_LOCATOR}
\]

ROWID

\[
\text{SQL TYPE IS ROWID}
\]

Table locator

\[
\text{SQL TYPE IS TABLE LIKE table-name AS LOCATOR}
\]

LOB file reference

\[
\text{SQL TYPE IS BLOB_FILE \leftarrow CLOB_FILE \leftarrow DBCLOB_FILE}
\]
LOB locator

SQL TYPE IS BLOB_LOCATOR

CLOB_LOCATOR

DBCLOB_LOCATOR

LOB variable

SQL TYPE IS BLOB

CLOB

DBCLOB

You can also use BINARY LARGE OBJECT as an alternative for BLOB.

CLOB

You can also use either CHARACTER LARGE OBJECT or CHAR LARGE OBJECT as an alternative for CLOB.

XML LOB variable

SQL TYPE IS XML AS BLOB

CLOB

DBCLOB

You can also use BINARY LARGE OBJECT as an alternative for BLOB.

CLOB

You can also use either CHARACTER LARGE OBJECT or CHAR LARGE OBJECT as an alternative for CLOB.

XML file reference

SQL TYPE IS XML AS BLOB_FILE

CLOB_FILE

DBCLOB_FILE

The following constant declarations are generated by the SQL preprocessor. You can use them to set the file option variable when you use the file reference host variables:

DCL SQL_FILE_READ FIXED BIN(31) VALUE(2);
DCL SQL_FILE_CREATE FIXED BIN(31) VALUE(8);
DCL SQL_FILE_OVERWRITE FIXED BIN(31) VALUE(16);
DCL SQL_FILE_APPEND FIXED BIN(31) VALUE(32);

Determining equivalent SQL and PL/I data types: The base SQLTYPE and SOLLEN of host variables are determined according to Table 8 on page 127 and Table 9 on page 127. If a host variable appears with an indicator variable, the SQLTYPE is the base SQLTYPE plus one.

To determine the PL/I data type that is equivalent to a given SQL data type, you can use Table 10 on page 128 and Table 11 on page 129.
### Table 8. SQL data types generated from PL/I declarations

<table>
<thead>
<tr>
<th>PL/I Data Type</th>
<th>SQLTYPE of Host Variable</th>
<th>SQLLEN of Host Variable</th>
<th>SQL Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>BIN FIXED(p), 7 &lt; p &lt;= 15</td>
<td>500</td>
<td>2</td>
<td>SMALLINT</td>
</tr>
<tr>
<td>BIN FIXED(p), 15 &lt; p &lt;= 31</td>
<td>496</td>
<td>4</td>
<td>INTEGER</td>
</tr>
<tr>
<td>BIN FIXED(p), 31 &lt; p &lt;= 63</td>
<td>492</td>
<td>8</td>
<td>BIGINT</td>
</tr>
<tr>
<td>DEC FIXED(p,s), 0&lt;=p&lt;=15, 0&lt;=s&lt;=p</td>
<td>484</td>
<td>p (byte 1) s (byte 2)</td>
<td>DECIMAL(p,s)</td>
</tr>
<tr>
<td>BIN FLOAT(p), 1 &lt;= p &lt;= 21</td>
<td>480</td>
<td>4</td>
<td>REAL or FLOAT(n) 1&lt;=n&lt;=21</td>
</tr>
<tr>
<td>BIN FLOAT(p), 22 &lt;= p &lt;= 53</td>
<td>480</td>
<td>8</td>
<td>DOUBLE PRECISION or FLOAT(n) 22&lt;=n&lt;=53</td>
</tr>
</tbody>
</table>

Under FLOAT(NODFP):

| DEC FLOAT(p), 1 <= p <= 6               | 480                      | 4                       | FLOAT (single precision) |
| DEC FLOAT(p), 7 <= p <= 16              | 480                      | 8                       | FLOAT (double precision) |

Under FLOAT(DFP):

| DEC FLOAT(p), 1 <= p <= 7               | 996                      | 4                       | DECFLOAT (single precision) |
| DEC FLOAT(p), 7 <= p <= 16              | 996                      | 8                       | DECFLOAT (double precision) |
| DEC FLOAT(p), 16 <= p <= 34             | 996                      | 16                      | DECFLOAT (extended decimal) |

| CHAR(n)                                 | 452                      | n                       | CHAR(n)                |
| CHAR(n) VARYING                         | 448                      | n                       | VARCHAR(n)            |
| GRAPHIC(n), 1 <= n <= 127              | 468                      | n                       | GRAPHIC(n)            |
| GRAPHIC(n) VARYING                     | 464                      | n                       | VARGRAPHIC(n)         |

### Table 9. SQL data types generated from SQL TYPE declarations

<table>
<thead>
<tr>
<th>PL/I Data Type</th>
<th>SQLTYPE of Host Variable</th>
<th>SQLLEN of Host Variable</th>
<th>SQL Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>SQL TYPE IS BLOB(n) 1&lt;n&lt;2147483647</td>
<td>404</td>
<td>n</td>
<td>BLOB(n)</td>
</tr>
<tr>
<td>SQL TYPE IS CLOB(n) 1&lt;n&lt;2147483647</td>
<td>408</td>
<td>n</td>
<td>CLOB(n)</td>
</tr>
<tr>
<td>SQL TYPE IS DBCLOB(n) 1&lt;n&lt;1073741823 (2)</td>
<td>412</td>
<td>n</td>
<td>DBCLOB(n) (2)</td>
</tr>
<tr>
<td>SQL TYPE IS ROWID</td>
<td>904</td>
<td>40</td>
<td>ROWID</td>
</tr>
<tr>
<td>SQL TYPE IS VARBINARY(n) 1&lt;n&lt;32704</td>
<td>908</td>
<td>n</td>
<td>VARBINARY(n)</td>
</tr>
<tr>
<td>SQL TYPE IS BINARY(n) 1&lt;n&lt;255</td>
<td>912</td>
<td>n</td>
<td>BINARY(n)</td>
</tr>
<tr>
<td>SQL TYPE IS BLOB_FILE</td>
<td>916</td>
<td>267</td>
<td>BLOB File Reference (1)</td>
</tr>
<tr>
<td>SQL TYPE IS CLOB_FILE</td>
<td>920</td>
<td>267</td>
<td>CLOB File Reference (1)</td>
</tr>
<tr>
<td>SQL TYPE IS DBCLOB_FILE</td>
<td>924</td>
<td>267</td>
<td>DBCLOB File Reference (1)</td>
</tr>
<tr>
<td>SQL TYPE IS BLOB_LOCATOR</td>
<td>960</td>
<td>4</td>
<td>BLOB Locator (1)</td>
</tr>
<tr>
<td>SQL TYPE IS CLOB_LOCATOR</td>
<td>964</td>
<td>4</td>
<td>CLOB Locator (1)</td>
</tr>
<tr>
<td>SQL TYPE IS DBCLOB_LOCATOR</td>
<td>968</td>
<td>4</td>
<td>DBCLOB Locator (1)</td>
</tr>
</tbody>
</table>
### Table 9. SQL data types generated from SQL TYPE declarations (continued)

<table>
<thead>
<tr>
<th>PL/I Data Type</th>
<th>SQLTYPE of Host Variable</th>
<th>SQLLEN of Host Variable</th>
<th>SQL Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>SQL TYPE IS RESULT_SET_LOCATOR</td>
<td>972</td>
<td>4</td>
<td>Result Set Locator</td>
</tr>
<tr>
<td>SQL TYPE IS TABLE LIKE table-name AS LOCATOR</td>
<td>976</td>
<td>4</td>
<td>Table Locator (1)</td>
</tr>
</tbody>
</table>

**Notes:**
1. Do not use this data type as a column type.
2. \( n \) is the number of double-byte characters.

### Table 10. SQL data types mapped to PL/I declarations

<table>
<thead>
<tr>
<th>SQL Data Type</th>
<th>PL/I Equivalent</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>SMALLINT</td>
<td>BIN FIXED(15)</td>
<td></td>
</tr>
<tr>
<td>INTEGER</td>
<td>BIN FIXED(31)</td>
<td></td>
</tr>
<tr>
<td>BIGINT</td>
<td>BIN FIXED(63)</td>
<td></td>
</tr>
<tr>
<td>DECIMAL((p,s))</td>
<td>DEC FIXED((p)) or DEC FIXED((p,s))</td>
<td>( 1 \leq p \leq 31 ) and ( 0 \leq s \leq p )</td>
</tr>
</tbody>
</table>

**Under FLOAT(NODFP):**

<table>
<thead>
<tr>
<th>SQL Data Type</th>
<th>PL/I Equivalent</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>REAL or FLOAT((n))</td>
<td>BIN FLOAT((p)) or DEC FLOAT((m))</td>
<td>( 1 \leq p \leq 21 ) and ( 1 \leq m \leq 6 )</td>
</tr>
<tr>
<td>DOUBLE PRECISION, DOUBLE, or FLOAT((n))</td>
<td>BIN FLOAT((p)) or DEC FLOAT((m))</td>
<td>( 22 \leq p \leq 53 ) and ( 7 \leq m \leq 16 )</td>
</tr>
</tbody>
</table>

**Under FLOAT(DFP):**

<table>
<thead>
<tr>
<th>SQL Data Type</th>
<th>PL/I Equivalent</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>DECFLOAT</td>
<td>DEC FLOAT((m))</td>
<td>Short Decimal Float</td>
</tr>
<tr>
<td></td>
<td></td>
<td>( 1 \leq m \leq 7 )</td>
</tr>
<tr>
<td>DECFLOAT</td>
<td>DEC FLOAT((m))</td>
<td>Long Decimal Float</td>
</tr>
<tr>
<td></td>
<td></td>
<td>( 7 \leq m \leq 16 )</td>
</tr>
<tr>
<td>DECFLOAT</td>
<td>DEC FLOAT((m))</td>
<td>Extended Decimal Float</td>
</tr>
<tr>
<td></td>
<td></td>
<td>( 16 \leq m \leq 34 )</td>
</tr>
<tr>
<td>CHAR((n))</td>
<td>CHAR((n))</td>
<td>( 1 \leq n \leq 32767 )</td>
</tr>
<tr>
<td>VARCHAR((n))</td>
<td>CHAR((n)) VAR</td>
<td></td>
</tr>
<tr>
<td>GRAPHIC((n))</td>
<td>GRAPHIC((n))</td>
<td>( n ) is a positive integer that refers to the number of double-byte characters, not to the number of bytes; ( 1 \leq n \leq 16383 )</td>
</tr>
<tr>
<td>VARGRAPHIC((n))</td>
<td>GRAPHIC((n)) VAR</td>
<td>( n ) is a positive integer that refers to the number of double-byte characters, not to the number of bytes; ( 1 \leq n \leq 16383 )</td>
</tr>
<tr>
<td>DATE</td>
<td>CHAR((n))</td>
<td>( n ) must be at least 10</td>
</tr>
<tr>
<td>TIME</td>
<td>CHAR((n))</td>
<td>( n ) must be at least 8</td>
</tr>
<tr>
<td>TIMESTAMP</td>
<td>CHAR((n))</td>
<td>( n ) must be at least 26</td>
</tr>
</tbody>
</table>
### Table 11. SQL data types mapped to SQL TYPE declarations

<table>
<thead>
<tr>
<th>SQL Data Type</th>
<th>PL/I Equivalent</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>Result set locator</td>
<td>SQL TYPE IS RESULT_SET_LOCATOR</td>
<td>Use this data type only for receiving result sets. Do not use this data type as a column type.</td>
</tr>
<tr>
<td>Table locator</td>
<td>SQL TYPE IS TABLE LIKE table-name AS LOCATOR</td>
<td>Use this data type only in a user-defined function or stored procedure to receive rows of a transition table. Do not use this data type as a column type.</td>
</tr>
<tr>
<td>BLOB locator</td>
<td>SQL TYPE IS BLOB_LOCATOR</td>
<td>Use this data type only to manipulate data in BLOB columns. Do not use this data type as a column type.</td>
</tr>
<tr>
<td>CLOB locator</td>
<td>SQL TYPE IS CLOB_LOCATOR</td>
<td>Use this data type only to manipulate data in CLOB columns. Do not use this data type as a column type.</td>
</tr>
<tr>
<td>DBCLOB locator</td>
<td>SQL TYPE IS DBCLOB_LOCATOR</td>
<td>Use this data type only to manipulate data in DBCLOB columns. Do not use this data type as a column type.</td>
</tr>
<tr>
<td>BLOB file reference</td>
<td>SQL TYPE IS BLOB_FILE</td>
<td>Use this data type only as a reference to a BLOB file. Do not use this data type as a column type.</td>
</tr>
<tr>
<td>CLOB file reference</td>
<td>SQL TYPE IS CLOB_FILE</td>
<td>Use this data type only as a reference to a CLOB file. Do not use this data type as a column type.</td>
</tr>
<tr>
<td>DBCLOB file reference</td>
<td>SQL TYPE IS DBCLOB_FILE</td>
<td>Use this data type only as a reference to a DBCLOB file. Do not use this data type as a column type.</td>
</tr>
<tr>
<td>BLOB(n)</td>
<td>SQL TYPE IS BLOB(n)</td>
<td>1 &lt; n &lt; 2147483647</td>
</tr>
<tr>
<td>CLOB(n)</td>
<td>SQL TYPE IS CLOB(n)</td>
<td>1 &lt; n &lt; 2147483647</td>
</tr>
<tr>
<td>DBCLOB(n)</td>
<td>SQL TYPE IS DBCLOB(n)</td>
<td>n is the number of double-byte characters. 1 &lt; n &lt; 1073741823</td>
</tr>
<tr>
<td>ROWID</td>
<td>SQL TYPE IS ROWID</td>
<td></td>
</tr>
<tr>
<td>XML AS</td>
<td>SQL TYPE IS XML AS ...</td>
<td>Used to describe an XML version of a BLOB, CLOB, DBCLOB, BLOB_FILE, CLOB_FILE, or DBCLOB_FILE</td>
</tr>
</tbody>
</table>

**Determining compatibility of SQL and PL/I data types:** PL/I host variables in SQL statements must be type compatible with the columns which use them:

- Numeric data types are compatible with each other. A SMALLINT, INTEGER, DECIMAL, or FLOAT column is compatible with a PL/I host variable of BIN FIXED(15), BIN FIXED(31), DECIMAL(p,s), BIN FLOAT(n) where n is from 22 to 53, or DEC FLOAT(m) where m is from 7 to 16.

- Character data types are compatible with each other. A CHAR or VARCHAR column is compatible with a fixed-length or varying-length PL/I character host variable.

- Datetime data types are compatible with character host variables. A DATE, TIME, or TIMESTAMP column is compatible with a fixed-length or varying-length PL/I character host variable.

When necessary, the Database Manager automatically converts a fixed-length character string to a varying-length string or a varying-length string to a fixed-length character string.
Using host structures

A PL/I host structure name can be a structure name with members that are not structures or unions. For example:

```pli
dcl 1 A,
  2 B,
    3 C1 char(...),
    3 C2 char(...);
```

In this example, B is the name of a host structure consisting of the scalars C1 and C2.

Host structures are limited to two levels. A host structure can be thought of as a named collection of host variables.

Each leaf element of a host structure must have one of the following valid host data attributes as discussed in "Declaring host variables" on page 124:

- CHARACTER, GRAPHIC, or WIDECHAR
- FIXED BINARY, FIXED DECIMAL, or FLOAT
- SQL TYPE

Using indicator variables

An indicator variable is a two-byte integer (BIN FIXED(15)), an array of two-byte integers, or a structure that contains only two-byte integers (or arrays thereof).

On retrieval, an indicator variable is used to show whether its associated host variable has been assigned a null value. On assignment to a column, a negative indicator variable is used to indicate that a null value should be assigned.

A structure can be used as an indicator variable only when the associated host variable is also a structure.

Indicator variables are declared in the same way as host variables and the declarations of the two can be mixed in any way that seems appropriate to the programmer.

The SQL Preprocessor does not require that indicator arrays have a lower bound of one.

An indicator variable must have the attribute REAL NATIVE SIGNED FIXED BIN(15).

Given the statement:

```sql
exec sql fetch Cls_Cursor into :Cls_Cd,
    :Day :Day_Ind,
    :Bgn :Bgn_Ind,
    :End :End_Ind;
```

Variables can be declared as follows:

```pli
exec sql begin declare section;
dcl Cls_Cd char(7);
dcl Day bin fixed(15);
dcl Bgn char(8);
dcl End char(8);
dcl (Day_Ind, Bgn_Ind, End_Ind) bin fixed(15);
exec sql end declare section;
```
Host structure example

The following example shows the declaration of a host structure and an indicator array followed by an SQL statement that could be used to retrieve the data into the host structure:

```pli
  dcl 1 games,
      5 sunday,
      10 opponents char(30),
      10 gtime char(10),
      10 tv char(6),
      10 comments char(120) var;
  dcl indicator(4) fixed bin (15);
  exec sql
    fetch cursor_a
    into :games.sunday:indicator;
```

General information about LOBs

LOBs, CLOBs, and BLOBs can be as large as 2,147,483,647 bytes long (2 Gigabytes). Double Byte CLOBs can be 1,073,741,823 characters long (1 Gigabyte).

Manipulating LOB data

To use large object (LOB) data from a DB2 table, use techniques like LOB locators and LOB file references to manipulate the data while the data is still in the database. Declaring a host variable to hold all of the LOB data can be inefficient or impractical, because this requires your program to allocate large amounts of storage and requires DB2 to move large amounts of data.

This information describes the using of LOB locators and LOB file references. A PL/I sample program, the pliclob sample program is also provided to show how to manipulate CLOBs in a PL/I and DB2 environment. For more information about these techniques to minimize the moving around of large pieces of data, see Ways to manipulate LOB data in the Writing DB2 application chapter in the DB2 for z/OS Application Programming and SQL Guide, or the Redbooks, LOBs with DB2 for z/OS: Stronger and Faster.

LOB locators

You can use LOB Locators to avoid materialization of the LOB data and all the underlying activities associated with it. The benefits of using LOB locators are listed as follows:

- Saving storage when manipulating LOBs with LOB locators
- Manipulating data without retrieving it from the data base
- Avoiding the use of large amounts of storage to hold the LOB
- Avoiding the time and resource expenditures for moving large pieces of data thereby improving performance

LOB locators are especially useful under the following circumstances:

- When you need only a small part of a LOB
- When you do not have enough memory for the entire LOB
- When performance is important
- In a client or server environment to avoid moving data over the network from one system to another

The following code example is from the pliclob sample program. The sample program uses LOB locators to identify and manipulate sections of the resume.
CLOB from the dsn8a10.emp_photo_resume DB2 V10 table. (The numbers that precede each line are not part of the program, but are used in the explanation after the program.)

1. dcl hv_loc_resume sql type is clob_locator;
2. exec sql
3. select resume into :hv_loc_resume
4. from dsn8a10.emp_photo_resume
5. where empno = :hv_empno;
6. exec sql
7. set :start_resume = (posstr(:hv_loc_resume, 'Resume:'));

In lines 2 - 5, LOB locator hv_loc_resume is set to the location of the resume of the employee number hv_empno in the emp_photo_resume table. In lines 7 - 8, the start_resume host variable is set to the beginning of the 'Resume:' section of the resume. Then you can start manipulating the resume data while the resume is still in the database.

LOB file reference variables

You can use LOB file reference variables to import or export data between a LOB column and an external file outside the DB2 system. The benefits of using LOB file reference variables are listed as follows:

- Uses less processing time than moving LOB data with a host variable. The movement of the data would not be overlapped with any DB2 processing or network transfer time.
- Uses less application storage. The LOB data is moved directly from DB2 to a file and is not materialized in the memory of the application.

The following code example is from the pliclob sample program. The sample program uses LOB file references to create a new, trimmed down version of the resume in an external file. (The numbers that precede each line are not part of the program, but are used in the explanation after the program.)

1. dcl hv_clob_file sql type is clob_file;
2. name_string = '/SYSTEM/tmp/pliclob2.txt';
3. hv_clob_file.sql_lob_file_name_len = length(name_string);
4. hv_clob_file.sql_lob_file_name = name_string;
5. hv_clob_file.sql_lob_file_options = ior(sql_file_overwrite);
6. exec sql
7. values ( substr(:hv_loc_resume,:start_resume,
8. :start_pers_info-:start_resume)
9. || substr(:hv_loc_resume,:start_work_hist,
10. :start_interests-:start_work_hist)
11. )
12. )
13. into :hv_clob_file;

The host variable hv_clob_file is declared as a LOB file reference. In lines 2 - 4, the file name field of the LOB file reference is set to the fully qualified file name and file name length is set to its length. The overwrite option is set so any existing file is overwritten (line 5). For details of these and other file options, see the DB2 for z/OS Application programming and SQL Guide.

Next the SQL VALUES statement is used to concatenate the resume name and work history sections of the resume directly into the hv_clob_file LOB file reference, as in lines 8 - 13.
Example

The following PL/I sample program shows how to manipulate CLOBs in a PL/I and DB2 environment.

You must have the DB2 supplied sample database installed for this program to run properly. This sample assumes DB2 V10 and table dsn8a10.emp_photo_resume. If you use a different version of DB2, you must change the table reference.

Notes:

- When you use the LOB locators and LOB file reference variables, the resume CLOB is still within the data base and not in memory or storage.
- The format of the resume CLOB in the data base remains unchanged, with the reformatting of the resume taking place only in the second file that was written out.
pliclob: procedure options(main);
  display('begin pliclob');
  exec sql include sqlca;

  dcl hv_empno char(06);
  dcl name_string char(256) var;
  dcl hv_resume sql type is clob(50k);
  dcl hv_clob_file sql type is clob_file;
  dcl hv_loc_resume sql type is clob_locator;
  dcl start_resume fixed bin(31);
  dcl start_pers_info fixed bin(31);
  dcl start_dept_info fixed bin(31);
  dcl start_education fixed bin(31);
  dcl start_work_hist fixed bin(31);
  dcl start_interests fixed bin(31);

  /* Extract resume CLOB for employee '000130' into a file in z/OS */
  /* UNIX file system. The contents of this file shows the initial */
  /* format of the resume CLOB in the data base. */
  /* Note: this program must have 'write' access to the directory */
  /* designated in the 'name_string' variable. */
  name_string = '/SYSTEM/tmp/pliclob1.txt';
  hv_clob_file.sql_lob_file_name_len = length(name_string);
  hv_clob_file.sql_lob_file_name = name_string;
  hv_clob_file.sql_lob_file_options = ior(sql_file_overwrite);

  hv_empno = '000130';
  exec sql
    select resume into :hv_clob_file
    from dsn8a10.emp_photo_resume
    where empno = :hv_empno;
  display('file1 sqlca.sqlcode = ' || sqlca.sqlcode);

  /* Next, a CLOB locator is used to locate the resume CLOB for */
  /* employee number '000130' in the data base. Then a series of */
  /* DB2 SET statements using the posstr DB2 function finds the */
  /* beginning position of each section within the resume. */
  exec sql
    select resume into :hv_loc_resume
    from dsn8a10.emp_photo_resume
    where empno = :hv_empno;
  display('select resume sqlcode = ' || sqlca.sqlcode);

  exec sql set :start_resume =
    (posstr(:hv_loc_resume, 'Resume:'));
  display('first set sqlcode = ' || sqlca.sqlcode);

  exec sql set :start_pers_info =
    (posstr(:hv_loc_resume, 'Personal Information'));
  display('second set sqlcode = ' || sqlca.sqlcode);

Figure 8. pliclob sample program
/* Finally, by using the CLOB locator and the start references */
/* of each section in the resume, along with the DB2 substr and */
/* concatenate (||) functions, the resume CLOB is written out to */
/* a second file in a slightly different format: */
/* 1. the Personal Information section is omitted due to */
/* privacy concerns. */
/* 2. the sections within the resume are written out in this */
/* order: Resume, Work History, Education then Department */
/* Information. */
/* */
/* After the second file is written out, the changes to the */
/* resume CLOB can be verified by comparing the contents of the */
/* two files pliclob1.txt and pliclob2.txt. */
/* */
/* Note: this program must have 'write' access to the directory */
/* designated in the 'name_string' variable. */
name_string = '/SYSTEM/tmp/pliclob2.txt';
hv_clob_file.sql_lob_file_name_len = length(name_string);
hv_clob_file.sql_lob_file_name = name_string;
hv_clob_file.sql_lob_file_options = ior(sql_file_overwrite);

exec sql
values ( substr(:hv_loc_resume,:start_resume,
    :start_pers_info:-:start_resume)
    || substr(:hv_loc_resume,:start_work_hist,
    :start_interests:-:start_work_hist)
    || substr(:hv_loc_resume,:start_education,
    :start_work_hist:-:start_education)
    || substr(:hv_loc_resume,:start_dept_info,
    :start_education:-:start_dept_info)
)
into :hv_clob_file;

display('file2 sqlca.sqlcode = ' || sqlca.sqlcode);
display('End pliclob');
end;

pliclob sample program (continued)

Suppressing SQL preprocessor messages

You can use the IBM supplied compiler user exit (IBMUEXIT) to suppress a message or to change the severity of a message. See "Example of suppressing SQL messages" on page 468 for an example of suppressing the preprocessor messages by modifying the user exit.
CICS Preprocessor

You can use EXEC CICS statements in PL/I applications that run as transactions under CICS.

If you do not specify the PP(CICS) option, EXEC CICS statements are parsed and variable references in them are validated. If they are correct, no messages are issued as long as the NOCOMPILE option is in effect. If you do not invoke the CICS translator and the COMPILE option is in effect, the compiler will issue S-level messages.

The compiler will invoke the CICS preprocessor if you specify the CICS suboption of the PP option. For compatibility, it will also invoke the CICS preprocessor if the compiler sees any of these options: CICS, XOPT or XOPTS. However, you should not specify one of these options and also the PP(CICS) options.

Programming and compilation considerations

When you are developing programs for execution under CICS, all the EXEC CICS commands must be translated in one of two ways:

- By the command language translator provided by CICS in a job step prior to the PL/I compilation
- By the PL/I CICS preprocessor as part of the PL/I compilation (this requires CICS TS 2.2 or later)

To use the CICS preprocessor, you must also specify the PP(CICS) and DFT(EBCDIC) compile time options. All data passed to CICS must be in the NATIVE format.

Unless you specify CICS as one of the suboptions of the PP(CICS) option, the compiler will flag any EXEC CICS statements in the source. Similarly, it will flag any EXEC CPSM or EXEC DLI statements if you do not specify CPSM or DLI respectively as a suboption of the PP(CICS) option.

If your CICS program is a MAIN procedure, you must also compile it with the SYSTEM(CICS) or the SYSTEM(MVS) option. If compiled with SYSTEM(MVS), the PTFs for runtime APAR PQ91318 must be applied. NOEXECOPS is implied with this option and all parameters passed to the MAIN procedure must be POINTERS. For a description of the SYSTEM compile time option, see “SYSTEM” on page 81.

If your want your CICS program to be reentrant and if your program uses FILEs or CONTROLLED variables, then you must compile it with the NOWRITABLE as well.

If your CICS program includes any files or uses any macros that contain EXEC CICS statements, you must also run the MACRO preprocessor before your code is translated (in either of the ways described above). If you are using the CICS preprocessor, you can specify this with one PP option as illustrated in the following example:

```
pp (macro(...) cics(...) )
```

The CICS preprocessor will add a set of declares for CICS variables and APIs to all non-nested procedures. Consequently, it keeps tracks of all statements that require a matching END statement, and if some of these statements are missing or
incorrect, the preprocessor may be misled and not insert these declares. The preprocessor will also terminate with a severe message if such statements are nested more than 150 deep.

Finally, in order to use the CICS preprocessor, you must have the CICS SDFHLOAD dataset as part of the STEPLIB DD for the PL/I compiler.

**CICS preprocessor options**

There are many options supported by CICS translator. For a description of these options, see the CICS Application Programming Guide.

Note that these options should be enclosed in quotes (single or double, as long as they match). For instance, to invoke the CICS preprocessor with the EDF option, you should specify the option PP(CICS('EDF')).

**Coding CICS statements in PL/I applications**

You can code CICS statements in your PL/I applications using the language defined in CICS on Open Systems Application Programming Guide. Specific requirements for your CICS code are described in the sections that follow.

**Embedding CICS statements**

If you use the CICS translator, rather than the integrated preprocessor, then the first statement of your PL/I program must be a PROCEDURE statement. You can add CICS statements to your program wherever executable statements can appear. Each CICS statement must begin with EXEC (or EXECUTE) CICS and end with a semicolon (;).

For example, the GETMAIN statement might be coded as follows:

```
EXEC CICS GETMAIN SET(BLK_PTR) LENGTH(STG(BLK));
```

**Comments:** In addition to the CICS statements, PL/I comments can be included in embedded CICS statements wherever a blank is allowed.

**Continuation for CICS statements:** Line continuation rules for CICS statements are the same as those for other PL/I statements.

**Including code:** If included code contains EXEC CICS statements or your program uses PL/I macros that generate EXEC CICS statements, you must use one of the following:

- The MACRO compile-time option
- The MACRO option of the PP option (before the CICS option of the PP option)

**Margins:** CICS statements must be coded within the columns specified in the MARGINS compile-time option.

**Statement labels:** EXEC CICS statements, like PL/I statements, can have a label prefix.

**Writing CICS transactions in PL/I**

You can use PL/I with CICS facilities to write application programs (transactions) for CICS subsystems. If you do this, CICS provides facilities to the PL/I program that would normally be provided directly by the operating system. These facilities include most data management facilities and all job and task management facilities.
You must observe the following restrictions on PL/I CICS programs:

- Macro-level CICS is not supported.
- PL/I input or output cannot be used except for:
  - PUT FILE(SYSPRINT)
  - DISPLAY
  - CALL PLIDUMP
- The PLISRTx built-in subroutines cannot be used.
- Routines written in a language other than PL/I cannot be called from a PL/I CICS program if those routines contain their own EXEC CICS statements. If you want to communicate with a non-PL/I program that contains EXEC CICS statements, you must use EXEC CICS LINK or EXEC CICS XCTL to do so.

Although PUT FILE(SYSPRINT) is permitted under CICS, you should generally not use it in production programs as it will degrade performance.

Since the CICS EIB address is only generated by either the CICS translator or the PL/I CICS preprocessor for an OPTIONS(MAIN) program, it is the user’s responsibility to establish the addressability to the EIB for the OPTIONS(FETCHABLE) routine.

Addressability can be achieved either by using this command:
EXEC CICS ADDRESS EIB(DFHEIPTR)

or by passing the EIB address as an argument to the CALL statement that invokes the external procedure.

**Error-handling**

Language Environment prohibits the use of the following EXEC CICS commands in any PL/I ON-unit or in any code called from a PL/I ON-unit.

- EXEC CICS ABEND
- EXEC CICS HANDLE AID
- EXEC CICS HANDLE ABEND
- EXEC CICS HANDLE CONDITION
- EXEC CICS IGNORE CONDITION
- EXEC CICS POP HANDLE
- EXEC CICS PUSH HANDLE

All other EXEC CICS commands are allowed within an ON-unit. However, they must be coded using the NOHANDLE option, the RESP option, or the RESP2 option.
Chapter 3. Using PL/I cataloged procedures

This chapter describes the standard cataloged procedures supplied by IBM for use with the IBM Enterprise PL/I for z/OS compiler. It explains how to invoke them, and how to temporarily or permanently modify them. The Language Environment SCEERUN data set must be located in STEPLIB and accessible to the compiler when you use any of the cataloged procedures described in this chapter.

A cataloged procedure is a set of job control statements, stored in a library, that includes one or more EXEC statements, each of which can be followed by one or more DD statements. You can retrieve the statements by naming the cataloged procedure in the PROC parameter of an EXEC statement in the input stream.

You can use cataloged procedures to save time and reduce Job Control Language (JCL) errors. If the statements in a cataloged procedure do not match your requirements exactly, you can easily modify them or add new statements for the duration of a job. You should review these procedures and modify them to obtain the most efficient use of the facilities available and to allow for your own conventions.

IBM-supplied cataloged procedures

The PL/I cataloged procedures supplied for use with Enterprise PL/I for z/OS are:

IBMZC
  Compile only

IBMZCB
  Compile and bind

IBMZCPL
  Compile, prelink, and link-edit

IBMZCBG
  Compile, bind, and run

IBMZCPLG
  Compile, prelink, link-edit, and run

IBMZCPLG
  Compile, prelink, load, and run

Cataloged procedures IBMZCB and IBMZCBG use features of the program management binder introduced in DFSMS/MVS® 1.4 in place of the prelinker supplied with Language Environment. These procedures produce a program object in a PDSE.

Cataloged procedures IBMZCPL, IBMZCPLG and IBMZCPLG use the prelinker supplied with Language Environment and produce a load module in PDS. Use these procedures if you do not want to use a PDSE. The information in this section describes the procedure steps of the different cataloged procedures. For a description of the individual statements for compiling and link editing, see “Invoking the compiler under z/OS using JCL” on page 157 and z/OS Language Environment Programming Guide. These cataloged procedures do not include a DD statement for the input data set; you must always provide one. The example shown in Figure 9 on page 140 illustrates the JCL statements you might use to invoke the cataloged procedure IBMZCBG to compile, bind, and run a PL/I program.
Enterprise PL/I requires a minimum REGION size of 32M. Large programs require more storage. If you do not specify REGION on the EXEC statement that invokes the cataloged procedure you are running, the compiler uses the default REGION size for your site. The default size might or might not be adequate, depending on the size of your PL/I program.

If you compile your programs with optimization turned on, the REGION size (and time) required may be much, much larger.

For an example of specifying REGION on the EXEC statement, see Figure 9.

```plaintext
//COLEGOC JOB
//STEP1 EXEC IBMZCBG, REGION.PLI=32M
//PLI.SYSLIN DD *
  .
  .
  .
  (insert PL/I program to be compiled here)
  .
  .
/*

Figure 9. Invoking a cataloged procedure

Compile only (IBMZC)

The IBMZC cataloged procedure, shown in Figure 10 on page 141 includes only one procedure step, in which the options specified for the compilation are OBJECT and OPTIONS. (IBMZPLI is the symbolic name of the compiler.) In common with the other cataloged procedures that include a compilation procedure step, IBMZC does not include a DD statement for the input data set; you must always supply an appropriate statement with the qualified ddname PLI.SYSIN.

The OBJECT compile-time option causes the compiler to place the object module, in a syntax suitable for input to the linkage editor, in the standard data set defined by the DD statement with the name SYSLIN. This statement defines a temporary data set named &&LOADSET on a sequential device; if you want to retain the object module after the end of your job, you must substitute a permanent name for &&LOADSET (that is, a name that does not start with &&) and specify KEEP in the appropriate DISP parameter for the last procedure step that used the data set. You can do this by providing your own SYSLIN DD statement, as shown below. The data set name and disposition parameters on this statement will override those on the IBMZC procedure SYSLIN DD statement. In this example, the compile step is the only step in the job.

```
//PLICOMP EXEC IBMZC
//PLI.SYSLIN DD DSN=MYPROG,DISP=SHR
//PLI.SYSIN DD ...
```

The term MOD in the DISP parameter in Figure 10 on page 141 allows the compiler to place more than one object module in the data set, and PASS ensures that the data set is available to a later procedure step providing a corresponding DD statement is included there.
The SYSLIN SPACE parameter allows an initial allocation of 1 cylinder and, if necessary, 15 further allocations of 1 cylinder (a total of 16 cylinders).

Compile and bind (IBMZCB)

The IBMZCB cataloged procedure, shown in Figure 11 on page 142, includes two procedure steps: PLI, which is identical to cataloged procedure IBMZC, and BIND, which invokes the Program Management binder (symbolic name IEWBLINK) to bind the object module produced in the first procedure step.

Input data for the compilation procedure step requires the qualified ddname PLI.SYSIN. The COND parameter in the EXEC statement BIND specifies that this procedure step should be bypassed if the return code produced by the compiler is greater than 8 (that is, if a severe or unrecoverable error occurs during compilation).
The Program Management binder always places the program objects it creates in the standard data set defined by the DD statement with the name SYSLMOD. This statement in the cataloged procedure specifies a new temporary library &&GOSET, in which the program object will be placed and given the member name GO. In specifying a temporary library, the cataloged procedure assumes that you will run the program object in the same job; if you want to retain the program object, you must substitute your own statement for the DD statement with the name SYSLMOD.

Figure 11. Cataloged Procedure IBMZCB

The Program Management binder always places the program objects it creates in the standard data set defined by the DD statement with the name SYSLMOD. This statement in the cataloged procedure specifies a new temporary library &&GOSET, in which the program object will be placed and given the member name GO. In specifying a temporary library, the cataloged procedure assumes that you will run the program object in the same job; if you want to retain the program object, you must substitute your own statement for the DD statement with the name SYSLMOD.
Compile, bind, and run (IBMZCBG)

The IBMZCBG cataloged procedure, shown in Figure 12 on page 144, includes three procedure steps: PLI, BIND, and GO. PLI and BIND are identical to the two procedure steps of IBMZCB, and GO runs the program object created in the step BIND. The GO step is executed only if no severe or unrecoverable errors occurred in the preceding procedure steps.

Input data for the compilation procedure step should be specified in a DD statement with the name PLI.SYSIN, and for the GO step in a DD statement with the name GO.SYSIN.
Figure 12. Cataloged Procedure IBMZCBG
Compile, prelink, and link-edit (IBMZCPL)

The IBMZCPL cataloged procedure, shown in Figure 13, includes three procedure steps: PLI, which is identical to cataloged procedure IBMZC; PLKED, which invokes the Language Environment prelinker; and LKED, which invokes the linkage editor (symbolic name IEWL) to link-edit the object module produced in the first procedure step.

Input data for the compilation procedure step requires the qualified ddname PLI.SYSIN. The COND parameter in the EXEC statement LKED specifies that this procedure step should be bypassed if the return code produced by the compiler is greater than 8 (that is, if a severe or unrecoverable error occurs during compilation).

```plaintext
//IBMZCPL PROC LNGPRFX='IBMZ.V4R5M0',LIBPRFX='CEE',
  // SYSLBLK=3200,PLANG=EDCPMSGE,GOPGM=GO
/*
//*************************************************************
//= LICENSED MATERIALS - PROPERTY OF IBM
//= *
//= * 5655-H31 (C) COPYRIGHT IBM CORP. 1999, 2015
//= * ALL RIGHTS RESERVED.
//= *
//= US GOVERNMENT USERS RESTRICTED RIGHTS - USE,
//= DUPICATION OR DISCLOSURE RESTRICTED BY GSA
//= ADP SCHEDULE CONTRACT WITH IBM CORP.
//= *
//=*************************************************************
//= */
//*************************************************************
//= IBM ENTERPRISE PL/I FOR Z/OS
//= Version 4 Release 5 MODIFICATION 0
 //= */
 //= * COMPILE, PRELINK, LINK-EDIT A PL/I PROGRAM
 //= */
 //= PARAMETER DEFAULT VALUE USAGE
 //= LNGPRFX IBMZ.V4R5M0 PREFIX FOR LANGUAGE DATA SET NAMES
 //= LIBPRFX CEE PREFIX FOR LIBRARY DATA SET NAMES
 //= SYSLBLK 3200 BLKSIZE FOR OBJECT DATA SET
 //= PLANG EDCPMSGE PRELINKER MESSAGES MEMBER NAME
 //= GOPGM GO MEMBER NAME FOR LOAD MODULE
 //= */
//*************************************************************
//= * COMPILE STEP
//=*************************************************************
//PLI EXEC PGM=IBMZPLI,PARM='OBJECT,OPTIONS'
//STEPLIB DD DSN=&LNGPRFX.;SIBMZCMP,DISP=SHR
// DD DSN=&LIBPRFX.;SCEE_RUN,DISP=SHR
//SYSPRINT DD SYSOUT=*
//SYSOUT DD SYSOUT=*
//SYSLIN DD DSN=&&LOADSET,DISP=(MOD,PASS),UNIT=SYSALLDA,
// SPACE=(CYL,(1,1)),DCB=(LRECL=80,BLKSIZE=&SYSLBLK)
//SYSUT1 DD DSN=&&SYSUT1,UNIT=SYSALLDA,
// SPACE=(1024,(200,50)),CONTIG,ROUND),DCB=BLKSIZE=1024
```
Cataloged Procedure IBMZCPL (continued)

The linkage editor always places the load modules it creates in the standard data set defined by the DD statement with the name SYSLMOD. This statement in the cataloged procedure specifies a new temporary library &&GOSET, in which the load module will be placed and given the member name GO. In specifying a temporary library, the cataloged procedure assumes that you will run the load module in the same job; if you want to retain the module, you must substitute your own statement for the DD statement with the name SYSLMOD.

The SYSLIN DD statement in Figure 13 on page 145 shows how to concatenate a data set defined by a DD statement named SYSIN with the primary input (SYSLIN) to the linkage editor. You could place linkage editor control statements in the input stream by this means, as described in the z/OS Language Environment Programming Guide.

Compile, prelink, link-edit, and run (IBMZCPLG)

The IBMZCPLG cataloged procedure, shown in Figure 14 on page 147 includes four procedure steps: PLI, PLKED, LKED, and GO. PLI, PLKED, and LKED are identical to the three procedure steps of IBMZCPL, and GO runs the load module created in the step LKED. The GO step is executed only if no severe or unrecoverable errors occurred in the preceding procedure steps.

Input data for the compilation procedure step should be specified in a DD statement with the name PLISYSIN, and for the GO step in a DD statement with the name GO.SYSIN.
Figure 14. Cataloged Procedure IBMZCPLG
Cataloged Procedure IBMZCPLG (continued)

Compile, prelink, load and run (IBMZCPG)

The IBMZCPG cataloged procedure, shown in Figure 15 on page 149, achieves the same results as IBMZCPLG but uses the loader instead of the linkage editor. Instead of using four procedure steps (compile, prelink, link-edit, and run), it has only three (compile, prelink, and load-and-run). The third procedure step runs the loader program. The loader program processes the object module produced by the compiler and runs the resultant executable program immediately. You must provide input data for the compilation step by supplying a qualified ddname PLISYSIN.

The use of the loader imposes certain restrictions on your PL/I program; before using this cataloged procedure, see z/OS Language Environment Programming Guide, which explains how to use the loader.
//IBMZCPG PROC LNGPRFX='IBMZ.V4R5M0',LIBPRFX='CEE',
// SYSLBLK=3200,PLANG=EDCPMSGE
//
//********************************************************************
// LICENSED MATERIALS - PROPERTY OF IBM
//
// 5655-H31 (C) COPYRIGHT IBM CORP. 1999, 2015
// ALL RIGHTS RESERVED.
//
// US GOVERNMENT USERS RESTRICTED RIGHTS - USE,
// DUPLICATION OR DISCLOSURE RESTRICTED BY GSA
// AOP SCHEDULE CONTRACT WITH IBM CORP.
//
//********************************************************************
//
// IBM ENTERPRISE PL/I FOR Z/OS
// Version 4 Release 5 MODIFICATION 0
//
// COMPILER, PRELINK, LOAD AND RUN A PL/I PROGRAM
//
// PARAMETER DEFAULT VALUE USAGE
// LNGPRFX IBMZ.V4R5M0 PREFIX FOR LANGUAGE DATA SET NAMES
// LIBPRFX CEE PREFIX FOR LIBRARY DATA SET NAMES
// SYSLBLK 3200 BLKSIZE FOR OBJECT DATA SET
// PLANG EDCPMSGE PRELINKER MESSAGES MEMBER NAME
//
//*********************************************************************
// COMPILE STEP
//*********************************************************************
//PLI EXEC PGM=IBMZPLI,PARM='OBJECT,OPTIONS'
//STEPLIB DD DSN=&LNGPRFX;.SIBMZCMP,DISP=SHR
// DD DSN=&LIBPRFX;.SCEERUN,DISP=SHR
//SYSPRINT DD SYSOUT=* 
//SYSOUT DD SYSOUT=* 
//SYSLIN DD DSN=&&LOADSET,DISP=(MOD,PASS),UNIT=SYSALLDA,
// SPACE=(CYL,(1,1)),DCB=(LRECL=80,BLKSIZE=&SYSLBLK)
//SYSUT1 DD DSN=&&SYSUT1,UNIT=SYSALLDA,
// SPACE=(1024,(200,50),,CONTIG,ROUND),DCB=BLKSIZE=1024
//*********************************************************************
// PRE-LINK-EDIT STEP
//*********************************************************************
//PLKED EXEC PGM=EDCPRLK,COND=(8,LT,PLI)
//STEPLIB DD DSN=&LIBPRFX;.SCEERUN,DISP=SHR
//SYSSMSGS DD DSN=&LIBPRFX;.SCEEMSGP(&PLANG),DISP=SHR
//SYSLIB DD DUMMY  
//SYSMOD DD DSN=&&PLNK,DISP=(,PASS),
// UNIT=SYSALLDA,SPACE=(CYL,(1,1)),
// DCB=(RECFM=FB,LRECL=80,BLKSIZE=&SYSLBLK)
//SYSIN DD DSN=*.PLI.SYSLIN,DISP=(OLD,DELETE)
//SYSPRINT DD SYSOUT=* 
//SYSOUT DD SYSOUT=* 

Figure 15. Cataloged Procedure IBMZCPG
Invoking a cataloged procedure

To invoke a cataloged procedure, specify its name in the PROC parameter of an EXEC statement. For example, to use the cataloged procedure IBMZC, you could include the following statement in the appropriate position among your other job control statements in the input stream:

```
//stepname EXEC PROC=IBMZC
```

You do not need to code the keyword PROC. If the first operand in the EXEC statement does not begin PGM= or PROC=, the job scheduler interprets it as the name of a cataloged procedure. The following statement is equivalent to that given above:

```
//stepname EXEC IBMZC
```

If you include the parameter MSGLEVEL=1 in your JOB statement, the operating system will include the original EXEC statement in its listing, and will add the statements from the cataloged procedure. In the listing, cataloged procedure statements are identified by XX or X/ as the first two characters; X/ signifies a statement that was modified for the current invocation of the cataloged procedure.

You might be required to modify the statements of a cataloged procedure for the duration of the job step in which it is invoked, either by adding DD statements or by overriding one or more parameters in the EXEC or DD statements. For example, cataloged procedures that invoke the compiler require the addition of a DD statement with the name SYSIN to define the data set containing the source statements. Also, whenever you use more than one standard link-edit procedure step in a job, you must modify all but the first cataloged procedure that you invoke if you want to run more than one of the load modules.

Specifying multiple invocations

You can invoke different cataloged procedures, or invoke the same cataloged procedure several times, in the same job. No special problems are likely to arise unless more than one of these cataloged procedures involves a link-edit procedure step, in which case you must take the following precautions to ensure that all your load modules can be run.

When the linkage editor creates a load module, it places the load module in the standard data set defined by the DD statement with the name SYSLMOD. When
the binder creates a program object, it places the program object in the PDSE defined by the DD statement with the name SYSLMOD. In the absence of a linkage editor NAME statement, the linkage editor or the binder uses the member name specified in the DSNAME parameter as the name of the module. In the standard cataloged procedures, the DD statement with the name SYSLMOD always specifies a temporary library &GOSET with the member name GO.

If you use the cataloged procedure IBMZCBG twice within the same job to compile, bind, and run two PL/I programs, and do not name each of the two program objects that the binder creates, the first program object runs twice, and the second one not at all.

To prevent this, use one of the following methods:

• Delete the library &GOSET at the end of the GO step. In the first invocation of the cataloged procedure at the end of the GO step, add a DD statement with the syntax:
  
  //GO.SYSLMOD DD DSN=&GOSET,
  // DISP=(OLD,DELETE)

• Modify the DD statement with the name SYSLMOD in the second and subsequent invocations of the cataloged procedure so as to vary the names of the load modules. For example:

  //BIND.SYSLMOD DD DSN=&GOSET(GO1)

  and so on.

• Use the NAME linkage editor option to give a different name to each program object and change each job step EXEC statement to specify the running of the program object with the name for that job step.

To assign a membername to the program object, you can use the linkage editor NAME option with the DSNAME parameter on the SYSLMOD DD statement. When you use this procedure, the membername must be identical to the name on the NAME option if the EXEC statement that runs the program refers to the SYSLMOD DD statement for the name of the module to be run.

Another option is to give each program a different name by using GOPGM on the EXEC procedure statement. For example:

  // EXEC IBMZCBG,GOPGM=GO2

### Modifying the PL/I cataloged procedures

You can modify a cataloged procedure temporarily by including parameters in the EXEC statement that invokes the cataloged procedure, or by placing additional DD statements after the EXEC statement. Temporary modifications apply only for the duration of the job step in which the procedure is invoked. They do not affect the master copy of the cataloged procedure in the procedure library.

Temporary modifications can apply to EXEC or DD statements in a cataloged procedure. To change a parameter of an EXEC statement, you must include a corresponding parameter in the EXEC statement that invokes the cataloged procedure. To change one or more parameters of a DD statement, you must include a corresponding DD statement after the EXEC statement that invokes the cataloged procedure. Although you cannot add a new EXEC statement to a cataloged procedure, you can always include additional DD statements.
**EXEC statement**

If a parameter of an EXEC statement that invokes a cataloged procedure has an unqualified name, the parameter applies to all the EXEC statements in the cataloged procedure. The effect on the cataloged procedure depends on the parameters, as follows:

- **PARM** applies to the first procedure step and nullifies any other PARM parameters.
- **COND** and **ACCT** apply to all the procedure steps.
- **TIME** and **REGION** apply to all the procedure steps and override existing values.

For example, the statement:

//stepname EXEC IBMZCBG,PARM='OFFSET',REGION=32M

- Invokes the cataloged procedure IBMZCBG.
- Substitutes the option **OFFSET** for **OBJECT** and **OPTIONS** in the EXEC statement for procedure step PLI.
- Nullifies the PARM parameter in the EXEC statement for procedure step BIND.
- Specifies a region size of 32M for all three procedure steps.

To change the value of a parameter in only one EXEC statement of a cataloged procedure, or to add a new parameter to one EXEC statement, you must identify the EXEC statement by qualifying the name of the parameter with the name of the procedure step. For example, to alter the region size for procedure step PLI only in the preceding example, code:

//stepname EXEC PROC=IBMZCBG,PARM='OFFSET',REGION.PLI=90M

A new parameter specified in the invoking EXEC statement overrides completely the corresponding parameter in the procedure EXEC statement.

You can nullify all the options specified by a parameter by coding the keyword and equal sign without a value. For example, to suppress the bulk of the linkage editor listing when invoking the cataloged procedure IBMZCBG, code:

//stepname EXEC IBMZCBG,PARM.BIND=

**DD statement**

To add a DD statement to a cataloged procedure, or to modify one or more parameters of an existing DD statement, you must include a DD statement with the form **procstepname.ddname** in the appropriate position in the input stream. If **ddname** is the name of a DD statement already present in the procedure step identified by **procstepname**, the parameters in the new DD statement override the corresponding parameters in the existing DD statement; otherwise, the new DD statement is added to the procedure step. For example, the statement:

//PLI.SYSIN DD *

adds a DD statement to the procedure step PLI of cataloged procedure IBMZC and the effect of the statement:

//PLI.SYSPRINT DD SYSOUT=C

is to modify the existing DD statement SYSPRINT (causing the compiler listing to be transmitted to the system output device of class C).
Overriding DD statements must appear after the procedure invocation and in the same order as they appear in the cataloged procedure. Additional DD statements can appear after the overriding DD statements are specified for that step.

To override a parameter of a DD statement, code either a revised form of the parameter or a replacement parameter that performs a similar function (for example, SPLIT for SPACE). To nullify a parameter, code the keyword and equal sign without a value. You can override DCB subparameters by coding only those you want to modify; that is, the DCB parameter in an overriding DD statement does not necessarily override the entire DCB parameter of the corresponding statement in the cataloged procedures.
Chapter 4. Compiling your program

This chapter describes how to invoke the compiler under z/OS UNIX System Services (z/OS UNIX) and the job control statements used for compiling under z/OS. The Language Environment SCEERUN data set must be accessible to the compiler when you compile your program.

Invoking the compiler under z/OS UNIX

To compile your program under the z/OS UNIX environment, use the pli command.

```
pli command_line_option input_file
```

command_line_option
You can specify a command_line_option in the following ways:
- option
- Option flag (usually a single letter preceded by -)

If you choose to specify compile-time options on the command line, the format differs from either setting them in your source file using %PROCESS statements. See “Specifying compile-time options under z/OS UNIX” on page 156.

input_file
The z/OS UNIX file specification for your program files. If you omit the extension from your file specification, the compiler assumes an extension of .pli. If you omit the complete path, the current directory is assumed.

Input files
The pli command compiles PL/I source files, links the resulting object files with any object files and libraries specified on the command line in the order indicated, and produces a single executable file.

The pli command accepts the following types of files:

Source files—.pli
All .pli files are source files for compilation. The pli command sends source files to the compiler in the order they are listed. If the compiler cannot find a specified source file, it produces an error message and the pli command proceeds to the next file if one exists.

All HFS source files must be line-delimited and encoded in EBCDIC.

Object files—.o
All .o files are object files. The pli command sends all object files along with library files to the linkage editor at link-edit time unless you specify the -c option. After it compiles all the source files, the compiler invokes the linkage editor to link-edit the resulting object files with any object files specified in the input file list, and produces a single executable output file.
Library files—.a

The pli command sends all of the library files (.a files) to the linkage editor at link-edit time.

Specifying compile-time options under z/OS UNIX

Enterprise PL/I provides compile-time options to change any of the compiler’s default settings. You can specify options on the command line, and they remain in effect for all compilation units in the file, unless %PROCESS statements in your source program override them.

See “Compile-time option descriptions” on page 5 for a description of these options.

When you specify options on the command line, they override the default settings of the option. They are overridden by options set in the source file.

You can specify compile-time options on the command line in three ways:

- -q option_keyword (compiler-specific)
- Single and multiletter flags
- -q+/u/myopts.txt

-q option_keyword

You can specify options on the command line using the -q option format.

You can have multiple -q options on the same command line, but they must be separated by blanks. Option keywords can appear in either uppercase or lowercase, but you must specify the -q in lowercase.

Some compile-time options allow you to specify suboptions. These suboptions are indicated on the command line with an equal sign following the -q option_keyword. Multiple suboptions must be separated with a colon(;) and no intervening blanks.

An option, for example, that contains multiple suboptions is RULES ("RULES" on page 69). To specify RULES(LAXDCL) on the command line, you would enter:

-q rules=ibm:laxdcl

The LIMITS option ("LIMITS" on page 46) is slightly more complex since each of its suboptions also has an argument. You would specify LIMITS(EXTNAME(31),FIXEDDEC(15)) on the command line as shown in the following example:

-q limits=extname=31:fixeddec=15

Single and multiletter flags

The z/OS UNIX family of compilers uses a number of common conventional flags. Each language has its own set of additional flags.

Some flag options have arguments that form part of the flag, for example:
Specifying compile-time options

pli samp.pli -I/home/test3/include

In this case, /home/test3/include is an include directory to be searched for INCLUDE files.

Each flag option should be specified as a separate argument.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>-c</td>
<td>Compile only.</td>
</tr>
<tr>
<td>-e</td>
<td>Create names and entries for a fetchable load module.</td>
</tr>
<tr>
<td>-I&lt;dir&gt;'</td>
<td>Add path &lt;dir&gt; to the directories to be searched for INCLUDE files. -I must be followed by a path and only a single path is allowed per -I option. To add multiple paths, use multiple -I options. There shouldn't be any spaces between -I and the path name.</td>
</tr>
<tr>
<td>-O, -O2</td>
<td>Optimize generated code. This option is equivalent to -qOPT=2.</td>
</tr>
<tr>
<td>-q&lt;option&gt;'</td>
<td>Pass it to the compiler. &lt;option&gt; is a compile-time option. Each option should be delimited by a comma and each suboption should be delimited by an equal sign or colon. There shouldn't be any spaces between -q and &lt;option&gt;.</td>
</tr>
<tr>
<td>-v</td>
<td>Display compile and link steps and execute them.</td>
</tr>
<tr>
<td>-#</td>
<td>Display compile and link steps, but do not execute them.</td>
</tr>
</tbody>
</table>

Note: You must specify an argument where indicated; otherwise, the results are unpredictable.

Invoking the compiler under z/OS using JCL

Although you will probably use cataloged procedures rather than supply all the JCL required for a job step that invokes the compiler, you should be familiar with these statements so that you can make the best use of the compiler and, if necessary, override the statements of the cataloged procedures.

So-called "batch compilation", whereby one compilation produced more than one object deck, is not supported.

Invoking the compiler via BPXBATCH is also not supported.

The following section describes the JCL needed for compilation. The IBM-supplied cataloged procedures described in “IBM-supplied cataloged procedures” on page 139 contain these statements. You need to code them yourself only if you are not using the cataloged procedures.

EXEC statement

The basic EXEC statement is:

```
//stepname EXEC PGM
```

512K is required for the REGION parameter of this statement.

If you compile your programs with optimization turned on, the REGION size (and time) required may be much, much larger.
The PARM parameter of the EXEC statement can be used to specify one or more of the optional facilities provided by the compiler. These facilities are described under “Specifying options in the EXEC statement” on page 161. See Chapter 1, “Using compiler options and facilities,” on page 5 for a description of the options.

### DD statements for the standard data sets

The compiler requires several standard data sets, the number of data sets depends on the optional facilities specified. You must define these data sets in DD statements with the standard ddnames shown, together with other characteristics of the data sets, in Table 13. The DD statements SYSIN, SYSUT1, and SYSPRINT are always required.

You can store any of the standard data sets on a direct access device, but you must include the SPACE parameter in the DD statement. This parameter defines the data set to specify the amount of auxiliary storage required. The amount of auxiliary storage allocated in the IBM-supplied cataloged procedures should suffice for most applications.

**Table 13. Compiler standard data sets**

<table>
<thead>
<tr>
<th>Standard DDNAME</th>
<th>Contents of data set</th>
<th>Possible device classes</th>
<th>Record format (RECFM)</th>
<th>Record size (LRECL)</th>
</tr>
</thead>
<tbody>
<tr>
<td>SYSDEBUG</td>
<td>TEST(SEPARATE) output</td>
<td>SYSDA, F,FB</td>
<td>&gt;=80 and &lt;=1024</td>
<td></td>
</tr>
<tr>
<td>SYSDEFSD</td>
<td>XINFO(DEF) output</td>
<td>SYSDA, F,FB</td>
<td>128</td>
<td></td>
</tr>
<tr>
<td>SYSIN</td>
<td>Input to the compiler</td>
<td>SYSSQ, F,FB,U, VB, V</td>
<td>&lt;101(100) or &lt;105(104)</td>
<td></td>
</tr>
<tr>
<td>SYSLIB</td>
<td>Source statements for INCLUDE files</td>
<td>SYSDA, F,FB,U, V</td>
<td>&lt;101 or &lt;105</td>
<td></td>
</tr>
<tr>
<td>SYSLIN</td>
<td>Object module</td>
<td>SYSSQ, FB</td>
<td>80</td>
<td></td>
</tr>
<tr>
<td>SYSPRINT</td>
<td>Listing, including messages</td>
<td>SYSSQ, VBA</td>
<td>137</td>
<td></td>
</tr>
<tr>
<td>SYSUT1</td>
<td>Temporary workfile</td>
<td>SYSDA, F</td>
<td>4051</td>
<td></td>
</tr>
<tr>
<td>SYSXMI</td>
<td>XINFO(XMI) output</td>
<td>SYSDA, VB</td>
<td>16383</td>
<td></td>
</tr>
<tr>
<td>SYSXMLSD</td>
<td>XINFO(XML) output</td>
<td>SYSDA, VB</td>
<td>16383</td>
<td></td>
</tr>
<tr>
<td>SYSADATA</td>
<td>XINFO(MSG) output</td>
<td>SYSDA, U</td>
<td>1024</td>
<td></td>
</tr>
</tbody>
</table>

1. The possible device classes are:
   - **SYSSQ**: Sequential device
   - **SYSDA**: Direct access device

   Block size can be specified except for SYSUT1. The block size and logical record length for SYSUT1 is chosen by the compiler.

**Note:**

The only value for compile-time SYSPRINT that can be overridden is BLKSIZE.
Input (SYSIN)
Input to the compiler must be a data set defined by a DD statement with the name SYSIN. This data set must have CONSECUTIVE organization. The input must be one or more external PL/I procedures. If you want to compile more than one external procedure in a single job or job step, precede each procedure, except possibly the first, with a %PROCESS statement.

80-byte records are commonly used as the input medium for PL/I source programs. The input data set can be on a direct access device or some other sequential media. The input data set can contain either fixed-length records (blocked or unblocked), variable-length records (coded or uncoded), or undefined-length records. The maximum record size is 100 bytes.

The maximum number of lines in the input file is 999,999.

When data sets are concatenated for input to the compiler, the concatenated data sets must have similar characteristics (for example, block size and record format).

Output (SYSLIN, SYSPUNCH)
Output in the form of one or more object modules from the compiler will be stored in the data set SYSLIN if you specify the OBJECT compile-time option. This data set is defined by the DD statement.

The object module is always in the form of 80-byte fixed-length records, blocked or unblocked. If the BLKSIZE is specified for SYSLIN and is something other than 80, then the LRECL must be specified as 80.

The SYSLIN DD must name either a temporary dataset or a permanent dataset: it cannot specify a concatenation of datasets of any type.

The SYSLIN DD must specify a sequential dataset, not a PDS or PDSE.

The data set defined by the DD statement with the name SYSPUNCH is also used to store the output from the preprocessor if you specify the MDECK compile-time option.

Temporary workfile (SYSUT1)
The compiler requires a data set for use as a temporary workfile. It is defined by a DD statement with the name SYSUT1, and is known as the spill file. It must be on a direct access device, and must not be allocated as a multi-volume data set.

The spill file is used as a logical extension to main storage and is used by the compiler and by the preprocessor to contain text and dictionary information. The LRECL and BLKSIZE for SYSUT1 is chosen by the compiler based on the amount of storage available for spill file pages.

The DD statements given in this publication and in the cataloged procedures for SYSUT1 request a space allocation in blocks of 1024 bytes. This is to insure that adequate secondary allocations of direct access storage space are acquired.

Listing (SYSPRINT)
The compiler generates a listing that includes all the source statements that it processed, information relating to the object module, and, when necessary, messages. Most of the information included in the listing is optional, and you can specify those parts that you require by including the appropriate compile-time
options. The information that can appear, and the associated compile-time options, are described under “Using the compiler listing” on page 96.

You must define the data set, in which you want the compiler to store its listing, in a DD statement with the name SYSPRINT. This data set must have CONSECUTIVE organization. Although the listing is usually printed, it can be stored on any sequential or direct access device. For printed output, the following statement will suffice if your installation follows the convention that output class A refers to a printer:

//SYSPRINT DD SYSOUT=A

Source Statement Library (SYSLIB)

If you use the %INCLUDE statement to introduce source statements into the PL/I program from a library, you can either define the library in a DD statement with the name SYSLIB, or you can choose your own ddname (or ddnames) and specify a ddname in each %INCLUDE statement.

The DD should specify a PDS or PDSE, but not the actual member. For example to include the file HEADER from the library SYSLIB using the dataset INCLUDE.PLI, the %INCLUDE statement would like

%INCLUDE HEADER;

or

%INCLUDE SYSLIB( HEADER );

The DD statement should be

SYSLIB DD DISP=SHR,DSN=INCLUDE.PLI

but it should not be

SYSLIB DD DISP=SHR,DSN=INCLUDE.PLI(HEADER)

All %INCLUDE files must have the same record format (fixed, variable, undefined), the same logical record length, and the same left and right margins as the SYSIN source file.

The BLOCKSIZE of the library must be less than or equal to 32,760 bytes.

The maximum number of lines in any one include file is 999,999.

Specifying options

For each compilation, the IBM-supplied or installation default for a compile-time option applies unless it is overridden by specifying the option in a %PROCESS statement or in the PARM parameter of an EXEC statement.

An option specified in the PARM parameter overrides the default value, and an option specified in a %PROCESS statement overrides both that specified in the PARM parameter and the default value.

Note: When conflicting attributes are specified either explicitly or implicitly by the specification of other options, the latest implied or explicit option is accepted. No diagnostic message is issued to indicate that any options are overridden in this way.
Specifying options in the EXEC statement

To specify options in the EXEC statement, code `PARM=` followed by the list of options, in any order separating the options with commas and enclosing the list within single quotation marks, for example:

```
//STEP1 EXEC PGM=IBMZPLI,PARM='OBJECT,LIST'
```

Any option that has quotation marks, for example `MARGINI('c')`, must have the quotation marks duplicated. The length of the option list must not exceed 100 characters, including the separating commas. However, many of the options have an abbreviated syntax that you can use to save space. If you need to continue the statement onto another line, you must enclose the list of options in parentheses (instead of in quotation marks) enclose the options list on each line in quotation marks, and ensure that the last comma on each line except the last line is outside of the quotation marks. An example covering all the above points is as follows:

```
//STEP1 EXEC PGM=IBMZPLI,PARM=('AG,A',
   'C,F(I)',
   'M,MI(\'X\'),NEST,STG,X')
```

If you are using a cataloged procedure, and want to specify options explicitly, you must include the `PARM` parameter in the EXEC statement that invokes it, qualifying the keyword `PARM` with the name of the procedure step that invokes the compiler. For example:

```
//STEP1 EXEC nnnnnnn,PARM.PLI='A,LIST'
```

Specifying options in the EXEC statement using an options file

Another way to specify options in the EXEC statement is by declaring all your options in an options file and coding the following:

```
//STEP1 EXEC PGM=IBMZPLI,PARM='+DD:OPTIONS'
```

This method allows you to provide a consistent set of options that you frequently use. This is especially effective if you want other programmers to use a common set of options. It also gets you past the 100-character limit.

The `MARGINS` option does not apply to options files: the data in column 1 will be read as part of the options. Also, if the file is F-format, any data after column 72 will be ignored.

The parm string can contain "normal" options and can point to more than one options file. For instance, to specify the option `LIST` as well as options from both the file in the `GROUP DD` and in the `PROJECT DD`, you could specify

```
PARM='LIST +DD:GROUP +DD:PROJECT'
```

The options in the `PROJECT DD` file would have precedence over options in the `GROUP DD` file.

Also, in this case, the `LIST` option might be turned off by a `NOLIST` option specified in either of the options files. To insure that the `LIST` option is on, you could specify

```
PARM='+DD:GROUP +DD:PROJECT LIST'
```

Options files may also be used under z/OS UNIX. For example, in z/OS UNIX, to compile sample.pli with options from the file `/u/pli/group.opt`, you would specify

```
pli -q+/u/pli/group.opt sample.pli
```
Earlier releases of the compiler used the character '@' as the trigger character that preceded the options file specification. This character is not part of the invariant set of EBCDIC code points, and for that reason the character '+', which is invariant, is preferred. However, the '@' character may be still be used as long as it is specified with the hex value '7C'x.
Chapter 5. Link-editing and running

After compilation, your program consists of one or more object modules that contain unresolved references to each other, as well as references to the Language Environment run-time library. These references are resolved during link-editing (statically) or during execution (dynamically). There are two ways to link-edit statically:

1. Use the prelinker prior to the traditional link step
2. Link without the prelinker, which is similar to linking with PL/I for MVS & VM except that depending on which compile-time options you use, you may now need to use a PDSE to hold the resultant load module.

After you compile your PL/I program, the next step is to link and run your program with test data to verify that it produces the results you expect. When using Enterprise PL/I we recommend you select the method of linking without the prelinker (as described in Item 2 above).

Language Environment provides the run-time environment and services you need to execute your program. For instructions on linking and running PL/I and all other Language Environment-conforming language programs, see z/OS Language Environment Programming Guide. For information about migrating your existing PL/I programs to Language Environment, see Enterprise PL/I for z/OS Compiler and Run-Time Migration Guide.

Link-edit considerations

If you compile with the option RENT or the option LIMITS(EXTNAME(n)) with n > 8, then you must use the prelinker or use a PDSE for your linker output.

Using the binder

You must place the binder output into a PDSE.

When linking a DLL, you must specify any needed definition side-decks during the bind step.

You can use the binder in place of the prelinker and linkage-editor, with the following exceptions:

- Releases of CICS prior to CICS Transaction Server 1.3 do not support PDSEs.
  From CICS Transaction Server 1.3 onwards, there is support in CICS for PDSEs.
  Refer to the CICS Transaction Server for z/OS Release Guide, GC34-5701 for a list of prerequisite APAR fixes.
- MTF does not support PDSEs.

Using the prelinker

If you use the prelinker, you must prelink together in one job step all object decks defining external references in any of the input object decks.

For instance, if A and B are separately compiled programs and A statically calls B, then you cannot prelink A and B separately and then later link them together. Instead, you must link A and B together in one prelink job.
Using the ENTRY card

If you are building a module that will be fetched and which has an Enterprise PL/I routine as its entry point, then the ENTRY card should specify the name of that PL/I entry point. If the module is to be fetched from Enterprise PL/I, you may specify CEESTART on the ENTRY card, although this is strongly not recommended. However, if the module is to be fetched from COBOL or assembler, then the ENTRY card absolutely must specify the name of the PL/I entry point into the module and not CEESTART.

Run-time considerations

You can specify run-time options as parameters passed to the program initialization routine. You can also specify run-time options in the PLIXOPT variable. It might also prove beneficial, from a performance standpoint, if you alter your existing programs by using the PLIXOPT variable to specify your run-time options and recompiling your programs. For a description of using PLIXOPT, see z/OS Language Environment Programming Guide.

To simplify input/output at the terminal, various conventions have been adopted for stream files that are assigned to the terminal. Three areas are affected:

1. Formatting of PRINT files
2. The automatic prompting feature
3. Spacing and punctuation rules for input.

Note: No prompting or other facilities are provided for record I/O at the terminal, so you are strongly advised to use stream I/O for any transmission to or from a terminal.

Formatting conventions for PRINT files

When a PRINT file is assigned to the terminal, it is assumed that it will be read as it is being printed. Spacing is therefore reduced to a minimum to reduce printing time. The following rules apply to the PAGE, SKIP, and ENDPAGE keywords:

- PAGE options or format items result in three lines being skipped.
- SKIP options or format items larger than SKIP (2) result in three lines being skipped. SKIP (2) or less is treated in the usual manner.
- The ENDPAGE condition is never raised.

Changing the format on PRINT files

If you want normal spacing to apply to output from a PRINT file at the terminal, you must supply your own tab table for PL/I. This is done by declaring an external structure called PLITABS in the main program or in a program linked with the main program and initializing the element PAGELENGTH to the number of lines that can fit on your page. This value differs from PAGESIZE, which defines the number of lines you want to print on the page before ENDPAGE is raised (see Figure 17 on page 165). If you require a PAGELENGTH of 64 lines, declare PLITABS as shown in Figure 16 on page 165. For information on overriding the tab table, see “Overriding the tab control table” on page 234.

If your code contains a declare for PLITABS, not only must the values be valid, but the first field in the PLITABS structure must also be valid. This field is supposed to hold the offset to the field specifying the number of tabs set by the structure, and the Enterprise PL/I library code will not work correctly if this is not true.
Automatic prompting

When the program requires input from a file that is associated with a terminal, it issues a prompt. This takes the form of printing a colon on the next line and then skipping to column 1 on the line following the colon. This gives you a full line to enter your input, as follows:

```
: (space for entry of your data)
```

---

**Figure 16. Declaration of PLITABS.** This declaration gives the standard page size, line size and tabulating positions.

```
DCL 1 PLITABS STATIC EXTERNAL,
( 2 OFFSET INIT (14),
  2 PAGESIZE INIT (60),
  2 LINESIZE INIT (120),
  2 PAGELENGTH INIT (64),
  2 FILL1 INIT (0),
  2 FILL2 INIT (0),
  2 FILL3 INIT (0),
  2 NUMBER_OF_TABS INIT (5),
  2 TAB1 INIT (25),
  2 TAB2 INIT (49),
  2 TAB3 INIT (73),
  2 TAB4 INIT (97),
  2 TAB5 INIT (121)) FIXED BIN (15,0);
```

---

**Figure 17. PAGELENGTH and PAGESIZE.** PAGELENGTH defines the size of your paper, PAGESIZE the number of lines in the main printing area.

---

**Automatic prompting**

When the program requires input from a file that is associated with a terminal, it issues a prompt. This takes the form of printing a colon on the next line and then skipping to column 1 on the line following the colon. This gives you a full line to enter your input, as follows:

```
: (space for entry of your data)
```
This type of prompt is referred to as a primary prompt.

**Overriding automatic prompting**
You can override the primary prompt by making a colon the last item in the request for the data. You cannot override the secondary prompt. For example, the two PL/I statements:

```pli
PUT SKIP EDIT ('ENTER TIME OF PERIHELION') (A);
GET EDIT (PERITIME) (A(10));
```

result in the terminal displaying:

```
ENTER TIME OF PERIHELION
    :   (automatic prompt)
(space for entry of data)
```

However, if the first statement has a colon at the end of the output, as follows:

```pli
PUT EDIT ('ENTER TIME OF PERIHELION:') (A);
```

the sequence is:

```
ENTER TIME OF PERIHELION: (space for entry of data)
```

**Note:** The override remains in force for only one prompt. You will be automatically prompted for the next item unless the automatic prompt is again overridden.

**Punctuating long input lines**

**Line continuation character**
To transmit data that requires 2 or more lines of space at the terminal as one data-item, type an SBCS hyphen as the last character in each line except the last line. For example, to transmit the sentence “this data must be transmitted as one unit.” you enter:

```pli
:'this data must be transmitted -
+:as one unit.'
```

Transmission does not occur until you press ENTER after “unit.”. The hyphen is removed. The item transmitted is called a “logical line.”

**Note:** To transmit a line whose last data character is a hyphen or a PL/I minus sign, enter two hyphens at the end of the line, followed by a null line as the next line. For example:

```pli
xyz--
(press ENTER only, on this line)
```

**Punctuating GET LIST and GET DATA statements**
For GET LIST and GET DATA statements, a comma is added to the end of each logical line transmitted from the terminal, if the programmer omits it. Thus there is no need to enter blanks or commas to delimit items if they are entered on separate logical lines. For the PL/I statement GET LIST(A,B,C); you can enter at the terminal:

```pli
:1
+:2
+:3
```
This rule also applies when entering character-string data. Therefore, a character string must transmit as one logical line. Otherwise, commas are placed at the break points. For example, if you enter:

```
: 'COMMAS SHOULD NOT BREAK
+:UP A CLAUSE.'
```

the resulting string is: “COMMAS SHOULD NOT BREAK, UP A CLAUSE.” The comma is not added if a hyphen was used as a line continuation character.

### Automatic padding for GET EDIT

For a GET EDIT statement, there is no need to enter blanks at the end of the line. The data will be padded to the specified length. Thus, for the PL/I statement:

```
GET EDIT (NAME) (A(15));
```

you can enter the 5 characters SMITH. The data will be padded with ten blanks so that the program receives the fifteen characters:

```
'SMITH
```

**Note:** A single data item must transmit as a logical line. Otherwise, the first line transmitted will be padded with the necessary blanks and taken as the complete data item.

### Use of SKIP for terminal input

All uses of SKIP for input are interpreted as SKIP(1) when the file is allocated to the terminal. SKIP(1) is treated as an instruction to ignore all unused data on the currently available logical line.

### ENDFILE

The end-of-file can be entered at the terminal by keying in a logical line that consists of the two characters “/*”. Any further attempts to use the file without closing it result in the ENDFILE condition being raised.

### SYSPRINT considerations

The PL/I standard SYSPRINT file is shared by multiple enclaves within an application. You can issue I/O requests, for example STREAM PUT, from the same or different enclaves. These requests are handled using the standard PL/I SYSPRINT file as a file which is common to the entire application. The SYSPRINT file is implicitly closed only when the application terminates, not at the termination of the enclave.

The standard PL/I SYSPRINT file contains user-initiated output only, such as STREAM PUTs. Run-time library messages and other similar diagnostic output are directed to the Language Environment MSGFILE. See the *z/OS Language Environment Programming Guide* for details on redirecting SYSPRINT file output to the Language Environment MSGFILE.

To be shared by multiple enclaves within an application, the PL/I SYSPRINT file must be declared as an EXTERNAL FILE constant with a file name of SYSPRINT and also have the attributes STREAM and OUTPUT as well as the (implied) attribute of PRINT, when OPENed. This is the standard SYSPRINT file as defaulted by the compiler.

There exists only one standard PL/I SYSPRINT FILE within an application and this file is shared by all enclaves within the application. For example, the SYSPRINT file can be shared by multiple nested enclaves within an application or
by a series of enclaves that are created and terminated within an application by the
Language Environment preinitialization function. To be shared by an enclave
within an application, the PL/I SYSPRINT file must be declared in that enclave.
The standard SYSPRINT file cannot be shared by passing it as a file argument
between enclaves. The declared attributes of the standard SYSPRINT file should be
the same throughout the application, as with any EXTERNALy declared constant.
PL/I does not enforce this rule. Both the TITLE option and the
MSGFILE(SYSPRINT) option attempt to route SYSPRINT to another data set. As
such, if the two options were used together, there will be a conflict and the TITLE
option will be ignored.

Having a common SYSPRINT file within an application can be an advantage to
applications that utilize enclaves that are closely tied together. However, since all
enclaves in an application write to the same shared data set, this might require
some coordination among the enclaves.

The SYSPRINT file is opened (implicitly or explicitly) when first referenced within
an enclave of the application. When the SYSPRINT file is CLOSED, the file
resources are released (as though the file had never been opened) and all enclaves
are updated to reflect the closed status.

If SYSPRINT is utilized in a multiple enclave application, the LINENO built-in
function only returns the current line number until after the first PUT or OPEN in
an enclave has been issued. This is required in order to maintain full compatibility
with old programs.

The COUNT built-in function is maintained at an enclave level. It always returns a
value of zero until the first PUT in the enclave is issued. If a nested child enclave
is invoked from a parent enclave, the value of the COUNT built-in function is
undefined when the parent enclave regains control from the child enclave.

The TITLE option can be used to associate the standard SYSPRINT file with
different operating system data sets, keeping in mind that a particular open
association has to be closed before another one is opened. This association is
retained across enclaves for the duration of the open.

PL/I condition handling associated with the standard PL/I SYSPRINT file retains
its current semantics and scope. For example, an ENDPAGE condition raised
within a child enclave will only invoke an established ON-unit within that child
enclave. It does not cause invocation of an ON-unit within the parent enclave.

The tabs for the standard PL/I SYSPRINT file can vary when PUTs are done from
different enclaves, if the enclaves contain a user PLITABS table.

If the PL/I SYSPRINT file is utilized as a RECORD file or as a STREAM INPUT
file, PL/I supports it at an individual enclave or task level, but not as a shareable
file among enclaves. If the PL/I SYSPRINT file is open at the same time with
different file attributes (e.g. RECORD and STREAM) in different enclaves of the
same application, results are unpredictable.

SYSPRINT may also be shared between code compiled by Enterprise PL/I and by
older PL/I compilers, but the following must all apply:
• SYSPRINT must be declared as STREAM OUTPUT
• the application must not be running under TSO

Enterprise PL/I for z/OS Programming Guide
• if the runtime option MSGFILE(SYSPRINT) is in effect, then there must be no preinitialized programs and no stored procedures in the application

**Using MSGFILE(SYSPRINT)**

Any file attributes that are specified in the ENVIRONMENT option of the file declaration for SYSPRINT STREAM PRINT are ignored.

Any attributes that are specified on the OPEN statement for SYSPRINT are ignored.

When you use the OPEN statement to open the PL/I SYSPRINT STREAM PRINT file, the file is marked as opened in the PL/I control blocks, but it is actually opened by the Language Environment.

When you use the CLOSE statements to close the PL/I SYSPRINT STREAM PRINT file, the file is marked as closed in the PL/I control blocks, but Language Environment still keeps it open.

The synchronization between the Language Environment messages and PL/I user-specified output is not provided, so the order of the output is unpredictable.

The use of MSGFILE(SYSPRINT) restricts the line size specified by the LINESIZE option to a maximum of 225 characters.

---

### Using FETCH in your routines

In Enterprise PL/I, you can fetch routines compiled by PL/I, C, COBOL or assembler.

**Fetching Enterprise PL/I routines**

Almost all the restrictions imposed by the older PL/I compilers on fetched modules have been removed. So a FETCHed module can now:

- Fetch other modules
- Perform any I/O operations on any PL/I file. The file can be opened either by the fetched module, by the main module, or by some other fetched module.
- ALLOCATE and FREE its own CONTROLLED variables

There are, however, a few restrictions on a Enterprise PL/I module that is to be fetched. These restrictions are:

1. **OPTIONS(FETCHABLE)** should be specified on the PROCEDURE statement of the fetched routine if there is no ENTRY card provided during the link-edit step.

2. The ENTRY card should specify the name of that PL/I entry point
   - If the module is to be fetched from Enterprise PL/I, you may specify CEESTART on the ENTRY card, although this is strongly not recommended.
   - However, if the module is to be fetched from COBOL or assembler, then the ENTRY card absolutely must specify the name of the PL/I entry point into the module and not CEESTART.

3. If the RENT compiler option was used to compile any of the fetched code, then the module must be linked as a DLL.

4. If the NORENT compiler option was used to compile the fetching code, then the fetched module must satisfy at least one of the following:
• It is a MAIN module
• It consists only of NORENT code
• It has as its entry point code compiled with the C or Enterprise PL/I compiler with the NORENT option in effect. In this case, the module may also contain code compiled with the RENT option, but calling that code is not supported

5. If the RENT compiler option was used to compile the fetching code, then the ENTRY that is fetched must not be declared in the fetching module as OPTIONS(COBOL) or OPTIONS(ASM). If you want to avoid passing descriptors in this situation, you should specify the OPTIONS(NODESCRIPTOR) attribute on the ENTRY declare.

6. An Enterprise PL/I routine can not fetch itself.

In the case of a nonreentrant and nonreusable module that is loaded multiple times, the order of processing occurs in a last-in first-out order.
For example, if Program A loads module LOADMODA, then calls Program B, which also loads LOADMODA, then issues a DELETE against LOADMODA, the copy of LOADMODA to be deleted is the one associated with Program B. At this point, the copy of LOADMODA associated with Program A still exists.
In other words, the DELETE requested against LOADMODA will release the last copy that was loaded, regardless of which program issues the request.

NORENT WRITABLE code is serially usable, and for that reason, the pointer that is used to represent a FETCHABLE constant is zeroed out in the prologue code of any NORENT WRITABLE routine. While this insures that the code is serially re usable while also providing the correct PL/I semantics, it does impose a restriction on the use of FETCH with TITLE in NORENT WRITABLE code: if a routine that did a FETCH A TITLE('B') is exited and reentered, then it must re-execute the FETCH A TITLE('B'), before executing any CALL A statements (otherwise it would do an implicit FETCH of A (but without any TITLE) before making the CALL).

As an illustration of these restrictions, consider the compiler user exit. If you specify the EXIT compile-time option, the compiler will fetch and call a Enterprise PL/I module named IBMUEXIT.

First note that the compiler user exit must be compiled with the RENT option since the compiler expects it to be a DLL.
In accordance with Item 1 above, the PROCEDURE statement for this routine looks like:

```pli
ibmuexit:
    proc ( addr_Userexit_Interface_Block,
           addr_Request_Area )
    options( fetchable );
    dcl addr_Userexit_Interface_Block pointer byvalue;
    dcl addr_Request_Area pointer byvalue;
```

In accordance with Item 3 above, the linker option DYNAM=DLL must be specified when linking the user exit into a DLL. The DLL must be linked either into a PDSE or into a temporary dataset (in which case DSNTYPE=LIBRARY must be specified on the SYSLMOD DD statement).
All the JCL to compile, link, and invoke the user exit is given in the JCL below in Figure 18 on page 172. The one significant difference between the sample below and the code excerpts above is that, in the code below, the fetched user exit does not receive two BYVALUE pointers to structures, but instead it receives the two structures BYADDR. In order to make this change work, the code specifies OPTIONS(NODESCRIPTOR) on each of its PROCEDURE statements.
/*
** compile the user exit
*****************************************************************************/
//PLEXIT EXEC PGM=IBMZPLI,
//STEPLIB DD DSN=IBMZ.V4R5MO.SIBMZCMP,DISP=SHR
//SYSPRINT DD SYSOUT=*
//SYSLIN DD DSN=&&LOADSET,DISP=(MOD,PASS),UNIT=SYSSQ,
//   SPACE=(CYL,(3,1))
//SYSUT1 DD DSN=&&SYSUT1,UNIT=SYSDA,
//   SPACE=(*,*,CONTIG,ROUND),DCB=BLKSIZE=1024
//SYSIN DD *
*Process or('!') not('!');
*Process limits(extname(31));
*Process RENT;
***************************************************************************/

/******************************************************************************/
/*
** NAME - IBMUEXIT.PLI
**
** User-exit sample program.
**
** Licensed Materials - Property of IBM
** 5639-A83, 5639-A24 (C) Copyright IBM Corp. 1992,2015.
** All Rights Reserved.
** US Government Users Restricted Rights-- Use, duplication or
** disclosure restricted by GSA ADP Schedule Contract with
** IBM Corp.
**
** DISCLAIMER OF WARRANTIES
** The following "enclosed" code is sample code created by IBM
** Corporation. This sample code is not part of any standard
** IBM product and is provided to you solely for the purpose of
** assisting you in the development of your applications. The
** code is provided "AS IS", without warranty of any kind.
** IBM shall not be liable for any damages arising out of your
** use of the sample code, even if IBM has been advised of the
** possibility of such damages.
**
******************************************************************************/

/******************************************************************************/
/*
** During initialization, IBMUEXIT is called. It reads
** information about the messages being screened from a text
** file and stores the information in a hash table. IBMUEXIT
** also sets up the entry points for the message filter service
** and termination service.
**
** For each message generated by the compiler, the compiler
** calls the message filter registered by IBMUEXIT. The filter
** looks the message up in the hash table previously created.
**
** The termination service is called at the end of the compile
** but does nothing. It could be enhanced to generates reports
** or do other cleanup work.
**
******************************************************************************/

Figure 18. Sample JCL to compile, link, and invoke the user exit
pack: package exports(*);

Dcl
1 Uex UIB native Based( null() ),
 2 Uex UIB Length fixed bin(31),
2 Uex UIB Exit_token pointer, /* for user exit's use*/
2 Uex UIB User_char_str pointer, /* to exit option str */
2 Uex UIB_User_char_len fixed bin(31),
2 Uex UIB Filename_str pointer, /* to source filename */
2 Uex UIB Filename_len fixed bin(31),
2 Uex UIB return_code fixed bin(31), /* set by exit procs */
2 Uex UIB reason_code fixed bin(31), /* set by exit procs */
2 Uex UIB Exit_Routs, /* exit entries setat initialization */
3 ( Uex UIB Termination,
  Uex UIB Message_Filter, /* call for each msg */
    *, *, *, *
) limited entry ( *
  *, /* to Uex UIB */
  * /* to a request area */
);

/***************************************************************************/
/* */
/* Request Area for Initialization exit */
/* */
/***************************************************************************/
Dcl 1 Uex ISA native based( null() ),
 2 Uex ISA Length fixed bin(31);

/***************************************************************************/
/* */
/* Request Area for Message_Filter exit */
/* */
/***************************************************************************/
Dcl 1 Uex MFA native based( null() ),
 2 Uex MFA Length fixed bin(31),
 2 Uex MFA_Facility_Id char(3),
 2 * char(1),
 2 Uex MFA_Message_no fixed bin(31),
 2 Uex MFA_Severity fixed bin(15),
 2 Uex MFA_New_Severity fixed bin(15); /* set by exit proc */

/***************************************************************************/
/* */
/* Request Area for Terminate exit */
/* */
/***************************************************************************/
Dcl 1 Uex TSA native based( null() ),
 2 Uex TSA Length fixed bin(31);

Sample JCL to compile, link, and invoke the user exit (continued)

Chapter 5. Link-editing and running 173
Severity Codes

dcl uex_Severity_Normal fixed bin(15) value(0);
dcl uex_Severity_Warning fixed bin(15) value(4);
dcl uex_Severity_Error fixed bin(15) value(8);
dcl uex_Severity_Severe fixed bin(15) value(12);
dcl uex_Severity_Unrecoverable fixed bin(15) value(16);

Return Codes

dcl uex_Return_Normal fixed bin(15) value(0);
dcl uex_Return_Warning fixed bin(15) value(4);
dcl uex_Return_Error fixed bin(15) value(8);
dcl uex_Return_Severe fixed bin(15) value(12);
dcl uex_Return_Unrecoverable fixed bin(15) value(16);

Reason Codes

dcl uex_Reason_Output fixed bin(15) value(0);
dcl uex_Reason_Suppress fixed bin(15) value(1);

dcl hashsive fixed bin(15) value(97);
dcl hashtable(0:hashsize-1) ptr init((hashsize) null());

dcl 1 message_item native based,
    2 message_Info,
    3 facid char(3),
    3 msgno fixed bin(31),
    3 newsev fixed bin(15),
    3 reason fixed bin(31),
    2 link pointer;

Sample JCL to compile, link, and invoke the user exit (continued)
ibmuexit: proc ( ue, ia )
  options( fetchable nodescriptor );

dcl 1 ue like uex_UIb byaddr;
dcl 1 ia like uex_Isa byaddr;

dcl sysuexit file stream input env(recsize(80));
dcl p pointer;
dcl bucket fixed bin(31);
dcl based_Chars char(8) based;
dcl title_Str char(8) var;

ue.uex_UIb_Message_Filter = message_Filter;
ue.uex_UIb_Termination = exitterm;

on undefinedfile(sysuexit)
begin;
  put edit ('** User exit unable to open exit file ') (A) skip;
  put skip;
  signal error;
end;

if ue.uex_UIb_User_Char_Len = 0 then
  do;
    open file(sysuexit);
  end;
else
  do;
    title_Str = substr( ue.uex_UIb_User_Char_Str->based_Chars,
                       1, ue.uex_UIb_User_Char_Len );
    open file(sysuexit) Title(title_Str);
  end;

on error, endfile(sysuexit)
goto done;

allocate message_item set(p);

/***************************************************************/
/* */
/* Skip header lines and read first data line */
/***************************************************************/

get file(sysuexit) list(p->message_info) skip(3);

Sample JCL to compile, link, and invoke the user exit (continued)
do loop;

/*************************************************************/
/* */
/* Put message information in hash table */
/* */
/*************************************************************/

bucket = mod(p->msgno, hashsize);
p->link = hashtable(bucket);
hashtable(bucket) = p;

/*************************************************************/
/* */
/* Read next data line */
/* */
/*************************************************************/

allocate message_item set(p);
get file(sysuexit) skip;
get file(sysuexit) list(p->message_info);
end;

/*************************************************************/
/* */
/* Clean up */
/* */
/*************************************************************/

done:
free p->message_Item;
close file(sysuexit);
end;

message_Filter:
proc ( ue, mf )
options( nodescriptor );
dcl 1 ue like uex_Uib byaddr;
dcl 1 mf like uex_Mfa byaddr;
dcl p pointer;
dcl bucket fixed bin(15);
on error snap system;
ue.uex_Uib_Reason_Code = uex_Reason_Output;
ue.uex_Uib_Return_Code = 0;

mf.uex_Mfa_New_Severity = mf.uex_Mfa_Severity;

/*************************************************************/
/* */
/* Calculate bucket for error message */
/* */
/*************************************************************/

bucket = mod(mf.uex_Mfa_Message_No, hashsize);

Sample JCL to compile, link, and invoke the user exit (continued)
do p = hashtable(bucket) repeat (p->link) while(p!=null())
  until (p->msgno = mf.uex_Mfa_Message_No &
           p->facid = mf.Uex_Mfa_Facility_Id);
end;

if p = null() then;
else
  do;
    /***********************************************************/
    /* */
    /* Filter error based on information in has table */
    /* */
    /***********************************************************/
    ue.uex_Uib_Reason_Code = p->reason;
    if p->newsev < 0 then;
      else
        mf.uex_Mfa_New_Severity = p->newsev;
      end;
    end;
  end;
end pack;

//*********************************************************************
//* link the user exit
//*********************************************************************
//LKEDEXIT EXEC PGM=IEWL,PARM='XREF,LIST,LET,DYNAM=DLL',
// COND=(9,LT,PLIEXIT),REGION=5000K
//SYSLIB DD DSN=CEE.SCEELKED,DISP=SHR
//SYSLMOD DD DSN=&&EXITLIB(IBMUEXIT),DISP=(NEW,PASS),UNIT=SYSDA,
//         SPACE=(TRK,(7,1,1)),DSNTYPE=LIBRARY
//SYSUT1 DD DSN=&&SYSUT1,UNIT=SYSDA,SPACE=(CYL,(3,1)),
//         DCB=BLKSIZE=1024
//SYSPRINT DD SYSOUT=X
//SYSLIN DD DDNAME=SYSIN
//LKD.SYSIN DD *
ENTRY IBMUEXIT

Sample JCL to compile, link, and invoke the user exit (continued)
Fetching PL/I MAIN routines

In an Enterprise PL/I application, you can also fetch a PL/I MAIN program. When a FETCH of a PL/I MAIN occurs, a child enclave is created. Here are a few of the rules you need to follow:

- You cannot pass any runtime options to a fetched MAIN program in the parameter string.
- If the SYSTEM(MVS) compiler option is specified, you can pass an arbitrary parameter list, but if the parameter is anything other than a single CHAR VARYING string, the compiler flags the MAIN routine with a warning message.
- You must not specify OPTIONS(ASM) or OPTIONS(NODESCRIPTOR) in the ENTRY declaration for the fetched MAIN routine in the fetching program.
- Attempts to pass runtime options should be avoided, because attempts to parse them might produce LE informational messages regarding invalid runtime options. If NOEXECOPS is specified in the fetched MAIN routine, the passed char varying string is not parsed for the runtime options.
- If no parameters are passed to the fetched MAIN program, the fetching program should:
  - either specify OPTIONS(LINKAGE(SYSTEM)) in its ENTRY declaration for the fetched MAIN routine.
  - or be compiled with DEFAULT(LINKAGE(SYSTEM)).

Examples

Here is the sample of the PL/I fetched MAIN program:

FMAIN: proc(parm) options(main,noexecops);
  DCL parm char(*) var;
  DCL SYSPRINT print;
  DCL PLIXOPT CHAR(11) VAR INIT('RPTOPTS(ON)')
Here is the sample of the PL/I MAIN program which fetches another PL/I MAIN:

```
MainFet: Proc Options(main);
  Dcl Parm char(1000) var;
  Dcl SYSPRINT print;
  Dcl Fmain entry(char(*) var) ;
  Put skip list("MainFet: start ");
  Parm = 'local-parm';
  Put skip list("MainFet parm: "|| Parm);
  Fetch Fmain;
  Call Fmain(Parm);
  Release Fmain;
  Put skip list("MainFet: testcase finished ");
End;
```

### Fetching z/OS C routines

Unless the NORENT option has been specified, the ENTRY declaration in the routine that fetches an z/OS C routine must not specify OPTIONS(COBOL) or OPTIONS(ASM)—these should be specified only for COBOL or ASM routines not linked as DLLs.

The z/OS C documentation provides instructions on how to compile and link an z/OS C DLL.

### Fetching assembler routines

Unless the NORENT option has been specified, the ENTRY declaration in the routine that fetches an assembler routine must specify OPTIONS(ASM).

### Invoking MAIN under TSO/E

If you compile your MAIN program with the SYSTEM(MVS) option, you can invoke the program using the TSO CALL command or as a TSO command processor. Both the run-time options and parameters can be passed the same way as under MVS batch.

For example, if the MAIN program TSOARG1 was link-edited as a member of a data set named `userid.TEST.load`, it can be invoked by:

- CALL TEST(TSOARG1) 'RPTSTG(ON),TRAP(ON)/THIS IS MY ARGUMENT'
- or
- CALL TEST(TSOARG1) '/THIS IS MY ARGUMENT'
- or
- TSOARG1 TRAP(ON)/THIS IS MY ARGUMENT
- or
- TSOARG1 /THIS IS MY ARGUMENT

**NOTE:** The data set containing TSOARG1 (`userid.TEST.load`) must be in the standard TSO program search list to run the program without using the CALL statement. This can be accomplished by issuing the TSO command

```
TSOLIB ACTIVATE DSN('userid.TEST.load')
```

However, if you compile your MAIN program with the SYSTEM(TSO) option, the program will be passed a pointer to the Command Processor Parameter List (CPPL). In this case, NOEXECOPS is in effect. Your program can be invoked as a TSO command, but it cannot be invoked by a TSO CALL statement. For example:
The program in Figure 19 uses the SYSTEM(TSO) interface to address and display the program arguments from the CPPL.

```
*process system(tso);
tsoarg2: proc (cppl_ptr) options(main);
dcl cppl_ptr pointer;
dcl 1 cppl based(cppl_ptr),
  2 cpplcbuf pointer,
  2 cpplupt pointer,
  2 cpplpscb pointer,
  2 cpplect pointer;
dcl 1 cpplbuf based(cpplcbuf),
  2 len fixed bin(15),
  2 offset fixed bin(15),
  2 argstr char(1000);
dcl my_argument char(1000) varying;
dcl my_argument_len fixed bin(31);
dcl length builtin;

  my_argument_len  =  len - offset - 4;
  if my_argument_len  =  0 then
    my_argument = ' ';
  else
    my_argument = substr(argstr,offset + 1, my_argument_len);
  display('Program args: ' || my_argument);
end tsoarg2;
```

Figure 19. Sample program to display program arguments from the CPPL under TSO when using SYSTEM(STD) option

Whether your MAIN program is invoked as a command or via CALL, you can always specify run-time options via the PLIXOPT string.

**Invoking MAIN under z/OS UNIX**

Under z/OS UNIX, if you compile a MAIN program with the SYSTEM(MVS) option, the program will be passed, as usual, one CHARACTER VARYING string containing the parameters specified when it was invoked.

However, under z/OS UNIX, if you compile a MAIN program with the SYSTEM(OS) option, the program will be passed 7 parameters as specified in the z/OS UNIX manuals. These 7 parameters include:

- the argument count (which includes the name of the executable as the first "argument"
- the address of an array of addresses of the arguments
- the address of an array of addresses of the arguments as null-terminated character strings
- the count of environment variables set
- the address of an array of addresses of the lengths of the environment variables
- the address of an array of addresses of the environment variables as null-terminated character strings

The program in Figure 20 on page 181 uses the SYSTEM(OS) interface to address and display the individual arguments and environment variables.
*process display(std) system(os);

sayargs:
proc(argc, pArgLen, pArgStr, envc, pEnvLen, pEnvStr, pParmSelf)
options( main, noexecops );

dcl argc fixed bin(31) nonasgn byaddr;
dcl pArgLen pointer nonasgn byvalue;
dcl pArgStr pointer nonasgn byvalue;
dcl envc fixed bin(31) nonasgn byaddr;
dcl pEnvLen pointer nonasgn byvalue;
dcl pEnvStr pointer nonasgn byvalue;
dcl pParmSelf pointer nonasgn byvalue;

dcl q(4095) pointer based;
dcl bxb fixed bin(31) based;
dcl bcz char(31) varz based;

display( 'argc = ' | | argc );
do jx = 1 to argc;
   display( 'pargStr(jx) = ' || pArgStr->q(jx)->bcz );
end;
display( 'envc = ' | | envc );
do jx = 1 to envc;
   display( 'pEnvStr(jx) = ' || pEnvStr->q(jx)->bcz );
end;
end;

Figure 20. Sample program to display z/OS UNIX args and environment variables
Part 2. Using I/O facilities

Chapter 6. Using data sets and files ........ 185
   Allocating files ................................ 185
   Associating data sets with files under z/OS ....... 186
      Associating several files with one data set .... 188
      Associating several data sets with one file .... 189
      Concatenating several data sets .............. 189
   Accessing HFS files under z/OS .................. 189
   Associating data sets with files under z/OS UNIX ... 191
      Using environment variables .................. 191
   Using the TITLE option of the OPEN statement .... 192
   Attempting to use files not associated with data sets .... 193
   How PL/I finds data sets ...................... 193
   Specifying characteristics using DD_DDNAME environment variables .................................. 194
      APPEND ..................................... 194
      BUFSIZE ................................... 195
      CHARSET for record I/O ...................... 195
      CHARSET for stream I/O ...................... 195
      DELAY ..................................... 195
      DELIMIT ................................... 196
      LRECL ..................................... 196
      LRMBSKIP .................................. 196
      PROMPT .................................... 196
      PUTPAGE ................................... 196
      RECCOUNT .................................. 197
      RECSIZE ................................... 197
      SAMELINE .................................. 197
      SKIP0 ..................................... 198
      TYPE ...................................... 198
   Establishing data set characteristics .......... 199
   Blocks and records ................................ 200
   Information interchange codes .................. 200
   Record formats ................................ 200
      Fixed-length records ......................... 200
      Variable-length records ..................... 201
      Undefined-length records .................... 202
   Data set organization ................................ 202
   Labels ....................................... 203
   Data Definition (DD) statement ............... 203
      Use of the conditional subparameters .......... 203
      Data set characteristics ...................... 204
   Using the TITLE option of the OPEN statement .... 204
   Associating PL/I files with data sets .......... 205
      Opening a file ................................ 205
      Closing a file ................................ 206
   Specifying characteristics in the
      ENVIRONMENT attribute ...................... 207
      The ENVIRONMENT attribute .................. 207
      Record formats for record-oriented data
         transmission ................................ 209
      Record formats for stream-oriented data
         transmission ................................ 209
      RECSIZE option ................................ 210
      BLKSIZE option ................................ 210
   Record format, BLKSIZE, and RECSIZE
      defaults .................................... 212
   GENKEY option — key classification ............. 212
   SCALARVARYING option — varying-length strings ................. 213
   KEYLENGTH option ................................ 214
   ORGANIZATION option ........................... 214
   Data set types used by PL/I record I/O .......... 214
   Setting environment variables under z/OS UNIX .................. 215
   PL/I standard files (SYSPRINT and SYSIN)
      under z/OS UNIX ............................. 215
   Redirecting standard input, output, and error
      devices under z/OS UNIX ...................... 216

Chapter 7. Using libraries ..................... 217
   Types of libraries ................................ 217
   How to use a library ............................ 217
   Creating a library ................................ 218
   SPACE parameter ................................ 218
   Creating and updating a library member .......... 219
   Examples ..................................... 219
   Extracting information from a library directory .... 221

Chapter 8. Defining and using consecutive data
   sets ......................................... 223
   Using stream-oriented data transmission .......... 223
   Defining files using stream I/O ................... 223
   Defining stream files using PL/I dynamic
      allocation ................................... 224
   Specifying ENVIRONMENT options .................. 224
      CONSECUTIVE ................................ 224
      Record format options ......................... 225
      RECSIZE ................................... 225
   Defaults for record format, BLKSIZE, and
      RECSIZE .................................... 225
   GRAPHIC option ................................ 226
   Creating a data set with stream I/O ................ 226
   Essential information ............................ 227
   Examples ..................................... 227
   Accessing a data set with stream I/O .............. 230
   Essential information ............................ 230
   Record format .................................. 230
   Example ...................................... 230
   Using PRINT files with stream I/O .................. 231
   Controlling printed line length .................... 232
   Overriding the tab control table .................. 234
   Using SYSPRINT and SYSPRN files .................. 235
   Controlling input from the terminal .............. 236
   Format of data .................................. 237
   Stream and record files ........................... 238
   Defining QSAM files using PL/I dynamic
      allocation ................................... 238
   Capital and lowercase letters ...................... 239
   End-of-file .................................... 239
   COPY option of GET statement .................... 239
**Chapter 9. Controlling output to the terminal** 241
- Format of PRINT files 241
- Stream and record files 241
- Output from the PUT EDIT command 241

**Chapter 10. Using record-oriented data transmission** 243
- Specifying record format 244
- Defining files using record I/O 244
- Specifying ENVIRONMENT options 244
  - CONSECUTIVE 245
  - ORGANIZATION(CONSECUTIVE) 245
  - CTLASA|CTL360 245
  - LEAVE|REREAD 246
- Creating a data set with record I/O 247
  - Essential information 248
- Accessing and updating a data set with record I/O 248
  - Essential information 249
  - Example of consecutive data sets 249

**Chapter 11. Defining and using regional data sets** 253
- Defining REGIONAL(1) data sets using PL/I dynamic allocation 255
- Defining files for a regional data set 255
  - Specifying ENVIRONMENT options 255
    - REGIONAL option 256
  - Using keys with REGIONAL data sets 256
- Using REGIONAL(1) data sets 256
- Dummy Records 257
- Creating a REGIONAL(1) data set 257
  - Example 257
- Accessing and updating a REGIONAL(1) data set 258
  - Sequential access 259
  - Direct access 259
  - Example 259
- Essential information for creating and accessing regional data sets 261

**Chapter 12. Defining and using VSAM data sets** 265
- Defining VSAM file using PL/I dynamic allocation 265
- Using VSAM data sets 265
  - How to run a program with VSAM data sets 265
  - Pairing an Alternate Index Path with a File 266
- VSAM organization 266
  - Keys for VSAM data sets 268
    - Keys for indexed VSAM data sets 268
    - Relative byte addresses (RBA) 269
    - Relative record numbers 269
  - Choosing a data set type 269
- Defining files for VSAM data sets 271
  - Specifying ENVIRONMENT options 271
    - BKWD option 272
    - BUFND option 272
    - BUFNI option 273
    - BUFSP option 273
    - GENKEY option 273
    - PASSWORD option 273
    - REUSE option 273
    - SKIP option 274

  - VSAM option 274
  - Performance options 274
  - Defining Files for Alternate Index Paths 274
  - Defining VSAM data sets 275
  - Entry-sequenced data sets 276
  - Loading an ESDS 276
  - Using a SEQUENTIAL file to access an ESDS 277
  - Defining and loading an ESDS 277
  - Updating an ESDS 278
  - Key-sequenced and indexed entry-sequenced data sets 279
  - Alternate Indexes for KSDSs or Indexed ESDSs 285
  - Relative-record data sets 292
  - Using Files Defined for non-VSAM Data Sets 299
  - Using Shared Data Sets 299

184 Enterprise PL/I for z/OS Programming Guide
Chapter 6. Using data sets and files

Your PL/I programs process and transmit units of information called *records*. A collection of records is called a *data set*. Data sets are physical collections of information external to PL/I programs; they can be created, accessed, or modified by programs written in PL/I or other languages or by the utility programs of the operating system.

Your PL/I program recognizes and processes information in a data set by using a symbolic or logical representation of the data set called a *file*. This chapter describes how to allocate files and associate data sets with the files known within your program. It introduces the five major types of data sets, how they are organized and accessed, and some of the file and data set characteristics you need to know how to specify.

Note: INDEXED implies VSAM and is supported only under batch.

Allocating files

For any type of file, including sequential, VSAM, REGIONAL(1), and HFS, you can define the external name by using the following methods:

- In the MVS or TSO environment:
  - A *ddname* in the JCL
  - An environment variable name
  - The TITLE option of the OPEN statement
- In the UNIX System Services environment:
  - An environment variable name
  - The TITLE option of the OPEN statement

For HFS files under the batch environment, the following rules apply:

- If FILEDATA is not specified, the default setting is TEXT.
- If the record format or the ENVIRONMENT option is not specified, the default is V with the LF type.
- If FILEDATA is BINARY, only fixed-length files are valid. Error message MSGIBM0210S is issued when FILEDATA=BINARY is specified with varying-length files.
- For varying-length files, files are assumed to be TYPE=LF; that is, the records are delimited with the LF (x'15') character.
- For data files that contain delimiters, specify FILEDATA=TEXT, so the PL/I library can handle the delimiters properly.

Dynamic allocation

A PL/I program dynamically allocates the file by using the attributes that are specified by the environment variable or TITLE option.

To use PL/I dynamic allocation, you must specify the file names by using the DSN() format for MVS data sets or the PATH() format for HFS files. All options and attributes must be in uppercase, except for the *pathname* suboption of the
PATH option, which is case-sensitive. Do not use temporary data set names in the
DSN() option. See the following examples to specify file names:

OPEN FILE(FILEIN) TITLE('DSN(USER.FILE.EXT),SHR');
OPEN FILE(FILEIN) TITLE('PATH(/usr/FILE.EXT)');

EXPORT DD_FILE="DSN(USER.FILE.EXT),SHR"
EXPORT DD_FILE="PATH(/usr/FILE.EXT)"

The following rules of precedence apply in determining when the dynamic
allocation takes place:

1. If one of the following DD statements exists for the file, it is used. This rule is not valid in the UNIX System Services environment.
   - JCL DD
   - TSO ALLOCATE
   - User-initiated dynamic allocation

2. If a DD statement does not exist for the file and the TITLE option is specified in the OPEN statement, the TITLE option is used by associating it with the external name for the file.

3. If a DD statement does not exist for the file and the TITLE option is not specified, but an environment variable exists for the file, the environment variable is used by associating it with the external name for the file.

For MVS data sets, the Enterprise PL/I run time checks the contents of the environment variable or TITLE option at each OPEN statement:

- If a file with the same external name was dynamically allocated by a previous OPEN statement, and the contents of the environment variable or TITLE option have changed since that OPEN statement, the run time dynamically deallocates the previous allocation and reallocates the file by using the options that are currently set in the environment variable or specified in the TITLE option.
- If the contents of the environment variable or TITLE option have not changed, the run time uses the current allocation without deallocating or reallocating.

For HFS files, the DD statement is deallocated and reallocated at each subsequent OPEN statement.

Note:

1. When you use the DSN() or the PATH() format specification for PL/I dynamic allocation, if a DD statement and the TITLE option of the OPEN statement are both specified for a file, the DD statement is used and the TITLE option is ignored.

2. Under the UNIX System Services environment, user-initiated dynamic allocation is not supported. If it is attempted, the UNDEFINEDFILE condition is raised, even when the TITLE or environment variable option is used, because the external name is in use.

---

**Associating data sets with files under z/OS**

A file used within a PL/I program has a **PL/I file name**. The physical data set external to the program has a name by which it is known to the operating system: a **data set name** or **dsname**. In some cases the data set has no name; it is known to the system by the device on which it exists.

The operating system needs a way to recognize which physical data set is referred to by your program, so you must write a **data definition** or **DD** statement, external
to your program, that associates the PL/I file name with a dsname. For example, if
you have the following file declaration in your program:

```pli
dcl stock file stream input;
```

you should create a DD statement with a *data definition name (ddname)* that matches
the name of the PL/I file. The DD statement specifies a physical data set name
dsname) and gives its characteristics:

```
//go.stock dd dsn=parts.instock, ...
```

You’ll find some guidance in writing DD statements in this manual, but for more
detail refer to the job control language (JCL) manuals for your system.

There is more than one way to associate a data set with a PL/I file. You associate a
data set with a PL/I file by ensuring that the ddname of the DD statement that
defines the data set is the same as one of the following:

- The declared PL/I file name
- The character-string value of the expression specified in the TITLE option of the
  associated OPEN statement.

You must choose your PL/I file names so that the corresponding ddnames
conform to the following restrictions:

- If a file is opened implicitly, or if no TITLE option is included in the OPEN
  statement that explicitly opens the file, the ddname defaults to the file name. If
  the file name is longer than 8 characters, the default ddname is composed of the
  first 8 characters of the file name.
- The character set of the JCL does not contain the break character (_).
  Consequently, this character cannot appear in ddnames. Do not use break
  characters among the first 8 characters of file names, unless the file is to be
  opened with a TITLE option with a valid ddname as its expression. The
  alphabetic extender characters $, @, and #, however, are valid for ddnames, but
  the first character must be one of the letters A through Z.

Since external names are limited to 7 characters, an external file name of more than
7 characters is shortened into a concatenation of the first 4 and the last 3 characters
of the file name. Such a shortened name is not, however, the name used as the
ddname in the associated DD statement.

Consider the following statements:

1. `OPEN FILE(MASTER);`
2. `OPEN FILE(OLDMASTER);`
3. `READ FILE(DETAIL) ...;`

When statement number 1 is run, the file name MASTER is taken to be the same
as the ddname of a DD statement in the current job step. When statement number
2 is run, the name OLDMASTER is taken to be the same as the ddname of a DD
statement in the current job step. (The first 8 characters of a file name form the
ddname. If OLDMASTER is an external name, it will be shortened by the compiler
to OLDMTER for use within the program.) If statement number 3 causes implicit
opening of the file DETAIL, the name DETAIL is taken to be the same as the
ddname of a DD statement in the current job step.

In each of the above cases, a corresponding DD statement must appear in the job
stream; otherwise, the UNDEFINEDFILE condition is raised. The three DD
statements could start as follows:
If the file reference in the statement which explicitly or implicitly opens the file is not a file constant, the DD statement name must be the same as the value of the file reference. The following example illustrates how a DD statement should be associated with the value of a file variable:

```
DCL PRICES FILE VARIABLE,
   RPRICE FILE;
   PRICES = RPRICE;
   OPEN FILE(PRICES);
```

The DD statement should associate the data set with the file constant RPRICE, which is the value of the file variable PRICES, thus:

```
//RPRICE DD DSNAME=...
```

Use of a file variable also allows you to manipulate a number of files at various times by a single statement. For example:

```
DECLARE F FILE VARIABLE,
   A FILE,
   B FILE,
   C FILE;

   DO F=A,B,C;
      READ FILE (F) ...;

   END;
```

The READ statement reads the three files A, B, and C, each of which can be associated with a different data set. The files A, B, and C remain open after the READ statement is executed in each instance.

The following OPEN statement illustrates use of the TITLE option:

```
OPEN FILE(DETAIL) TITLE('DETAIL1');
```

For this statement to be executed successfully, you must have a DD statement in the current job step with DETAIL1 as its ddname. It could start as follows:

```
//DETAIL1 DD DSNAME=DETAILA,...
```

Thus, you associate the data set DETAILA with the file DETAIL through the ddname DETAIL1.

### Associating several files with one data set

You can use the TITLE option to associate two or more PL/I files with the same external data set, provided the first file association is closed before opening a second file association against the same TITLE name. The following example, where INVNTTRY is the name of a DD statement defining a data set to be associated with two files:
OPEN FILE (FILE1) TITLE('INVNTRY');

.....

CLOSE FILE (FILE1);

.....

OPEN FILE (FILE2) TITLE('INVNTRY');

If the file is not closed first before the second open is done, the UNDEFINEDFILE condition will be raised with a subcode1 value of 59, stating that an open was attempted against a file that was already open.

**Associating several data sets with one file**

The file name can, at different times, represent entirely different data sets. In the above example of the OPEN statement, the file DETAIL1 is associated with the data set named in the DSNAME parameter of the DD statement DETAIL1. If you closed and reopened the file, you could specify a different ddname in the TITLE option to associate the file with a different data set.

Use of the TITLE option allows you to choose dynamically, at open time, one among several data sets to be associated with a particular file name. Consider the following example:

```
DO IDENT='A','B','C';
   OPEN FILE(MASTER)
       TITLE('MASTER1'||IDENT);
   ..
   ..
   CLOSE FILE(MASTER);
END;
```

In this example, when MASTER is opened during the first iteration of the do-group, the associated ddname is taken to be MASTER1A. After processing, the file is closed, dissociating the file name and the ddname. During the second iteration of the do-group, MASTER is opened again. This time, MASTER is associated with the ddname MASTER1B. Similarly, during the final iteration of the do-group, MASTER is associated with the ddname MASTER1C.

**Concatenating several data sets**

For input only, you can concatenate two or more sequential or regional data sets (that is, link them so that they are processed as one continuous data set) by omitting the ddname from all but the first of the DD statements that describe them. For example, the following DD statements cause the data sets LIST1, LIST2, and LIST3 to be treated as a single data set for the duration of the job step in which the statements appear:

```
//GO.LIST DD DSNAME=LIST1,DISP=OLD
//   DD DSNAME=LIST2,DISP=OLD
//   DD DSNAME=LIST3,DISP=OLD
```

When read from a PL/I program, the concatenated data sets need not be on the same volume. You cannot process concatenated data sets backward.

**Accessing HFS files under z/OS**

You can access HFS files from a batch program by specifying the HFS file name in the DD statement or in the TITLE option of the OPEN statement.

For example, to access the HFS file `/u/USER/sample.txt` by using the DD HFS, you would code the DD statement as follows:
To access the same file by using the TITLE option of the OPEN statement, you would code:

```plaintext
OPEN FILE(HFS) TITLE('//u/USER/sample.txt');
```

Note the two forward slashes in the TITLE option: the first indicates that what follows is a file name (rather than a DD name), and the second is the start of the fully qualified HFS file name (and fully qualified names have to be used when HFS files are referenced under batch, because there is no current directory that could be used to complete a file specification).

You can access HFS files from a batch program using PL/I dynamic allocation by specifying the HFS file name using one of the following methods:

- The DD statement
- The TITLE option of the OPEN statement
- The PUTENV built-in function
- The PLIXOPT string using the ENVAR option

The following example shows how to access HFS files by using these methods:

```plaintext
//HFS DD PATH='/u/USER/sample.txt',PATHOPTS=ORDONLY...
OPEN FILE(HFS) TITLE('PATH(/u/USER/sample.txt)');
xx = putenv('DD_HFS=/u/USER/sample.txt');
Dcl plixopt char(100) var ext static
init('ENVAR("DD_HFS=PATH(/u/USER/sample.txt)")');
```

Note: To use PL/I dynamic allocation, specify the file names by using the DSN() format for MVS data sets, or the PATH() format for HFS files.

PL/I decides how to treat HFS files under batch in the following order:

1. If ENV(F) is specified in the file declaration, the file is assumed to consist of fixed length records.
2. If ENV(V) is specified in the file declaration, the file is assumed to consist of lf-delimited records.
3. If FILEDATA=BINARY is specified on the file's DD statement, the file is assumed to consist of fixed length records.
4. Otherwise the file is assumed to consist of lf-delimited records.

To access a fixed length z/OS file from UNIX, the record size of the file must be specified either in the ENVIRONMENT attribute of the file or in the TITLE option on the OPEN statement. So,

- If the file declaration does not contain ENV( F RECSIZE(...)), you must specify the data set name and these attributes in the TITLE option. For example, for a file of fixed length 80-byte records, you could specify a TITLE option as follows:
  `/dataset.name,type(fixed),recesize(80)`
- If the ENVIRONMENT attribute specifies only one of F or RECSIZE, you must specify the data set name and the omitted attribute in the TITLE option.
- If the ENVIRONMENT attribute specifies both F and RECSIZE, you need to specify only the data set name in the TITLE option.
Associating data sets with files under z/OS UNIX

A file used within a PL/I program has a PL/I file name. A data set also has a name by which it is known to the operating system.

PL/I needs a way to recognize the data sets to which the PL/I files in your program refer, so you must provide an identification of the data set to be used, or allow PL/I to use a default identification.

You can identify the data set explicitly using either an environment variable or the TITLE option of the OPEN statement.

To use PL/I dynamic allocation, you must specify the file names by using the DSN() format for MVS data sets, or the PATH() format for HFS files.

To access HFS files from z/OS UNIX program using PL/I dynamic allocation, you can specify the HFS filename by using one of the following methods:

- The TITLE option of the OPEN statement
- The PUTENV built-in function
- The PLIXOPT string using the ENVAR option
- The EXPORT statement

Using environment variables

Use the export command to establish an environment variable that identifies the data set to be associated with a PL/I file, and, optionally, to specify the characteristics of that data set. The information provided by the environment variable is called data definition (or DD) information.

These environment variable names have the form DD_DDNAME where the DDNAME is the name of a PL/I file constant (or an alternate DDNAME, as defined below). If the filename refers to an HFS file, the filename has to be properly qualified. Otherwise, the PL/I library assumes the filename refers to an MVS data set.

Examples:

- declare MyFile stream output;
  export DD_MYFILE=/datapath/mydata.dat
  where /datapath/mydata.dat refers to an HFS file. The filename is fully-qualified.
- export DD_MYFILE=./mydata.dat
  where ./mydata.dat refers to an HFS file residing in the current directory.
- export DD_MYFILE=mydata.dat
  where mydata.dat refers to an MVS data set.

The following examples show that the HFS files are accessed when using PL/I dynamic allocation:

export DD_HFS="PATH(/u/USER/sample.txt)"
export DD_FILE="DSN(USER.FILE.EXT),SHR"

If you are familiar with the IBM mainframe environment, you can think of the environment variable much like:

DD statement in z/OS
ALLOCATE statement in TSO
For more information about the syntax and options you can use with the DD_DDNAME environment variable, see “Specifying characteristics using DD_DDNAME environment variables” on page 194.

Under z/OS UNIX, where more types of varying length HFS files are supported than under batch, PL/I treats an HFS file as follows:

1. If ENV(F) is specified in the file declaration, the file is assumed to consist of fixed length records.
2. If a TYPE is specified in the file's EXPORT statement, the file is assumed to consist of records of that type.
3. Otherwise the file is assumed to consist of LF-delimited records.

Using the TITLE option of the OPEN statement

You can use the TITLE option of the OPEN statement to identify the data set to be associated with a PL/I file, and, optionally, to provide additional characteristics of that data set.

```
TITLE('expression')
```

The expression must yield a character string with the following syntax:

```
alternate_ddname
/filespec,
,dd_option

/DSN(->dsn)
/PATH(->pathname)
```

**alternate_ddname**

The name of an alternate DD_DDNAME environment variable. An alternate DD_DDNAME environment variable is not named after a file constant. For example, if you have a file named INVENTORY in your program, and you establish two DD_DDNAME environment variables—the first named INVENTORY and the second named PARTS—you can associate the file with the second one using this statement:

```
open file(INVENTORY) title('PARTS');
```

**filespec**

Any valid file specification on the system you are using. The maximum length of filespec is 1023 characters.

**dd_option**

One or more options allowed in a DD_DDNAME environment variable

For more information about the options of the DD_DDNAME environment variable, see “Specifying characteristics using DD_DDNAME environment variables” on page 194.

**dsn**

The fully qualified MVS data set name

**dsn_option**

One or more DSN options
For more information about the DSN options, see “Defining QSAM files using PL/I dynamic allocation” on page 238, “Defining REGIONAL() data sets using PL/I dynamic allocation” on page 255, and “Defining VSAM file using PL/I dynamic allocation” on page 265.

pathname

The fully qualified HFS path name

Here is an example of using the OPEN statement in this manner with a z/OS DSN:

```pli
open file(Payroll) title('/June.Dat,append(n),recsize(52)');
```

Note the required leading forward slash in the TITLE option. This leading forward slash indicates that what follows is a file name (rather than a DD name). In this case, June.Dat refers to an MVS data set.

If June.Dat is an HFS file, the example looks like this:

```pli
open file(Payroll) title('//u/USER/June.Dat,append(n),recsize(52)');
```

Note the two forward slashes in the TITLE option: the first indicates that what follows is a file name (rather than a DD name), and the second is the start of the fully qualified HFS file name.

Relative HFS file names can also be specified in place of fully qualified names, for example:

```pli
open file(Payroll) title('./June.Dat,append(n),recsize(52)');
```

The dataset name June.Dat will be prefixed with the pathname of the current z/OS UNIX directory.

With this form, PL/I obtains all DD information either from the TITLE expression or from the ENVIRONMENT attribute of a file declaration - a DD_DDNAME environment variable is not referenced.

The following examples show how to use the OPEN statement with PL/I dynamic allocation support:

- For a z/OS data set:
  ```pli
  OPEN FILE(QSAM01) title('DSN(USER.FILE.EXT),SHR');
  ```

- For an HFS file:
  ```pli
  OPEN FILE(QSAM01) title('PATH(//u/USER/sample.txt)');
  ```

**Attempting to use files not associated with data sets**

If you attempt to use a file that has not been associated with a data set, (either through the use of the TITLE option of the OPEN statement or by establishing a DD_DDNAME environment variable), the UNDEFINEDFILE condition is raised. The only exceptions are the files SYSIN and SYSPRINT; these default to stdin and stdout, respectively.

**How PL/I finds data sets**

PL/I establishes the path for creating new data sets or accessing existing data sets in one of the following ways:

- The current directory.
- The paths as defined by the export DD_DDNAME environment variable.
 Specifying characteristics using DD_DDNAME environment variables

You use the export command to establish an environment variable that identifies the data set to be associated with a PL/I file, and, optionally, provide additional characteristics of that data set. This information provided by the environment variable is called data definition (or DD) information.

The syntax of the DD_DDNAME environment variable is:

```
DD_DDNAME="filespec", option

-DSN(-dsname-), -dsn_option

-PATH(-pathname-)
```

Blanks are acceptable within the syntax. In addition, the syntax of the statement is not checked at the time the command is entered. It is verified when the data set is opened. If the syntax is wrong, UNDEFINEDFILE is raised with the oncode 96.

**DD_DDNAME**
The name of the environment variable. The DDNAME must be in uppercase and can be either the name of a file constant or an alternate DDNAME that you specify in the TITLE option of your OPEN statement. The TITLE option is described in [“Using the TITLE option of the OPEN statement” on page 192](#).

If you use an alternate DDNAME, and it is longer than 31 characters, only the first 31 characters are used in forming the environment variable name.

**filespec**
The specifications of a file or the name of a device to be associated with the PL/I file

**option**
The options that you can specify as DD information

**dsname**
The fully-qualified MVS data set name

**pathname**
The fully-qualified HFS path name

The options that you can specify as DD information are described in the following topics, beginning with [“APPEND”](#) and ending with [“TYPE” on page 198](#). Note that these options do not apply to the DSN() or PATH() formats.

**APPEND**
The APPEND option specifies whether an existing data set is to be extended or re-created.

```
APPEND(Y-N)
```

Y  Specifies that new records are to be added to the end of a sequential data set, or inserted in a relative or indexed data set.
N  Specifies that, if the file exists, it is to be recreated.

The APPEND option applies only to OUTPUT files. APPEND is ignored if:
  • The file does not exist
  • The file does not have the OUTPUT attribute
  • The organization is REGIONAL(1)

**BUFSIZE**
The BUFSIZE option specifies the number of bytes for a buffer.

```
BUFSIZE=(n)
```

RECORD output is buffered by default and has a default value for BUFSIZE of 64k. STREAM output is buffered, but not by default, and has a default value for BUFSIZE of zero.

If the value of zero is given to BUFSIZE, the number of bytes for buffering is equal to the value specified in the RECSIZE or LRECL option.

The BUFSIZE option is valid only for a consecutive binary file. If the file is used for terminal input, you should assign the value of zero to BUFSIZE for increased efficiency.

**CHARSET for record I/O**
This version of the CHARSET option applies only to consecutive files using record I/O. It gives the user the capability of using ASCII data files as input files, and specifying the character set of output files.

```
CHARSET=(ASIS, EBCDIC, ASCII)
```

Choose a suboption of CHARSET based on what form the file has (input) or what form you want the file have (output).

**CHARSET for stream I/O**
This version of the CHARSET option applies for stream input and output files. It gives the user the capability of using ASCII data files as input files, and specifying the character set of output files. If you attempt to specify ASIS when using stream I/O, no error is issued and character sets are treated as EBCDIC.

```
CHARSET=(EBCDIC, ASCII)
```

Choose a suboption of CHARSET based on what form the file has (input) or what form you want the file to have (output).

**DELAY**
The DELAY option specifies the number of milliseconds to delay before retrying an operation that fails when a file or record lock cannot be obtained by the system.
This option is applicable only to VSAM files.

**DELDIT**
The DELIMIT option specifies whether the input file contains field delimiters or not. A field delimiter is a blank or a user-defined character that separates the fields in a record. This is applicable for sort input files only.

The sort utility distinguishes text files from binary files with the presence of field delimiters. Input files that contain field delimiters are processed as text files; otherwise, they are considered to be binary files. The library needs this information in order to pass the correct parameters to the sort utility.

**LRECL**
The LRECL option is the same as the RECSIZE option.

If LRECL is not specified and not implied by a LINESIZE value (except for TYPE(FIXED) files, the default is 1024.

**LRMSKIP**
The LRMSKIP option allows output to commence on the nth (n refers to the value specified with the SKIP option of the PUT or GET statement) line of the first page for the first SKIP format item to be executed after a file is opened.

If n is zero or 1, output commences on the first line of the first page.

**PROMPT**
The PROMPT option specifies whether or not colons should be visible as prompts for stream input from the terminal.

**PUTPAGE**
The PUTPAGE option specifies whether or not the form feed character should be followed by a carriage return character. This option applies only to printer-destined files. Printer-destined files are stream output files declared with the PRINT attribute, or record output files declared with the CTLASA environment option.
NOCR
Indicates that the form feed character ('0C'x) is not followed by a carriage return character ('0D'x).

CR
Indicates that the carriage return character is appended to the form feed character. This option should be specified if output is sent to non-IBM printers.

RECCOUNT
The RECCOUNT option specifies the maximum number of records that can be loaded into a relative or regional data set that is created during the PL/I file opening process.

RECSIZE
The RECSIZE option specifies the length, n, of records in the data set.

SAMELINE
The SAMELINE option specifies whether the system prompt occurs on the same line as the statement that prompts for input.

Note: Under z/OS, it is recommended to omit the TITLE option with both the /filespec parameter and RECCOUNT parameter for improved functionality and performance of REGIONAL(1) data sets. In such a case, the number of records that will be loaded into the file depends on the space allocated to the first extent of the data set. See Chapter 11, “Defining and using regional data sets,” on page 253 for additional information.
Example 2
Given the statement PUT SKIP LIST('ENTER');, output is as follows:

```
prompt(y) sameline(y)  ENTER: (cursor)
prompt(y) sameline(n)  ENTER: (cursor)
prompt(n) sameline(n)  ENTER: (cursor)
prompt(n) sameline(n)  ENTER: (cursor)
```

**SKIP0**
The SKIP0 option specifies where the line cursor moves when SKIP(0) statement is coded in the source program. SKIP0 applies to terminal files that are not linked as PM applications.

```
Example 2
Given the statement PUT SKIP LIST('ENTER');, output is as follows:
```

```
```

**SKIP0(N)**
Specifies that the cursor is to be moved to the beginning of the next line.

**SKIP0(Y)**
Specifies that the cursor to be moved to the beginning of the current line.

The following example shows how you could make the output to the terminal skip zero lines so that the cursor moves to the beginning of the current output line:

```
export DD_SYSPRINT='stdout:,SKIP0(Y)'
```

**TYPE**
The TYPE option specifies the format of records in a native file.

```
Example 2
Given the statement PUT SKIP LIST('ENTER');, output is as follows:
```

```
```

**CRLF**
Specifies that records are delimited by the CR - LF character combination. ('CR' and 'LF' represent the ASCII values of carriage return and line feed, '0D'x and '0A'x, respectively. For an output file, PL/I places the characters at the end of each record; for an input file, PL/I discards the characters. For both input and output, the characters are not counted in consideration for RECSIZE.

The data set must not contain any record that is longer than the value determined for the record length of the data set.
**LF** Specifies that records are delimited by the LF character combination. (LF' represents the ASCII values of feed or '0A'.) For an output file, PL/I places the characters at the end of each record; for an input file, PL/I discards the characters. For both input and output, the characters are not counted in consideration for RECSIZE.

The data set must not contain any record that is longer than the value determined for the record length of the data set.

**TEXT**
Equivalent to LF.

**FIXED**
Specifies that each record in the data set has the same length. The length determined for records in the data set is used to recognize record boundaries.

All characters in a TYPE(FIXED) file are considered as data, including control characters if they exist. Make sure the record length you specify reflects the presence of these characters or make sure the record length you specify accounts for all characters in the record.

**CRLFEOF**
 Except for output files, this suboption specifies the same information as CRLF. When one of these files is closed for output, an end-of-file marker is appended to the last record.

**U** Indicates that records are unformatted. These unformatted files cannot be used by any record or stream I/O statements except OPEN and CLOSE. You can read from a TYPE(U) file only by using the FILEREAD built-in function. You can write to a TYPE(U) file only by using the FILEWRITE built-in function.

The TYPE option applies only to CONSECUTIVE files, except that it is ignored for printer-destined files with ASA(N) applied.

If your program attempts to access an existing data set with TYPE(FIXED) in effect and the length of the data set is not a multiple of the logical record length you specify, PL/I raises the UNDEFINEDFILE condition.

When using nonprint files with the TYPE(FIXED) attribute, SKIP is replaced by trailing blanks to the end of the line. If TYPE(LF) is being used, SKIP is replaced by LF with no trailing blanks.

---

**Establishing data set characteristics**

A data set consists of records stored in a particular format which the operating system data management routines understand. When you declare or open a file in your program, you are describing to PL/I and to the operating system the characteristics of the records that file will contain. You can also use JCL or an expression in the TITLE option of the OPEN statement to describe to the operating system the characteristics of the data in data sets or in the PL/I files associated with them.

You do not always need to describe your data both within the program and outside it; often one description will serve for both data sets and their associated PL/I files. There are, in fact, advantages to describing your data's characteristics in only one place. These are described later in this chapter and in following chapters.

To effectively describe your program data and the data sets you will be using, you need to understand something of how the operating system moves and stores data.
Blocks and records

The items of data in a data set are arranged in blocks separated by interblock gaps (IBG). (Some manuals refer to these as interrecord gaps.)

A block is the unit of data transmitted to and from a data set. Each block contains one record, part of a record, or several records. You can specify the block size in the BLKSIZE parameter of the DD statement or in the BLKSIZE option of the ENVIRONMENT attribute.

A record is the unit of data transmitted to and from a program. You can specify the record length in the LRECL parameter of the DD statement, in the TITLE option of the OPEN statement, or in the RECSIZE option of the ENVIRONMENT attribute.

When writing a PL/I program, you need consider only the records that you are reading or writing; but when you describe the data sets that your program will create or access, you must be aware of the relationship between blocks and records.

Blocking conserves storage space in a magnetic storage volume because it reduces the number of interblock gaps, and it can increase efficiency by reducing the number of input/output operations required to process a data set. Records are blocked and deblocked by the data management routines.

Information interchange codes

The normal code in which data is recorded is the Extended Binary Coded Decimal Interchange Code (EBCDIC).

Each character in the ASCII code is represented by a 7-bit pattern and there are 128 such patterns. The ASCII set includes a substitute character (the SUB control character) that is used to represent EBCDIC characters having no valid ASCII code. The ASCII substitute character is translated to the EBCDIC SUB character, which has the bit pattern 00111111.

Record formats

The records in a data set have one of the following formats:

- Fixed-length
- Variable-length
- Undefined-length.

Records can be blocked if required. The operating system will deblock fixed-length and variable-length records, but you must provide code in your program to deblock undefined-length records.

You specify the record format in the RECFM parameter of the DD statement, in the TITLE option of the OPEN statement, or as an option of the ENVIRONMENT attribute.

Fixed-length records

You can specify the following formats for fixed-length records:

- F Fixed-length, unblocked
- FB Fixed-length, blocked
- FS Fixed-length, unblocked, standard
- FBS Fixed-length, blocked, standard.

In a data set with fixed-length records, as shown in Figure 21 on page 201, all records have the same length. If the records are blocked, each block usually
contains an equal number of fixed-length records (although a block can be truncated). If the records are unblocked, each record constitutes a block.

---

**Unblocked records (F format):**

```
Record | IBG | Record | IBG | Record
```

**Blocked records (FB format):**

```
Block
Record | Record | Record | IBG | Record | Record | Record | Record | ...
```

*Figure 21. Fixed-length records*

Because it bases blocking and deblocking on a constant record length, the operating system processes fixed-length records faster than variable-length records.

**Variable-length records**

You can specify the following formats for variable-length records:

- **V**  Variable-length, unblocked
- **VB** Variable-length, blocked
- **VS** Variable-length, unblocked, spanned
- **VBS** Variable-length, blocked, spanned

V-format allows both variable-length records and variable-length blocks. A 4-byte prefix of each record and the first 4 bytes of each block contain control information for use by the operating system (including the length in bytes of the record or block). Because of these control fields, variable-length records cannot be read backward.

V-format signifies unblocked variable-length records. Each record is treated as a block containing only one record. The first 4 bytes of the block contain block control information, and the next 4 contain record control information.

VB-format signifies blocked variable-length records. Each block contains as many complete records as it can accommodate. The first 4 bytes of the block contain block control information, and a 4-byte prefix of each record contains record control information.

**Spanned Records:** A spanned record is a variable-length record in which the length of the record can exceed the size of a block. If this occurs, the record is divided into segments and accommodated in two or more consecutive blocks by specifying the record format as either VS or VBS. Segmentation and assembly are handled by the operating system. The use of spanned records allows you to select a block size, independently of record length, that will combine optimum use of auxiliary storage with maximum efficiency of transmission.

VS-format is similar to V-format. Each block contains only one record or segment of a record. The first 4 bytes of the block contain block control information, and the
next 4 contain record or segment control information (including an indication of whether the record is complete or is a first, intermediate, or last segment).

VBS-format differs from VS-format in that each block contains as many complete records or segments as it can accommodate; each block is, therefore, approximately the same size (although there can be a variation of up to 4 bytes, since each segment must contain at least 1 byte of data).

**Undefined-length records**

U-format allows the processing of records that do not conform to F- and V-formats. The operating system and the compiler treat each block as a record; your program must perform any required blocking or deblocking.

**Data set organization**

The data management routines of the operating system can handle a number of types of data sets, which differ in the way data is stored within them and in the allowed means of access to the data. The three main types of non-VSAM data sets and the corresponding keywords describing their PL/I organization are as follows:

<table>
<thead>
<tr>
<th>Type of data set</th>
<th>PL/I organization</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sequential</td>
<td>CONSECUITVE or ORGANIZATION(consecutive)</td>
</tr>
<tr>
<td>Indexed</td>
<td>INDEXED or ORGANIZATION(indexed)</td>
</tr>
<tr>
<td>Direct</td>
<td>REGIONAL or ORGANIZATION(relative)</td>
</tr>
</tbody>
</table>

A fourth type, *partitioned*, has no corresponding PL/I organization.

PL/I also provides support for three types of VSAM data organization: ESDS, KSDS, and RRDS. For more information about VSAM data sets, see Chapter 12, “Defining and using VSAM data sets,” on page 265.

In a sequential (or CONSECUITVE) data set, records are placed in physical sequence. Given one record, the location of the next record is determined by its physical position in the data set. Sequential organization can be selected for direct access devices.

An indexed sequential (or INDEXED) data set must reside on a direct access volume. An index or set of indexes maintained by the operating system gives the location of certain principal records. This allows direct retrieval, replacement, addition, and deletion of records, as well as sequential processing.

A direct (or REGIONAL) data set must reside on a direct access volume. The data set is divided into regions, each of which contains one or more records. A key that specifies the region number allows direct access to any record; sequential processing is also possible.

In a partitioned data set, independent groups of sequentially organized data, each called a member, reside in a direct access data set. The data set includes a directory that lists the location of each member. Partitioned data sets are often called libraries. The compiler includes no special facilities for creating and accessing partitioned data sets. Each member can be processed as a CONSECUITVE data set by a PL/I program. The use of partitioned data sets as libraries is described under Chapter 7, “Using libraries,” on page 217.

---
1. Do not confuse the terms “sequential” and “direct” with the PL/I file attributes SEQUENTIAL and DIRECT. The attributes refer to how the file is to be processed, and not to the way the corresponding data set is organized.
Labels

The operating system uses internal labels to identify direct access volumes and to store data set attributes (for example, record length and block size). The attribute information must originally come from a DD statement or from your program.

IBM standard labels have two parts: the initial volume label and header labels. The initial volume label identifies a volume and its owner; the header labels precede and follow each data set on the volume. Header labels contain system information, device-dependent information (for example, recording technique), and dataset characteristics.

direct access volumes have IBM standard labels. Each volume is identified by a volume label, which is stored on the volume. This label contains a volume serial number and the address of a volume table of contents (VTOC). The table of contents, in turn, contains a label, termed a data set control block (DSCB), for each data set stored on the volume.

Data Definition (DD) statement

A data definition (DD) statement is a job control statement that defines a data set to the operating system, and is a request to the operating system for the allocation of input/output resources. If the data sets are not dynamically allocated, each job step must include a DD statement for each data set that is processed by the step.

Your z/OS JCL User's Guide describes the syntax of job control statements. The operand field of the DD statement can contain keyword parameters that describe the location of the data set (for example, volume serial number and identification of the unit on which the volume will be mounted) and the attributes of the data itself (for example, record format).

The DD statement enables you to write PL/I source programs that are independent of the data sets and input/output devices they will use. You can modify the parameters of a data set or process different data sets without recompiling your program.

The following paragraphs describe the relationship of some operands of the DD statement to your PL/I program.

The LEAVE and REREAD options of the ENVIRONMENT attribute allow you to use the DISP parameter to control the action taken when the end of a magnetic-tape volume is reached or when a magnetic-tape data set is closed. The LEAVE and REREAD options are described under "LEAVE|REREAD" on page 246.

Write validity checking, which was standard in PL/I Version 1, is no longer performed. Write validity checking can be requested through the OPTCD subparameter of the DCB parameter of the JCL DD statement. See OS/VS2 Job Control Language manual.

Use of the conditional subparameters

If you use the conditional subparameters of the DISP parameter for data sets processed by PL/I programs, the step abend facility must be used. The step abend facility is obtained as follows:

1. The ERROR condition should be raised or signaled whenever the program is to terminate execution after a failure that requires the application of the conditional subparameters.
2. The PL/I user exit must be changed to request an ABEND.

Data set characteristics
The DCB (data control block) parameter of the DD statement allows you to describe the characteristics of the data in a data set, and the way it will be processed, at run time. Whereas the other parameters of the DD statement deal chiefly with the identity, location, and disposal of the data set, the DCB parameter specifies information required for the processing of the records themselves. The subparameters of the DCB parameter are described in your z/OS JCL User’s Guide.

The DCB parameter contains subparameters that describe:
- The organization of the data set and how it will be accessed (CYLOFL, DSORG, LIMCT, NTM, and OPTCD subparameters)
- Device-dependent information such as the line spacing for a printer (CODE, FUNC, MODE, OPTCD=J, PRTSP, and STACK subparameters)
- The record format (BLKSIZE, KEYLEN, LRECL, and RECFM subparameters)
- The ASA control characters (if any) that will be inserted in the first byte of each record (RECFM subparameter).

You can specify BLKSIZE, LRECL, KEYLEN, and RECFM (or their equivalents) in the ENVIRONMENT attribute of a file declaration in your PL/I program instead of in the DCB parameter.

You cannot use the DCB parameter to override information already established for the data set in your PL/I program (by the file attributes declared and the other attributes that are implied by them). DCB subparameters that attempt to change information already supplied are ignored.

For a new dataset, the attributes of the file defined in the program will be used if there is a conflict with the DD statement.

You may see message IEC225I with RC=4 issued when closing PDS files. This message can be safely ignored.

An example of the DCB parameter is:
```
DCB=(RECFM=FB,BLKSIZE=400,LRECL=40)
```

which specifies that fixed-length records, 40 bytes in length, are to be grouped together in a block 400 bytes long.

Using the TITLE option of the OPEN statement
You can use the TITLE option of the OPEN statement to identify the data set to be associated with a PL/I file and, optionally, to provide additional characteristics of the data set.

For more information about using the TITLE option of the OPEN statement, see "Using the TITLE option of the OPEN statement" on page 192.
Associating PL/I files with data sets

Opening a file
The execution of a PL/I OPEN statement associates a file with a data set. This requires merging of the information describing the file and the data set. If any conflict is detected between file attributes and data set characteristics, the UNDEFINEDFILE condition is raised.

Subroutines of the PL/I library create a skeleton data control block for the data set. They use the file attributes from the DECLARE and OPEN statements and any attributes implied by the declared attributes, to complete the data control block as far as possible. (See Figure 22 on page 206.) They then issue an OPEN macro instruction, which calls the data management routines to check that the correct volume is mounted and to complete the data control block.

The data management routines examine the data control block to see what information is still needed and then look for this information, first in the DD statement, and finally, if the data set exists and has standard labels, in the data set labels. For new data sets, the data management routines begin to create the labels (if they are required) and to fill them with information from the data control block.

For INPUT data sets, the PL/I program can override the DCB attributes as long as there is no conflict in attributes. For OUTPUT data sets, the PL/I program cannot override the DCB attributes. However, if any DCB attributes are missing from the data set when it is opened, they will be obtained from the PL/I program, if provided.

When the DCB fields are filled in from these sources, control returns to the PL/I library subroutines. If any fields still are not filled in, the PL/I OPEN subroutine provides default information for some of them. For example, if LRECL is not specified, it is provided from the value given for BLKSIZE.
Using system-determined block size: If you use the system-determined block size function of the Data Facility Product on z/OS, DFP determines the optimal block size for the device type that is assigned. When you create a new DASD data set, the system derives the optimum block size and saves it in the data set label when all of the following conditions are true:

- Block size is not available or specified from any source. BLKSIZE=0 can be specified.
- You specify LRECL or it is in the data class. The data set does not have to be SMS managed.
- You specify RECFM or it is in the data class. It must be fixed or variable.
- You specify DSORG as PS or PO or you omit DSORG and it is PS or PO in the data class.

When system-determined block size is active, DFP determines the block size and places it in the data set label before PL/I opens the file. Therefore, if the PL/I program specifies a block size with its ENVIRONMENT option, it must not be in conflict with the value from system-determined block size.

For detailed information about system-determined block size, see related topics in z/OS DFSMS Using Data Sets.

Closing a file

The execution of a PL/I CLOSE statement dissociates a file from the data set with which it was associated. The PL/I library subroutines first issue a CLOSE macro instruction and, when control returns from the data management routines, release the data control block that was created when the file was opened. The data management routines complete the writing of labels for new data sets and update the labels of existing data sets.
Specifying characteristics in the ENVIRONMENT attribute

You can use various options in the ENVIRONMENT attribute. Each type of file has different attributes and environment options, which are listed below.

The ENVIRONMENT attribute

You use the ENVIRONMENT attribute of a PL/I file declaration file to specify information about the physical organization of the data set associated with a file, and other related information. The format of this information must be a parenthesized option list.

```
ENVIRONMENT (option-list)
```

Abbreviation: ENV

You can specify the options in any order, separated by blanks or commas.

The following example illustrates the syntax of the ENVIRONMENT attribute in the context of a complete file declaration (the options specified are for VSAM and are discussed in Chapter 12, “Defining and using VSAM data sets,” on page 265).

```
DCL FILENAME FILE RECORD SEQUENTIAL
  INPUT ENV(VSAM GENKEY);
```

Table 14 summarizes the ENVIRONMENT options and file attributes. Certain qualifications on their use are presented in the notes and comments for the figure. Those options that apply to more than one data set organization are described in the remainder of this chapter. In addition, in the following chapters, each option is described with each data set organization to which it applies.

<table>
<thead>
<tr>
<th>Data set type</th>
<th>Stream</th>
<th>Record</th>
<th>Legend</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td>C</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>D</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>I</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>N</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>O</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>S</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>V</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>VSAM</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>R</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>I</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>C</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>D</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>I</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>N</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>V</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>VSAM</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>R</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>I</td>
</tr>
</tbody>
</table>

**Table 14. Attributes of PL/I file declarations.**

<table>
<thead>
<tr>
<th>File Type</th>
<th>Consecutive</th>
<th>Sequential</th>
<th>Direct</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

**Legend:**

- C: Checked for VSAM
- D: Default
- I: Must be specified or implied
- N: Ignored for VSAM
- O: Optional
- S: Must be specified
- V: Invalid
- VSAM: Relevant for VSAM
- R: Relevant for record
- I: Implied
- C: Checked
- D: Default

**File attributes:**

<table>
<thead>
<tr>
<th>File attribute</th>
<th>Attributes implied</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Chapter 6. Using data sets and files 207
### Data set organization options

The options that specify data set organization are:

- **CONSECUTIVE**
- **INDEXED**
- **REGIONAL**
- **VSAM**

Each option is described in the discussion of the data set organization to which it applies.

---

**Notes:**
1. A file with the INPUT attribute cannot have the PRINT attribute.
2. Keyed is required for INDEXED and REGIONAL output.
Other ENVIRONMENT options: You can use a constant or variable with those ENVIRONMENT options that require integer arguments, such as block sizes and record lengths. The variable must not be subscripted or qualified, and must have attributes FIXED BINARY(31,0) and STATIC.

The list of equivalents for ENVIRONMENT options and DCB parameters are:

<table>
<thead>
<tr>
<th>ENVIRONMENT option</th>
<th>DCB subparameter</th>
</tr>
</thead>
<tbody>
<tr>
<td>Record format</td>
<td>RECFM1</td>
</tr>
<tr>
<td>RECSIZE</td>
<td>LRECL</td>
</tr>
<tr>
<td>BLKSIZE</td>
<td>BLKSIZE</td>
</tr>
<tr>
<td>CTLASA</td>
<td>CTL360</td>
</tr>
<tr>
<td>KEYLENGTH</td>
<td>KEYLEN</td>
</tr>
</tbody>
</table>

**Note:** VS must be specified as an ENVIRONMENT option, not in the DCB.

Record formats for record-oriented data transmission

Record formats supported depend on the data set organization.

Records can have one of the following formats:

<table>
<thead>
<tr>
<th>Format</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>F</td>
<td>Fixed-length unblocked</td>
</tr>
<tr>
<td>FB</td>
<td>Fixed-length blocked</td>
</tr>
<tr>
<td>FS</td>
<td>Fixed-length unblocked, standard</td>
</tr>
<tr>
<td>FBS</td>
<td>Fixed-length blocked, standard</td>
</tr>
<tr>
<td>V</td>
<td>Variable-length unblocked</td>
</tr>
<tr>
<td>VB</td>
<td>Variable-length blocked</td>
</tr>
<tr>
<td>VS</td>
<td>Variable-length spanned</td>
</tr>
<tr>
<td>VBS</td>
<td>Variable-length blocked, spanned</td>
</tr>
<tr>
<td>U</td>
<td>Undefined-length (cannot be blocked)</td>
</tr>
</tbody>
</table>

When U-format records are read into a varying-length string, PL/I sets the length of the string to the block length of the retrieved data.

These record format options do not apply to VSAM data sets. If you specify a record format option for a file associated with a VSAM data set, the option is ignored.

You can only specify VS-format records for data sets with consecutive organization.

Record formats for stream-oriented data transmission

The record format options for stream-oriented data transmission are discussed in "Using stream-oriented data transmission” on page 223."
**RECSIZE option**

The **RECSIZE** option specifies the record length.

```
  RECSIZE(record-length)
```

For files associated with VSAM data sets, **record-length** is the sum of:
1. The length required for data. For variable-length and undefined-length records, this is the maximum length.
2. Any control bytes required. Variable-length records require 4 (for the record-length prefix); fixed-length and undefined-length records do not require any.

For VSAM data sets, the maximum and average lengths of the records are specified to the Access Method Services utility when the data set is defined. If you include the **RECSIZE** option in the file declaration for checking purposes, you should specify the maximum record size. If you specify **RECSIZE** and it conflicts with the values defined for the data set, the UNDEFINEDFILE condition is raised.

You can specify **record-length** as an integer or as a variable with attributes FIXED BINARY(31,0) STATIC.

The value is subject to the following conventions:

**Maximum:**
- Fixed-length, and undefined (except ASCII data sets): 32760
- V-format, and VS- and VBS-format with UPDATE files: 32756
- VS- and VBS-format with INPUT and OUTPUT files: 16777215
- ASCII data sets: 9999
- VSAM data sets: 32761

*Note:* For VS- and VBS-format records longer than 32,756 bytes, you must specify the length in the **RECSIZE** option of ENVIRONMENT, and for the DCB subparameter of the DD statement you must specify LRECL=X. If **RECSIZE** exceeds the allowed maximum for INPUT or OUTPUT, either a record condition occurs or the record is truncated. UPDATE files are not supported with LRECL=X.

**Zero value:**
- A search for a valid value is made first:
  - In the DD statement for the data set associated with the file, and second
  - In the data set label.

  If neither of these provides a value, default action is taken (see "Record format, BLKSIZE, and RECSIZE defaults" on page 212).

**Negative Value:**
The UNDEFINEDFILE condition is raised.

**BLKSIZE option**

The **BLKSIZE** option specifies the maximum block size on the data set.

```
  BLKSIZE(block-size)
```

**block-size** is the sum of:
1. The total length(s) of one of the following:
   - A single record
   - A single record and either one or two record segments
   - Several records
   - Several records and either one or two record segments
   - Two record segments
   - A single record segment.

   For variable-length records, the length of each record or record segment includes the 4 control bytes for the record or segment length.

   The preceding list summarizes all the possible combinations of records and record segments options: fixed- or variable-length blocked or unblocked, spanned or unspanned. When specifying a blocksize for spanned records, note that each record and each record segment requires 4 control bytes for the record length and that these quantities are in addition to the 4 control bytes required for each block.

2. Any further control bytes required.
   - Variable-length blocked records require 4 (for the block size).
   - Fixed-length and undefined-length records do not require any further control bytes.

3. Any block prefix bytes required (ASCII data sets only).

*block-size* can be specified as an integer, or as a variable with attributes FIXED BINARY(31,0) STATIC.

The value is subject to the following conventions:

*Maximum:* 32760

*Zero value:* If you set BLKSIZE to 0, under z/OS the Data Facility Product sets the block size. For an elaboration of this topic, see “Record format, BLKSIZE, and RECSIZE defaults” on page 212.

*Negative value:* The UNDEFINEDFILE condition is raised.

The relationship of block size to record length depends on the record format:

*FB-format or FBS-format*:
   - The block size must be a multiple of the record length.

*VB-format*:
   - The block size must be equal to or greater than the sum of:
     1. The maximum length of any record
     2. Four control bytes.

*VS-format or VBS-format*:
   - The block size can be less than, equal to, or greater than the record length.

*Notes:*
   - Use the BLKSIZE option with unblocked (F- or V-format) records in either of the following ways:
     - Specify the BLKSIZE option, but not the RECSIZE option. Set the record length equal to the block size (minus any control or prefix bytes), and leave the record format unchanged.
Specify both BLKSIZE and RECSIZE and ensure that the relationship of the two values is compatible with blocking for the record format you use. Set the record format to FB or VB, whichever is appropriate.

- If for FB-format or FBS-format records the block size equals the record length, the record format is set to F.
- The BLKSIZE option does not apply to VSAM data sets, and is ignored if you specify it for one.

**Record format, BLKSIZE, and RECSIZE defaults**

If you do not specify either the record format, block size, or record length for a non-VSAM data set, the following default action is taken:

**Record format:**

A search is made in the associated DD statement or data set label. If the search does not provide a value, the UNDEFINEDFILE condition is raised, except for files associated with dummy data sets or the foreground terminal, in which case the record format is set to U.

**Block size or record length:**

If one of these is specified, a search is made for the other in the associated DD statement or data set label. If the search provides a value, and if this value is incompatible with the value in the specified option, the UNDEFINEDFILE condition is raised. If the search is unsuccessful, a value is derived from the specified option (with the addition or subtraction of any control or prefix bytes).

If neither is specified, the UNDEFINEDFILE condition is raised, except for files associated with dummy data sets, in which case BLKSIZE is set to 121 for F-format or U-format records and to 129 for V-format records. For files associated with the foreground terminal, RECSIZE is set to 120.

If you are using z/OS with the Data Facility Product system-determined block size, DFP determines the optimum block size for the device type assigned. If you specify BLKSIZE(0) in either the DD assignment or the ENVIRONMENT statement, DFP calculates BLKSIZE using the record length, record format, and device type.

**GENKEY option — key classification**

The GENKEY (generic key) option applies only to INDEXED and VSAM key-sequenced data sets. It enables you to classify keys recorded in a data set and to use a SEQUENTIAL KEYED INPUT or SEQUENTIAL KEYED UPDATE file to access records according to their key classes.

A generic key is a character string that identifies a class of keys; all keys that begin with the string are members of that class. For example, the recorded keys “ABCD”, “ABCE”, and “ABDF” are all members of the classes identified by the generic keys “A” and “AB”, and the first two are also members of the class “ABC”; and the three recorded keys can be considered to be unique members of the classes “ABCD”, “ABCE”, and “ABDF”, respectively.

The GENKEY option allows you to start sequential reading or updating of a VSAM data set from the first record that has a key in a particular class, and for an INDEXED data set from the first nondummy record that has a key in a particular class. You identify the class by including its generic key in the KEY option of a
READ statement. Subsequent records can be read by READ statements without the KEY option. No indication is given when the end of a key class is reached.

Although you can retrieve the first record having a key in a particular class by using a READ with the KEY option, you cannot obtain the actual key unless the records have embedded keys, since the KEYTO option cannot be used in the same statement as the KEY option.

In the following example, a key length of more than 3 bytes is assumed:

```
DCL IND FILE RECORD SEQUENTIAL KEYED
    UPDATE ENV (GENKEY);

    READ FILE(IND) INTO(INFIELD)
        KEY ('ABC');

    NEXT: READ FILE (IND) INTO (INFIELD);

    GO TO NEXT;
```

The first READ statement causes the first nondummy record in the data set whose key begins with “ABC” to be read into INFIELD; each time the second READ statement is executed, the nondummy record with the next higher key is retrieved. Repeated execution of the second READ statement could result in reading records from higher key classes, since no indication is given when the end of a key class is reached. It is your responsibility to check each key if you do not wish to read beyond the key class. Any subsequent execution of the first READ statement would reposition the file to the first record of the key class “ABC”.

If the data set contains no records with keys in the specified class, or if all the records with keys in the specified class are dummy records, the KEY condition is raised. The data set is then positioned either at the next record that has a higher key or at the end of the file.

The presence or absence of the GENKEY option affects the execution of a READ statement which supplies a source key that is shorter than the key length specified in the KEYLEN subparameter. This KEYLEN subparameter is found in the DD statement that defines the indexed data set. If you specify the GENKEY option, it causes the source key to be interpreted as a generic key, and the data set is positioned to the first nondummy record in the data set whose key begins with the source key. If you do not specify the GENKEY option, a READ statement’s short source key is padded on the right with blanks to the specified key length, and the data set is positioned to the record that has this padded key (if such a record exists). For a WRITE statement, a short source key is always padded with blanks.

Use of the GENKEY option does not affect the result of supplying a source key whose length is greater than or equal to the specified key length. The source key, truncated on the right if necessary, identifies a specific record (whose key can be considered to be the only member of its class).

**SCALARVARYING option — varying-length strings**

You use the SCALARVARYING option in the input/output of varying-length strings; you can use it with records of any format.
When storage is allocated for a varying-length string, the compiler includes a 2-byte prefix that specifies the current length of the string. For an element varying-length string, this prefix is included on output, or recognized on input, only if SCALARVARYING is specified for the file.

When you use locate mode statements (LOCATE and READ SET) to create and read a data set with element varying-length strings, you must specify SCALARVARYING to indicate that a length prefix is present, since the pointer that locates the buffer is always assumed to point to the start of the length prefix.

When you specify SCALARVARYING and element varying-length strings are transmitted, you must allow two bytes in the record length to include the length prefix.

A data set created using SCALARVARYING should be accessed only by a file that also specifies SCALARVARYING.

You must not specify SCALARVARYING and CTLASA/CTL360 for the same file, as this causes the first data byte to be ambiguous.

**KEYLENGTH option**

Use the KEYLENGTH option to specify the length of the recorded key for KEYED files where \( n \) is the length. You can specify KEYLENGTH for INDEXED files.

If you include the KEYLENGTH option in a VSAM file declaration for checking purposes, and the key length you specify in the option conflicts with the value defined for the data set, the UNDEFINEDFILE condition is raised.

**ORGANIZATION option**

The ORGANIZATION option specifies the organization of the data set associated with the PL/I file.

Data set types used by PL/I record I/O

Data sets with the RECORD attribute are processed by record-oriented data transmission in which data is transmitted to and from auxiliary storage exactly as it appears in the program variables; no data conversion takes place. A record in a data set corresponds to a variable in the program.
Table 15 shows the facilities that are available with the various types of data sets that can be used with PL/I Record I/O.

Table 15. A comparison of data set types available to PL/I record I/O

<table>
<thead>
<tr>
<th></th>
<th>VSAM KSDS</th>
<th>VSAM ESDS</th>
<th>VSAM RRDS</th>
<th>INDEXED</th>
<th>CONSECUTIVE</th>
<th>REGIONAL</th>
</tr>
</thead>
<tbody>
<tr>
<td>SEQUENCE</td>
<td>Key order</td>
<td>Entry order</td>
<td>Numbered</td>
<td>Key order</td>
<td>Entry order</td>
<td>By region</td>
</tr>
<tr>
<td>DEVICES</td>
<td>DASD</td>
<td>DASD</td>
<td>DASD</td>
<td>DASD</td>
<td>DASD, card, etc.</td>
<td>DASD</td>
</tr>
</tbody>
</table>

**ACCESS**

<table>
<thead>
<tr>
<th></th>
<th></th>
<th>123</th>
<th>123</th>
<th>123</th>
<th>12</th>
<th>2</th>
<th>12</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>By key</td>
<td>123</td>
<td>123</td>
<td>123</td>
<td>12</td>
<td>2</td>
<td>12</td>
</tr>
<tr>
<td>2</td>
<td>Sequential</td>
<td>123</td>
<td>123</td>
<td>123</td>
<td>12</td>
<td>2</td>
<td>12</td>
</tr>
<tr>
<td>3</td>
<td>Backward</td>
<td>123</td>
<td>123</td>
<td>123</td>
<td>12</td>
<td>2</td>
<td>12</td>
</tr>
</tbody>
</table>

Alternate index access as above:

<table>
<thead>
<tr>
<th></th>
<th></th>
<th>No</th>
<th>No</th>
<th>No</th>
<th>No</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>With new keys</td>
<td>No</td>
<td>No</td>
<td>No</td>
<td>No</td>
</tr>
<tr>
<td>2</td>
<td>At end</td>
<td>In empty slots</td>
<td>With new keys</td>
<td>At end</td>
<td>In empty slots</td>
</tr>
<tr>
<td>3</td>
<td>In empty slots</td>
<td>No</td>
<td>No</td>
<td>Yes, 1</td>
<td>No</td>
</tr>
</tbody>
</table>

**DELETION**

<table>
<thead>
<tr>
<th></th>
<th></th>
<th>Yes, 1</th>
<th>No</th>
<th>Yes, 1</th>
<th>Yes, 2</th>
<th>No</th>
<th>Yes, 1</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Space reusable</td>
<td>No</td>
<td>Yes, 1</td>
<td>Yes, 2</td>
<td>No</td>
<td>Yes, 1</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>Space not reusable</td>
<td>No</td>
<td>Yes, 1</td>
<td>Yes, 2</td>
<td>No</td>
<td>Yes, 1</td>
<td></td>
</tr>
</tbody>
</table>

The following chapters describe how to use Record I/O data sets for different types of data sets:

- Chapter 8, “Defining and using consecutive data sets,” on page 223
- Chapter 11, “Defining and using regional data sets,” on page 253
- Chapter 12, “Defining and using VSAM data sets,” on page 265

Setting environment variables under z/OS UNIX

There are a number of environment variables that can be set and exported for use with z/OS UNIX.

To set the environment variables system wide so all users have access to them, add the lines suggested in the subsections to the file `/etc/profile`. To set them for a specific user only, add them to the file `.profile` in the user's home directory. The variables are set the next time the user logs on.

The following example illustrates how to set environment variables:

```
LANG=ja_JP
NLSPATH=/usr/lib/nls/msg/%L/%N:/usr/lib/nls/msg/prime/%N
LIBPATH=/home/joe/usr/lib:/home/joe/mylib:/usr/lib
export LANG NLSPATH LIBPATH
```

Rather than using the last statement in the previous example, you could have added `export` to each of the preceding lines (export LANG=ja_JP...).
You can use the ECHO command to determine the current setting of an environment variable. To define the value of BYPASS, you can use either of the following two examples:

```bash
echo $LANG

echo $LIBPATH
```

**PL/I standard files (SYSPRINT and SYSIN) under z/OS UNIX**

SYSIN is read from stdin and SYSPRINT is directed to stdout by default. If you want either to be associated differently, you must use the TITLE option of the OPEN statement, or establish a DD_DDNAME environment variable naming a data set or another device. Environment variables are discussed above in “Setting environment variables under z/OS UNIX” on page 215.

**Redirecting standard input, output, and error devices under z/OS UNIX**

You can also redirect standard input, standard output, and standard error devices to a file. You can use redirection in the following program:

```pli
Hello2: proc options(main);
    put list('Hello!');
end;
```

After compiling and linking the program, you can invoke it from the command line by entering:

```bash
hello2 > hello2.out
```

If you want to combine stdout and stderr in a single file, enter the following command:

```bash
hello2 > hello2.out 2>&1
```

As is true with display statements, the `>` sign redirects the output to the file that is specified after it, in this case `hello2.out`. This means that the word ‘Hello’ is written in the file `hello2.out`. Note also that the output includes printer control characters since the PRINT attribute is applied to SYSPRINT by default.

READ statements can also access data from stdin.
Chapter 7. Using libraries

Within the z/OS operating system, the terms “partitioned data set”, “partitioned data set/extension”, and “library” are synonymous and refer to a type of data set that can be used for the storage of other data sets (usually programs in the form of source, object or load modules). A library must be stored on direct access storage and be wholly contained in one volume. It contains independent, consecutively organized data sets, called members. Each member has a unique name, not more than 8 characters long, which is stored in a directory that is part of the library. All the members of one library must have the same data characteristics because only one data set label is maintained.

You can create members individually until there is insufficient space left for a new entry in the directory, or until there is insufficient space for the member itself. You can access members individually by specifying the member name.

Use DD statements or their conversational mode equivalent to create and access members.

You can delete members by means of the IBM utility program IEHPROGM. This deletes the member name from the directory so that the member can no longer be accessed, but you cannot use the space occupied by the member itself again unless you recreate the library or compress the unused space using, for example, the IBM utility program IEBCOPY. If you attempt to delete a member by using the DISP parameter of a DD statement, it causes the whole data set to be deleted.

Types of libraries

You can use the following types of libraries with a PL/I program:

- The system program library SYS1.LINKLIB or its equivalent. This can contain all system processing programs such as compilers and the linkage editor.
- Private program libraries. These usually contain user-written programs. It is often convenient to create a temporary private library to store the load module output from the linkage editor until it is executed by a later job step in the same job. The temporary library will be deleted at the end of the job. Private libraries are also used for automatic library call by the linkage editor and the loader.
- The system procedure library SYS1.PROCLIB or its equivalent. This contains the job control procedures that have been cataloged for your installation.

How to use a library

A PL/I program can use a library directly. If you are adding a new member to a library, its directory entry will be made by the operating system when the associated file is closed, using the member name specified as part of the data set name.

If you are accessing a member of a library, its directory entry can be found by the operating system from the member name that you specify as part of the data set name.
More than one member of the same library can be processed by the same PL/I program, but only one such file can be open as output at any one time. You access different members by giving the member name in a DD statement.

Creating a library

To create a library include in your job step a DD statement containing the information given in Table 16. The information required is similar to that for a consecutively organized data set (see “Defining files using record I/O” on page 244) except for the SPACE parameter.

Table 16. Information required when creating a library

<table>
<thead>
<tr>
<th>Information Required</th>
<th>Parameter of DD statement</th>
</tr>
</thead>
<tbody>
<tr>
<td>Type of device that will be used</td>
<td>UNIT=</td>
</tr>
<tr>
<td>Serial number of the volume that will contain the library</td>
<td>VOLUME=SER</td>
</tr>
<tr>
<td>Name of the library</td>
<td>DSNAME=</td>
</tr>
<tr>
<td>Amount of space required for the library</td>
<td>SPACE=</td>
</tr>
<tr>
<td>Disposition of the library</td>
<td>DISP=</td>
</tr>
</tbody>
</table>

SPACE parameter

The SPACE parameter in a DD statement that defines a library must always be of the form:

\[ \text{SPACE=}(\text{units},(\text{quantity},\text{increment},\text{directory})) \]

Although you can omit the third term (increment), indicating its absence by a comma, the last term, specifying the number of directory blocks to be allocated, must always be present.

The amount of auxiliary storage required for a library depends on the number and sizes of the members to be stored in it and on how often members will be added or replaced. (Space occupied by deleted members is not released.) The number of directory blocks required depends on the number of members and the number of aliases. You can specify an incremental quantity in the SPACE parameter that allows the operating system to obtain more space for the data set, if such is necessary at the time of creation or at the time a new member is added; the number of directory blocks, however, is fixed at the time of creation and cannot be increased.

For example, the DD statement:

\[
\begin{align*}
// & \text{ PDS DD UNIT=SYSDA, VOL=SER=3412,} \\
// & \text{ DSNAME=ALIB,} \\
// & \text{ SPACE=(CYL,(5,,10),} \\
// & \text{ DISP=(,CATLG)}
\end{align*}
\]

requests the job scheduler to allocate 5 cylinders of the DASD with a volume serial number 3412 for a new library name ALIB, and to enter this name in the system catalog. The last term of the SPACE parameter requests that part of the space allocated to the data set be reserved for ten directory blocks.
Creating and updating a library member

The members of a library must have identical characteristics. Otherwise, you might later have difficulty retrieving them. Identical characteristics are necessary because the volume table of contents (VTOC) will contain only one data set control block (DSCB) for the library and not one for each member. When using a PL/I program to create a member, the operating system creates the directory entry; you cannot place information in the user data field.

When creating a library and a member at the same time, your DD statement must include all the parameters listed under "Creating a library" on page 218 (although you can omit the DISP parameter if the data set is to be temporary). The DSNAME parameter must include the member name in parentheses. For example, DSNAME=ALIB(MEM1) names the member MEM1 in the data set ALIB. If the member is placed in the library by the linkage editor, you can use the linkage editor NAME statement or the NAME compile-time option instead of including the member name in the DSNAME parameter. You must also describe the characteristics of the member (record format, etc.) either in the DCB parameter or in your PL/I program. These characteristics will also apply to other members added to the data set.

When creating a member to be added to an existing library, you do not need the SPACE parameter. The original space allocation applies to the whole of the library and not to an individual member. Furthermore, you do not need to describe the characteristics of the member, since these are already recorded in the DSCB for the library.

To add two more members to a library in one job step, you must include a DD statement for each member, and you must close one file that refers to the library before you open another.

Examples

The use of the cataloged procedure IBMZC to compile a simple PL/I program and place the object module in a new library named EXLIB is shown in Figure 23 on page 220. The DD statement that defines the new library and names the object module overrides the DD statement SYSLIN in the cataloged procedure. (The PL/I program is a function procedure that, given two values in the form of the character string produced by the TIME built-in function, returns the difference in milliseconds.)

The use of the cataloged procedure IBMZCL to compile and link-edit a PL/I program and place the load module in the existing library HPU8.CCLM is shown in Figure 24 on page 220.
To use a PL/I program to add or delete one or more records within a member of a library, you must rewrite the entire member in another part of the library. This is rarelly an economic proposition, since the space originally occupied by the member cannot be used again. You must use two files in your PL/I program, but both can be associated with the same DD statement. The program shown in Figure 26 on page 221 updates the member created by the program in Figure 25 on page 221. It copies all the records of the original member except those that contain only blanks.

Figure 23. Creating new libraries for compiled object modules

Figure 24. Placing a load module in an existing library
Extracting information from a library directory

The directory of a library is a series of records (entries) at the beginning of the data set. There is at least one directory entry for each member. Each entry contains a member name, the relative address of the member within the library, and a variable amount of user data.
User data is information inserted by the program that created the member. An entry that refers to a member (load module) written by the linkage editor includes user data in a standard format, described in the systems manuals.

If you use a PL/I program to create a member, the operating system creates the directory entry for you and you cannot write any user data. However, you can use assembler language macro instructions to create a member and write your own user data. The method for using macro instructions to do this is described in the data management manuals.
### Chapter 8. Defining and using consecutive data sets

This chapter covers consecutive data set organization and the ENVIRONMENT options that define consecutive data sets for stream and record-oriented data transmission. It then covers how to create, access, and update consecutive data sets for each type of transmission.

In a data set with consecutive organization, records are organized solely on the basis of their successive physical positions; when the data set is created, records are written consecutively in the order in which they are presented. You can retrieve the records only in the order in which they were written. See Table 14 on page 207 for valid file attributes and ENVIRONMENT options for consecutive data sets.

### Using stream-oriented data transmission

This section covers how to define data sets for use with PL/I files that have the STREAM attribute. It covers the ENVIRONMENT options you can use and how to create and access data sets. The essential parameters of the DD statements you use in creating and accessing these data sets are summarized in tables, and several examples of PL/I programs are included to illustrate the text.

Data sets with the STREAM attribute are processed by stream-oriented data transmission, which allows your PL/I program to ignore block and record boundaries and treat a data set as a continuous stream of data values in character or graphic form.

You create and access data sets for stream-oriented data transmission using the list-, data-, and edit-directed input and output statements described in the PL/I Language Reference.

For output, PL/I converts the data items from program variables into character form if necessary, and builds the stream of characters or graphics into records for transmission to the data set.

For input, PL/I takes records from the data set and separates them into the data items requested by your program, converting them into the appropriate form for assignment to program variables.

You can use stream-oriented data transmission to read or write graphic data. There are terminals, printers, and data-entry devices that, with the appropriate programming support, can display, print, and enter graphics. You must be sure that your data is in a format acceptable for the intended device, or for a print utility program.

### Defining files using stream I/O

You define files for stream-oriented data transmission by a file declaration with the following attributes:

```pli
DCL filename FILE STREAM
    INPUT |
    {OUTPUT [PRINT]}
    ENVIRONMENT(options);
```

© Copyright IBM Corp. 1999, 2015

223
Default file attributes are shown in Table 14 on page 207. The FILE attribute is described in the PL/I Language Reference. The PRINT attribute is described further in "Using PRINT files with stream I/O" on page 231. Options of the ENVIRONMENT attribute are discussed below.

### Defining stream files using PL/I dynamic allocation

To define the stream files, you can use a DD statement, an environment variable, or the TITLE option of the OPEN statement.

When an environment variable or the TITLE option is used, the name must be in uppercase. Specify the MVS data set in one of the following ways:

- **DSN(data-set-name)**
- **DSN(data-set-name(member-name))**

*data-set-name* must be fully qualified and cannot be a temporary data set; for example, it must not start with &.

Specify the HFS file as follows:

**PATH(absolute-path-name)**

You can specify the following attributes in any order after the DSN keyword:

- NEW, OLD, SHR, or MOD
- TRACKS or CYL
- SPACE\((n,m)\)
- VOL(volser)
- UNIT(type)
- KEEP, DELETE, CATALOG, or UNCATALOG
- STORCLAS(storageclass)
- MGMTCLAS(managementclass)
- DATACLAS(dataclass)

**Note:** You cannot create a PDS or PDSE using an environment variable or the TITLE option of the OPEN statement, but you can create a new member in an existing PDS or PDSE.

### Specifying ENVIRONMENT options

Table 14 on page 207 summarizes the ENVIRONMENT options. The options applicable to stream-oriented data transmission are:

- **CONSECUTIVE** or **ORGANIZATION(CONSECUTIVE)**
- **F|FB|FS|V|VB|VS|VBS|U**
- **RECSIZE(record-length)**
- **BLKSIZE(block-size)**
- **GRAPHIC**
- **LEAVE|REREAD**

**BLKSIZE** is described in Chapter 6, “Using data sets and files,” on page 185, beginning on page "BLKSIZE option" on page 210. **LEAVE** and **REREAD** are described later in this chapter, beginning at "LEAVE|REREAD" on page 246. Descriptions of the rest of these options follow immediately below.

**CONSECUTIVE**

STREAM files must have CONSECUTIVE data set organization; however, it is not necessary to specify this in the ENVIRONMENT options since CONSECUTIVE is the default data set organization. The CONSECUTIVE option for STREAM files is the same as that described in "Data set organization” on page 202."
Record format options

Although record boundaries are ignored in stream-oriented data transmission, record format is important when creating a data set. This is not only because record format affects the amount of storage space occupied by the data set and the efficiency of the program that processes the data, but also because the data set can later be processed by record-oriented data transmission.

Having specified the record format, you need not concern yourself with records and blocks as long as you use stream-oriented data transmission. You can consider your data set a series of characters or graphics arranged in lines, and you can use the SKIP option or format item (and, for a PRINT file, the PAGE and LINE options and format items) to select a new line.

Records can have one of the following formats, which are described in details in “Record formats” on page 200.

<table>
<thead>
<tr>
<th>Format</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>F</td>
<td>fixed-length, unblocked</td>
</tr>
<tr>
<td>FB</td>
<td>fixed-length, blocked</td>
</tr>
<tr>
<td>FS</td>
<td>fixed-length, standard</td>
</tr>
<tr>
<td>FBS</td>
<td>fixed-length, blocked, standard</td>
</tr>
<tr>
<td>V</td>
<td>variable-length, unblocked</td>
</tr>
<tr>
<td>VB</td>
<td>variable-length, blocked</td>
</tr>
<tr>
<td>VS</td>
<td>variable-length, standard</td>
</tr>
<tr>
<td>VBS</td>
<td>variable-length, blocked</td>
</tr>
<tr>
<td>U</td>
<td>undefined-length</td>
</tr>
</tbody>
</table>

Blocking and deblocking of records are performed automatically.

**RECSIZE**

RECSIZE for stream-oriented data transmission is the same as that described in “Specifying characteristics in the ENVIRONMENT attribute” on page 207. Additionally, a value specified by the LINESIZE option of the OPEN statement overrides a value specified in the RECSIZE option. LINESIZE is discussed in the PL/I Language Reference.

Additional record-size considerations for list- and data-directed transmission of graphics are given in the PL/I Language Reference.

**Defaults for record format, BLKSIZE, and RECSIZE**

If you do not specify the record format, BLKSIZE, or RECSIZE option in the ENVIRONMENT attribute, or in the associated DD statement or data set label, the following action is taken:
Input files:
Defaults are applied as for record-oriented data transmission, described in "Record format, BLKSIZE, and RECSIZE defaults" on page 212.

Output files:

Record format
Set to VB-format

Record length
The specified or default LINESIZE value is used:
• PRINT files:
  – F, FB, FBS, or U: line size + 1
  – V or VB: line size + 5
• Non-PRINT files:
  – F, FB, FBS, or U: linesize
  – V or VB: linesize + 4

Block Size
Files associated with SYSOUT:
• F, FB, or FBS: record length
• V or VB: record length + 4

New or Temporary data set:
• Optimum block size determined by DFP

GRAPHIC option
Specify the GRAPHIC option for edit-directed I/O.

The ERROR condition is raised for list- and data-directed I/O if you have graphics in input or output data and do not specify the GRAPHIC option.

For edit-directed I/O, the GRAPHIC option specifies that left and right delimiters are added to DBCS variables and constants on output, and that input graphics will have left and right delimiters. If you do not specify the GRAPHIC option, left and right delimiters are not added to output data, and input graphics do not require left and right delimiters. When you do specify the GRAPHIC option, the ERROR condition is raised if left and right delimiters are missing from the input data.

For information on the graphic data type, and on the G-format item for edit-directed I/O, see the PL/I Language Reference.

Creating a data set with stream I/O
To create a data set, you must give the operating system certain information either in your PL/I program or in the DD statement that defines the data set. For z/OS UNIX, use one of the following to provide the additional information:
• TITLE option of the OPEN statement
• DD_DDNAME environment variable
• ENVIRONMENT attribute

The following paragraphs indicate the essential information, and discuss some of the optional information you can supply.
Essential information
When your application creates a STREAM file, PL/I will derive a line-size value for that file from one of the following sources in order of declining precedence.
• LINESIZE option of the OPEN statement
• RECSIZE option of the ENVIRONMENT attribute
• RECSIZE option of the TITLE option of the OPEN statement
• RECSIZE option of the DD_DDNAME environment variable
• PL/I-supplied default value

If a LINESIZE value is supplied, but a RECSIZE value is not, then PL/I derives the record-length value as follows:
• For a V-format PRINT file, the value is LINESIZE + 5
• For a V-format non-PRINT file, the value is LINESIZE + 4
• For a F-format PRINT file, the value is LINESIZE + 1
• In all other cases, the value is LINESIZE

If a LINESIZE value is not supplied, but a RECSIZE value is, then PL/I derives the line-size value from RECSIZE as follows:
• For a V-format PRINT file, the value is RECSIZE - 5
• For a V-format non-PRINT file, the value is RECSIZE - 4
• For a F-format PRINT file, the value is RECSIZE - 1
• In all other cases, the value is RECSIZE

If neither LINESIZE nor RECSIZE is supplied, then PL/I determines a default line-size value based on the attributes of the file and the type of associated data set. In cases where PL/I cannot supply an appropriate default line size, the UNDEFINEDFILE condition is raised.

A default line-size value is supplied for an OUTPUT file when:
• The file has the PRINT attribute. In this case, the value is obtained from the tab control table.
• The associated data set is the terminal (stdout: or stderr:). In this case the value is 120.

Note that if the LINESIZE option is specified (on the OPEN statement) and RECSIZE is also specified (in the ENVIRONMENT attribute, the TITLE option or the DD statement), if the record size value is too small to hold the LINESIZE (taking into account the record format and appropriate control byte overhead), then
• For users who have LE for z/OS 1.9 or earlier releases
  For DD SYSOUT= files, the LINESIZE option will be used to determine a new record size that matches the given LINESIZE. For DD DSN= files and all other files, the UNDEFINEDFILE condition will be raised.
• For users with LE for z/OS releases subsequent to 1.9
  The UNDEFINEDFILE condition will be raised for all files

Examples
The use of edit-directed stream-oriented data transmission to create a data set on a direct access storage device is shown in Figure 27 on page 228. The data read from the input stream by the file SYSIN includes a field VREC that contains five unnamed 7-character subfields; the field NUM defines the number of these
subfields that contain information. The output file WORK transmits to the data set the whole of the field FREC and only those subfields of VREC that contain information.

```pli
//EX7#2 JOB
//STEP1 EXEC IBMZCBG
//PLI.SYSIN DD *
PEOPLE: PROC OPTIONS(MAIN);
  DCL WORK FILE STREAM OUTPUT,
  1 REC,
    2 FREC,
    3 NAME CHAR(19),
    3 NUM CHAR(1),
    3 PAD CHAR(25),
    2 VREC CHAR(35),
    EOF BIT(1) INIT('0'B),
    IN CHAR(80) DEF REC;
  ON ENDFILE(SYSIN) EOF='1'B;
  OPEN FILE(WORK) LINESIZE(400);
  GET FILE(SYSIN) EDIT(IN)(A(80));
  DO WHILE (¬EOF);
    PUT FILE(WORK) EDIT(IN)(A(45+7*NUM));
    GET FILE(SYSIN) EDIT(IN)(A(80));
  END;
  CLOSE FILE(WORK);
END PEOPLE;

/*@ 
//GO.WORK DD DSN=HPU8.PEOPLE,DISP=(NEW,CATLG),UNIT=SYSDA, 
// SPACE=(TRK,(1,1))
//GO.SYSIN DD *
R.C.ANDERSON 0 202848 DOCTOR
B.F.BENNETT 2 771239 PLUMBER VICTOR HAZEL
R.E.COLE 5 698635 COOK ELLEN VICTOR JOAN ANN OTTO
J.F.COOPER 5 418915 LAWYER FRANK CAROL DONALD NORMAN BRENDA
A.J.CORNELL 3 237837 BARBER ALBERT ERIC JANET
E.F.FERRIS 4 158636 CARPENTER GERALD ANNA MARY HAROLD
/*@ 

Figure 27. Creating a data set with stream-oriented data transmission
```

Figure 28 on page 229 shows an example of a program using list-directed output to write graphics to a stream file. It assumes that you have an output device that can print graphic data. The program reads employee records and selects persons living in a certain area. It then edits the address field, inserting one graphic blank between each address item, and prints the employee number, name, and address.
//EX7#3 JOB
//STEP1 EXEC IBMZCBG
//PLI.SYSIN DD *
% PROCESS GRAPHIC;

XAMPLE1: PROC OPTIONS(MAIN);
DCL INFILE FILE INPUT RECORD,
OUTFILE FILE OUTPUT STREAM ENV(GRAPHIC);
/* GRAPHIC OPTION MEANS DELIMITERS WILL BE INSERTED ON OUTPUT FILES. */
DCL 1 IN,
3 EMPNO CHAR(6),
3 SHIFT1 CHAR(1),
3 NAME,
5 LAST G(7),
5 FIRST G(7),
3 SHIFT2 CHAR(1),
3 ADDRESS,
5 ZIP CHAR(6),
5 SHIFT3 CHAR(1),
5 DISTRICT G(5),
5 CITY G(5),
5 OTHER G(8),
5 SHIFT4 CHAR(1);
DCL EOF BIT(1) INIT('0'B);
DCL ADDRWK G(20);
ON ENDFILE (INFILE) EOF = '1'B;
READ FILE(INFILE) INTO(IN);
DO WHILE(¬EOF);
  DO;
    IF SUBSTR(ZIP,1,3)¬='300'
      THEN LEAVE;
    L=0;
    ADDRWK=DISTRICT;
    DO I=1 TO 5;
      IF SUBSTR(DISTRICT,I,1)¬= <
        THEN LEAVE; /* SUBSTR BIF PICKS 3P */
      END; /* THE ITH GRAPHIC CHAR */
      L=L+I+1; /* IN DISTRICT */
      SUBSTR(ADDRWK,L,5)=CITY;
      DO I=1 TO 5;
        IF SUBSTR(CITY,I,1)= <
          THEN LEAVE;
      END;
      L=L+I;
      SUBSTR(ADDRWK,L,8)=OTHER;
    PUT FILE(OUTFILE) SKIP /* THIS DATA SET */
      EDIT(EMPNO,IN.LAST,FIRST,ADDRWK) /* REQUIRES UTILITY */
        (A(8),G(7),G(7),X(4),G(20)); /* TO PRINT GRAPHIC */
        /* DATA */
      END; /* END OF NON-ITERATIVE DO */
  READ FILE(INFILE) INTO (IN);
END; /* END OF DO WHILE(¬EOF) */
END XAMPLE1;

/*
GO.OUTFILE DD SYSOUT=A,DCB=(RECFM=VB,LRECL=121,BLKSIZE=129)
GO.INFILE DD *
ABCODEF<
>300099< 3 3 3 3 3 3 3>
ABCD <
>300011< 3 3 3 3>
*/

Figure 28. Writing graphic data to a stream file

Chapter 8. Defining and using consecutive data sets 229
Accessing a data set with stream I/O

A data set accessed using stream-oriented data transmission need not have been created by stream-oriented data transmission, but it must have CONSECUTIVE organization, and all the data in it must be in character or graphic form. You can open the associated file for input, and read the records the data set contains; or you can open the file for output, and extend the data set by adding records at the end.

To access a data set, you must use one of the following to identify it:
- ENVIRONMENT attribute
- DD_DDNAME environment variable
- TITLE option of the OPEN statement

The following paragraphs describe the essential information you must include in the DD statement, and discuss some of the optional information you can supply. The discussions do not apply to data sets in the input stream.

Essential information

When your application accesses an existing STREAM file, PL/I must obtain a record-length value for that file. If the data set does not have a record-length, the value can come from one of the following sources:
- The LINESIZE option of the OPEN statement
- The RECSIZE option of the ENVIRONMENT attribute
- The RECSIZE option of the DD_DDNAME environment variable
- The RECSIZE option of the TITLE option of the OPEN statement
- PL/I-supplied default value

If you are using an existing OUTPUT file, or if you supply a RECSIZE value, PL/I determines the record-length value as described in “Creating a data set with stream I/O” on page 226.

PL/I uses a default record-length value for an INPUT file when:
- The file is SYSIN, value = 80
- The file is associated with the terminal (stdout: or stderr:), value = 120

Record format

When using stream-oriented data transmission to access a data set, you do not need to know the record format of the data set (except when you must specify a block size); each GET statement transfers a discrete number of characters or graphics to your program from the data stream.

If you do give record-format information, it must be compatible with the actual structure of the data set. For example, if a data set is created with F-format records, a record size of 600 bytes, and a block size of 3600 bytes, you can access the records as if they are U-format with a maximum block size of 3600 bytes; but if you specify a block size of 3500 bytes, your data will be truncated.

Example

The program in Figure 29 on page 231 reads the data set created by the program in Figure 27 on page 228 and uses the file SYSPRINT to list the data it contains. (For details on SYSPRINT, see “Using SYSIN and SYSPRINT files” on page 235.) Each set of data is read, by the GET statement, into two variables: FREC, which always contains 45 characters; and VREC, which always contains 35 characters. At each execution of the GET statement, VREC consists of the number of characters generated by the expression 7*NUM, together with sufficient blanks to bring the total number of characters to 35. The DISP parameter of the DD statement could
read simply DISP=OLD; if DELETE is omitted, an existing data set will not be deleted.

```
//EX7#5 JOB
//STEP1 EXEC IBMZCBG
//PLI.SYSIN DD *
PEOPLE: PROC OPTIONS(MAIN);
  DCL WORK FILE STREAM INPUT,
    1 REC,
    2 FREC,
      3 NAME CHAR(19),
      3 NUM CHAR(1),
      3 SERNO CHAR(7),
      3 PROF CHAR(18),
    2 VREC CHAR(35),
    IN CHAR(80) DEF REC,
    EOF BIT(1) INIT('0'B);
  ON ENDFILE(WORK) EOF='1'B;
  OPEN FILE(WORK);
  GET FILE(WORK) EDIT(IN,VREC)(A(45),A(7*NUM));
  DO WHILE (¬EOF);
    PUT FILE(SYSPRINT) SKIP EDIT(IN)(A);
    GET FILE(WORK) EDIT(IN,VREC)(A(45),A(7*NUM));
  END;
  CLOSE FILE(WORK);
END PEOPLE;

/*
//GO.WORK DD DSN=HPU8.PEOPLE,DISP=(OLD,DELETE)
```

Figure 29. Accessing a data set with stream-oriented data transmission

Using PRINT files with stream I/O

Both the operating system and the PL/I language include features that facilitate the formatting of printed output. The operating system allows you to use the first byte of each record for a print control character. The control characters, which are not printed, cause the printer to skip to a new line or page. (Tables of print control characters are given in Figure 32 on page 246 and Table 18 on page 246.)

In a PL/I program, the use of a PRINT file provides a convenient means of controlling the layout of printed output from stream-oriented data transmission. The compiler automatically inserts print control characters in response to the PAGE, SKIP, and LINE options and format items.

You can apply the PRINT attribute to any STREAM OUTPUT file, even if you do not intend to print the associated data set directly. When a PRINT file is associated with a direct access data set, the print control characters have no effect on the layout of the data set, but appear as part of the data in the records.

PRINT files opened with FB or VB will cause the UNDEFINEDFILE condition to be raised. PRINT files should be opened with the "A" option; i.e., FBA or VBA.

The compiler reserves the first byte of each record transmitted by a PRINT file for an American National Standard print control character, and inserts the appropriate characters automatically.

A PRINT file uses only the following five print control characters:

<table>
<thead>
<tr>
<th>Character</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>//EX7#5</td>
<td>JOB</td>
</tr>
<tr>
<td>//STEP1</td>
<td>EXEC</td>
</tr>
<tr>
<td>//PLI.SYSIN</td>
<td>DD *</td>
</tr>
<tr>
<td>PEOPLE: PROC OPTIONS(MAIN);</td>
<td></td>
</tr>
<tr>
<td>DCL WORK FILE STREAM INPUT,</td>
<td></td>
</tr>
<tr>
<td>1 REC,</td>
<td></td>
</tr>
<tr>
<td>2 FREC,</td>
<td></td>
</tr>
<tr>
<td>3 NAME CHAR(19),</td>
<td></td>
</tr>
<tr>
<td>3 NUM CHAR(1),</td>
<td></td>
</tr>
<tr>
<td>3 SERNO CHAR(7),</td>
<td></td>
</tr>
<tr>
<td>3 PROF CHAR(18),</td>
<td></td>
</tr>
<tr>
<td>2 VREC CHAR(35),</td>
<td></td>
</tr>
<tr>
<td>IN CHAR(80) DEF REC,</td>
<td></td>
</tr>
<tr>
<td>EOF BIT(1) INIT('0'B);</td>
<td></td>
</tr>
<tr>
<td>ON ENDFILE(WORK) EOF='1'B;</td>
<td></td>
</tr>
<tr>
<td>OPEN FILE(WORK);</td>
<td></td>
</tr>
<tr>
<td>GET FILE(WORK) EDIT(IN,VREC)(A(45),A(7*NUM));</td>
<td></td>
</tr>
<tr>
<td>DO WHILE (¬EOF);</td>
<td></td>
</tr>
<tr>
<td>PUT FILE(SYSPRINT) SKIP EDIT(IN)(A);</td>
<td></td>
</tr>
<tr>
<td>GET FILE(WORK) EDIT(IN,VREC)(A(45),A(7*NUM));</td>
<td></td>
</tr>
<tr>
<td>END;</td>
<td></td>
</tr>
<tr>
<td>CLOSE FILE(WORK);</td>
<td></td>
</tr>
<tr>
<td>END PEOPLE;</td>
<td></td>
</tr>
<tr>
<td>/*</td>
<td></td>
</tr>
<tr>
<td>//GO.WORK DD DSN=HPU8.PEOPLE,DISP=(OLD,DELETE)</td>
<td></td>
</tr>
</tbody>
</table>
The compiler handles the PAGE, SKIP, and LINE options or format items by padding the remainder of the current record with blanks and inserting the appropriate control character in the next record. If SKIP or LINE specifies more than a 3-line space, the compiler inserts sufficient blank records with appropriate control characters to accomplish the required spacing. In the absence of a print control option or format item, when a record is full the compiler inserts a blank character (single line space) in the first byte of the next record.

If a PRINT file is being transmitted to a terminal, the PAGE, SKIP, and LINE options will never cause more than 3 lines to be skipped, unless formatted output is specified.

**Controlling printed line length**

You can limit the length of the printed line produced by a PRINT file either by specifying a record length in your PL/I program (ENVIRONMENT attribute) or in a DD statement, or by giving a line size in an OPEN statement (LINESIZE option). The record length must include the extra byte for the print control character, that is, it must be 1 byte larger than the length of the printed line (5 bytes larger for V-format records). The value you specify in the LINESIZE option refers to the number of characters in the printed line; the compiler adds the print control character.

The blocking of records has no effect on the appearance of the output produced by a PRINT file, but it does result in more efficient use of auxiliary storage when the file is associated with a data set on a direct access device. If you use the LINESIZE option, ensure that your line size is compatible with your block size. For F-format records, block size must be an exact multiple of (line size+1); for V-format records, block size must be at least 9 bytes greater than line size.

Although you can vary the line size for a PRINT file during execution by closing the file and opening it again with a new line size, you must do so with caution if you are using the PRINT file to create a data set on a direct access device. You cannot change the record format that is established for the data set when the file is first opened. If the line size you specify in an OPEN statement conflicts with the record format already established, the UNDEFINEDFILE condition is raised. To prevent this, either specify V-format records with a block size at least 9 bytes greater than the maximum line size you intend to use, or ensure that the first OPEN statement specifies the maximum line size. (Output destined for the printer can be stored temporarily on a direct access device, unless you specify a printer by using UNIT=, even if you intend it to be fed directly to the printer.)

Since PRINT files have a default line size of 120 characters, you need not give any record format information for them. In the absence of other information, the compiler assumes V-format records. The complete default information is:

```
BLKSIZE=129
LRECL=125
RECFM=VBA.
```
Example: Figure 30 on page 234 illustrates the use of a PRINT file and the printing options of stream-oriented data transmission statements to format a table and write it onto a direct access device for printing on a later occasion. The table comprises the natural sines of the angles from 0° to 359° 54' in steps of 6'. The statements in the ENDPAGE ON-unit insert a page number at the bottom of each page, and set up the headings for the following page. The DD statement defining the data set created by this program includes no record-format information. The compiler infers the following from the file declaration and the line size specified in the statement that opens the file TABLE:

- **Record format** = V (the default for a PRINT file).
- **Record size** = 98 (line size + 1 byte for print control character + 4 bytes for record control field).
- **Block size** = 102 (record length + 4 bytes for block control field).

The program in Figure 34 on page 252 uses record-oriented data transmission to print the table created by the program in Figure 30 on page 234.
Overriding the tab control table

Data-directed and list-directed output to a PRINT file are aligned on preset tabulator positions. See Figure 16 on page 165 and Figure 31 on page 235 for examples of declaring a tab table. The definitions of the fields in the table are as follows:

OFFSET OF TAB COUNT:
Halfword binary integer that gives the offset of “Tab count,” the field that indicates the number of tabs to be used.

PAGESIZE:
Halfword binary integer that defines the default page size. This page size is used for dump output to the PLIDUMP data set as well as for stream output.

%PROCESS INT F(I) AG A(F) OP STG NEST X(F) SOURCE ;
%PROCESS LIST;

SINE: PROC OPTIONS(MAIN);
  DCL TABLE FILE STREAM OUTPUT PRINT;
  DCL DEG FIXED DEC(5,1) INIT(0); /* INIT(0) FOR ENDPAGE */
  DCL MIN FIXED DEC(3,1);
  DCL PGNO FIXED DEC(2) INIT(0);
  DCL ONCODE BUILTIN;

  ON ERROR
  BEGIN;
    ON ERROR SYSTEM;
    DISPLAY ('ONCODE = '|| ONCODE);
  END;

  ON ENDPAGE(TABLE)
  BEGIN;
    DCL I;
    IF PGNO = 0 THEN
      PUT FILE(TABLE) EDIT ('PAGE',PGNO)
      (LINE(55),COL(80),A,F(3));
    IF DEG = 360 THEN
      PUT FILE(TABLE) PAGE EDIT ('NATURAL SINES') (A);
      IF PGNO = 0 THEN
        PUT FILE(TABLE) EDIT ((I DO I = 0 TO 54 BY 6))
        (SKIP(3),10 F(9));
      PGNO = PGNO + 1;
    ELSE
      PUT FILE(TABLE) PAGE;
    END;
  END;

  OPEN FILE(TABLE) PAGESIZE(52) LINESIZE(93);
  SIGNAL ENDPAGE(TABLE);

  PUT FILE(TABLE) EDIT
  ((DEG,(SIND(DEG+MIN) DO MIN = 0 TO .9 BY .1) DO DEG = 0 TO 359))
  (SKIP(2), 5 (COL(1), F(3), 10 F(9,4 )));

Figure 30. Creating a print file via stream data transmission. The example in Figure 34 on page 252 will print the resultant file.

Overriding the tab control table

Data-directed and list-directed output to a PRINT file are aligned on preset tabulator positions. See Figure 16 on page 165 and Figure 31 on page 235 for examples of declaring a tab table. The definitions of the fields in the table are as follows:

OFFSET OF TAB COUNT:
Halfword binary integer that gives the offset of “Tab count,” the field that indicates the number of tabs to be used.

PAGESIZE:
Halfword binary integer that defines the default page size. This page size is used for dump output to the PLIDUMP data set as well as for stream output.
**LINESIZE:**
Halfword binary integer that defines the default line size.

**PAGELENGTH:**
Halfword binary integer that defines the default page length for printing at a terminal.

**FILLERS:**
Three halfword binary integers; reserved for future use.

**TAB COUNT:**
Halfword binary integer that defines the number of tab position entries in the table (maximum 255). If tab count = 0, any specified tab positions are ignored.

**Tab1–Tabn:**
n halfword binary integers that define the tab positions within the print line. The first position is numbered 1, and the highest position is numbered 255. The value of each tab should be greater than that of the tab preceding it in the table; otherwise, it is ignored. The first data field in the printed output begins at the next available tab position.

You can override the default PL/I tab settings for your program by causing the linkage editor to resolve an external reference to PLITABS. To cause the reference to be resolved, supply a table with the name PLITABS, in the format described above.

To supply this tab table, include a PL/I structure in your source program with the name PLITABS, which you must declare to be STATIC EXTERNAL in your MAIN procedure or in a program linked with your MAIN procedure. An example of the PL/I structure is shown in Figure 31. This example creates three tab settings, in positions 30, 60, and 90, and uses the defaults for page size and line size. Note that TAB1 identifies the position of the second item printed on a line; the first item on a line always starts at the left margin. The first item in the structure is the offset to the NO_OF_TABS field. The FILL fields must not be omitted.

```pli
DCL 1 PLITABS STATIC EXT,
  2 (OFFSET INIT(14),
      PAGESIZE INIT(60),
      LINESIZE INIT(120),
      PAGELENGTH INIT(0),
      FILL1 INIT(0),
      FILL2 INIT(0),
      FILL3 INIT(0),
      NO_OF_TABS INIT(3),
      TAB1 INIT(30),
      TAB2 INIT(60),
      TAB3 INIT(90)) FIXED BIN(15,0);
```

*Figure 31. PL/I structure PLITABS for modifying the preset tab settings*

### Using SYSIN and SYSPRINT files

If you code a GET statement without the FILE option in your program, the compiler inserts the file name SYSIN. If you code a PUT statement without the FILE option, the compiler inserts the name SYSPRINT.

If you do not declare SYSPRINT, the compiler gives the file the attribute PRINT in addition to the normal default attributes. The complete set of attributes will be:
FILE STREAM OUTPUT PRINT EXTERNAL

Since SYSPRINT is a PRINT file, the compiler also supplies a default line size of 120 characters and a V-format record. You need give only a minimum of information in the corresponding DD statement; if your installation uses the usual convention that the system output device of class A is a printer, the following is sufficient:

//SYSPRINT DD SYSOUT=A

Note: SYSin and SYSPRINT are established in the User Exit during initialization. IBM-supplied defaults for SYSin and SYSPRINT are directed to the terminal.

You can override the attributes given to SYSPRINT by the compiler by explicitly declaring or opening the file. For more information about the interaction between SYSPRINT and the z/OS Language Environment message file option, see the z/OS Language Environment Programming Guide.

The compiler does not supply any special attributes for the input file SYSin; if you do not declare it, it receives only the default attributes. The data set associated with SYSin is usually in the input stream; if it is not in the input stream, you must supply full DD information.

For more information about SYSPRINT, see “SYSPRINT considerations” on page 167.

Controlling input from the terminal

You can enter data at the terminal for an input file in your PL/I program if you do the following:

1. Declare the input file explicitly or implicitly with the CONSECUTIVE environment option (all stream files meet this condition)
2. Allocate the input file to the terminal.

You can usually use the standard default input file SYSin because it is a stream file and can be allocated to the terminal.

You are prompted for input to stream files by a colon (:). You will see the colon each time a GET statement is executed in the program. The GET statement causes the system to go to the next line. You can then enter the required data. If you enter a line that does not contain enough data to complete execution of the GET statement, a further prompt, which is a plus sign followed by a colon (+:), is displayed.

By adding a hyphen to the end of any line that is to continue, you can delay transmission of the data to your program until you enter two or more lines.

If you include output statements that prompt you for input in your program, you can inhibit the initial system prompt by ending your own prompt with a colon. For example, the GET statement could be preceded by a PUT statement such as:

PUT SKIP LIST(‘ENTER NEXT ITEM:’);

To inhibit the system prompt for the next GET statement, your own prompt must meet the following conditions:

1. It must be either list-directed or edit-directed, and if list-directed, must be to a PRINT file.
2. The file transmitting the prompt must be allocated to the terminal. If you are merely copying the file at the terminal, the system prompt is not inhibited.

Under TSO, there is support of an environment variable called TSO_INPUT_OPT to help control the input from a terminal in a more flexible way than the default way described above.

The syntax for the TSO_INPUT_OPT environment variable, which has to be in uppercase, is:

```
TSO_INPUT_OPT = option
```

When specifying the options, which could be in uppercase or lowercase, blanks are not allowed. If more than one option is specified, they have to be separated by a comma. In addition, the syntax of the statement is not checked at the time the command is entered. It is verified when the data set is opened. If the syntax is wrong, the UNDEFINEDFILE condition is raised with the oncode 96.

The options you can specify are:

**PROMPT**

The PROMPT option specifies whether or not a colon should be visible as prompts for stream input from the terminal.

```
PROMPT (N) (Y)
```

**SAMELINE**

The SAMELINE option specifies whether the system prompt occurs on the same line as the statement that prompts for input.

```
SAMELINE (N) (Y)
```

For more information on how these options will affect your input and output to the terminal, see Example 1 and Example 2 under the SAMELINE option in Chapter 6.

One way to specify this environment variable under TSO is via the PLIXOPT string. For example:

```
DCL PLIXOPT char(50) var ext static
init('ENVAR("TSO_INPUT_OPT=PROMPT(Y),SAMELINE(Y)")');
```

**Format of data**

The data you enter at the terminal should have exactly the same format as stream input data in batch mode, except for the following variations:

- Simplified punctuation for input: If you enter separate items of input on separate lines, there is no need to enter intervening blanks or commas; the compiler will insert a comma at the end of each line.
  For instance, in response to the statement:
  ```
  GET LIST(I,J,K);
  ```
your terminal interaction could be as follows:

```
; 1 +:2 +:3
```

with a carriage return following each item. It would be equivalent to:

```
; 1,2,3
```

If you wish to continue an item onto another line, you must end the first line with a continuation character. Otherwise, for a GET LIST or GET DATA statement, a comma will be inserted, and for a GET EDIT statement, the item will be padded (see next paragraph).

- Automatic padding for GET EDIT: There is no need to enter blanks at the end of a line of input for a GET EDIT statement. The item you enter will be padded to the correct length.

For instance, for the PL/I statement:

```
GET EDIT(NAME)(A(15));
```

you could enter the five characters:

```
SMITH
```

followed immediately by a carriage return. The item will be padded with 10 blanks, so that the program receives a string 15 characters long. If you wish to continue an item on a second or subsequent line, you must add a continuation character to the end of every line except the last; the first line transmitted would otherwise be padded and treated as the complete data item.

- SKIP option or format item: A SKIP in a GET statement asks the program to ignore data not yet entered. All uses of SKIP(n) where \( n \) is greater than one are taken to mean SKIP(1). SKIP(1) is taken to mean that all unused data on the current line is ignored.

### Stream and record files

You can allocate both stream and record files to the terminal. However, no prompting is provided for record files. If you allocate more than one file to the terminal, and one or more of them is a record file, the output of the files will not necessarily be synchronized. The order in which data is transmitted to and from the terminal is not guaranteed to be the same order in which the corresponding PL/I I/O statements are executed.

Also, record file input from the terminal is received in upper case letters because of a TCAM restriction. To avoid problems you should use stream files wherever possible.

### Defining QSAM files using PL/I dynamic allocation

QSAM or HFS files can be defined using a DD statement, an environment variable, or the TITLE option of the OPEN statement.

When an environment variable or TITLE option is used, the name must be in uppercase. Specify the MVS data set in one of the following ways:

- `DSN(data-set-name)`
- `DSN(data-set-name (member-name))`
data-set-name must be fully qualified and cannot be a temporary data set; for example, it must not start with &.

Specify the HFS file as follows:

PATH (absolute-path-name)

The following attributes can be specified in any order after the DSN keyword:
- NEW, OLD, SHR, or MOD
- TRACKS or CYL
- SPACE(n,m)
- VOL(volser)
- UNIT(type)
- KEEP, DELETE, CATALOG, or UNCATALOG
- STORCLAS(storageclass)
- MGMTCLAS(managementclass)
- DATACLAS(dataclass)

Note: You cannot create a PDS or PDSE by using an environment variable or the TITLE option of the OPEN statement, but you can create a new member in an existing PDS or PDSE.

Capital and lowercase letters

For stream files, character strings are transmitted to the program as entered in lowercase or uppercase. For record files, all characters become uppercase.

End-of-file

The characters /* in positions one and two of a line that contains no other characters are treated as an end-of-file marker; that is, they raise the ENDFILE condition.

Under the USS environment, you can also use the key sequence ESC-D as the end-of-file marker. ESC is predefined as a specific character, and you need to use this predefined character in the sequence.

COPY option of GET statement

The GET statement can specify the COPY option; but if the COPY file, as well as the input file, is allocated to the terminal, no copy of the data will be printed.
Chapter 9. Controlling output to the terminal

At your terminal you can obtain data from a PL/I file that has been both:
1. Declared explicitly or implicitly with the CONSECUTIVE environment option. All stream files meet this condition.
2. Allocated to the terminal.

The standard print file SYSPRINT generally meets both these conditions.

Format of PRINT files

Data from SYSPRINT or other PRINT files is not normally formatted into pages at the terminal. Three lines are always skipped for PAGE and LINE options and format items. The ENDPAGE condition is normally never raised. SKIP(n), where n is greater than three, causes only three lines to be skipped. SKIP(0) is implemented by backspacing, and should therefore not be used with terminals that do not have a backspace feature.

You can cause a PRINT file to be formatted into pages by inserting a tab control table in your program. The table must be called PLITABS, and its contents are explained in "Overriding the tab control table" on page 234. You must initialize the element PAGELENGTH to the length of page you require—that is, the length of the sheet of paper on which each page is to be printed, expressed as the maximum number of lines that could be printed on it. You must initialize the element PAGESIZE to the actual number of lines to be printed on each page. After the number of lines in PAGESIZE has been printed on a page, ENDPAGE is raised, for which standard system action is to skip the number of lines equal to PAGELENGTH minus PAGESIZE, and then start printing the next page. For other than standard layout, you must initialize the other elements in PLITABS to the values shown in Figure 16 on page 165. You can also use PLITABS to alter the tabulating positions of list-directed and data-directed output. You can use PLITABS for SYSPRINT when you need to format page breaks in ILC applications. Set PAGESIZE to 32767 and use the PUT PAGE statement to control page breaks.

Although some types of terminals have a tabulating facility, tabulating of list-directed and data-directed output is always achieved by transmission of blank characters.

Stream and record files

You can allocate both stream and record files to the terminal. However, if you allocate more than one file to the terminal, and one of the files is SYSPRINT or a record file, the output of the files is not necessarily synchronized. There is no guarantee that the order in which data is transmitted between the program and the terminal is the same as the order in which the corresponding PL/I output statements are executed.

Output from the PUT EDIT command

The format of the output from a PUT EDIT command to a terminal is line mode TPUTs with “Start of field” and “end of field” characters appearing as blanks on the screen.
Chapter 10. Using record-oriented data transmission

PL/I supports various types of data sets with the RECORD attribute (see Table 20 on page 247). This section covers how to use consecutive data sets.

Table 17 lists the statements and options that you can use to create and access a consecutive data set using record-oriented data transmission.

Table 17. Statements and options allowed for creating and accessing consecutive data sets

<table>
<thead>
<tr>
<th>File declaration</th>
<th>Valid statements with Options you must specify</th>
<th>Other options you can specify</th>
</tr>
</thead>
<tbody>
<tr>
<td>SEQUENTIAL OUTPUT BUFFERED</td>
<td>WRITE FILE(file-reference) FROM(reference); LOCATE based-variable FILE(file-reference);</td>
<td>SET(pointer-reference)</td>
</tr>
<tr>
<td>SEQUENTIAL OUTPUT</td>
<td>WRITE FILE(file-reference) FROM(reference);</td>
<td></td>
</tr>
<tr>
<td>SEQUENTIAL INPUT BUFFERED</td>
<td>READ FILE(file-reference) INTO(reference); READ FILE(file-reference) SET(pointer-reference); READ FILE(file-reference) IGNORE(expression);</td>
<td></td>
</tr>
<tr>
<td>SEQUENTIAL INPUT</td>
<td>READ FILE(file-reference) INTO(reference); READ FILE(file-reference) IGNORE(expression);</td>
<td></td>
</tr>
<tr>
<td>SEQUENTIAL UPDATE BUFFERED</td>
<td>READ FILE(file-reference) INTO(reference); READ FILE(file-reference) SET(pointer-reference); READ FILE(file-reference) IGNORE(expression); REWRITE FILE(file-reference);</td>
<td>FROM(reference)</td>
</tr>
</tbody>
</table>
Table 17. Statements and options allowed for creating and accessing consecutive data sets (continued)

<table>
<thead>
<tr>
<th>File declaration</th>
<th>Valid statements with Options you must specify</th>
<th>Other options you can specify</th>
</tr>
</thead>
<tbody>
<tr>
<td>SEQUENTIAL UPDATE</td>
<td>READ FILE(file-reference) INTO(reference);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) IGNORE(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>REWRITE FILE(file-reference) FROM(reference);</td>
<td></td>
</tr>
</tbody>
</table>

Notes:
1. The complete file declaration would include the attributes FILE, RECORD and ENVIRONMENT.

2. The statement READ FILE (file-reference); is a valid statement and is equivalent to READ FILE(file-reference) IGNORE (1);

Specifying record format

If you give record-format information, it must be compatible with the actual structure of the data set. For example, if you create a data set with FB-format records, with a record size of 600 bytes and a block size of 3600 bytes, you can access the records as if they are U-format with a maximum block size of 3600 bytes. If you specify a block size of 3500 bytes, your data is truncated.

Defining files using record I/O

You define files for record-oriented data transmission by using a file declaration with the following attributes:

DCL filename FILE RECORD
  INPUT | OUTPUT | UPDATE
  SEQUENTIAL
  BUFFERED
  ENVIRONMENT(options);

Default file attributes are shown in Table 14 on page 207. The file attributes are described in the PL/I Language Reference. Options of the ENVIRONMENT attribute are discussed below.

Specifying ENVIRONMENT options

The ENVIRONMENT options applicable to consecutive data sets are:

F|FB|FS|FBS|V|VB|U
RECSIZE(record-length)
BLKSIZE(block-size)
SCALARVARYING
CONSECUTIVE or ORGANIZATION(CONSECUTIVE)
CTLASA|CTL360
LEAVE|REREAD
The options through SCALARVARYING are described in “Specifying characteristics in the ENVIRONMENT attribute” on page 207, and those after SCALARVARYING are described below.

See Table 14 on page 207 to find which options you must specify, which are optional, and which are defaults.

**CONSECUTIVE**

The CONSECUTIVE option defines a file with consecutive data set organization, which is described in this chapter and in “Data set organization” on page 202.

CONSECUTIVE is the default.

**ORGANIZATION(CONSECUTIVE)**

Specifies that the file is associated with a consecutive data set. The ORGANIZATION option is described in “ORGANIZATION option” on page 214.

The file can be either a native data set or a VSAM data set.

**CTLASACTL360**

The printer control options CTLASA and CTL360 apply only to OUTPUT files associated with consecutive data sets. They specify that the first character of a record is to be interpreted as a control character.

The CTLASA option specifies American National Standard Vertical Carriage Positioning Characters or American National Standard Pocket Select Characters (Level 1). The CTL360 option specifies IBM machine-code control characters.

The American National Standard control characters, listed in Figure 32 on page 246, cause the specified action to occur before the associated record is printed or punched.

The machine code control characters differ according to the type of device. The IBM machine code control characters for printers are listed in Table 18 on page 246.
### Table 18. IBM machine code print control characters (CTL360)

<table>
<thead>
<tr>
<th>Code byte</th>
<th>Action</th>
<th>Act immediately (no printing)</th>
</tr>
</thead>
<tbody>
<tr>
<td>00000001</td>
<td>Print only (no space)</td>
<td>—</td>
</tr>
<tr>
<td>00001001</td>
<td>Space 1 line</td>
<td>00001011</td>
</tr>
<tr>
<td>00010001</td>
<td>Space 2 lines</td>
<td>00010011</td>
</tr>
<tr>
<td>00011001</td>
<td>Space 3 lines</td>
<td>00011011</td>
</tr>
<tr>
<td>10001001</td>
<td>Skip to channel 1</td>
<td>10001011</td>
</tr>
<tr>
<td>10010001</td>
<td>Skip to channel 2</td>
<td>10010011</td>
</tr>
<tr>
<td>10011001</td>
<td>Skip to channel 3</td>
<td>10011011</td>
</tr>
<tr>
<td>10100001</td>
<td>Skip to channel 4</td>
<td>10100011</td>
</tr>
<tr>
<td>10101001</td>
<td>Skip to channel 5</td>
<td>10101011</td>
</tr>
<tr>
<td>10110001</td>
<td>Skip to channel 6</td>
<td>10110011</td>
</tr>
<tr>
<td>10111001</td>
<td>Skip to channel 7</td>
<td>10111011</td>
</tr>
<tr>
<td>11000001</td>
<td>Skip to channel 8</td>
<td>11000011</td>
</tr>
<tr>
<td>11001001</td>
<td>Skip to channel 9</td>
<td>11001011</td>
</tr>
<tr>
<td>11010001</td>
<td>Skip to channel 10</td>
<td>11010011</td>
</tr>
<tr>
<td>11011001</td>
<td>Skip to channel 11</td>
<td>11011011</td>
</tr>
<tr>
<td>11100001</td>
<td>Skip to channel 12</td>
<td>11100011</td>
</tr>
</tbody>
</table>

#### LEAVE|REREAD

The magnetic tape handling options LEAVE and REREAD allow you to specify the action to be taken when the end of a magnetic tape volume is reached, or when a data set on a magnetic tape volume is closed. The LEAVE option prevents the tape from being rewound. The REREAD option rewinds the tape to allow reprocessing of the data set. If you do not specify either of these, the action at end-of-volume or on closing of a data set is controlled by the DISP parameter of the associated DD statement.
If a data set is first read or written forward and then read backward in the same program, specify the LEAVE option to prevent rewinding when the file is closed (or, with a multivolume data set, when volume switching occurs).

The effects of the LEAVE and REREAD options are summarized in Table 19.

**Table 19. Effect of LEAVE and REREAD Options**

<table>
<thead>
<tr>
<th>ENVIRONMENT option</th>
<th>DISP parameter</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>REREAD</td>
<td>—</td>
<td>Positions the current volume to reprocess the data set.</td>
</tr>
<tr>
<td>LEAVE</td>
<td>—</td>
<td>Positions the current volume at the logical end of the data set.</td>
</tr>
<tr>
<td>Neither REREAD nor LEAVE</td>
<td>PASS</td>
<td>Positions the volume at the end of the data set.</td>
</tr>
<tr>
<td></td>
<td>DELETE</td>
<td>Rewinds the current volume.</td>
</tr>
<tr>
<td></td>
<td>KEEP, CATLG, UNCATLG</td>
<td>Rewinds and unloads the current volume.</td>
</tr>
</tbody>
</table>

### Creating a data set with record I/O

When you create a consecutive data set, you must open the associated file for SEQUENTIAL OUTPUT. You can use either the WRITE or LOCATE statement to write records. Table 17 on page 243 shows the statements and options for creating a consecutive data set.

When creating a data set, you must identify it to the operating system in a DD statement. The following paragraphs, summarized in Table 20, tell what essential information you must include in the DD statement and discuss some of the optional information you can supply.

**Table 20. Creating a consecutive data set with record I/O: essential parameters of the DD statement**

<table>
<thead>
<tr>
<th>Storage device</th>
<th>When required</th>
<th>What you must state</th>
<th>Parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td>All</td>
<td>Always</td>
<td>Output device</td>
<td>UNIT= or SYSOUT= or VOLUME=REF=</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Block size¹</td>
<td>DCB=(BLKSIZE=...)</td>
</tr>
<tr>
<td>Direct access only</td>
<td>Always</td>
<td>Storage space required</td>
<td>SPACE=</td>
</tr>
</tbody>
</table>
Table 20. Creating a consecutive data set with record I/O: essential parameters of the DD statement (continued)

<table>
<thead>
<tr>
<th>Storage device</th>
<th>When required</th>
<th>What you must state</th>
<th>Parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td>Direct access</td>
<td>Data set to be used by another job step but not required at end of job</td>
<td>Disposition</td>
<td>DISP=</td>
</tr>
<tr>
<td></td>
<td>Data set to be kept after end of job</td>
<td>Disposition</td>
<td>DISP=</td>
</tr>
<tr>
<td></td>
<td>Data set to be on particular device</td>
<td>Name of data set</td>
<td>DSNAME=</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Volume serial number</td>
<td>VOLUME=SER=</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>or VOLUME=REF=</td>
</tr>
</tbody>
</table>

Notes: ¹Or you could specify the block size in your PL/I program by using the ENVIRONMENT attribute.

Essential information

When you create a consecutive data set you must specify:

- The name of data set to be associated with your PL/I file. A data set with consecutive organization can exist on any type of device.
- The record length. You can specify the record length using the RECSIZE option of the ENVIRONMENT attribute, of the DD_DDNAME environment variable, or of the TITLE option of the OPEN statement.

For files associated with the terminal device (stdout: or stderr:), PL/I uses a default record length of 120 when the RECSIZE option is not specified.

Accessing and updating a data set with record I/O

Once you create a consecutive data set, you can open the file that accesses it for sequential input, for sequential output, or, for data sets on direct access devices, for updating. See Figure 33 on page 250 for an example of a program that accesses and updates a consecutive data set. If you open the file for output, and extend the data set by adding records at the end, you must specify DISP=MOD in the DD statement. If you do not, the data set will be overwritten. If you open a file for updating, you can update only records in their existing sequence, and if you want to insert records, you must create a new data set. Table 17 on page 243 shows the statements and options for accessing and updating a consecutive data set.

When you access a consecutive data set by a SEQUENTIAL UPDATE file, you must retrieve a record with a READ statement before you can update it with a REWRITE statement; however, every record that is retrieved need not be rewritten. A REWRITE statement will always update the last record read.

Consider the following:

```plaintext
READ FILE(F) INTO(A);
  .
  .
  .
READ FILE(F) INTO(B);
  .
  .
  .
REWRITE FILE(F) FROM(A);
```
The REWRITE statement updates the record that was read by the second READ statement. The record that was read by the first statement cannot be rewritten after the second READ statement has been executed.

You cannot update a consecutive data set on magnetic tape except by adding records at the end. To replace or insert records, you must read the data set and write the updated records into a new data set.

You can read a consecutive data set on magnetic tape forward only. Reading backward is not supported.

To access a data set, you must identify it to the operating system in a DD statement. Table 21 summarizes the DD statement parameters needed to access a consecutive data set.

Table 21. Accessing a consecutive data set with record I/O: essential parameters of the DD statement

<table>
<thead>
<tr>
<th>Parameters</th>
<th>What you must state</th>
<th>When required</th>
</tr>
</thead>
<tbody>
<tr>
<td>DSNAME=</td>
<td>Name of data set</td>
<td>Always</td>
</tr>
<tr>
<td>DISP=</td>
<td>Disposition of data set</td>
<td></td>
</tr>
<tr>
<td>UNIT= or VOLUME=REF=</td>
<td>Input device</td>
<td>If data set not cataloged (all devices)</td>
</tr>
<tr>
<td>VOLUME=SER=</td>
<td>Volume serial number</td>
<td>If data set not cataloged (direct access)</td>
</tr>
<tr>
<td>DCB=(BLKSIZE=</td>
<td>Block size</td>
<td>If data set does not have standard labels</td>
</tr>
</tbody>
</table>

Note: 1 Or you could specify the block size in your PL/I program by using the ENVIRONMENT attribute.

The following paragraphs indicate the essential information you must include in the DD statement, and discuss some of the optional information you can supply. The discussions do not apply to data sets in the input stream.

**Essential information**

If the data set is cataloged, you need to supply only the following information in the DD statement:

- The name of the data set (DSNAME parameter). The operating system will locate the information describing the data set in the system catalog, and, if necessary, will request the operator to mount the volume containing it.
- Confirmation that the data set exists (DISP parameter). If you open the data set for output with the intention of extending it by adding records at the end, code DISP=MOD; otherwise, opening the data set for output will result in it being overwritten.

If the data set is not cataloged, you must additionally specify the device that will read the data set and, direct access devices, give the serial number of the volume that contains the data set (UNIT and VOLUME parameters).

**Example of consecutive data sets**

Creating and accessing consecutive data sets are illustrated in the program in Figure 33 on page 250. The program merges the contents of two data sets, in the input stream, and writes them onto a new data set, &&TEMP; each of the original
data sets contains 15-byte fixed-length records arranged in EBCDIC collating sequence. The two input files, INPUT1 and INPUT2, have the default attribute BUFFERED, and locate mode is used to read records from the associated data sets into the respective buffers. Access of based variables in the buffers should not be attempted after the file has been closed.

```
//EXAMPLE JOB
//STEP1 EXEC IBMZCBG
//PLI.SYSIN DD *
%PROCESS INT F(I) AG A(F) OP STG X(F) SOURCE ;
%PROCESS LIST;

MERGE: PROC OPTIONS(MAIN);
  DCL (INPUT1, /* FIRST INPUT FILE */
    INPUT2, /* SECOND INPUT FILE */
    OUT ) FILE RECORD SEQUENTIAL; /* RESULTING MERGED FILE*/
  DCL SYSPRINT FILE PRINT; /* NORMAL PRINT FILE */
  DCL INPUT1_EOF BIT(1) INIT('0'B); /* EOF FLAG FOR INPUT1 */
  DCL INPUT2_EOF BIT(1) INIT('0'B); /* EOF FLAG FOR INPUT2 */
  DCL OUT_EOF BIT(1) INIT('0'B); /* EOF FLAG FOR OUT */
  DCL TRUE BIT(1) INIT('1'B); /* CONSTANT TRUE */
  DCL FALSE BIT(1) INIT('0'B); /* CONSTANT FALSE */
  DCL ITEM1 CHAR(15) BASED(A); /* ITEM FROM INPUT1 */
  DCL ITEM2 CHAR(15) BASED(B); /* ITEM FROM INPUT2 */
  DCL ITEM_LINE CHAR(15); /* ITEM FROM READ INTO */
  DCL A POINTER; /* POINTER VAR */
  DCL B POINTER; /* POINTER VAR */
  ON ENDFILE(INPUT1) INPUT1_EOF = TRUE;
  ON ENDFILE(INPUT2) INPUT2_EOF = TRUE;
  ON ENDFILE(OUT) OUT_EOF = TRUE;

  OPEN FILE(INPUT1) INPUT,
    FILE(INPUT2) INPUT,
    FILE(OUT) OUTPUT;

  READ FILE(INPUT1) SET(A); /* PRIMING READ */
  READ FILE(INPUT2) SET(B);

  DO WHILE (((INPUT1_EOF = FALSE) & (INPUT2_EOF = FALSE))
    IF ITEM1 > ITEM2 THEN
      DO;
        WRITE FILE(OUT) FROM(ITEM2);
        PUT FILE(SYSPRINT) SKIP EDIT('1>2', ITEM1, ITEM2)
          (A(5),A,A);
        READ FILE(INPUT2) SET(B);
      END;
      ELSE
      DO;
        WRITE FILE(OUT) FROM(ITEM1);
        PUT FILE(SYSPRINT) SKIP EDIT('1<2', ITEM1, ITEM2)
          (A(5),A,A);
        READ FILE(INPUT1) SET(A);
      END;
    END;

Figure 33. Merge Sort—creating and accessing a consecutive data set
```
Merge Sort—creating and accessing a consecutive data set (continued)

The program in Figure 34 on page 252 uses record-oriented data transmission to print the table created by the program in Figure 30 on page 234.
Figure 34. Printing record-oriented data transmission
Chapter 11. Defining and using regional data sets

This chapter covers regional data set organization, data transmission statements, and ENVIRONMENT options that define regional data sets. How to create and access regional data sets for each type of regional organization is then discussed.

A data set with regional organization is divided into regions, each of which is identified by a region number, and each of which can contain one record or more than one record, depending on the type of regional organization. The regions are numbered in succession, beginning with zero, and a record can be accessed by specifying its region number, and perhaps a key, in a data transmission statement.

Regional data sets are confined to direct access devices.

Regional organization of a data set allows you to control the physical placement of records in the data set, and to optimize the access time for a particular application. Such optimization is not available with consecutive or indexed organization, in which successive records are written either in strict physical sequence or in logical sequence depending on ascending key values; neither of these methods takes full advantage of the characteristics of direct access storage devices.

You can create a regional data set in a manner similar to a consecutive or indexed data set, presenting records in the order of ascending region numbers; alternatively, you can use direct access, in which you present records in random sequence and insert them directly into preformatted regions. Once you create a regional data set, you can access it by using a file with the attributes SEQUENTIAL or DIRECT as well as INPUT or UPDATE. You do not need to specify either a region number or a key if the data set is associated with a SEQUENTIAL INPUT or SEQUENTIAL UPDATE file. When the file has the DIRECT attribute, you can retrieve, add, delete, and replace records at random.

Records within a regional data set are either actual records containing valid data or dummy records.

The major advantage of regional organization over other types of data set organization is that it allows you to control the relative placement of records; by judicious programming, you can optimize record access in terms of device capabilities and the requirements of particular applications.

Direct access of regional data sets is quicker than that of indexed data sets, but regional data sets have the disadvantage that sequential processing can present records in random sequence; the order of sequential retrieval is not necessarily that in which the records were presented, nor need it be related to the relative key values.

Table 22 on page 254 lists the data transmission statements and options that you can use to create and access a regional data set.
<table>
<thead>
<tr>
<th>File declaration</th>
<th>Valid statements with options you must include</th>
<th>Other options you can also include</th>
</tr>
</thead>
<tbody>
<tr>
<td>SEQUENTIAL OUTPUT</td>
<td>WRITE FILE(file-reference) FROM(reference) KEYFROM(expression); LOCATE based-variable FROM(file-reference) KEYFROM(expression);</td>
<td>SET(pointer-reference)</td>
</tr>
<tr>
<td>SEQUENTIAL INPUT</td>
<td>READ FILE(file-reference) INTO(reference);</td>
<td>KEYTO(reference)</td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) SET(pointer-reference);</td>
<td>KEYTO(reference)</td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) IGNORE(expression);</td>
<td></td>
</tr>
<tr>
<td>SEQUENTIAL UPDATE</td>
<td>READ FILE(file-reference) INTO(reference);</td>
<td>KEYTO(reference)</td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) SET(pointer-reference);</td>
<td>KEYTO(reference)</td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) IGNORE(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>REWRITE FILE(file-reference);</td>
<td>FROM(reference)</td>
</tr>
<tr>
<td>DIRECT OUTPUT</td>
<td>WRITE FILE(file-reference) FROM(reference) KEYFROM(expression);</td>
<td></td>
</tr>
<tr>
<td>DIRECT INPUT</td>
<td>READ FILE(file-reference) INTO(reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td>DIRECT UPDATE</td>
<td>READ FILE(file-reference) INTO(reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>REWRITE FILE(file-reference) FROM(reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>WRITE FILE(file-reference) FROM(reference) KEYFROM(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>DELETE FILE(file-reference) KEY(expression);</td>
<td></td>
</tr>
</tbody>
</table>

Notes:

1. The complete file declaration would include the attributes FILE, RECORD, and ENVIRONMENT; if you use any of the options KEY, KEYFROM, or KEYTO, you must also include the attribute KEYED.
2. The statement READ FILE(file-reference); is equivalent to the statement READ FILE(file-reference) IGNORE(1);.
3. The file must not have the UPDATE attribute when creating new data sets.
Defining REGIONAL(1) data sets using PL/I dynamic allocation

REGIONAL(1) data sets can be defined by using either a DD statement, an environment variable, or the TITLE option of the OPEN statement.

When an environment variable or TITLE option is used, the name must be in uppercase. Specify the MVS data set in the following way:

DSN(data-set-name)

data-set-name must be fully qualified and cannot be a temporary data set; for example, it must not start with &.

You must specify one of the following attributes after the DSN keyword:
OLD
SHR

Defining files for a regional data set

Use a file declaration with the following attributes to define a sequential regional data set:

DCL filename FILE RECORD
  INPUT | OUTPUT | UPDATE
  SEQUENTIAL
  BUFFERED
  [KEYED]
  ENVIRONMENT(options);

Since BUFFERED and UNBUFFERED will be treated the same for REGIONAL(1) data sets, either option can be specified in the ENV option. For example, the FROM option is not required on a REWRITE for a SEQUENTIAL UNBUFFERED file and the LOCATE statement is allowed for OUTPUT SEQUENTIAL data sets even if UNBUFFERED is specified.

To define a direct regional data set, use a file declaration with the following attributes:

DCL filename FILE RECORD
  INPUT | OUTPUT | UPDATE
  DIRECT
  ENVIRONMENT(options);

Default file attributes are shown in Table 14 on page 207. The file attributes are described in the PL/I Language Reference. Options of the ENVIRONMENT attribute are discussed below.

Specifying ENVIRONMENT options

The ENVIRONMENT options applicable to regional data sets are:

REGIONAL((1))
F
RECSIZE(record-length)
BLKSIZE(block-size)
SCALARVARYING
**REGIONAL option**

Use the REGIONAL option to define a file with regional organization.

1 specifies REGIONAL(1)

**REGIONAL(1)**

specifies that the data set contains F-format records that do not have recorded keys. Each region in the data set contains only one record; therefore, each region number corresponds to a relative record within the data set (that is, region numbers start with 0 at the beginning of the data set).

Although REGIONAL(1) data sets have no recorded keys, you can use REGIONAL(1) DIRECT INPUT or UPDATE files to process data sets that do have recorded keys.

**RECSIZE(record-length)**

**BLKSIZE(block-size)**

If both RECSIZE and BLKSIZE are specified, they must specify the same value.

REGIONAL(1) organization is most suited to applications where there are no duplicate region numbers, and where most of the regions will be filled (reducing wasted space in the data set).

**Using keys with REGIONAL data sets**

There are two kinds of keys, recorded keys and source keys. A **recorded key** is a character string that immediately precedes each record in the data set to identify that record; its length cannot exceed 255 characters. A **source key** is the character value of the expression that appears in the KEY or KEYFROM option of a data transmission statement to identify the record to which the statement refers. When you access a record in a regional data set, the source key gives a region number, and can also give a recorded key.

Unlike the keys for indexed data sets, recorded keys in a regional data set are never embedded within the record.

**Using REGIONAL(1) data sets**

In a REGIONAL(1) data set, since there are no recorded keys, the region number serves as the sole identification of a particular record. The character value of the source key should represent an unsigned decimal integer that should not exceed 16777215 (although the actual number of records allowed can be smaller, depending on a combination of record size, device capacity, and limits of your access method. For direct regional(1) files with fixed format records, the maximum number of tracks which can be addressed by relative track addressing is 65,536.) If the region number exceeds this figure, it is treated as modulo 16777216; for instance, 16777226 is treated as 10. Only the characters 0 through 9 and the blank character are valid in the source key; leading blanks are interpreted as zeros. Embedded blanks are not allowed in the number; the first embedded blank, if any, terminates the region number. If more than 8 characters appear in the source key, only the rightmost 8 are used as the region number; if there are fewer than 8 characters, blanks (interpreted as zeros) are inserted on the left.
Dummy Records

Records in a REGIONAL(1) data set are either actual records containing valid data or dummy records. A dummy record in a REGIONAL(1) data set is identified by the constant (8)'1'B in its first byte. Although such dummy records are inserted in the data set either when it is created or when a record is deleted, they are not ignored when the data set is read; your PL/I program must be prepared to recognize them. You can replace dummy records with valid data. Note that if you insert (8)'1'B in the first byte, the record can be lost if you copy the file onto a data set that has dummy records that are not retrieved.

Creating a REGIONAL(1) data set

You can create a REGIONAL(1) data set either sequentially or by direct access. Table 22 on page 254 shows the statements and options for creating a regional data set.

When you use a SEQUENTIAL OUTPUT file to create the data set, the opening of the file causes all tracks on the data set to be cleared, and a capacity record to be written at the beginning of each track to record the amount of space available on that track. You must present records in ascending order of region numbers; any region you omit from the sequence is filled with a dummy record. If there is an error in the sequence, or if you present a duplicate key, the KEY condition is raised. When the file is closed, any space remaining at the end of the current extent is filled with dummy records.

If you use a DIRECT OUTPUT file to create the data set, the whole primary extent allocated to the data set is filled with dummy records when the file is opened. You can present records in random order; if you present a duplicate, the existing record will be overwritten.

For sequential creation, the data set can have up to 15 extents, which can be on more than one volume. For direct creation, the data set can have only one extent, and can therefore reside on only one volume.

Example

Creating a REGIONAL(1) data set is illustrated in Figure 35 on page 258. The data set is a list of telephone numbers with the names of the subscribers to whom they are allocated. The telephone numbers correspond with the region numbers in the data set, the data in each occupied region being a subscriber’s name.
Accessing and updating a REGIONAL(1) data set

Once you create a REGIONAL(1) data set, you can open the file that accesses it for SEQUENTIAL INPUT or UPDATE, or for DIRECT INPUT or UPDATE. You can
open it for OUTPUT only if the existing data set is to be overwritten. Table 22 on page 254 shows the statements and options for accessing a regional data set.

**Sequential access**

To open a SEQUENTIAL file that is used to process a REGIONAL(1) data set, use either the INPUT or UPDATE attribute. You must not include the KEY option in data transmission statements, but the file can have the KEYED attribute, since you can use the KEYTO option. If the target character string referenced in the KEYTO option has more than 8 characters, the value returned (the 8-character region number) is padded on the left with blanks. If the target string has fewer than 8 characters, the value returned is truncated on the left.

Sequential access is in the order of ascending region numbers. All records are retrieved, whether dummy or actual, and you must ensure that your PL/I program recognizes dummy records.

Using sequential input with a REGIONAL(1) data set, you can read all the records in ascending region-number sequence, and in sequential update you can read and rewrite each record in turn.

The rules governing the relationship between READ and REWRITE statements for a SEQUENTIAL UPDATE file that accesses a REGIONAL(1) data set are identical to those for a consecutive data set. Consecutive data sets are discussed in detail in Chapter 8, “Defining and using consecutive data sets,” on page 223.

**Direct access**

To open a DIRECT file that is used to process a REGIONAL(1) data set you can use either the INPUT or the UPDATE attribute. All data transmission statements must include source keys; the DIRECT attribute implies the KEYED attribute.

Use DIRECT UPDATE files to retrieve, add, delete, or replace records in a REGIONAL(1) data set according to the following conventions:

**Retrieval**

All records, whether dummy or actual, are retrieved. Your program must recognize dummy records.

**Addition**

A WRITE statement substitutes a new record for the existing record (actual or dummy) in the region specified by the source key.

**Deletion**

The record you specify by the source key in a DELETE statement is converted to a dummy record.

**Replacement**

The record you specify by the source key in a REWRITE statement, whether dummy or actual, is replaced.

**Example**

Updating a REGIONAL(1) data set is illustrated in Figure 36 on page 260. This program updates the data set and lists its contents. Before each new or updated record is written, the existing record in the region is tested to ensure that it is a dummy; this is necessary because a WRITE statement can overwrite an existing record in a REGIONAL(1) data set even if it is not a dummy. Similarly, during the sequential reading and printing of the contents of the data set, each record is tested and dummy records are not printed.
//EX10 JOB
//STEP2 EXEC IBMZCBG,PARM.PLI='NOP,MAR(1,72)',PARM.BIND='LIST'
//PLI.SYSIN DD *
ACR1: PROC OPTIONS(MAIN);
/* UPDATING A REGIONAL(1) DATA SET - PHONE DIRECTORY */
DCL NOS FILE RECORD KEYED ENV(REGIONAL(1));
DCL SYSIN FILE INPUT RECORD;
DCL (SYSIN_REC,NOS_REC) BIT(1) INIT('1'B);
DCL 1 CARD,
   2 NAME CHAR(20),
   2 (NEWNO,OLDNO) CHAR(2),
   2 CARD_1 CHAR(1),
   2 CODE CHAR(1),
   2 CARD_2 CHAR(54);
DCL IOFIELD CHAR(20);
DCL BYTE CHAR(1) DEF IOFIELD;
ON ENDFILE(SYSIN) SYSIN_REC = '0'B;
OPEN FILE (NOS) DIRECT UPDATE;
READ FILE(SYSIN) INTO(CARD);

DO WHILE(SYSIN_REC);
   SELECT(CODE);
   WHEN('A','C') DO;
      IF CODE = 'C' THEN
         DELETE FILE(NOS) KEY(OLDNO);
         READ FILE(NOS) KEY(NEWNO) INTO(IOFIELD);
         IF UNSPEC(BYTE) = (8)'1'B
            THEN WRITE FILE(NOS) KEYFROM(NEWNO) FROM(NAME);
         ELSE PUT FILE(SYSPRINT) SKIP LIST ('DUPLICATE:',NAME);
      END;
   WHEN('D') DELETE FILE(NOS) KEY(OLDNO);
   OTHERWISE PUT FILE(SYSPRINT) SKIP LIST ('INVALID CODE:',NAME);
END;
READ FILE(SYSIN) INTO(CARD);
END;

CLOSE FILE(SYSIN),FILE(NOS);
PUT FILE(SYSPRINT) PAGE;
OPEN FILE(NOS) SEQUENTIAL INPUT;
ON ENDFILE(NOS) NOS_REC = '0'B;
READ FILE(NOS) INTO(IOFIELD) KEYTO(NEWNO);

DO WHILE(NOS_REC);
   IF UNSPEC(BYTE) ~/ (8)'1'B
      THEN PUT FILE(SYSPRINT) SKIP EDIT (NEWNO,IOFIELD)(A(2),X(3),A);
      PUT FILE(SYSPRINT) SKIP EDIT (IOFIELD) (A);
      READ FILE(NOS) INTO(IOFIELD) KEYTO(NEWNO);
   END;
CLOSE FILE(NOS);
END ACR1;
/*
Figure 36. Updating a REGIONAL(1) data set
*/
Essential information for creating and accessing regional data sets

To create a regional data set, you must give the operating system certain information, either in your PL/I program or in the DD statement that defines the data set. The following paragraphs indicate the essential information, and discuss some of the optional information you can supply.

You must supply the following information when creating a regional data set:

- Device that will write your data set (UNIT or VOLUME parameter of DD statement).
- Block size: You can specify the block size either in your PL/I program (in the BLKSIZE option of the ENVIRONMENT attribute) or in the DD statement (BLKSIZE subparameter). If you do not specify a record length, unblocked records are the default and the record length is determined from the block size. If you do specify a record length, it must be equal to the block size.

If you want to keep a data set (that is, you do not want the operating system to delete it at the end of your job), the DD statement must name the data set and indicate how it is to be disposed of (DSNAME and DISP parameters). The DISP parameter alone will suffice if you want to use the data set in a later step but do not need it after the end of your job.

If you want your data set stored on a particular direct access device, you must indicate the volume serial number in the DD statement (SER or REF subparameter of VOLUME parameter). If you do not supply a serial number for a data set that you want to keep, the operating system allocates one, informs the operator, and prints the number on your program listing. All the essential parameters required in a DD statement for the creation of a regional data set are summarized in Table 23 on page 262 and Table 24 on page 262 lists the DCB subparameters needed. See your z/OS JCL User’s Guide for a description of the DCB subparameters.

You cannot place a regional data set on a system output (SYSOUT) device.

In the DCB parameter, if you specify the DSORG parameter, you must specify the data set organization as direct by coding DSORG=DA. You cannot specify the DUMMY or DSN=NULLFILE parameters in a DD statement for a regional data set. Using DSORG=DA may cause message IEC225I to be issued. This message can be safely ignored.
Table 23. Creating a regional data set: essential parameters of the DD statement

<table>
<thead>
<tr>
<th>Parameters</th>
<th>What you must state</th>
<th>When required</th>
</tr>
</thead>
<tbody>
<tr>
<td>UNIT= or VOLUME=REF=</td>
<td>Output device&lt;sup&gt;1&lt;/sup&gt;</td>
<td>Always</td>
</tr>
<tr>
<td>SPACE=</td>
<td>Storage space required&lt;sup&gt;2&lt;/sup&gt;</td>
<td></td>
</tr>
<tr>
<td>DCB=</td>
<td>Data control block information: see Table 24</td>
<td></td>
</tr>
<tr>
<td>DISP=</td>
<td>Disposition</td>
<td>Data set to be used in another job step but not required in another job</td>
</tr>
<tr>
<td>DISP=</td>
<td>Disposition</td>
<td>Data set to be kept after end of job</td>
</tr>
<tr>
<td>DSNAME=</td>
<td>Name of data set</td>
<td></td>
</tr>
<tr>
<td>VOLUME=SER= or VOLUME=REF=</td>
<td>Volume serial number</td>
<td>Data set to be on particular volume</td>
</tr>
</tbody>
</table>

<sup>1</sup>Regional data sets are confined to direct access devices.

<sup>2</sup>For sequential access, the data set can have up to 15 extents, which can be on more than one volume. For creation with DIRECT access, the data set can have only one extent.

To access a regional data set, you must identify it to the operating system in a DD statement. The following paragraphs indicate the minimum information you must include in the DD statement; this information is summarized in Table 25 on page 263.

If the data set is cataloged, you need to supply only the following information in your DD statement:

- The name of the data set (DSNAME parameter). The operating system locates the information that describes the data set in the system catalog and, if necessary, requests the operator to mount the volume that contains it.
- Confirmation that the data set exists (DISP parameter).

If the data set is not cataloged, you must, in addition, specify the device that will read the data set and give the serial number of the volume that contains the data set (UNIT and VOLUME parameters).

When opening a multiple-volume regional data set for sequential update, the ENDFILE condition is raised at the end of the first volume.

Table 24. DCB subparameters for a regional data set

<table>
<thead>
<tr>
<th>Subparameters</th>
<th>To specify</th>
<th>When required</th>
</tr>
</thead>
<tbody>
<tr>
<td>RECFM=F</td>
<td>Record format&lt;sup&gt;1&lt;/sup&gt;</td>
<td>These are always required</td>
</tr>
<tr>
<td>BLKSIZE=</td>
<td>Block size&lt;sup&gt;1&lt;/sup&gt;</td>
<td></td>
</tr>
<tr>
<td>DSORG=DA</td>
<td>Data set organization</td>
<td></td>
</tr>
</tbody>
</table>

<sup>1</sup>Or you can specify the block size in the ENVIRONMENT attribute.
Table 25. Accessing a regional data set: essential parameters of the DD statement

<table>
<thead>
<tr>
<th>Parameters</th>
<th>What you must state</th>
<th>When required</th>
</tr>
</thead>
<tbody>
<tr>
<td>DSNAME=</td>
<td>Name of data set</td>
<td>Always</td>
</tr>
<tr>
<td>DISP=</td>
<td>Disposition of data set</td>
<td></td>
</tr>
<tr>
<td>UNIT= or VOLUME=REF=</td>
<td>Input device</td>
<td>If data set not cataloged</td>
</tr>
<tr>
<td>VOLUME=SER=</td>
<td>Volume serial number</td>
<td></td>
</tr>
</tbody>
</table>
Chapter 12. Defining and using VSAM data sets

This chapter covers VSAM (the Virtual Storage Access Method) organization for record-oriented data transmission, VSAM ENVIRONMENT options, compatibility with other PL/I data set organizations, and the statements you use to load and access the three types of VSAM data sets that PL/I supports—entry-sequenced, key-sequenced, and relative record. The chapter is concluded by a series of examples showing the PL/I statements, Access Method Services commands, and JCL statements necessary to create and access VSAM data sets.

Enterprise PL/I provides no support for ISAM datasets.

For additional information about the facilities of VSAM, the structure of VSAM data sets and indexes, the way in which they are defined by Access Method Services, and the required JCL statements, see the VSAM publications for your system.

Defining VSAM file using PL/I dynamic allocation

VSAM data sets can be defined by using either a DD statement, an environment variable, or the TITLE option of the OPEN statement.

When an environment variable or TITLE option is used, the name must be in uppercase. Specify the MVS data set as follows:

```
DSN(data-set-name)
```

`data-set-name` must be fully qualified and cannot be a temporary data set; for example, it must not start with &.

You must specify one of the following attributes after the DSN keyword:

- OLD
- SHR

Using VSAM data sets

**How to run a program with VSAM data sets**

Before you execute a program that accesses a VSAM data set, you need to know:

- The name of the VSAM data set
- The name of the PL/I file
- Whether you intend to share the data set with other users

Then you can write the required DD statement to access the data set:

```
//filename DD DSNAME=dsname,DISP=OLD|SHR
```

For example, if your file is named PL1FILE, your data set named VSAMDS, and you want exclusive control of the data set, enter:

```
//PL1FILE DD DSNAME=VSAMDS,DISP=OLD
```

To share your data set, use DISP=SHR.
Enterprise PL/I has no support for ISAM data sets.

To optimize VSAM’s performance by controlling the number of VSAM buffers used for your data set, see the VSAM publications.

**Pairing an Alternate Index Path with a File**

When using an alternate index, you simply specify the name of the *path* in the DSNAME parameter of the DD statement associating the base data set/alternate index pair with your PL/I file. Before using an alternate index, you should be aware of the restrictions on processing; these are summarized in Table 27 on page 271.

Given a PL/I file called PL1FILE and the alternate index path called PERSALPH, the DD statement required would be:

```
//PL1FILE DD DSNAME=PERSALPH,DISP=OLD
```

**VSAM organization**

PL/I provides support for three types of VSAM data sets:

- Key-sequenced data sets (KSDS)
- Entry-sequenced data sets (ESDS)
- Relative record data sets (RRDS).

These correspond roughly to PL/I indexed, consecutive, and regional data set organizations, respectively. They are all ordered, and they can all have keys associated with their records. Both sequential and keyed access are possible with all three types.

Although only key-sequenced data sets have keys as part of their logical records, keyed access is also possible for entry-sequenced data sets (using relative-byte addresses) and relative record data sets (using relative record numbers).

All VSAM data sets are held on direct access storage devices, and a virtual storage operating system is required to use them.

The physical organization of VSAM data sets differs from those used by other access methods. VSAM does not use the concept of blocking, and, except for relative record data sets, records need not be of a fixed length. In data sets with VSAM organization, the data items are arranged in *control intervals*, which are in turn arranged in *control areas*. For processing purposes, the data items within a control interval are arranged in logical records. A control interval can contain one or more logical records, and a logical record can span two or more control intervals. Concern about blocking factors and record length is largely removed by VSAM, although records cannot exceed the maximum specified size. VSAM allows access to the control intervals, but this type of access is not supported by PL/I.

VSAM data sets can have two types of indexes—prime and alternate. A *prime index* is the index to a KSDS that is established when you define a data set; it always exists and can be the only index for a KSDS. You can have one or more *alternate indexes* on a KSDS or an ESDS. Defining an *alternate index* for an ESDS enables you to treat the ESDS, in general, as a KSDS. An alternate index on a KSDS enables a field in the logical record different from that in the prime index to be used as the key field. Alternate indexes can be either *nonunique*, in which duplicate keys are allowed, or *unique*, in which they are not. The prime index can never have duplicate keys.
Any change in a data set that has alternate indexes must be reflected in all the indexes if they are to remain useful. This activity is known as index upgrade, and is done by VSAM for any index in the index upgrade set of the data set. (For a KSDS, the prime index is always a member of the index upgrade set.) However, you must avoid making changes in the data set that would cause duplicate keys in the prime index or in a unique alternate index.

Before using a VSAM data set for the first time, you need to define it to the system with the DEFINE command of Access Method Services, which you can use to completely define the type, structure, and required space of the data set. This command also defines the data set’s indexes (together with their key lengths and locations) and the index upgrade set if the data set is a KSDS or has one or more alternate indexes. A VSAM data set is thus “created” by Access Method Services.

The operation of writing the initial data into a newly created VSAM data set is referred to as loading in this publication.

Use the three different types of data sets according to the following purposes:

- Use entry-sequenced data sets for data that you primarily access in the order in which it was created (or the reverse order).
- Use key-sequenced data sets when you normally access records through keys within the records (for example, a stock-control file where the part number is used to access a record).
- Use relative record data sets for data in which each item has a particular number, and you normally access the relevant record by that number (for example, a telephone system with a record associated with each number).

You can access records in all types of VSAM data sets either directly by means of a key, or sequentially (backward or forward). You can also use a combination of the two ways: Select a starting point with a key and then read forward or backward from that point.

You can create alternate indexes for key-sequenced and entry-sequenced data sets. You can then access your data in many sequences or by one of many keys. For example, you could take a data set held or indexed in order of employee number and index it by name in an alternate index. Then you could access it in alphabetic order, in reverse alphabetic order, or directly using the name as a key. You could also access it in the same kind of combinations by employee number.

Table 26 on page 268 shows how the same data could be held in the three different types of VSAM data sets and illustrates their respective advantages and disadvantages.
### Table 26. Types and advantages of VSAM data sets

<table>
<thead>
<tr>
<th>Data set type</th>
<th>Method of loading</th>
<th>Method of reading</th>
<th>Method of updating</th>
<th>Pros and cons</th>
</tr>
</thead>
<tbody>
<tr>
<td>Key-Sequenced</td>
<td>Sequentially in order or prime index which must be unique</td>
<td>KEYED by specifying key of record in prime index</td>
<td>KEYED specifying a unique key in any index</td>
<td><strong>Advantages:</strong> Complete access and updating</td>
</tr>
<tr>
<td></td>
<td></td>
<td>SEQUENTIAL backward or forward in order of any index</td>
<td>SEQUENTIAL following positioning by unique key</td>
<td><strong>Disadvantages:</strong> Records must be in order of prime index before loading</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Positioning by key followed by sequential reading either backward or forward</td>
<td>Record deletion allowed</td>
<td><strong>Uses:</strong> For uses where access will be related to key</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Record insertion allowed</td>
<td></td>
</tr>
<tr>
<td>Entry-Sequenced</td>
<td>Sequentially (forward only)</td>
<td>SEQUENTIAL backward or forward</td>
<td>New records at end only</td>
<td><strong>Advantages:</strong> Simple fast creation</td>
</tr>
<tr>
<td></td>
<td>The RBA of each record can be obtained and used as a key</td>
<td>KEYED using RBA</td>
<td>Existing records cannot have length changed</td>
<td><strong>Disadvantages:</strong> No requirement for a unique index</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Positioning by key followed by sequential either backward or forward</td>
<td>Record deletion not allowed</td>
<td><strong>Uses:</strong> For uses where data will primarily be accessed sequentially</td>
</tr>
<tr>
<td>Relative Record</td>
<td>Sequentially starting from slot 1</td>
<td>KEYED specifying numbers as key</td>
<td>Sequentially starting at a specified slot and continuing with next slot</td>
<td><strong>Advantages:</strong> Speedy access to record by number</td>
</tr>
<tr>
<td></td>
<td>KEYED specifying number of slot</td>
<td>Sequential forward or backward omitting empty records</td>
<td>Keyed specifying numbers as key</td>
<td><strong>Disadvantages:</strong> Structure tied to numbering sequences</td>
</tr>
<tr>
<td></td>
<td>Positioning by key followed by sequential writes</td>
<td></td>
<td>Record deletion allowed</td>
<td>Fixed length records</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Record insertion into empty slots allowed</td>
<td><strong>Uses:</strong> For use where records will be accessed by number</td>
</tr>
</tbody>
</table>

### Keys for VSAM data sets

All VSAM data sets can have keys associated with their records. For key-sequenced data sets, and for entry-sequenced data sets accessed via an alternate index, the key is a defined field within the logical record. For entry-sequenced data sets, the key is the relative byte address (RBA) of the record. For relative-record data sets, the key is a relative record number.

### Keys for indexed VSAM data sets

Keys for key-sequenced data sets and for entry-sequenced data sets accessed via an alternate index are part of the logical records recorded on the data set. You define the length and location of the keys when you create the data set.
The ways you can reference the keys in the KEY, KEYFROM, and KEYTO options are as described under “KEY(expression) Option,” “KEYFROM(expression) Option,” and “KEYTO(reference) Option” in Chapter 12 of the PL/I Language Reference.

**Relative byte addresses (RBA)**

Relative byte addresses allow you to use keyed access on an ESDS associated with a KEYED SEQUENTIAL file. The RBAs, or keys, are character strings of length 4, and their values are defined by VSAM. You cannot construct or manipulate RBAs in PL/I; you can, however, compare their values in order to determine the relative positions of records within the data set. RBAs are not normally printable.

You can obtain the RBA for a record by using the KEYTO option, either on a WRITE statement when you are loading or extending the data set, or on a READ statement when the data set is being read. You can subsequently use an RBA obtained in either of these ways in the KEY option of a READ or REWRITE statement.

Do not use an RBA in the KEYFROM option of a WRITE statement.

VSAM allows use of the relative byte address as a key to a KSDS, but this use is not supported by PL/I.

**Relative record numbers**

Records in an RRDS are identified by a relative record number that starts at 1 and is incremented by 1 for each succeeding record. You can use these relative record numbers as keys for keyed access to the data set.

Keys used as relative record numbers are character strings of length 8. The character value of a source key you use in the KEY or KEYFROM option must represent an unsigned integer. If the source key is not 8 characters long, it is truncated or padded with blanks (interpreted as zeros) on the left. The value returned by the KEYTO option is a character string of length 8, with leading zeros suppressed.

**Choosing a data set type**

When planning your program, the first decision to be made is which type of data set to use. There are three types of VSAM data sets and five types of non-VSAM data sets available to you. VSAM data sets can provide all the function of the other types of data sets, plus additional function available only in VSAM. VSAM can usually match other data set types in performance, and often improve upon it. However, VSAM is more subject to performance degradation through misuse of function.

The comparison of all eight types of data sets given in Table 15 on page 215 is helpful; however, many factors in the choice of data set type for a large installation are beyond the scope of this book.

When choosing between the VSAM data set types, you should base your choice on the most common sequence in which you will require your data. The following is a suggested procedure that you can use to help ensure a combination of data sets and indexes that provide the function you require.

1. Determine the type of data and how it will be accessed.
   a. Primarily sequentially — favors ESDS.
   b. Primarily by key — favors KSDS.
c. Primarily by number — favors RRDS.

2. Determine how you will load the data set. Note that you must load a KSDS in key sequence; thus an ESDS with an alternate index path can be a more practical alternative for some applications.

3. Determine whether you require access through an alternate index path. These are only supported on KSDS and ESDS. If you require an alternate index path, determine whether the alternate index will have unique or nonunique keys. Use of nonunique keys can limit key processing. However, it might also be impractical to assume that you will use unique keys for all future records; if you attempt to insert a record with a nonunique key in an index that you have created for unique keys, it will cause an error.

4. When you have determined the data sets and paths that you require, ensure that the operations you have in mind are supported. Figure 37 might be helpful.

Do not try to access a dummy VSAM data set, because you will receive an error message indicating that you have an undefined file.

Table 28 on page 278, Table 29 on page 279, and Table 30 on page 293 show the statements allowed for entry-sequenced data sets, indexed data sets, and relative record data sets, respectively.

<table>
<thead>
<tr>
<th></th>
<th>SEQUENTIAL</th>
<th>KEYED SEQUENTIAL</th>
<th>DIRECT</th>
</tr>
</thead>
<tbody>
<tr>
<td>INPUT</td>
<td>ESDS</td>
<td>ESDS</td>
<td>KSDS</td>
</tr>
<tr>
<td></td>
<td>KSDS</td>
<td>KSDS</td>
<td>RRDS</td>
</tr>
<tr>
<td></td>
<td>RRDS</td>
<td>RRDS</td>
<td>Path(U)</td>
</tr>
<tr>
<td></td>
<td>Path(N)</td>
<td>Path(N)</td>
<td></td>
</tr>
<tr>
<td></td>
<td>Path(U)</td>
<td>Path(U)</td>
<td></td>
</tr>
<tr>
<td>OUTPUT</td>
<td>ESDS</td>
<td>ESDS</td>
<td>KSDS</td>
</tr>
<tr>
<td></td>
<td>RRDS</td>
<td>RRDS</td>
<td>Path(U)</td>
</tr>
<tr>
<td>UPDATE</td>
<td>ESDS</td>
<td>ESDS</td>
<td>KSDS</td>
</tr>
<tr>
<td></td>
<td>KSDS</td>
<td>KSDS</td>
<td>RRDS</td>
</tr>
<tr>
<td></td>
<td>RRDS</td>
<td>RRDS</td>
<td>Path(U)</td>
</tr>
<tr>
<td></td>
<td>Path(N)</td>
<td>Path(N)</td>
<td></td>
</tr>
<tr>
<td></td>
<td>Path(U)</td>
<td>Path(U)</td>
<td></td>
</tr>
</tbody>
</table>

Key:  
ESDS  Entry-sequenced data set  
KSDS  Key-sequenced data set  
RRDS  Relative record data set  
Path(N) Alternate index path with nonunique keys  
Path(U) Alternate index path with unique keys

You can combine the attributes on the left with those at the top of the figure for the data sets and paths shown. For example, only an ESDS and an RRDS can be SEQUENTIAL OUTPUT.

PL/I does not support dummy VSAM data sets.

Figure 37. VSAM data sets and allowed file attributes
### Defining files for VSAM data sets

You define a sequential VSAM data set by using a file declaration with the following attributes:

```plaintext
DCL filename FILE RECORD
    INPUT | OUTPUT | UPDATE
    SEQUENTIAL
    BUFFERED
    [KEYED]
    ENVIRONMENT(options);
```

You define a direct VSAM data set by using a file declaration with the following attributes:

```plaintext
DCL filename FILE RECORD
    INPUT | OUTPUT | UPDATE
    DIRECT
    [KEYED]
    ENVIRONMENT(options);
```

Table 14 on page 207 shows the default attributes. The file attributes are described in the *PL/I Language Reference*. Options of the ENVIRONMENT attribute are discussed below.

Some combinations of the file attributes INPUT or OUTPUT or UPDATE and DIRECT or SEQUENTIAL or KEYED SEQUENTIAL are allowed only for certain types of VSAM data sets. Figure 37 on page 270 shows the compatible combinations.

### Specifying ENVIRONMENT options

Many of the options of the ENVIRONMENT attribute affecting data set structure are not needed for VSAM data sets. If you specify them, they are either ignored or are used for checking purposes. If those that are checked conflict with the values defined for the data set, the UNDEFINEDFILE condition is raised when an attempt is made to open the file.

The ENVIRONMENT options applicable to VSAM data sets are:

- `BKWD`
- `BUFND (n)`
- `BUFNI (n)`
- `BUFSP (n)`

---

**Table 27. Processing Allowed on Alternate Index Paths**

<table>
<thead>
<tr>
<th>Base Cluster Type</th>
<th>Alternate Index Key Type</th>
<th>Processing</th>
<th>Restrictions</th>
</tr>
</thead>
<tbody>
<tr>
<td>KSDS</td>
<td>Unique key</td>
<td>As normal KSDS</td>
<td>May not modify key of access.</td>
</tr>
<tr>
<td></td>
<td>Nonunique key</td>
<td>Limited keyed access</td>
<td>May not modify key of access.</td>
</tr>
<tr>
<td>ESDS</td>
<td>Unique key</td>
<td>As KSDS</td>
<td>No deletion.</td>
</tr>
<tr>
<td></td>
<td>Nonunique key</td>
<td>Limited keyed access</td>
<td>May not modify key of access.</td>
</tr>
</tbody>
</table>

---

Chapter 12. Defining and using VSAM data sets  271
GENKEY and SCALARVARYING options have the same effect as they do when you use them for non-VSAM data sets. Note that under VSAM RLS, options BUFND, BUFNI and BUFSP are ignored.

The options that are checked for a VSAM data set are RECSIZE and, for a key-sequenced data set, KEYLENGTH and KEYLOC. Table 14 on page 207 shows which options are ignored for VSAM. Table 14 on page 207 also shows the required and default options.

For VSAM data sets, you specify the maximum and average lengths of the records to the Access Method Services utility when you define the data set. If you include the RECSIZE option in the file declaration for checking purposes, specify the maximum record size. If you specify RECSIZE and it conflicts with the values defined for the data set, the UNDEFINEDFILE condition is raised.

**BKWD option**

Use the BKWD option to specify backward processing for a SEQUENTIAL INPUT or SEQUENTIAL UPDATE file associated with a VSAM data set.

Sequential reads (that is, reads without the KEY option) retrieve the previous record in sequence. For indexed data sets, the previous record is, in general, the record with the next lower key. However, if you are accessing the data set via a nonunique alternate index, records with the same key are recovered in their normal sequence. For example, if the records are:

A B C1 C2 C3 D E

where C1, C2, and C3 have the same key, they are recovered in the sequence:

E D C1 C2 C3 B A

When a file with the BKWD option is opened, the data set is positioned at the last record. ENDFILE is raised in the normal way when the start of the data set is reached.

Do not specify the BKWD option with either the REUSE option or the GENKEY option. Also, the WRITE statement is not allowed for files declared with the BKWD option.

**BUFND option**

Use the BUFND option to specify the number of data buffers required for a VSAM data set.

n specifies an integer, or a variable with attributes FIXED BINARY(31) STATIC.
Multiple data buffers help performance when the file has the SEQUENTIAL attribute and you are processing long group of contiguous records sequentially.

**BUFNI option**
Use the BUFNI option to specify the number of index buffers required for a VSAM key-sequence data set.

\[ \text{BUFNI} - (n) \]

- \( n \) specifies an integer, or a variable with attributes FIXED BINARY(31) STATIC.

Multiple index buffers help performance when the file has the KEYED attribute. Specify at least as many index buffers as there are levels in the index.

**BUFSP option**
Use the BUFSP option to specify, in bytes, the total buffer space required for a VSAM data set (for both the data and index components).

\[ \text{BUFSP} - (n) \]

- \( n \) specifies an integer, or a variable with attributes FIXED BINARY(31) STATIC.

It is usually preferable to specify BUFNI and BUFND options rather than BUFSP.

**GENKEY option**
For the description of this option, see “GENKEY option — key classification” on page 212.

**PASSWORD option**
When you define a VSAM data set to the system (using the DEFINE command of Access Method Services), you can associate READ and UPDATE passwords with it. From that point on, you must include the appropriate password in the declaration of any PL/I file that you use to access the data set.

\[ \text{PASSWORD} - (\text{password-specification}) \]

- **password-specification** is a character constant or character variable that specifies the password for the type of access your program requires. If you specify a constant, it must not contain a repetition factor; if you specify a variable, it must be level-1, element, static, and unsubscripted.

The character string is padded or truncated to 8 characters and passed to VSAM for inspection. If the password is incorrect, the system operator is given a number of chances to specify the correct password. You specify the number of chances to be allowed when you define the data set. After this number of unsuccessful tries, the UNDEFINEDFILE condition is raised.

**REUSE option**
Use the REUSE option to specify that an OUTPUT file associated with a VSAM data set is to be used as a work file.
The data set is treated as an empty data set each time the file is opened. Any secondary allocations for the data set are released, and the data set is treated exactly as if it were being opened for the first time.

Do not associate a file that has the REUSE option with a data set that has alternate indexes or the BKWD option, and do not open it for INPUT or UPDATE.

The REUSE option takes effect only if you specify REUSE in the Access Method Services DEFINE CLUSTER command.

**SKIP option**

Use the SKIP option of the ENVIRONMENT attribute to specify that the VSAM OPTCD "SKP" is to be used whenever possible. It is applicable to key-sequenced data sets that you access by means of a KEYED SEQUENTIAL INPUT or UPDATE file.

You should specify this option for the file if your program accesses individual records scattered throughout the data set, but does so primarily in ascending key order.

Omit this option if your program reads large numbers of records sequentially without the use of the KEY option, or if it inserts large numbers of records at specific points in the data set (mass sequential insert).

It is never an error to specify (or omit) the SKIP option; its effect on performance is significant only in the circumstances described.

**VSAM option**

You must specify the VSAM option for VSAM data sets.

**Performance options**

You can specify the buffer options in the AMP parameter of the DD statement; they are explained in your Access Method Services manual.

**Defining Files for Alternate Index Paths**

VSAM allows you to define alternate indexes on key sequenced and entry sequenced data sets. This enables you to access key sequenced data sets in a number of ways other than from the prime index. This also allows you to index and access entry sequenced data sets by key or sequentially in order of the keys. Consequently, data created in one form can be accessed in a large number of different ways. For example, an employee file might be indexed by personnel number, by name, and also by department number.

When an alternate index has been built, you actually access the data set through a third object known as an alternate index path that acts as a connection between the alternate index and the data set.
Two types of alternate indexes are allowed—unique key and nonunique key. For a unique key alternate index, each record must have a different alternate key. For a nonunique key alternate index, any number of records can have the same alternate key. In the example suggested above, the alternate index using the names could be a unique key alternate index (provided each person had a different name). The alternate index using the department number would be a nonunique key alternate index because more than one person would be in each department.

In most respects, you can treat a data set accessed through a unique key alternate index path like a KSDS accessed through its prime index. You can access the records by key or sequentially, you can update records, and you can add new records. If the data set is a KSDS, you can delete records, and alter the length of updated records. Restrictions and allowed processing are shown in Table 27 on page 271. When you add or delete records, all indexes associated with the data set are by default altered to reflect the new situation.

In data sets accessed through a nonunique key alternate index path, the record accessed is determined by the key and the sequence. The key can be used to establish positioning so that sequential access can follow. The use of the key accesses the first record with that key. When the data set is read backwards, only the order of the keys is reversed. The order of the records with the same key remains the same whichever way the data set is read.

### Defining VSAM data sets

Use the DEFINE CLUSTER command of Access Method Services to define and catalog VSAM data sets. To use the DEFINE command, you need to know:

- The name and password of the master catalog if the master catalog is password protected
- The name and password of the VSAM private catalog you are using if you are not using the master catalog
- Whether VSAM space for your data set is available
- The type of VSAM data set you are going to create
- The volume on which your data set is to be placed
- The average and maximum record size in your data set
- The position and length of the key for an indexed data set
- The space to be allocated for your data set
- How to code the DEFINE command
- How to use the Access Method Services program.

When you have the information, you are in a position to code the DEFINE command and then define and catalog the data set using Access Method Services.
Entry-sequenced data sets

The statements and options allowed for files associated with an ESDS are shown in Table 28.

Table 28. Statements and options allowed for loading and accessing VSAM entry-sequenced data sets

<table>
<thead>
<tr>
<th>File declaration</th>
<th>Valid statements, with options you must include</th>
<th>Other options you can also include</th>
</tr>
</thead>
<tbody>
<tr>
<td>SEQUENTIAL OUTPUT BUFFERED</td>
<td>WRITE FILE(file-reference) FROM(reference); LOCATE based-variable FILE(file-reference);</td>
<td>KEYTO(reference) SET(pointer-reference)</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
| SEQUENTIAL INPUT BUFFERED | READ FILE(file-reference) INTO(reference); READ FILE(file-reference) SET(pointer-reference); READ FILE(file-reference); | KEYTO(reference) or KEY(expression)
| | | |
| SEQUENTIAL UPDATE BUFFERED | READ FILE(file-reference) INTO(reference); READ FILE(file-reference) SET(pointer-reference); READ FILE(file-reference); WRITE FILE(file-reference) FROM(reference); REWRITE FILE(file-reference); | KEYTO(reference) FROM(reference) and/or KEY(expression)

Notes:
1. The complete file declaration would include the attributes FILE, RECORD, and ENVIRONMENT; if you use either of the options KEY or KEYTO, it must also include the attribute KEYED.
2. The statement “READ FILE(file-reference);” is equivalent to the statement “READ FILE(file-reference) IGNORE (1);”.
3. The expression used in the KEY option must be a relative byte address, previously obtained by means of the KEYTO option.

Loading an ESDS

When an ESDS is being loaded, the associated file must be opened for SEQUENTIAL OUTPUT. The records are retained in the order in which they are presented.

You can use the KEYTO option to obtain the relative byte address of each record as it is written. You can subsequently use these keys to achieve keyed access to the data set.
Using a SEQUENTIAL file to access an ESDS

You can open a SEQUENTIAL file that is used to access an ESDS with either the INPUT or the UPDATE attribute. If you use either of the options KEY or KEYTO, the file must also have the KEYED attribute.

Sequential access is in the order that the records were originally loaded into the data set. You can use the KEYTO option on the READ statements to recover the RBAs of the records that are read. If you use the KEY option, the record that is recovered is the one with the RBA you specify. Subsequent sequential access continues from the new position in the data set.

For an UPDATE file, the WRITE statement adds a new record at the end of the data set. With a REWRITE statement, the record rewritten is the one with the specified RBA if you use the KEY option; otherwise, it is the record accessed on the previous READ. You must not attempt to change the length of the record that is being replaced with a REWRITE statement.

The DELETE statement is not allowed for entry-sequenced data sets.

Defining and loading an ESDS

In Figure 38 on page 278 the data set is defined with the DEFINE CLUSTER command and given the name PLIVSAM.AJC1.BASE. The NONINDEXED keyword causes an ESDS to be defined.

The PL/I program writes the data set using a SEQUENTIAL OUTPUT file and a WRITE FROM statement. The DD statement for the file contains the DSNAME of the data set given in the NAME parameter of the DEFINE CLUSTER command.

The RBA of the records could have been obtained during the writing for subsequent use as keys in a KEYED file. To do this, a suitable variable would have to be declared to hold the key and the WRITE...KEYTO statement used. For example:

DCL CHARS CHAR(4);
WRITE FILE(FAMFILE) FROM (STRING)
  KEYTO(CHARS);

Note that the keys would not normally be printable, but could be retained for subsequent use.

The cataloged procedure IBMZCBG is used. Because the same program (in Figure 38 on page 278) can be used for adding records to the data set, it is retained in a library. This procedure is shown in the next example.
Updating an ESDS

Figure 39 on page 279 shows the addition of a new record on the end of an ESDS. This is done by executing again the program shown in Figure 38. A SEQUENTIAL OUTPUT file is used and the data set associated with it by use of the DSNAME parameter specifying the name PLIVSAM.AJC1.BASE specified in the DEFINE command shown in Figure 38.

```plaintext
//OPT9#7 JOB
//STEP1 EXEC PGM=IDCAMS,REGION=512K
//SYSPRINT DD SYSOUT=A
//SYSSIN DD *
DEFINE CLUSTER -
  (NAME(PLIVSAM.AJC1.BASE) -
  VOLUMES(nnnnnn) -
  NONINDEXED -
  RECORDSIZE(80 80) -
  TRACKS(2 2))
/*
//STEP2 EXEC IBMZCLG
//PLI.SYSIN DD *
CREATE: PROC OPTIONS(MAIN);
DCL
  FAMFILE FILE SEQUENTIAL OUTPUT ENV(VSAM),
  IN FILE RECORD INPUT,
  STRING CHAR(80),
  EOF BIT(1) INIT('0'B);
ON ENDFILE(IN) EOF='1'B;
READ FILE(IN) INTO (STRING);
DO I=1 BY 1 WHILE (-EOF);
  PUT FILE(SYSPRINT) SKIP EDIT (STRING) (A);
  WRITE FILE(FAMFILE) FROM (STRING);
  READ FILE(IN) INTO (STRING);
END;
  PUT SKIP EDIT(I-1,' RECORDS PROCESSED')(A);
END;
/*
//LKED.SYSLMOD DD DSN=HPU8.MYDS(PGMA),DISP=(NEW,CATLG),
// UNIT=SYSDA,SPACE=(CYL,(1,1,1))
//GO.FAMFILE DD DSNAME=PLIVSAM.AJC1.BASE,DISP=OLD
//GO.IN DD *
FRED 69 M
ANDY 70 M
SUZAN 72 F
/*
```

Figure 38. Defining and loading an entry-sequenced data set (ESDS)

**Updating an ESDS**

Figure 39 on page 279 shows the addition of a new record on the end of an ESDS. This is done by executing again the program shown in Figure 38. A SEQUENTIAL OUTPUT file is used and the data set associated with it by use of the DSNAME parameter specifying the name PLIVSAM.AJC1.BASE specified in the DEFINE command shown in Figure 38.
You can rewrite existing records in an ESDS, provided that the length of the record is not changed. You can use a SEQUENTIAL or KEYED SEQUENTIAL update file to do this. If you use keys, they can be the RBAs or keys of an alternate index path.

Delete is not allowed for ESDS.

**Key-sequenced and indexed entry-sequenced data sets**

The statements and options allowed for indexed VSAM data sets are shown in Table 29. An indexed data set can be a KSDS with its prime index, or either a KSDS or an ESDS with an alternate index. Except where otherwise stated, the following description applies to all indexed VSAM data sets.

**Table 29. Statements and options allowed for loading and accessing VSAM indexed data sets.**

<table>
<thead>
<tr>
<th>File declaration</th>
<th>Valid statements, with options you must include</th>
<th>Other options you can also include</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>SEQUENTIAL OUTPUT BUFFERED</strong></td>
<td>WRITE FILE(file-reference) FROM(reference) KEYFROM(expression);</td>
<td>LOCATE based-variable FILE(file-reference) KEYFROM(expression); SET(pointer-reference)</td>
</tr>
<tr>
<td></td>
<td>LOCATE based-variable FILE(file-reference) KEYFROM(expression);</td>
<td></td>
</tr>
<tr>
<td><strong>SEQUENTIAL INPUT BUFFERED</strong></td>
<td>READ FILE(file-reference) INTO(reference);</td>
<td>KEY(expression) or KEYTO(expression)</td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) SET(pointer-reference);</td>
<td>KEY(expression) or KEYTO(expression)</td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference)</td>
<td>IGNORE(expression)</td>
</tr>
</tbody>
</table>

---

Figure 39. Updating an ESDS

You can rewrite existing records in an ESDS, provided that the length of the record is not changed. You can use a SEQUENTIAL or KEYED SEQUENTIAL update file to do this. If you use keys, they can be the RBAs or keys of an alternate index path.

Delete is not allowed for ESDS.
<table>
<thead>
<tr>
<th>File declaration</th>
<th>Valid statements, with options you must include</th>
<th>Other options you can also include</th>
</tr>
</thead>
<tbody>
<tr>
<td>SEQUENTIAL UPDATE BUFFERED</td>
<td>READ FILE(file-reference) INTO(reference);</td>
<td>KEY(expression) or KEYTO(reference)</td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) SET(pointer-reference);</td>
<td>KEY(expression) or KEYTO(reference)</td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) [2]</td>
<td>IGNORE(expression)</td>
</tr>
<tr>
<td></td>
<td>WRITE FILE(file-reference) FROM(reference) KEYFROM(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>REWRITE FILE(file-reference);</td>
<td>FROM(reference) and/or KEY(expression)</td>
</tr>
<tr>
<td></td>
<td>DELETE FILE(file-reference)</td>
<td>KEY(expression)</td>
</tr>
<tr>
<td>DIRECT BUFFERED</td>
<td>READ FILE(file-reference) INTO(reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) SET(pointer-reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td>DIRECT OUTPUT BUFFERED</td>
<td>WRITE FILE(file-reference) FROM(reference) KEYFROM(expression);</td>
<td></td>
</tr>
<tr>
<td>DIRECT BUFFERED</td>
<td>READ FILE(file-reference) INTO(reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) SET(pointer-reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>REWRITE FILE(file-reference) FROM(reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>DELETE FILE(file-reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>WRITE FILE(file-reference) FROM(reference) KEYFROM(expression);</td>
<td></td>
</tr>
</tbody>
</table>
Table 29. Statements and options allowed for loading and accessing VSAM indexed data sets (continued).

<table>
<thead>
<tr>
<th>File declaration</th>
<th>Valid statements, with options you must include</th>
<th>Other options you can also include</th>
</tr>
</thead>
</table>

Notes:

1. The complete file declaration would include the attributes FILE and RECORD. If you use any of the options KEY, KEYFROM, or KEYTO, you must also include the attribute KEYED in the declaration.
2. The statement READ FILE(file-reference); is equivalent to the statement READ FILE(file-reference) IGNORE(1);
3. Do not associate a SEQUENTIAL OUTPUT file with a data set accessed via an alternate index.
4. Do not associate a DIRECT file with a data set accessed via a nonunique alternate index.
5. DELETE statements are not allowed for a file associated with an ESDS accessed via an alternate index.

Loading a KSDS or indexed ESDS: When a KSDS is being loaded, you must open the associated file for KEYED SEQUENTIAL OUTPUT. You must present the records in ascending key order, and you must use the KEYFROM option. Note that you must use the prime index for loading the data set; you cannot load a VSAM data set via an alternate index.

If a KSDS already contains some records, and you open the associated file with the SEQUENTIAL and OUTPUT attributes, you can add only records at the end of the data set. The rules given in the previous paragraph apply; in particular, the first record you present must have a key greater than the highest key present on the data set.

Figure 40 on page 282 shows the DEFINE command used to define a KSDS. The data set is given the name PLIVSAM.AJC2.BASE and defined as a KSDS because of the use of the INDEXED operand. The position of the keys within the record is defined in the KEYS operand.

Within the PL/I program, a KEYED SEQUENTIAL OUTPUT file is used with a WRITE...FROM...KEYFROM statement. The data is presented in ascending key order. A KSDS must be loaded in this manner.

The file is associated with the data set by a DD statement which uses the name given in the DEFINE command as the DSNNAME parameter.
DEFINE CLUSTER -
(NAME(PLIVSAM.AJC2.BASE) -
VOLUMES(nnnnnn) -
INDEXED -
TRACKS(3 1) -
KEYS(20 0) -
RECORDSIZE(23 80))
/*
EXEC IBMZCBG
//PLI.SYSIN DD *
TELNOS: PROC OPTIONS(MAIN);

DCL DIREC FILE RECORD SEQUENTIAL OUTPUT KEYED ENV(VSAM),
   CARD CHAR(80),
   NAME CHAR(20) DEF CARD POS(1),
   NUMBER CHAR(3) DEF CARD POS(21),
   OUTREC CHAR(23) DEF CARD POS(1),
   EOF BIT(1) INIT('0'B);
ON ENDFILE(SYSIN) EOF='1'B;
OPEN FILE(DIREC) OUTPUT;
GET FILE(SYSIN) EDIT(CARD)(A(80));
DO WHILE (¬EOF);
   WRITE FILE(DIREC) FROM(OUTREC) KEYFROM(NAME);
   GET FILE(SYSIN) EDIT(CARD)(A(80));
END;
CLOSE FILE(DIREC);
END TELNOS;
/*
GO.DIREC DD DSNAME=PLIVSAM.AJC2.BASE,DISP=OLD
GO.SYSIN DD *
ACTION,G. 162
BAKER,R. 152
BRAMLEY,O.H. 248
CHEESEMAN,D. 141
CORY,G. 336
ELLIOTT,D. 875
FIGGINS,S. 413
HARVEY,C.D.W. 205
HASTINGS,G.M. 391
KENDALL,J.G. 294
LANCASTER,W.R. 624
MILES,R. 233
NEWMAN,M.W. 450
PITT,W.H. 515
ROLF,D.E. 114
SHEERS,C.D. 241
SUTCLIFFE,M. 472
TAYLOR,G.C. 407
WILTON,L.W. 404
WINSTONE,E.M. 307
//
Using a SEQUENTIAL file to access a KSDS or indexed ESDS: You can open a
SEQUENTIAL file that is used to access a KSDS with either the INPUT or the
UPDATE attribute.

For READ statements without the KEY option, the records are recovered in
ascending key order (or in descending key order if the BKWD option is used). You
can obtain the key of a record recovered in this way by means of the KEYTO
option.

If you use the KEY option, the record recovered by a READ statement is the one
with the specified key. Such a READ statement positions the data set at the
specified record; subsequent sequential reads will recover the following records in
sequence.

WRITE statements with the KEYFROM option are allowed for KEYED
SEQUENTIAL UPDATE files. You can make insertions anywhere in the data set,
without respect to the position of any previous access. If you are accessing the data
set via a unique index, the KEY condition is raised if an attempt is made to insert a
record with the same key as a record that already exists on the data set. For a
nonunique index, subsequent retrieval of records with the same key is in the order
that they were added to the data set.

REWRITE statements with or without the KEY option are allowed for UPDATE
files. If you use the KEY option, the record that is rewritten is the first record with
the specified key; otherwise, it is the record that was accessed by the previous
READ statement. When you rewrite a record using an alternate index, do not
change the prime key of the record.

Using a DIRECT file to access a KSDS or indexed ESDS: You can open a
DIRECT file that is used to access an indexed VSAM data set with the INPUT,
OUTPUT, or UPDATE attribute. Do not use a DIRECT file to access the data set via
a nonunique index.

If you use a DIRECT OUTPUT file to add records to the data set, and if an attempt
is made to insert a record with the same key as a record that already exists, the
KEY condition is raised.

If you use a DIRECT INPUT or DIRECT UPDATE file, you can read, write, rewrite,
or delete records in the same way as for a KEYED SEQUENTIAL file.

Figure 41 on page 284 shows one method by which a KSDS can be updated using
the prime index.
DIRUPDT: PROC OPTIONS(MAIN);

DCL DIREC FILE RECORD KEYED ENV(VSAM),
  ONCODE BUILTIN,
  OUTREC CHAR(23),
  NUMBER CHAR(3) DEF OUTREC POS(21),
  NAME CHAR(20) DEF OUTREC,
  CODE CHAR(1),
  EOF BIT(1) INIT('0'B);

ON ENDFILE(SYSIN) EOF='1'B;

ON KEY(DIREC) BEGIN;
  IF ONCODE=51 THEN PUT FILE(SYSPRINT) SKIP EDIT
      ('NOT FOUND: ',NAME)(A(15),A);
  IF ONCODE=52 THEN PUT FILE(SYSPRINT) SKIP EDIT
      ('DUPLICATE: ',NAME)(A(15),A);
END;

OPEN FILE(DIREC) DIRECT UPDATE;

GET FILE(SYSIN) EDIT (NAME,NUMBER,CODE)
     (COLUMN(1),A(20),A(3),A(1));
DO WHILE (¬EOF);
  PUT FILE(SYSPRINT) SKIP EDIT (' ',NAME,'#',NUMBER,' ',CODE)
     (A(1),A(20),A(1),A(3),A(1),A(1));
  SELECT (CODE);
  WHEN('A') WRITE FILE(DIREC) FROM(OUTREC) KEYFROM(NAME);
  WHEN('C') REWRITE FILE(DIREC) FROM(OUTREC) KEY(NAME);
  WHEN('D') DELETE FILE(DIREC) KEY(NAME);
  OTHERWISE PUT FILE(SYSPRINT) SKIP EDIT
      ('INVALID CODE: ',NAME) (A(15),A);
END;

GET FILE(SYSIN) EDIT (NAME,NUMBER,CODE)
     (COLUMN(1),A(20),A(3),A(1));
END;

CLOSE FILE(DIREC);
PUT FILE(SYSPRINT) PAGE;
OPEN FILE(DIREC) SEQUENTIAL INPUT;

EOF='0'B;
ON ENDFILE(DIREC) EOF='1'B;

READ FILE(DIREC) INTO(OUTREC);
DO WHILE(~EOF);
  PUT FILE(SYSPRINT) SKIP EDIT(OUTREC)(A);
  READ FILE(DIREC) INTO(OUTREC);
  END;
CLOSE FILE(DIREC);
END DIRUPDT;

Figure 41. Updating a KSDS
Updating a KSDS (continued)

A DIRECT update file is used and the data is altered according to a code that is passed in the records in the file SYSIN:

A  Add a new record
C  Change the number of an existing name
D  Delete a record

At the label NEXT, the name, number, and code are read in and action taken according to the value of the code. A KEY ON-unit is used to handle any incorrect keys. When the updating is finished (at the label PRINT), the file DIREC is closed and reopened with the attributes SEQUENTIAL INPUT. The file is then read sequentially and printed.

The file is associated with the data set by a DD statement that uses the DSNAME PLIVSAM.AJC2.BASE defined in the Access Method Services DEFINE CLUSTER command in Figure 40 on page 282.

Methods of updating a KSDS: There are a number of methods of updating a KSDS. The method shown using a DIRECT file is suitable for the data as it is shown in the example. For mass sequential insertion, use a KEYED SEQUENTIAL UPDATE file. This gives faster performance because the data is written onto the data set only when strictly necessary and not after every write statement, and because the balance of free space within the data set is retained.

Statements to achieve effective mass sequential insertion are:

DCL DIREC KEYED SEQUENTIAL UPDATE  
  ENV(VSAM);
WRITE FILE(DIREC) FROM(OUTREC) 
  KEYFROM(NAME);

The PL/I input/output routines detect that the keys are in sequence and make the correct requests to VSAM. If the keys are not in sequence, this too is detected and no error occurs, although the performance advantage is lost.

Alternate Indexes for KSDSs or Indexed ESDSs
Alternate indexes allow you to access KSDSs or indexed ESDSs in various ways, using either unique or nonunique keys.
Unique Key Alternate Index Path: Figure 42 shows the creation of a unique key alternate index path for the ESDS defined and loaded in Figure 38 on page 278. Using this path, the data set is indexed by the name of the child in the first 15 bytes of the record.

Three Access Method Services commands are used. These are:

**DEFINE ALTERNATEINDEX:** defines the alternate index as a data set to VSAM.

**BLDINDEX:** places the pointers to the relevant records in the alternate index.

**DEFINE PATH:** defines an entity that can be associated with a PL/I file in a DD statement.

DD statements are required for the INFILE and OUTFILE operands of BLDINDEX and for the sort files. Care should be taken that the correct names are specified at the various points.

```plaintext
//OPT9#9 JOB
//STEP1 EXEC PGM=IDCAMS,REGION=512K
//SYSPRINT DD SYSOUT=A
//SYSIN DD *
DEFINE ALTERNATEINDEX -
  (NAME(PLIVSAM.AJC1.ALPHIND) -
  VOLUMES(nnnnnn) -
  TRACKS(4 1) -
  KEYS(15 0) -
  RECORDSIZE(20 40) -
  UNIQUEKEY -
  RELATE(PLIVSAM.AJC1.BASE))
/*
//STEP2 EXEC PGM=IDCAMS,REGION=512K
//DD1 DD DSNAME=PLIVSAM.AJC1.BASE,DISP=SHR
//DD2 DD DSNAME=PLIVSAM.AJC1.ALPHIND,DISP=SHR
//SYSPRINT DD SYSOUT=A
//SYSIN DD *
BLDINDEX INFILE(DD1) OUTFILE(DD2)
DEFINE PATH -
  (NAME(PLIVSAM.AJC1.ALPHPATH) -
  PATHENTRY(PLIVSAM.AJC1.ALPHIND))

Figure 42. Creating a Unique Key Alternate Index Path for an ESDS
```

Nonunique Key Alternate Index Path: Figure 43 on page 287 shows the creation of a nonunique key alternate index path for an ESDS. The alternate index enables the data to be selected by the sex of the children. This enables the girls or the boys to be accessed separately and every member of each group to be accessed by use of the key.

The three Access Method Services commands used are:

**DEFINE ALTERNATEINDEX:** defines the alternate index as a data set to VSAM.

**BLDINDEX:** places the pointers to the relevant records in the alternate index.

**DEFINE PATH:** defines an entity that can be associated with a PL/I file in a DD statement.
DD statements are required for the INFILE and OUTFILE operands of BLDINDEX and for the sort files. Care should be taken that the correct names are specified at the various points.

The fact that the index has nonunique keys is specified by the use of the NONUNIQUEKEY operand. When creating an index with nonunique keys, be careful to ensure that the RECORDSIZE you specify is large enough. In a nonunique alternate index, each alternate index record contains pointers to all the records that have the associated index key. The pointer takes the form of an RBA for an ESDS and the prime key for a KSDS. When a large number of records might have the same key, a large record is required.

Figure 44 on page 288 shows the creation of a unique key alternate index path for a KSDS. The data set is indexed by the telephone number, enabling the number to be used as a key to discover the name of the person on that extension. The fact that keys are to be unique is specified by UNIQUEKEY. Also, the data set will be able to be listed in numerical order to show which numbers are not used. The three Access Method Services commands used are:

DEFINE ALTERNATEINDEX: defines the data set that will hold the alternate index data.

BLDINDEX: places the pointer to the relevant records in the alternate index.

DEFINE PATH: defines the entity that can be associated with a PL/I file in a DD statement.

DD statements are required for the INFILE and OUTFILE of BLDINDEX and for the sort files. Be careful not to confuse the names involved.
When creating an alternate index with a unique key, you should ensure that no further records could be included with the same alternate key. In practice, a unique key alternate index would not be entirely satisfactory for a telephone directory as it would not allow two people to have the same number. Similarly, the prime key would prevent one person having two numbers. A solution would be to have an ESDS with two nonunique key alternate indexes, or to restructure the data format to allow more than one number per person and to have a nonunique key alternate index for the numbers.

Detecting Nonunique Alternate Index Keys: If you are accessing a VSAM data set by means of an alternate index path, the presence of nonunique keys can be detected by means of the SAMEKEY built-in function. After each retrieval, SAMEKEY indicates whether any further records exist with the same alternate index key as the record just retrieved. Hence, it is possible to stop at the last of a series of records with nonunique keys without having to read beyond the last record. SAMEKEY (file-reference) returns '1'B if the input/output statement has completed successfully and the accessed record is followed by another with the same key; otherwise, it returns '0'B.

Using Alternate Indexes with ESDSs: Figure 45 on page 290 shows the use of alternate indexes and backward reading on an ESDS. The program has four files:

BASEFLE
reads the base data set forward.

BACKFLE
reads the base data set backward.

ALPHFLE
is the alphabetic alternate index path indexing the children by name.

SEXFILE

is the alternate index path that corresponds to the sex of the children.

There are DD statements for all the files. They connect BASEFLE and BACKFLE to
the base data set by specifying the name of the base data set in the DSNAME
parameter, and connect ALPHFLE and SEXFILE by specifying the names of the
paths given in Figure 42 on page 286 and Figure 43 on page 287.

The program uses SEQUENTIAL files to access the data and print it first in the
normal order, then in the reverse order. At the label AGEQUERY, a DIRECT file is
used to read the data associated with an alternate index key in the unique alternate
index.

Finally, at the label SPRINT, a KEYED SEQUENTIAL file is used to print a list of
the females in the family, using the nonunique key alternate index path. The
SAMEKEY built-in function is used to read all the records with the same key. The
names of the females will be accessed in the order in which they were originally
entered. This will happen whether the file is read forward or backward. For a
nonunique key path, the BKWD option only affects the order in which the keys are
read; the order of items with the same key remains the same as it is when the file
is read forward.

Deletion: At the end of the example, the Access Method Services DELETE command
is used to delete the base data set. When this is done, the associated alternate
indexes and paths will also be deleted.

Using Alternate Indexes with KSDSs: Figure 46 on page 292 shows the use of a
path with a unique alternate index key to update a KSDS and then to access and
print it in the order of the alternate index.

The alternate index path is associated with the PL/I file by a DD statement that
specifies the name of the path (PLIVSAM_AJC2.NUMPATH, given in the DEFINE
PATH command in Figure 44 on page 288) as the DSNAME.

In the first section of the program, a DIRECT OUTPUT file is used to insert a new
record using the alternate index key. Note that any alteration made with an
alternate index must not alter the prime key or the alternate index key of access of
an existing record. Also, the alternation must not add a duplicate key in the prime
index or any unique key alternate index.

In the second section of the program (at the label PRINTIT), the data set is read in
the order of the alternate index keys using a SEQUENTIAL INPUT file. It is then
printed onto SYSPRINT.
READIT: PROC OPTIONS(MAIN);
   DCL BASEFLE FILE SEQUENTIAL INPUT ENV(VSAM), /*File to read base data set forward */
   BACKFLE FILE SEQUENTIAL INPUT ENV(VSAM BKWD), /*File to read base data set backward */
   ALPHFLE FILE DIRECT INPUT ENV(VSAM), /*File to access via unique alternate index path */
   SEXFILE FILE KEYED SEQUENTIAL INPUT ENV(VSAM), /*File to access via nonunique alternate index path */
   STRING CHAR(80), /*String to be read into */
   1 STRUC DEF (STRING),
   2 NAME CHAR(25),
   2 DATE_OF_BIRTH CHAR(2),
   2 FILL CHAR(10),
   2 SEX CHAR(1);
   DCL NAMEHOLD CHAR(25),SAMEKEY BUILTIN;
   DCL EOF BIT(1) INIT('0'B);
   /*Print out the family eldest first*/
   ON ENDFILE(BASEFLE) EOF='1'B;
   PUT EDIT('FAMILY ELDEST FIRST')(A);"
AGEQUERY:
EOF='0' B;
ON ENDFILE(SYSIN) EOF='1' B;
GET SKIP EDIT(NAMEHOLD)(A(25));
DO WHILE(¬EOF);
   READ FILE(ALPHFLE) INTO (STRING) KEY(NAMEHOLD);
   PUT SKIP (2) EDIT(NAMEHOLD,' WAS BORN IN ','
      DATE_OF_BIRTH)(A,X(1),A,X(1),A);
   GET SKIP EDIT(NAMEHOLD)(A(25));
END;
SPRINT:
CLOSE FILE(ALPHFLE);
PUT SKIP(1);
/*Use the alternate index to print out all the females in the family*/
ON ENDFILE(SEXFILE) GOTO FINITO;
PUT SKIP(2) EDIT('ALL THE FEMALES')(A);
READ FILE(SEXFILE) INTO (STRING) KEY('F');
PUT SKIP EDIT(STRING)(A);
DO WHILE(SAMEKEY(SEXFILE));
   READ FILE(SEXFILE) INTO (STRING);
   PUT SKIP EDIT(STRING)(A);
END;
FINITO:
   END;

/*
//GO.BASEFLE DD DSN=PLIVSAM.AJC1.BASE,DISP=SHR
//GO.BACKFLE DD DSN=PLIVSAM.AJC1.BASE,DISP=SHR
//GO.ALPHFLE DD DSN=PLIVSAM.AJC1.ALPATH,DISP=SHR
//GO.SEXFILE DD DSN=PLIVSAM.AJC1.SEXPATH,DISP=SHR
//GO.SYSIN DD *
ANDY
*/
/*
//STEP2 EXEC PGM=IDCAMS,REGION=512K
//SYSPRINT DD SYSOUT=A
//SYSPRINT DD *
   DELETE -
      PLIVSAM.AJC1.BASE
*/

Alternate Index Paths and Backward Reading with an ESDS (continued)
Relative-record data sets

The statements and options allowed for VSAM relative-record data sets (RRDS) are shown in Table 30 on page 293.
Table 30. Statements and options allowed for loading and accessing VSAM relative-record data sets.

<table>
<thead>
<tr>
<th>File declaration</th>
<th>Valid statements, with options you must include</th>
<th>Other options you can also include</th>
</tr>
</thead>
<tbody>
<tr>
<td>SEQUENTIAL OUTPUT BUFFERED</td>
<td>WRITE FILE(file-reference) FROM(reference); LOCATE based-variable FILE(file-reference);</td>
<td>KEYFROM(expression) or KEYTO(reference); SET(pointer-reference);</td>
</tr>
<tr>
<td>SEQUENTIAL INPUT BUFFERED</td>
<td>READ FILE(file-reference) INTO(reference); READ FILE(file-reference) SET(pointer-reference); READ FILE(file-reference)</td>
<td>KEY(expression) or KEYTO(reference);</td>
</tr>
<tr>
<td>SEQUENTIAL UPDATE BUFFERED</td>
<td>READ FILE(file-reference) INTO(reference); READ FILE(file-reference) SET(pointer-reference); READ FILE(file-reference)</td>
<td>KEY(expression) or KEYTO(expression); IGNORE(expression);</td>
</tr>
<tr>
<td>DIRECT OUTPUT BUFFERED</td>
<td>WRITE FILE(file-reference) FROM(reference) KEYFROM(expression);</td>
<td></td>
</tr>
<tr>
<td>DIRECT INPUT BUFFERED</td>
<td>READ FILE(file-reference) INTO(reference) KEY(expression); READ FILE(file-reference) SET(pointer-reference) KEY(expression);</td>
<td></td>
</tr>
</tbody>
</table>
Table 30. Statements and options allowed for loading and accessing VSAM relative-record data sets (continued).

<table>
<thead>
<tr>
<th>File declaration</th>
<th>Valid statements, with options you must include</th>
<th>Other options you can also include</th>
</tr>
</thead>
<tbody>
<tr>
<td>DIRECT UPDATE BUFFERED</td>
<td>READ FILE(file-reference) INTO(reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>READ FILE(file-reference) SET(pointer-reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>REWRITE FILE(file-reference) FROM(reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>DELETE FILE(file-reference) KEY(expression);</td>
<td></td>
</tr>
<tr>
<td></td>
<td>WRITE FILE(file-reference) FROM(reference) KEYFROM(expression);</td>
<td></td>
</tr>
</tbody>
</table>

Notes:
1. The complete file declaration would include the attributes FILE and RECORD. If you use any of the options KEY, KEYFROM, or KEYTO, your declaration must also include the attribute KEYED.
   The UNLOCK statement for DIRECT UPDATE files is ignored if you use it for files associated with a VSAM RRDS.
2. The statement READ FILE(file-reference); is equivalent to the statement READ FILE(file-reference) IGNORE(1);.

Loading an RRDS: When an RRDS is being loaded, you must open the associated file for OUTPUT. Use either a DIRECT or a SEQUENTIAL file.

For a DIRECT OUTPUT file, each record is placed in the position specified by the relative record number (or key) in the KEYFROM option of the WRITE statement (see “Keys for VSAM data sets” on page 268).

For a SEQUENTIAL OUTPUT file, use WRITE statements with or without the KEYFROM option. If you specify the KEYFROM option, the record is placed in the specified slot; if you omit it, the record is placed in the slot following the current position. There is no requirement for the records to be presented in ascending relative record number order. If you omit the KEYFROM option, you can obtain the relative record number of the written record by means of the KEYTO option.

If you want to load an RRDS sequentially, without use of the KEYFROM or KEYTO options, your file is not required to have the KEYED attribute.

It is an error to attempt to load a record into a position that already contains a record: if you use the KEYFROM option, the KEY condition is raised; if you omit it, the ERROR condition is raised.

In Figure 47 on page 296, the data set is defined with a DEFINE CLUSTER command and given the name PLIVSAM.AJC3.BASE. The fact that it is an RRDS is
determined by the NUMBERED keyword. In the PL/I program, it is loaded with a DIRECT OUTPUT file and a WRITE...FROM...KEYFROM statement is used.

If the data had been in order and the keys in sequence, it would have been possible to use a SEQUENTIAL file and write into the data set from the start. The records would then have been placed in the next available slot and given the appropriate number. The number of the key for each record could have been returned using the KEYTO option.

The PL/I file is associated with the data set by the DD statement, which uses as the DSNAME the name given in the DEFINE CLUSTER command.
Using a SEQUENTIAL file to access an RRDS: You can open a SEQUENTIAL file that is used to access an RRDS with either the INPUT or the UPDATE attribute. If you use any of the options KEY, KEYTO, or KEYFROM, your file must also have the KEYED attribute.
For **READ** statements without the **KEY** option, the records are recovered in ascending relative record number order. Any empty slots in the data set are skipped.

If you use the **KEY** option, the record recovered by a **READ** statement is the one with the relative record number you specify. Such a **READ** statement positions the data set at the specified record; subsequent sequential reads will recover the following records in sequence.

**WRITE** statements with or without the **KEYFROM** option are allowed for **KEYED SEQUENTIAL UPDATE** files. You can make insertions anywhere in the data set, regardless of the position of any previous access. For **WRITE** with the **KEYFROM** option, the **KEY** condition is raised if an attempt is made to insert a record with the same relative record number as a record that already exists on the data set. If you omit the **KEYFROM** option, an attempt is made to write the record in the next slot, relative to the current position. The **ERROR** condition is raised if this slot is not empty.

You can use the **KEYTO** option to recover the key of a record that is added by means of a **WRITE** statement without the **KEYFROM** option.

**REWRITE** statements, with or without the **KEY** option, are allowed for **UPDATE** files. If you use the **KEY** option, the record that is rewritten is the record with the relative record number you specify; otherwise, it is the record that was accessed by the previous **READ** statement.

**DELETE** statements, with or without the **KEY** option, can be used to delete records from the dataset.

**Using a DIRECT file to access an RRDS:** A **DIRECT** file used to access an RRDS can have the **OUTPUT**, **INPUT**, or **UPDATE** attribute. You can read, write, rewrite, or delete records exactly as though a **KEYED SEQUENTIAL** file were used.

Figure 48 on page 298 shows an RRDS being updated. A **DIRECT UPDATE** file is used and new records are written by key. There is no need to check for the records being empty, because the empty records are not available under VSAM.

In the second half of the program, starting at the label **PRINT**, the updated file is printed out. Again there is no need to check for the empty records as there is in **REGIONAL(1)**.

The **PL/I** file is associated with the data sets by a **DD** statement that specifies the **DSNAME** **PLIVSAM.AJC3.BASE**, the name given in the **DEFINE CLUSTER** command in Figure 48 on page 298.

At the end of the example, the **DELETE** command is used to delete the data set.
/* NOTE: WITH A WRITE STATEMENT AFTER THE DELETE FILE STATEMENT,
   * A DUPLICATE MESSAGE IS EXPECTED FOR CODE 'C' ITEMS
   * WHOSE NEWNO CORRESPONDS TO AN EXISTING NUMBER IN THE LIST,
   * FOR EXAMPLE, ELLIOT.
   * WITH A REWRITE STATEMENT AFTER THE DELETE FILE STATEMENT,
   * A NOT FOUND MESSAGE IS EXPECTED FOR CODE 'C' ITEMS
   * WHOSE NEWNO DOES NOT CORRESPOND TO AN EXISTING NUMBER IN
   * THE LIST, FOR EXAMPLE, NEWMAN AND BRAMLEY.
   */
//OPT9#18 JOB
//STEP1 EXEC IBMZCBG
//PLI.SYSIN DD *
ACRI: PROC OPTIONS(MAIN);
   DCL NOS FILE RECORD KEYED ENV(VSAM),NAME CHAR(20),
      (NEWNO,OLDNO) CHAR(2),CODE CHAR(1),IOFIELD CHAR(20),
      BYTE CHAR(1) DEF IOFIELD, EOF BIT(1) INIT('0'B),
      ONCODE BUILTIN;
   ON ENDFILE(SYSIN) EOF='1'B;
   OPEN FILE(NOS) DIRECT UPDATE;
   ON KEY(NOS) BEGIN;
      IF ONCODE=51 THEN PUT FILE(SYSPRINT) SKIP EDIT
         ('NOT FOUND:',NAME)(A(15),A);
      IF ONCODE=52 THEN PUT FILE(SYSPRINT) SKIP EDIT
         ('DUPLICATE:',NAME)(A(15),A);
   END;
   GET FILE(SYSIN) EDIT(NAME,NEWNO,OLDNO,CODE)
      (COLUMN(1),A(20),A(2),A(2),A(1));
   DO WHILE (¬EOF);
      PUT FILE(SYSPRINT) SKIP EDIT (' ',NAME,'#',NEWNO,OLDNO,' ',CODE)
         (A(1),A(20),A(1),2(A(2)),X(5),2(A(1)));
      SELECT(CODE);
      WHEN('A') WRITE FILE(NOS) KEYFROM(NEWNO) FROM(NAME);
      WHEN('C') DO;
         DELETE FILE(NOS) KEY(OLDNO);
         WRITE FILE(NOS) KEYFROM(NEWNO) FROM(NAME);
      END;
      WHEN('D') DELETE FILE(NOS) KEY(OLDNO);
      OTHERWISE PUT FILE(SYSPRINT) SKIP EDIT
         ('INVALID CODE: ',NAME)(A(15),A);
   END;

Figure 48. Updating an RRDS
Using Files Defined for non-VSAM Data Sets

Using Shared Data Sets

PL/I allows cross-region or cross-system sharing of data sets. The support for this type of sharing is provided by VSAM. For details about the support, see the following DFSMS manuals:

- DFSMS: Using Data Sets
- DFSMS: Access Method Services for Catalogs
Part 3. Improving your program

Chapter 13. Improving performance 303
Selecting compiler options for optimal performance 303
  OPTIMIZE 303
  GONUMBER 303
  ARCH 303
  REDUCE 304
  RULES 304
    IBM/ANS 304
    (NO)LAXCTL 305
  PREFIX 305
    CONVERSION 305
    FIXEDOVERFLOW 306
  DEFAULT 306
    BYADDR or BYVALUE 306
    (NON)CONNECTED 307
    (NO)DESCRIPTOR 307
    (NO)INLINE 308
    LINKAGE 308
    (RE)ORDER 308
  NOOVERLAP 309
  RETURNS(BYVALUE) or RETURNS(BYADDR) 309
Summary of compiler options that improve performance 309
Coding for better performance 309
  DATA-directed input and output 309
  Input-only parameters 310
  GOTO statements 310
  String assignments 310
  Loop control variables 311
  PACKAGEs versus nested PROCEDUREs 311
    Example with nested procedures 312
    Example with packaged procedures 312
  REDUCIBLE Functions 312
  DESCLOCATOR or DESCLIST 313
  DEFINED versus UNION 313
  Named constants versus static variables 313
    Example with optimal code but no meaningful names 314
    Example with meaningful names but not optimal code 314
    Example with optimal code AND meaningful names 314
  Avoiding calls to library routines 315
  Preloading library routines 315
Chapter 13. Improving performance

Many considerations for improving the speed of your program are independent of the compiler that you use and the platform on which it runs. This chapter, however, identifies those considerations that are unique to the PL/I compiler and the code it generates.

Selecting compiler options for optimal performance

The compiler options you choose can greatly improve the performance of the code generated by the compiler; however, like most performance considerations, there are trade-offs associated with these choices. Fortunately, you can weigh the trade-offs associated with compiler options without editing your source code because these options can be specified on the command line or in the configuration file.

If you want to avoid details, the least complex way to improve the performance of generated code is to specify the following (nondefault) compiler options:
- OPT(2) or OPT(3)
- DFT(REORDER)

The following sections describe, in more detail, performance improvements and trade-offs associated with specific compiler options.

**OPTIMIZE**

You can specify the OPTIMIZE option to improve the speed of your program; otherwise, the compiler makes only basic optimization efforts.

Choosing OPTIMIZE(2) directs the compiler to generate code for better performance. Usually, the resultant code is shorter than when the program is compiled under NOOPTIMIZE. Sometimes, however, a longer sequence of instructions runs faster than a shorter sequence. This occurs, for instance, when a branch table is created for a SELECT statement where the values in the WHEN clauses contain gaps. The increased number of instructions generated is usually offset by the execution of fewer instructions in other places.

Choosing OPTIMIZE(3) directs the compiler to generate even better code. However, when you specify the OPTIMIZE(3) option, the compilation will take longer (and sometimes much, much longer) than when you specify OPTIMIZE(2).

**GONUMBER**

Using this option results in a statement number table used for debugging. This added information can be extremely helpful when debugging, but including statement number tables increases the size of your executable file. Larger executable files can take longer to load.

**ARCH**

Using the highest value in the ARCH option allows the compiler to select from the largest set of instructions available under z/OS and thus permits it to generate the most optimal code.
REDUCE

The REDUCE option specifies that the compiler is permitted to reduce an assignment of a null string to a structure into a simple copy operation - even if that means padding bytes might be overwritten.

The REDUCE option will cause fewer lines of code to be generated for an assignment of a null string to a structure, and that will usually mean your compilation will be quicker and your code will run much faster. However, padding bytes may be zeroed out.

For instance, in the following structure, there is one byte of padding between field11 and field12.

```
dcl sample ext,
  5 field10   bin fixed(31),
  5 field11   dec fixed(13),
  5 field12   bin fixed(31),
  5 field13   bin fixed(31),
  5 field14   bit(32),
  5 field15   bin fixed(31),
  5 field16   bit(32),
  5 field17   bin fixed(31);
```

Now consider the assignment `sample = "";`

Under the NOREDUCE option, it will cause eight assignments to be generated, but the padding byte will be unchanged.

However, under REDUCE, the assignment would be reduced to three operations.

RULES

Most of the RULES suboptions affect only the severity with which certain coding practices, such as not declaring variables, are flagged and have no impact on performance. However, these suboptions do have an impact on performance.

IBM/ANS

When you use the RULES(IBM) option, the compiler supports scaled FIXED BINARY and, what is more important for performance, generates scaled FIXED BINARY results in some operations. Under RULES(ANS), scaled FIXED BINARY is not supported and scaled FIXED BINARY results are never generated. This means that the code generated under RULES(ANS) always runs at least as fast as the code generated under RULES(IBM), and sometimes runs faster.

For example, consider the following code fragment:
```
dcl (i,j,k) fixed bin(15);
  ...
  i = j / k;
```

Under RULES(IBM), the result of the division has the attributes FIXED BIN(31,16). This means that a shift instruction is required before the division and several more instructions are needed to perform the assignment.

Under RULES(ANS), the result of the division has the attributes FIXED BIN(15,0). This means that a shift is not needed before the division, and no extra instructions are needed to perform the assignment.
(NO)LAXCTL

Under RULES(LAXCTL), a CONTROLLED variable may be declared with constant extents and yet allocated with different extents.

For instance, under RULES(LAXCTL), you may declare a structure as follows:

```pli
dcl
  1 a controlled,
  2 b char(17),
  2 c char(29);
```

However, you could then allocate it as follows:

```pli
allocate
  1 a,
  2 b char(170),
  2 c char(290);
```

This has disastrous consequences for performance because it means that whenever the compiler sees a reference to the structure `A` or to any member of that structure, the compiler is forced to assume that it knows nothing about the lengths, dimensions or offsets of the fields in it.

However, the RULES(NOLAXCTL) option disallows this coding practice: under RULES(NOLAXCTL), if you then want to allocate a CONTROLLED variable with a variable extent, then that extents must be declared either with an asterisk or with a non-constant expression. Consequently, under RULES(NOLAXCTL), when a CONTROLLED variable is declared with constant extents, then the compiler can generate much better code for any reference to that variable.

PREFIX

This option determines if selected PL/I conditions are enabled by default. The default suboptions for PREFIX are set to conform to the PL/I language definition; however, overriding the defaults can have a significant effect on the performance of your program. The default suboptions are:

- CONVERSION
- INVALIDOP
- FIXEDOVERFLOW
- OVERFLOW
- INVALIDOP
- NOSIZE
- NOSTRINGRANGE
- NOSTRINGSIZE
- NOSUBSCRIPTRANGE
- UNDERFLOW
- ZERODIVIDE

By specifying the SIZE, STRINGRANGE, STRINGSIZE, or SUBSCRIPTRANGE suboptions, the compiler generates extra code that helps you pinpoint various problem areas in your source that would otherwise be hard to find. This extra code, however, can slow program performance significantly.

CONVERSION

When you disable the CONVERSION condition, some character-to-numeric conversions are done inline and without checking the validity of the source; therefore, specifying NOCONVERSION also affects program performance.
FIXEDOVERFLOW
On some platforms, the FIXEDOVERFLOW condition is raised by the hardware and the compiler does not need to generate any extra code to detect it.

DEFAULT
Using the DEFAULT option, you can select attribute defaults. As is true with the PREFIX option, the suboptions for DEFAULT are set to conform to the PL/I language definition. Changing the defaults in some instances can affect performance.

Some of the suboptions, such IBM/ANS and ASSIGNABLE/NONASSIGNABLE, have no effect on program performance. But other suboptions can affect performance to varying degrees and, if applied inappropriately, can make your program invalid. The more important of these suboptions are:

BYADDR or BYVALUE
When the DEFAULT(BYADDR) option is in effect, arguments are passed by reference (as required by PL/I) unless an attribute in an entry declaration indicates otherwise. As arguments are passed by reference, the address of the argument is passed from one routine (calling routine) to another (called routine) as the variable itself is passed. Any change made to the argument while in the called routine is reflected in the calling routine when it resumes execution.

Program logic often depends on passing variables by reference. Passing a variable by reference, however, can hinder performance in two ways:
1. Every reference to that parameter requires an extra instruction.
2. Since the address of the variable is passed to another routine, the compiler is forced to make assumptions about when that variable might change and generate very conservative code for any reference to that variable.

Consequently, you should pass parameters by value using the BYVALUE suboption whenever your program logic allows. Even if you use the BYADDR attribute to indicate that one parameter should be passed by reference, you can use the DEFAULT(BYVALUE) option to ensure that all other parameters are passed by value.

A BYVALUE argument should be one that could reasonably be passed in a register. Hence its type should be one of
- REAL FIXED BIN
- REAL FLOAT
- POINTER
- OFFSET
- HANDLE
- LIMITED ENTRY
- FILE
- ORDINAL
- CHAR(1)
- WCHAR(1)
- ALIGNED BIT(n) with n less than or equal to 8
Moreover, special care is required when using BYVALUE parameters with assembler code as any BYVALUE parameter that does not require a multiple of 4 bytes will be widened so that 4 bytes are used in the parameter list.

If a procedure receives and modifies only one parameter that is passed by BYADDR, consider converting the procedure to a function that receives that parameter by value. The function would then end with a RETURN statement containing the updated value of the parameter.

**Procedure with BYADDR parameter:**

```plaintext
da: proc( parm1, parm2, ..., parmN );
dcl parm1 byaddr ...;
dcl parm2 byvalue ...;
.
.
dcl parmN byvalue ...;
/* program logic */
end;
```

**Faster, equivalent function with BYVALUE parameter:**

```plaintext
da: proc( parm1, parm2, ..., parmN )
  returns( ... /* attributes of parm1 */ );
dcl parm1 byvalue ...;
dcl parm2 byvalue ...;
.
.
dcl parmN byvalue ...;
/* program logic */
return( parm1 );
end;
```

**(NON)CONNECTED**

The DEFAULT(NONCONNECTED) option indicates that the compiler assumes that any aggregate parameters are NONCONNECTED. References to elements of NONCONNECTED aggregate parameters require the compiler to generate code to access the parameter's descriptor, even if the aggregate is declared with constant extents.

The compiler does not generate these instructions if the aggregate parameter has constant extents and is CONNECTED. Consequently, if your application never passes nonconnected parameters, your code is more optimal if you use the DEFAULT(CONNECTED) option.

**(NO)DESCRIPTOR**

The DEFAULT(DESCRIPTOR) option indicates that, by default, a descriptor is passed for any string, area, or aggregate parameter; however, the descriptor is used only if the parameter has nonconstant extents or if the parameter is an array with the NONCONNECTED attribute. In this case, the instructions and space required to pass the descriptor provide no benefit and incur substantial cost (the size of a structure descriptor is often greater than size of the structure itself). Consequently,
by specifying DEFAULT(NODESCRIPTOR) and using OPTIONS(DESCRIPTOR) only as needed on PROCEDURE statements and ENTRY declarations, your code runs more optimally.

(NO)INLINE
The suboption NOINLINE indicates that procedures and begin blocks should not be inlined.

Inlining occurs only when you specify optimization.

Inlining user code eliminates the overhead of the function call and linkage, and also exposes the function's code to the optimizer, resulting in faster code performance. Inlining produces the best results when the overhead for the function is nontrivial, for example, when functions are called within nested loops. Inlining is also beneficial when the inlined function provides additional opportunities for optimization, such as when constant arguments are used.

For programs containing many procedures that are not nested:
- If the procedures are small and only called from a few places, you can increase performance by specifying INLINE.
- If the procedures are large and called from several places, inlining duplicates code throughout the program. This increase in the size of the program might offset any increase of speed. In this case, you might prefer to leave NOINLINE as the default and specify OPTIONS(INLINE) only on individually selected procedures.

When you use inlining, you need more stack space. When a function is called, its local storage is allocated at the time of the call and freed when it returns to the calling function. If that same function is inlined, its storage is allocated when the function that calls it is entered, and is not freed until that calling function ends. Ensure that you have enough stack space for the local storage of the inlined functions.

LINKAGE
This suboption tells the compiler the default linkage to use when the LINKAGE suboption of the OPTIONS attribute or option for an entry has not been specified.

The compiler supports various linkages, each with its unique performance characteristics. When you invoke an ENTRY provided by an external entity (such as an operating system), you must use the linkage previously defined for that ENTRY.

As you create your own applications, however, you can choose the linkage convention. The OPTLINK linkage is strongly recommended because it provides significantly better performance than other linkage conventions.

(RE)ORDER
The DEFAULT(ORDER) option indicates that the ORDER option is applied to every block, meaning that variables in that block referenced in ON-units (or blocks dynamically descendant from ON-units) have their latest values. This effectively prohibits almost all optimization on such variables. Consequently, if your program logic allows, use DEFAULT(REORDER) to generate superior code.
The DEFAULT(NOOVERLAP) option lets the compiler assume that the source and target in an assignment do not overlap, and it can therefore generate smaller and faster code.

However, if you use this option, you must ensure that the source and target in assignment do not overlap. For example, under the DEFAULT( NOOVERLAP ) option, the assignment in this example would be invalid:

```
dcl c char(20);
substr(c,2,5) = substr(c,1,5);
```

**RETURNS(BYVALUE) or RETURNS(BYADDR)**

When the DEFAULT(RETURNS(BYVALUE)) option is in effect, the BYVALUE attribute is applied to all RETURNS description lists that do not specify BYADDR. This means that these functions return values in registers, when possible, in order to produce the most optimal code.

### Summary of compiler options that improve performance

In summary, the following options (if appropriate for your application) can improve performance:

- OPTIMIZE(3)
- ARCH(10)
- REDUCE
- RULES(ANS NOLAXCTL)
- DEFAULT with the following suboptions
  - BYVALUE
  - CONNECTED
  - NODESCRIPTOR
  - INLINE
  - LINKAGE(OPTLINK)
  - REORDER
  - NOOVERLAP
  - RETURNS(BYVALUE)

### Coding for better performance

As you write code, there is generally more than one correct way to accomplish a given task. Many important factors influence the coding style you choose, including readability and maintainability. The following sections discuss choices that you can make while coding that potentially affect the performance of your program.

**DATA-directed input and output**

Using GET DATA and PUT DATA statements for debugging can prove very helpful. When you use these statements, however, you generally pay the price of decreased performance. This cost to performance is usually very high when you use either GET DATA or PUT DATA without a variable list.

Many programmers use PUT DATA statements in their ON ERROR code as illustrated in the following example:

```
on error
  begin;
    on error system;
  .
  .
```

---

Chapter 13. Improving performance 309
In this case, the program would perform more optimally by including a list of selected variables with the PUT DATA statement.

The ON ERROR block in the previous example contained an ON ERROR system statement before the PUT DATA statement. This prevents the program from getting caught in an infinite loop if an error occurs in the PUT DATA statement (which could occur if any variables to be listed contained invalid FIXED DECIMAL values) or elsewhere in the ON ERROR block.

**Input-only parameters**

If a procedure has a BYADDR parameter which it uses as input only, it is best to declare that parameter as NONASSIGNABLE (rather than letting it get the default attribute of ASSIGNABLE). If that procedure is later called with a constant for that parameter, the compiler can put that constant in static storage and pass the address of that static area.

This practice is particularly useful for strings and other parameters that cannot be passed in registers (input-only parameters that can be passed in registers are best declared as BYVALUE).

In the following declaration, for instance, the first parameter to getenv is an input-only CHAR VARYINGZ string:

```
dcl getenv entry( char(*) varyingz nonasgn byaddr,  
                 pointer byaddr )  
    returns( native fixed bin(31) optional )  
  options( nodescriptor );
```

If this function is invoked with the string 'IBM_OPTIONS', the compiler can pass the address of that string rather than assigning it to a compiler-generated temporary storage area and passing the address of that area.

**GOTO statements**

A GOTO statement that uses either a label in another block or a label variable severely limits optimizations that the compiler might perform. If a label array is initialized and declared AUTOMATIC, either implicitly or explicitly, any GOTO to an element of that array will hinder optimization. However, if the array is declared as STATIC, the compiler assumes the CONSTANT attribute for it and no optimization is hindered.

**String assignments**

When one string is assigned to another, the compiler ensures that:
- The target has the correct value even if the source and target overlap.
- The source string is truncated if it is longer than the target.

This assurance comes at the price of some extra instructions. The compiler attempts to generate these extra instructions only when necessary, but often you, as the programmer, know they are not necessary when the compiler cannot be sure. For instance, if the source and target are based character strings and you know
they cannot overlap, you could use the PLIMOVE built-in function to eliminate the extra code the compiler would otherwise be forced to generate.

In the example which follows, faster code is generated for the second assignment statement:

```plaintext
dcl based_Str char(64) based( null() );
dcl target_Addr pointer;
dcl source_Addr pointer;

target_Addr->based_Str = source_Addr->based_Str;
call plimove( target_Addr, source_Addr, stg(based_Str) );
```

If you have any doubts about whether the source and target might overlap or whether the target is big enough to hold the source, you should not use the PLIMOVE built-in.

**Loop control variables**

Program performance improves if your loop control variables are one of the types in the following list. You should rarely, if ever, use other types of variables.

- FIXED BINARY with zero scale factor
- FLOAT
- ORDINAL
- HANDLE
- POINTER
- OFFSET

Performance also improves if loop control variables are not members of arrays, structures, or unions. The compiler issues a warning message when they are. Loop control variables that are AUTOMATIC and not used for any other purpose give you the optimal code generation.

If a loop control variable is a FIXED BINARY, performance is best if it has precision 31 and is SIGNED.

Performance is decreased if your program depends not only on the value of a loop control variable, but also on its address. For example, the performance is decreased if the ADDR built-in function is applied to the variable or if the variable is passed by BYADDR to another routine.

**PACKAGEs versus nested PROCEDUREs**

Calling nested procedures requires that an extra *hidden parameter* (the backchain pointer) is passed. As a result, the fewer nested procedures that your application contains, the faster it runs.

To improve the performance of your application, you can convert a mother-daughter pair of nested procedures into level-1 sister procedures inside of a package. This conversion is possible if your nested procedure does not rely on any of the automatic and internal static variables declared in its parent procedures.

If procedure b in “Example with nested procedures” on page 312 does not use any of the variables declared in a, you can improve the performance of both procedures by reorganizing them into the package illustrated in “Example with packaged procedures” on page 312.
Example with nested procedures

```
a: proc;

  dcl (i,j,k) fixed bin;
  dcl ib based fixed bin;
  .
  .
  .
  call b( addr(i) );
  .
  .

b: proc( px );
  dcl px pointer;
  display( px->ib );
  end;
end;
```

Example with packaged procedures

```
p: package exports( a );

  dcl ib based fixed bin;

a: proc;

  dcl (i,j,k) fixed bin;
  .
  .
  .
  call b( addr(i) );
  .
  .

end;

b: proc( px );
  dcl px pointer;
  display( px->ib );
  end;
end p;
```

REDUCIBLE Functions

REDUCIBLE indicates that a procedure or entry need not be invoked multiple times if the argument(s) stays unchanged, and that the invocation of the procedure has no side effects.

For example, a user-written function that computes a result based on unchanging data should be declared REDUCIBLE. A function that computes a result based on changing data, such as a random number or time of day, should be declared IRREDUCIBLE.

In the following example, $f$ is invoked only once since REDUCIBLE is part of the declaration. If IRREDUCIBLE had been used in the declaration, $f$ would be invoked twice.

```
dcl (f) entry options( reducible ) returns( fixed bin );

select;
  when( f(x) < 0 )
    .
    .
```
when( f(x) > 0 )
  .
  .
otherwise
  .
  .
end;

**DESCLOCATOR or DESCLIST**

When the DEFAULT(DESCLOCATOR) option is in effect, the compiler passes arguments requiring descriptors (such as strings and structures) via a descriptor locator in much the same way that the old compiler did. More information on descriptors and how they are passed is available in Chapter 23, “PL/I descriptors,” on page 475.

This option allows you to invoke an entry point that is not always passed all of the arguments that it declares.

This option also allows you to continue the somewhat unwise programming practice of passing a structure and receiving it as a pointer.

However, the code generated by the compiler for DEFAULT(DESCLOCATOR) may, in some situations, perform less well than that for DEFAULT(DESCLIST).

**DEFINED versus UNION**

The UNION attribute is more powerful than the DEFINED attribute and provides more function. In addition, the compiler generates better code for union references.

In the following example, the pair of variables b3 and b4 perform the same function as b1 and b2, but the compiler generates more optimal code for the pair in the union.

dcl b1 bit(32);
dcl b2 bit(16) def b1;

dcl
  1 * union,
  2 b3 bit(32),
  2 b4 bit(16);

Code that uses UNIONs instead of the DEFINED attribute is subject to less misinterpretation. Variable declarations in unions are in a single location making it easy to realize that when one member of the union changes, all of the others change also. This dynamic change is less obvious in declarations that use DEFINED variables since the declare statements can be several lines apart.

**Named constants versus static variables**

You can define named constants by declaring a variable with the VALUE attribute. If you use static variables with the INITIAL attribute and you do not alter the variable, you should declare the variable a named constant using the VALUE attribute. The compiler does not treat NONASSIGNABLE scalar STATIC variables as true named constants.

The compiler generates better code whenever expressions are evaluated during compilation, so you can use named constants to produce efficient code with no loss
Coding for better performance

in readability. For example, identical object code is produced for the two usages of the VERIFY built-in function in the following example:

dcl numeric char value('0123456789');

jx = verify( string, numeric );
jx = verify( string, '0123456789' );

The following examples illustrate how you can use the VALUE attribute to get optimal code without sacrificing readability.

**Example with optimal code but no meaningful names**

```pli
select( x );
when( '01'b4 )
  .
  .
  when( '02'b4 )
  .
  .
when( '03'b4 )
  .
end;
```

**Example with meaningful names but not optimal code**

```pli
dcl ( a1 init( '01'b4), a2 init( '02'b4), a3 init( '03'b4), a4 init( '04'b4), a5 init( '05'b4) ) bit(8) aligned static nonassignable;

dcl x bit(8) aligned;
select( x );
when( a1 )
  .
  .
when( a2 )
  .
when( a3 )
  .
end;
```

**Example with optimal code AND meaningful names**

```pli
dcl ( a1 value( '01'b4), a2 value( '02'b4), a3 value( '03'b4), a4 value( '04'b4), a5 value( '05'b4) ) bit(8);

dcl x bit(8) aligned;
```
Avoiding calls to library routines

The bitwise operations (prefix NOT, infix AND, infix OR, and infix EXCLUSIVE OR) are often evaluated by calls to library routines. These operations are, however, handled without a library call if either of the following conditions is true:

- Both operands are bit(1)
- Both operands are aligned and have the same constant length.

For certain assignments, expressions, and built-in function references, the compiler generates calls to library routines. If you avoid these calls, your code generally runs faster.

To help you determine when the compiler generates such calls, the compiler generates a message whenever a conversion is done using a library routine.

When your code refers to a member of a BASED structure using REFER, the compiler often has to generate one or more calls to a library routine to map the structure at run-time. These calls can be expensive, and so when the compiler makes these calls, it will issue a message so that you can locate these potential hot-spots in your code.

If you do have code that uses BASED structures with REFER and which the compiler flags with this message, you may be able to get better performance by passing the structure to a subroutine that declares a corresponding structure with * extents. This will cause the structure to be mapped once at the CALL statement, but there will be no further remappings when it is accessed in the called subroutine.

Preloading library routines

The PL/I library contains one RMODE 24 routine that is used for low-level system i/o functions: IBMPOIOA. If your code does RECORD i/o or uses SYSPRINT as a STREAM OUTPUT file (without compiling with the STDSYS option), then you will significantly improve your performance if you preload this routine or put it into the (E)LPA.
Coding for better performance
Part 4. Using interfaces to other products

Chapter 14. Using the Sort program ..... 319
Preparing to use Sort ..... 319
Choosing the type of Sort ..... 320
Specifying the sorting field ..... 323
Specifying the records to be sorted ..... 324
Maximum record lengths ..... 325
Determining storage needed for Sort
Main storage ..... 325
Auxiliary storage ..... 326
Calling the Sort program ..... 326
Example 1 ..... 327
Example 2 ..... 328
Example 3 ..... 328
Example 4 ..... 328
Example 5 ..... 328
Determining whether the Sort was successful ..... 328
Establishing data sets for Sort
Sort work data sets ..... 329
Input data set ..... 329
Output data set ..... 330
Checkpoint data set ..... 330
Sort data input and output ..... 330
Data input and output handling routines ..... 330
E15—Input handling routine (Sort Exit E15) ..... 331
E35—Output handling routine (Sort Exit E35) ..... 334
Calling PLISRKT example ..... 335
Calling PLISRKT2 example ..... 336
Calling PLISRKC example ..... 337
Calling PLISRKT3 example ..... 338
Sorting variable-length records example ..... 339

Chapter 15. ILC with C ..... 341
Equivalent data types ..... 341
Simple type equivalence ..... 341
Struct type equivalence ..... 342
Enum type equivalence ..... 342
File type equivalence ..... 343
Using C functions ..... 343
Matching simple parameter types ..... 344
Matching string parameter types ..... 347
Functions returning ENTRYs ..... 348
Linkages ..... 349
Sharing output and input ..... 351
Sharing output ..... 351
Sharing input ..... 352
Using the ATTACH statement ..... 352
Redirecting C standard streams ..... 352
Summary. ..... 352

Chapter 16. Interfacing with Java ..... 355
Java Native Interface (JNI) ..... 355
Calling PL/I program from Java ..... 356
JNI Sample Program #1 - 'Hello World' ..... 356
Writing Java Sample Program #1 ..... 356
Step 1: Writing the Java Program ..... 356
Step 2: Compiling the Java Program ..... 357
Step 3: Writing the PL/I Program ..... 357
Step 4: Compiling and Linking the PL/I Program ..... 359
Step 5: Running the Sample Program ..... 359
JNI Sample Program #2 - Passing a String ..... 359
Writing Java Sample Program #2 ..... 359
Step 1: Writing the Java Program ..... 359
Step 2: Compiling the Java Program ..... 361
Step 3: Writing the PL/I Program ..... 362
Step 4: Compiling and Linking the PL/I Program ..... 363
Step 5: Running the Sample Program ..... 364
JNI Sample Program #3 - Passing an Integer ..... 364
Writing Java Sample Program #3 ..... 364
Step 1: Writing the Java Program ..... 364
Step 2: Compiling the Java Program ..... 366
Step 3: Writing the PL/I Program ..... 366
Step 4: Compiling and Linking the PL/I Program ..... 367
Step 5: Running the Sample Program ..... 368
JNI Sample Program #4 - Java Invocation API ..... 368
Writing Java Sample Program #4 ..... 368
Step 1: Writing the Java Program ..... 368
Step 2: Compiling the Java Program ..... 368
Step 3: Writing the PL/I Program ..... 368
Step 4: Compiling and Linking the PL/I Program ..... 372
Step 5: Running the Sample Program ..... 372
Attaching programs to an existing Java VM ..... 372
Determining equivalent Java and PL/I data types ..... 373

© Copyright IBM Corp. 1999, 2015
Chapter 14. Using the Sort program

The compiler provides an interface called PLISRTx (x = A, B, C, or D) that allows you to make use of the IBM-supplied Sort programs.

To use the Sort program with PLISRTx, you must:
1. Include a call to one of the entry points of PLISRTx, passing it the information on the fields to be sorted. This information includes the length of the records, the maximum amount of storage to use, the name of a variable to be used as a return code, and other information required to carry out the sort.
2. Specify the data sets required by the Sort program in JCL DD statements.

When used from PL/I, the Sort program sorts records of all normal lengths on a large number of sorting fields. Data of most types can be sorted into ascending or descending order. The source of the data to be sorted can be either a data set or a user-written PL/I procedure that the Sort program will call each time a record is required for the sort. Similarly, the destination of the sort can be a data set or a PL/I procedure that handles the sorted records.

Using PL/I procedures allows processing to be done before or after the sort itself, thus allowing a complete sorting operation to be handled completely by a call to the sort interface. It is important to understand that the PL/I procedures handling input or output are called from the Sort program itself and will effectively become part of it.

PL/I can operate with DFSORT™ or a program with the same interface. DFSORT is a release of the program product 5740-SM1. DFSORT has many built-in features you can use to eliminate the need for writing program logic (for example, INCLUDE, OMIT, OUTREC and SUM statement plus the many ICETOOL operators). See DFSORT Application Programming Guide for details and Getting Started with DFSORT for a tutorial.

The following material applies to DFSORT. Because you can use programs other than DFSORT, the actual capabilities and restrictions vary. For these capabilities and restrictions, see DFSORT Application Programming Guide, or the equivalent publication for your sort product.

To use the Sort program you must include the correct PL/I statements in your source program and specify the correct data sets in your JCL.

Preparing to use Sort

Before using Sort, you must determine the type of sort you require, the length and format of the sorting fields in the data, the length of your data records, and the amount of auxiliary and main storage you will allow for sorting.

To determine the PLISRTx entry point that you will use, you must decide the source of your unsorted data, and the destination of your sorted data. You must choose between data sets and PL/I subroutines. Using data sets is simpler to understand and gives faster performance. Using PL/I subroutines gives you more flexibility and more function, enabling you to manipulate or print the data before it
is sorted, and to make immediate use of it in its sorted form. If you decide to use an input or output handling subroutine, you will need to read "Data input and output handling routines" on page 330.

The entry points and the source and destination of data are as follows:

<table>
<thead>
<tr>
<th>Entry point</th>
<th>Source</th>
<th>Destination</th>
</tr>
</thead>
<tbody>
<tr>
<td>PLISRTA</td>
<td>Data set</td>
<td>Data set</td>
</tr>
<tr>
<td>PLISRTB</td>
<td>Subroutine</td>
<td>Data set</td>
</tr>
<tr>
<td>PLISRTC</td>
<td>Data set</td>
<td>Subroutine</td>
</tr>
<tr>
<td>PLISRTD</td>
<td>Subroutine</td>
<td>Subroutine</td>
</tr>
</tbody>
</table>

Having determined the entry point you are using, you must now determine the following things about your data set:

- The position of the sorting fields; these can be either the complete record or any part or parts of it
- The type of data these fields represent, for example, character or binary
- Whether you want the sort on each field to be in ascending or descending order
- Whether you want equal records to be retained in the order of the input, or whether their order can be altered during sorting

Specify these options on the SORT statement, which is the first argument to PLISRTx. After you have determined these, you must determine two things about the records to be sorted:

- Whether the record format is fixed or varying
- The length of the record, which is the maximum length for varying

Specify these on the RECORD statement, which is the second argument to PLISRTx.

Finally, you must decide on the amount of main and auxiliary storage you will allow for the Sort program. For further details, see "Determining storage needed for Sort" on page 325.

### Choosing the type of Sort

If you want to make the best use of the Sort program, you must understand something of how it works. In your PL/I program you specify a sort by using a CALL statement to the sort interface subroutine PLISRTx. This subroutine has four entry points: x=A, B, C, and D. Each specifies a different source for the unsorted data and destination for the data when it has been sorted. For example, a call to PLISRTA specifies that the unsorted data (the input to sort) is on a data set, and that the sorted data (the output from sort) is to be placed on another data set. The CALL PLISRTx statement must contain an argument list giving the Sort program information about the data set to be sorted, the fields on which it is to be sorted, the amount of space available, the name of a variable into which Sort will place a return code indicating the success or failure of the sort, and the name of any output or input handling procedure that can be used.

The sort interface routine builds an argument list for the Sort program from the information supplied by the PLISRTx argument list and the choice of PLISRTx entry point. Control is then transferred to the Sort program. If you have specified an output- or input-handling routine, this will be called by the Sort program as many times as is necessary to handle each of the unsorted or sorted records. When
the sort operation is complete, the Sort program returns to the PL/I calling
procedure communicating its success or failure in a return code, which is placed in
one of the arguments passed to the interface routine. The return code can then be
tested in the PL/I routine to discover whether processing should continue.
Figure 49 on page 322 is a simplified flowchart showing this operation.
Within the Sort program itself, the flow of control between the Sort program and input- and output-handling routines is controlled by return codes. The Sort program...

Figure 49. Flow of control for Sort program
program calls these routines at the appropriate point in its processing. (Within the
Sort program, and its associated documentation, these routines are known as user
exits. The routine that passes input to be sorted is the E15 sort user exit. The
routine that processes sorted output is the E35 sort user exit.) From the routines,
Sort expects a return code indicating either that it should call the routine again, or
that it should continue with the next stage of processing.

The important points to remember about Sort are: (1) it is a self-contained program
that handles the complete sort operation, and (2) it communicates with the caller,
and with the user exits that it calls, by means of return codes.

The remainder of this chapter gives detailed information on how to use Sort from
PL/I. First the required PL/I statements are described, and then the data set
requirements. The chapter finishes with a series of examples showing the use of
the four entry points of the sort interface routine.

Specifying the sorting field

The SORT statement is the first argument to PLISRTx. The syntax of the SORT
statement must be a character string expression that takes the form:

'bSORTbFIELDS=(start1,length1,form1,seq1,
...startn,lengthn,formn,seqn)[,other options]b'

For example:

' SORT FIELDS=(1,10,CH,A) '  

b represents one or more blanks. Blanks shown are mandatory. No other blanks
are allowed.

start, length, form, seq
defines a sorting field. You can specify any number of sorting fields, but there
is a limit on the total length of the fields. If more than one field is to be sorted
on, the records are sorted first according to the first field, and then those that
are of equal value are sorted according to the second field, and so on. If all the
sorting values are equal, the order of equal records will be arbitrary unless you
use the EQUALS option. (See later in this definition list.) Fields can overlay
each other.

For DFSORT (5740-SM1), the maximum total length of the sorting fields is
restricted to 4092 bytes and all sorting fields must be within 4092 bytes of the
start of the record. Other sort products might have different restrictions.

start is the starting position within the record. Give the value in bytes
except for binary data where you can use a “byte.bit” notation. The
first byte in a string is considered to be byte 1, the first bit is bit 0.
(Thus the second bit in byte 2 is referred to as 2.1.) For varying length
records you must include the 4-byte length prefix, making 5 the first
byte of data.

length is the length of the sorting field. Give the value in bytes except for
binary where you can use “byte.bit” notation. The length of sorting
fields is restricted according to their data type.

form is the format of the data. This is the format assumed for the purpose of
sorting. All data passed between PL/I routines and Sort must be in the
form of character strings. The main data types and the restrictions on
their length are shown below. Additional data types are available for
special-purpose sorts. See the DFSORT Application Programming Guide,
or the equivalent publication for your sort product.
Code     Data type and length
CH       character 1–4096
ZD       zoned decimal, signed 1–32
PD       packed decimal, signed 1–32
FI       fixed point, signed 1–256
BI       binary, unsigned 1 bit to 4092 bytes
FL       floating-point, signed 1–256

The sum of the lengths of all fields must not exceed 4092 bytes.

seq is the sequence in which the data will be sorted as follows:
A     ascending (that is, 1,2,3,...)
D     descending (that is, ...,3,2,1)

Note: You cannot specify E, because PL/I does not provide a method of passing a user-supplied sequence.

other options
You can specify a number of other options, depending on your Sort program. You must separate them from the FIELDS operand and from each other by commas. Do not place blanks between operands.

FILSZ=y specifies the number of records in the sort and allows for optimization by Sort. If y is only approximate, E should precede y.

SKIPREC=y specifies that y records at the start of the input file are to be ignored before sorting the remaining records.

CKPT or CHKPT specifies that checkpoints are to be taken. If you use this option, you must provide a SORTCKPT data set and DFSORT's 16NCKPT=NO installation option must be specified.

EQUALS|NOEQUALS specifies whether the order of equal records will be preserved as it was in the input (EQUALS) or will be arbitrary (NOEQUALS). You could improve sort performance by using the NOEQUALS. The default option is chosen when Sort is installed. The IBM-supplied default is NOEQUALS.

DYNALLOC=(d,n) (OS/VS Sort only) specifies that the program dynamically allocates intermediate storage.

d     is the device type (3380, etc.).
n     is the number of work areas.

Specifying the records to be sorted
Use the RECORD statement as the second argument to PLISRTx. The syntax of the RECORD statement must be a character string expression which, when evaluated, takes the syntax shown below:

'bRECORDbTYPE=rectype[,LENGTH=(L1,[,L4,L5]]b'

For example:

' RECORD TYPE=F,LENGTH=(80) '
represents one or more blanks. Blanks shown are mandatory. No other blanks are allowed.

**TYPE**
specifies the type of record as follows:
- **F** fixed length
- **V** varying length EBCDIC
- **D** varying length ASCII

Even when you use input and output routines to handle the unsorted and sorted data, you must specify the record type as it applies to the work data sets used by Sort.

If varying length strings are passed to Sort from an input routine (E15 exit), you should normally specify V as a record format. However, if you specify F, the records are padded to the maximum length with blanks.

**LENGTH**
specifies the length of the record to be sorted. You can omit LENGTH if you use PLISRTA or PLISRTC, because the length will be taken from the input data set. Note that there is a restriction on the maximum and minimum length of the record that can be sorted (see below). For varying length records, you must include the four-byte prefix.

- **L1** is the length of the record to be sorted. For VSAM data sets sorted as varying records it is the maximum record size+4.
- **"** represent two arguments that are not applicable to Sort when called from PL/I. You must include the commas if the arguments that follow are used.
- **L4** specifies the minimum length of record when varying length records are used. If supplied, it is used by Sort for optimization purposes.
- **L5** specifies the modal (most common) length of record when varying length records are used. If supplied, it is used by Sort for optimization purposes.

**Maximum record lengths**
The length of a record can never exceed the maximum length specified by the user. The maximum record length for variable length records is 32756 bytes and for fixed length records it is 32760 bytes.

**Determining storage needed for Sort**

**Main storage**
Sort requires both main and auxiliary storage. The minimum main storage for DFSORT is 88K bytes, but for best performance, more storage (on the order of 1 megabyte or more) is recommended. DFSORT can take advantage of storage above 16M virtual or extended architecture processors. Under z/OS, DFSORT can also take advantage of expanded storage. You can specify that Sort use the maximum amount of storage available by passing a storage parameter in the following manner:

```
DCL MAXSTOR FIXED BINARY (31,0);
UNSPEC(MAXSTOR)='00000000'B||UNSPEC('MAX');
CALL PLISRRTA
(' SORT FIELDS=(1,80,CH,A) ',
 ' RECORD TYPE=F,LENGTH=(80) ',
 MAXSTOR,
 RETCODE,
 'TASK');
```
If files are opened in E15 or E35 exit routines, enough residual storage should be allowed for the files to open successfully.

**Auxiliary storage**
Calculating the minimum auxiliary storage for a particular sorting operation is a complicated task. To achieve maximum efficiency with auxiliary storage, use direct access storage devices (DASDs) whenever possible. For more information on improving program efficiency, see the [DFSORT Application Programming Guide](#), particularly the information about dynamic allocation of workspace which allows DFSORT to determine the auxiliary storage needed and allocate it for you.

If you are interested only in providing enough storage to ensure that the sort will work, make the total size of the SORTWK data sets large enough to hold three sets of the records being sorted. (You will not gain any advantage by specifying more than three if you have enough space in three data sets.)

However, because this suggestion is an approximation, it might not work, so you should check the sort manuals. If this suggestion does work, you will probably have wasted space.

### Calling the Sort program
When you have determined the points described above, you are in a position to write the CALL PLISRTx statement. You should do this with some care; for the entry points and arguments to use, see Table 31.

<table>
<thead>
<tr>
<th>Entry points</th>
<th>Arguments</th>
</tr>
</thead>
<tbody>
<tr>
<td>PLISRTA</td>
<td>(sort statement, record statement, storage, return code, data set prefix, message level, sort technique)</td>
</tr>
<tr>
<td>PLISRTB</td>
<td>(sort statement, record statement, storage, return code, input routine, data set prefix, message level, sort technique)</td>
</tr>
<tr>
<td>PLISRTC</td>
<td>(sort statement, record statement, storage, return code, output routine, data set prefix, message level, sort technique)</td>
</tr>
<tr>
<td>PLISRTD</td>
<td>(sort statement, record statement, storage, return code, input routine, output routine, data set prefix, message level, sort technique)</td>
</tr>
</tbody>
</table>

**Sort statement**
Character string expression containing the Sort program SORT statement. Describes sorting fields and format. See “Specifying the sorting field” on page 323.

**Record statement**
Character string expression containing the Sort program RECORD statement. Describes the length and record format of data. See “Specifying the records to be sorted” on page 324.

**Storage**
Fixed binary expression giving maximum amount of main storage to be used by the Sort program. Must be $>88K$ bytes for DFSORT. See also “Determining storage needed for Sort” on page 325.

**Return code**
Fixed binary variable of precision $(31,0)$ in which Sort places a return code when it has completed. The meaning of the return code is:

- 0=Sort successful
- 16=Sort failed
- 20=Sort message data set missing

---

326  
Enterprise PL/I for z/OS Programming Guide
Table 31. The entry points and arguments to PLISRTx (x = A, B, C, or D) (continued)

<table>
<thead>
<tr>
<th>Entry points</th>
<th>Arguments</th>
</tr>
</thead>
<tbody>
<tr>
<td>Input routine</td>
<td>(PLISRTB and PLISRTD only.) Name of the PL/I external or internal procedure used to supply the records for the Sort program at sort exit E15.</td>
</tr>
<tr>
<td>Output routine</td>
<td>(PLISRTC and PLISRTD only.) Name of the PL/I external or internal procedure to which Sort passes the sorted records at sort exit E35.</td>
</tr>
<tr>
<td>Data set prefix</td>
<td>Character string expression of four characters that replaces the default prefix of 'SORT' in the names of the sort data sets SORTIN, SORTOUT, SORTWKnn and SORTCNTL, if used. Thus if the argument is “TASK”, the data sets TASKIN, TASKOUT, TASKWKnn, and TASKCNTL can be used. This facility enables multiple invocations of Sort to be made in the same job step. The four characters must start with an alphabetic character and must not be one of the reserved names PEER, BALN, CRCX, OSCl, POLY, DIAG, SYSC, or LIST. You must code a null string for this argument if you require either of the following arguments but do not require this argument.</td>
</tr>
<tr>
<td>Message level</td>
<td>Character string expression of two characters indicating how Sort’s diagnostic messages are to be handled, as follows: NO No messages to SYSOUT AP All messages to SYSOUT CP Critical messages to SYSOUT SYSOUT will normally be allocated to the printer, hence the use of the mnemonic letter “P”. Other codes are also allowed for certain of the Sort programs. For further details on these codes, see DFSORT Application Programming Guide. You must code a null string for this argument if you require the following argument but you do not require this argument.</td>
</tr>
<tr>
<td>Sort technique</td>
<td>(This is not used by DFSORT; it appears for compatibility reasons only.) Character string of length 4 that indicates the type of sort to be carried out, as follows: PEER Peerage sort BALN Balanced CRCX Criss-cross sort OSCl Oscillating POLY Polyphase sort Normally the Sort program will analyze the amount of space available and choose the most effective technique without any action from you. You should use this argument only as a bypass for sorting problems or when you are certain that performance could be improved by another technique. See DFSORT Application Programming Guide for further information.</td>
</tr>
</tbody>
</table>

The examples below indicate the form that the CALL PLISRTx statement normally takes.

**Example 1**

A call to PLISRTA sorting 80-byte records from SORTIN to SORTOUT using 1048576 (1 megabyte) of storage, and a return code, RETCODE, declared as FIXED BINARY (31,0).

```plaintext
CALL PLISRTA (' SORT FIELDS=(1,80,CH,A) ',
' RECORD TYPE=F,LENGTH=(80) ',
1048576,
RETCODE);
```
Example 2
This example is the same as example 1 except that the input, output, and work data sets are called TASKIN, TASKOUT, TASKWK01, and so forth. This might occur if Sort was being called twice in one job step.

```
CALL PLISRTA (' SORT FIELDS=(1,80,CH,A) ',
               ' RECORD TYPE=F,LENGTH=(80) ',
               1048576,
               RETCODE,
               'TASK');
```

Example 3
This example is the same as example 1 except that the sort is to be undertaken on two fields. First, bytes 1 to 10 which are characters, and then, if these are equal, bytes 11 and 12 which contain a binary field, both fields are to be sorted in ascending order.

```
CALL PLISRTA (' SORT FIELDS=(1,10,CH,A,11,2,BI,A) ',
               ' RECORD TYPE=F,LENGTH=(80) ',
               1048576,
               RETCODE);
```

Example 4
This example shows a call to PLISRTB. The input is to be passed to Sort by the PL/I routine PUTIN, the sort is to be carried out on characters 1 to 10 of an 80 byte fixed length record. Other information as above.

```
CALL PLISRTB (' SORT FIELDS=(1,10,CH,A) ',
               ' RECORD TYPE=F,LENGTH=(80) ',
               1048576,
               RETCODE,
               PUTIN);
```

Example 5
This example shows a call to PLISRTD. The input is to be supplied by the PL/I routine PUTIN and the output is to be passed to the PL/I routine PUTOUT. The record to be sorted is 84 bytes varying (including the length prefix). It is to be sorted on bytes 1 through 5 of the data in ascending order, then if these fields are equal, on bytes 6 through 10 in descending order. (Note that the 4-byte length prefix is included so that the actual values used are 5 and 10 for the starting points.) If both these fields are the same, the order of the input is to be retained. (The EQUALS option does this.)

```
CALL PLISRTD (' SORT FIELDS=(5,5,CH,A,10,5,CH,D),EQUALS ',
               ' RECORD TYPE=V,LENGTH=(84) ',
               1048576,
               RETCODE,
               PUTIN, /*input routine (sort exit E15)*/
               PUTOUT); /*output routine (sort exit E35)*/
```

Determining whether the Sort was successful
When the sort is completed, Sort sets a return code in the variable named in the fourth argument of the call to PLISRTx. It then returns control to the statement that follows the CALL PLISRTx statement. The value returned indicates the success or failure of the sort as follows:

```
0     Sort successful
16    Sort failed
20    Sort message data set missing
```
You must declare the variable to which the return code is passed as FIXED BINARY (31,0). It is standard practice to test the value of the return code after the CALL PLISRTx statement and take appropriate action according to the success or failure of the operation.

For example (assuming the return code was called RETCODE):

IF RETCODE=0 THEN DO;
  PUT DATA(RETCODE);
  SIGNAL ERROR;
END;

If the job step that follows the sort depends on the success or failure of the sort, you should set the value returned in the Sort program as the return code from the PL/I program. This return code is then available for the following job step. The PL/I return code is set by a call to PLIRETC. You can call PLIRETC with the value returned from Sort thus:

CALL PLIRETC(RETCODE);

You should not confuse this call to PLIRETC with the calls made in the input and output routines, where a return code is used for passing control information to Sort.

Establishing data sets for Sort

If DFSORT was installed in a library not known to the system, you must specify the DFSORT library in a JOBLIB or STEPLIB DD statement.

When you call Sort, certain sort data sets must not be open. These are:

SYSOUT

A data set (normally the printer) on which messages from the Sort program will be written.

Sort work data sets

SORTWK01–SORTWK32

Note: If you specify more than 32 sort work data sets, DFSORT will only use the first 32.

****WK01–****WK32

From 1 to 32 working data sets used in the sorting process. These must be direct access. For a discussion of space required and number of data sets, see “Determining storage needed for Sort” on page 325.

**** represents the four characters that you can specify as the data set prefix argument in calls to PLISRTx. This allows you to use data sets with prefixes other than SORT. They must start with an alphabetic character and must not be the names PEER, BALN, CRCX, OSCL, POLY, SYSC, LIST, or DIAG.

Input data set

SORTIN

****IN

The input data set used when PLISRTA and PLISRTC are called.

See ****WK01–****WK32 above for a detailed description.
Output data set

SORTOUT

****OUT
The output data set used when PLISRTA and PLISRTB are called.
See ****WK01–****WK32 above for a detailed description.

Checkpoint data set

SORTCKPT

Data set used to hold checkpoint data, if CKPT or CHKPT option was used in the SORT statement argument and DFSORT’s 16NCKPT=NO installation option was specified. For information on this program DD statement, see DFSORT Application Programming Guide.

DFSPARM SORTCNTL

Data set from which additional or changed control statements can be read (optional). For additional information on this program DD statement, see DFSORT Application Programming Guide.

See ****WK01–****WK32 above for a detailed description.

Sort data input and output

The source of the data to be sorted is provided either directly from a data set or indirectly by a routine (Sort Exit E15) written by the user. Similarly, the destination of the sorted output is either a data set or a routine (Sort Exit E35) provided by the user.

PLISRTA is the simplest of all of the interfaces because it sorts from data set to data set. An example of a PLISRTA program is in Figure 53 on page 335. Other interfaces require either the input handling routine or the output handling routine, or both.

Data input and output handling routines

The input handling and output handling routines are called by Sort when PLISRTB, PLISRTC, or PLISRTD is used. They must be written in PL/I, and can be either internal or external procedures. If they are internal to the routine that calls PLISRTx, they behave in the same way as ordinary internal procedures in respect of scope of names. The input and output procedure names must themselves be known in the procedure that makes the call to PLISRTx.

The routines are called individually for each record required by Sort or passed from Sort. Therefore, each routine must be written to handle one record at a time. Variables declared as AUTOMATIC within the procedures will not retain their values between calls. Consequently, items such as counters, which need to be retained from one call to the next, should either be declared as STATIC or be declared in the containing block.

The E15 and E35 sort exits must not be MAIN procedures.
Input routines are normally used to process the data in some way before it is sorted. You can use input routines to print the data, as shown in the Figure 54 on page 336 and Figure 56 on page 338, or to generate or manipulate the sorting fields to achieve the correct results.

The input handling routine is used by Sort when a call is made to either PLISRTB or PLISRTD. When Sort requires a record, it calls the input routine which should return a record in character string format, with a return code of 12. This return code means that the record passed is to be included in the sort. Sort continues to call the routine until a return code of 8 is passed. A return code of 8 means that all records have already been passed, and that Sort is not to call the routine again. If a record is returned when the return code is 8, it is ignored by Sort.

The data returned by the routine must be a character string. It can be fixed or varying. If it is varying, you should normally specify V as the record format in the RECORD statement which is the second argument in the call to PLISRTx. However, you can specify F, in which case the string will be padded to its maximum length with blanks. The record is returned with a RETURN statement, and you must specify the RETURNS attribute in the PROCEDURE statement. The return code is set in a call to PLIRETC. A flowchart for a typical input routine is shown in Figure 50 on page 332.
Figure 50. Flowcharts for input and output handling subroutines

Skeletal code for a typical input routine is shown in Figure 51 on page 333.
Examples of an input routine are given in Figure 54 on page 336 and Figure 56 on page 338.

In addition to the return codes of 12 (include current record in sort) and 8 (all records sent), Sort allows the use of a return code of 16. This ends the sort and causes Sort to return to your PL/I program with a return code of 16—Sort failed.

**Note:** A call to PLIRETC sets a return code that will be passed by your PL/I program, and will be available to any job steps that follow it. When an output handling routine has been used, it is good practice to reset the return code with a call to PLIRETC after the call to PLISRTx to avoid receiving a nonzero completion code. By calling PLIRETC with the return code from Sort as the argument, you can make the PL/I return code reflect the success or failure of the sort. This practice is shown in Figure 55 on page 337.
Output handling routines are normally used for any processing that is necessary after the sort. This could be to print the sorted data, as shown in Figure 55 on page 337 and Figure 56 on page 338, or to use the sorted data to generate further information. The output handling routine is used by Sort when a call is made to PLISRTC or PLISRTD. When the records have been sorted, Sort passes them, one at a time, to the output handling routine. The output routine then processes them as required. When all the records have been passed, Sort sets up its return code and returns to the statement after the CALL PLISRTx statement. There is no indication from Sort to the output handling routine that the last record has been reached. Any end-of-data handling must therefore be done in the procedure that calls PLISRTx.

The record is passed from Sort to the output routine as a character string, and you must declare a character string parameter in the output handling subroutine to receive the data. The output handling subroutine must also pass a return code of 4 to Sort to indicate that it is ready for another record. You set the return code by a call to PLIRETC.

The sort can be stopped by passing a return code of 16 to Sort. This will result in Sort returning to the calling program with a return code of 16–Sort failed.

The record passed to the routine by Sort is a character string parameter. If you specified the record type as F in the second argument in the call to PLISRTx, you should declare the parameter with the length of the record. If you specified the record type as V, you should declare the parameter as adjustable, as in the following example:

```
DCL STRING CHAR(*);
```

Figure 57 on page 339 shows a program that sorts varying length records.

A flowchart for a typical output handling routine is given in Figure 50 on page 332. Skeletal code for a typical output handling routine is shown in Figure 52.

```
E35: PROC(STRING); /*The procedure must have a character string parameter to receive the record from Sort*/

DCL STRING CHAR(80); /*Declaration of parameter*/

(Your code goes here)

CALL PLIRETC(4); /*Pass return code to Sort indicating that the next sorted record is to be passed to this procedure*/
END E35; /*End of procedure returns control to Sort*/
```

Figure 52. Skeletal code for an output handling procedure

You should note that a call to PLIRETC sets a return code that will be passed by your PL/I program, and will be available to any job steps that follow it. When you have used an output handling routine, it is good practice to reset the return code with a call to PLIRETC after the call to PLISRTx to avoid receiving a nonzero completion code. By calling PLIRETC with the return code from Sort as the argument, you can make the PL/I return code reflect the success or failure of the sort. This practice is shown in the examples at the end of this chapter.
Calling PLISRTA example

After each time that the PL/I input- and output-handling routines communicate
the return-code information to the Sort program, the return-code field is reset to
zero; therefore, it is not used as a regular return code other than its specific use for
the Sort program.

For details on handling conditions, especially those that occur during the input-
and output-handling routines, see z/OS Language Environment Programming Guide.

//OPT14#7 JOB ...
//STEP1 EXEC IBMZCBG
//PLI.SYSIN DD *
EX106: PROC OPTIONS(MAIN);
   DCL RETURN_CODE FIXED BIN(31,0);
   CALL PLISRTA (' SORT FIELDS=(7,74,CH,A) ',
                  ' RECORD TYPE=F,LENGTH=(80) ',
                   1048576
                   RETURN_CODE);
   SELECT (RETURN_CODE);
      WHEN(0) PUT SKIP EDIT ('SORT COMPLETE RETURN_CODE 0') (A);
      WHEN(16) PUT SKIP EDIT ('SORT FAILED, RETURN_CODE 16') (A);
      WHEN(20) PUT SKIP EDIT ('SORT MESSAGE DATASET MISSING ') (A);
      OTHER PUT SKIP EDIT ('INVALID SORT RETURN_CODE = ', RETURN_CODE) (A,F(2));
   END /* select */;
   CALL PLIRETC(RETURN_CODE);
   /*set PL/I return code to reflect success of sort*/
   END EX106;

//GO.SORTIN DD *
//GO.SYSPRINT DD SYSOUT=A
//GO.SORTOUT DD SYSOUT=A
//GO.SYSOUT DD SYSOUT=A
//GO.SORTWK01 DD UNIT=SYSDA,SPACE=(CYL,2)
/*
Figure 53. PLISRTA—sorting from input data set to output data set

Chapter 14. Using the Sort program 335
Calling PLISRTB example

/*OPT14#8 JOB ...
//STEP1 EXEC IBMZCBG
//PLI.SYSIN DD *
EX107: PROC OPTIONS(MAIN);
   DCL RETURN_CODE FIXED BIN(31,0);
   CALL PLISRTB (' SORT FIELDS=(7,74,CH,A) ',
     ' RECORD TYPE=F,LENGTH=(80) ',
     1048576
     RETURN_CODE,
     E15X);
   SELECT(RETURN_CODE);
   WHEN(0) PUT SKIP EDIT
     ('SORT COMPLETE RETURN_CODE 0') (A);
   WHEN(16) PUT SKIP EDIT
     ('SORT FAILED, RETURN_CODE 16') (A);
   WHEN(20) PUT SKIP EDIT
     ('SORT MESSAGE DATASET MISSING ') (A);
   OTHER PUT SKIP EDIT
     ('INVALID RETURN_CODE = ',RETURN_CODE)(A,F(2));
   END /* select */;
   CALL PLIRETC(RETURN_CODE);
   /*set PL/I return code to reflect success of sort*/
E15X: /* INPUT HANDLING ROUTINE GETS RECORDS FROM THE INPUT
      STREAM AND PUTS THEM BEFORE THEY ARE SORTED*/
   PROC RETURNS (CHAR(80));
   DCL SYSIN FILE RECORD INPUT,
      INFIELD CHAR(80);
   ON ENDFILE(SYSIN) BEGIN;
   PUT SKIP(3) EDIT ('END OF SORT PROGRAM INPUT')(A);
   CALL PLIRETC(8); /* signal that last record has
   already been sent to sort*/
   INFIELD = '';
   GOTO ENDE15;
   END;
   READ FILE (SYSIN) INTO (INFIELD);
   PUT SKIP EDIT (INFIELD)(A(80)); /*PRINT INPUT*/
   CALL PLIRETC(12); /* request sort to include current
   record and return for more*/
   ENDE15:
   RETURN(INFIELD);
   END EX107;
/*
//GO.SYSIN DD *
003329HOOKER S.W. RIVERDALE, SATCHWELL LANE, BACONSFIELD
002886BOOKER R.R. ROTORUA, LINKEDGE LANE, TOBLEY
003077ROOKER & SON, LITTLETON NURSERIES, SHOLTSPAR
059334HOOK E.H. 109 ELMTREE ROAD, GANNET PARK, NORTHAMPTON
073872HOME TAVERN, WESTLEIGH
000931FOREST, IVER, BUCKS
/*
//GO.SYSPRINT DD SYSOUT=A
//GO.SORTOUT DD SYSOUT=A
//GO.SYSOUT DD SYSOUT=A
//GO.SORTCNTL DD *
   OPTION DYNALLOC=(3380,2),SKIPREC=2
/*

Figure 54. PLISRTB—sorting from input handling routine to output data set
Calling PLISRTC example

//OPT14#9 JOB ...
//STEP1 EXEC IBMZCBG
//PLI.SYSIN DD *
EX108: PROC OPTIONS(MAIN);
  DCL RETURN_CODE FIXED BIN(31,0);
  CALL PLISRTC (' SORT FIELDS=(7,74,CH,A) ',
                 ' RECORD TYPE=F,LENGTH=(80) ',
                 1048576
                 RETURN_CODE,
                 E35X);
  SELECT(RETURN_CODE);
    WHEN(0) PUT SKIP EDIT
       ('SORT COMPLETE RETURN_CODE 0') (A);
    WHEN(16) PUT SKIP EDIT
       ('SORT FAILED, RETURN_CODE 16') (A);
    WHEN(20) PUT SKIP EDIT
       ('SORT MESSAGE DATASET MISSING ') (A);
    OTHER PUT SKIP EDIT
       ('INVALID RETURN_CODE = ', RETURN_CODE) (A,F(2));
  END /* select */;
  CALL PLIRETC (RETURN_CODE);
  /*set PL/I return code to reflect success of sort*/
E35X: /* output handling routine prints sorted records*/
  PROC (INREC);
    DCL INREC CHAR(80);
    PUT SKIP EDIT (INREC) (A);
    CALL PLIRETC(4); /*request next record from sort*/
  END E35X;
END EX108;
/*
//GO.STEPLIB DD DSN=SYS1.SORTLINK,DISP=SHR
//GO.SYSPRINT DD SYSOUT=A
//GO.SYSOUT DD SYSOUT=A
//GO.SORTIN DD *
003329HOOKER S.W. RIVERDALE, SATCHWELL LANE, BACONSFIELD
002886BOOKER R.R. ROTORUA, LINKEDGE LANE, TOBLEY
003077ROOKER & SON, LITTLETON NURSERIES, SHOLTSPAR
059334HOOK E.H. 109 ELMTREE ROAD, GANNET PARK, NORTHAMPTON
073872HOME TAVERN, WESTLEIGH
000931FOREST, IVER, BUCKS
/*
//GO.SORTCNTL DD *
  OPTION DYNALLOC=(3380,2),SKIPREC=2
*/
Figure 55. PLISRTC—sorting from input data set to output handling routine
Calling PLISRTD example

```
//OPT14#10 JOB ...
//STEP1 EXEC IBMZCBG
//PLI.SYSIN DD *
EX109: PROC OPTIONS(MAIN);
  DCL RETURN_CODE FIXED BIN(31,0);
  CALL PLISRTD (' SORT FIELDS=(7,74,CH,A) ',
  ' RECORD TYPE=F,LENGTH=(80) ',
  1048576
  RETURN_CODE,
  E15X,
  E35X);
SELECT(RETURN_CODE);
  WHEN(0) PUT SKIP EDIT
    ('SORT COMPLETE RETURN_CODE 0') (A);
  WHEN(20) PUT SKIP EDIT
    ('SORT MESSAGE DATASET MISSING ') (A);
  OTHER PUT SKIP EDIT
    ('INVALID RETURN_CODE = ', RETURN_CODE) (A,F(2));
END /* select */;
CALL PLIRETC(RETURN_CODE);
/*set PL/I return code to reflect success of sort*/
E15X: /* Input handling routine prints input before sorting*/
  PROC RETURNS(CHAR(80));
    DCL INFIELD CHAR(80);
    ON ENDFILE(SYSIN) BEGIN;
      PUT SKIP(3) EDIT ('END OF SORT PROGRAM INPUT. ',
      'SORTED OUTPUT SHOULD FOLLOW') (A);
      CALL PLIRETC(8); /* Signal end of input to sort*/
      INFIELD = '';
      GOTO ENDE15;
    END;
    GET FILE (SYSIN) EDIT (INFIELD) (A(80));
    PUT SKIP EDIT (INFIELD) (A);
    CALL PLIRETC(12); /*Input to sort continues*/
  END E15X;
ENDE15: RETURN(INFIELD);
END E15X;
E35X: /* Output handling routine prints the sorted records*/
  PROC (INREC);
    DCL INREC CHAR(80);
    NEXT: CALL PLIRETC(4); /* Request next record from sort*/
  END E35X;
END EX109;
/*
//GO.SYSOUT DD SYSOUT=A
//GO.SYSPRINT DD SYSOUT=A
//GO.SORTWK01 DD UNIT=SYSDA,SPACE=(CYL,1)
//GO.SORTWK02 DD UNIT=SYSDA,SPACE=(CYL,1)
//GO.SORTWK03 DD UNIT=SYSDA,SPACE=(CYL,1)
//GO.SYSIN DD *
003329HOOKER S.W. RIVERDALE, SATCHWELL LANE, BACONSFIELD
002886BOOKER R.R. ROTORUA, LINKEDGE LANE, TOBLEY
003077ROOKER & SON, LITTLETON NURSERIES, SHOLTSPAR
059334HOOK E.H. 109 ELMTREE ROAD, GANNE PARK, NORTHAMPTON
073872HOME TAVERN, WESTLEIGH
000931FOREST, IVER, BUCKS
*/
```

Figure 56. PLISRTD—sorting from input handling routine to output handling routine
Sorting variable-length records example

```pli
//OPT14#11 JOB ...
//STEPI EXEC IBMZCBG
//PLI.SYSIN DD *
/* PL/I EXAMPLE USING PLISRTD TO SORT VARIABLE-LENGTH RECORDS */

EX1306: PROC OPTIONS(MAIN);
    DCL RETURN_CODE FIXED BIN(31,0);
    CALL PLISRTD (' SORT FIELDS=(11,14,CH,A) ',
        ' RECORD TYPE=V,LENGTH=(84,,,24,44) ',
        /*NOTE THAT LENGTH IS MAX AND INCLUDES 4 BYTE LENGTH PREFIX*/
        1048576
        RETURN_CODE,
        PUTIN,
        PUTOUT);

    SELECT(RETURN_CODE);
    WHEN(0) PUT SKIP EDIT ('SORT COMPLETE RETURN_CODE 0') (A);
    WHEN(16) PUT SKIP EDIT ( 'SORT FAILED, RETURN_CODE 16') (A);
    WHEN(20) PUT SKIP EDIT ( 'SORT MESSAGE DATASET MISSING ') (A);
    OTHER PUT SKIP EDIT ( 'INVALID RETURN_CODE = ', RETURN_CODE)
        (A,F(2));
END /* SELECT */;

CALL PLIRETC(RETURN_CODE);
/*SET PL/I RETURN CODE TO REFLECT SUCCESS OF SORT*/

PUTIN: PROC RETURNS (CHAR(80) VARYING);
/*OUTPUT HANDLING ROUTINE*/
/*NOTE THAT VARYING MUST BE USED ON RETURNS ATTRIBUTE WHEN USING VARYING LENGTH RECORDS*/
DCL STRING CHAR(80) VAR;
ON ENDFILE(SYSIN) BEGIN;
    PUT SKIP EDIT ('END OF INPUT')(A);
    CALL PLIRETC(8);
    STRING = '';
    GOTO ENDPUT;
END;
GET EDIT(STRING)(A(80));
I=INDEX(STRING || ' ' ' ')-1;/*RESET LENGTH OF THE*/
STRING = SUBSTR(STRING,1,I); /* STRING FROM 80 TO */
/* LENGTH OF TEXT IN */
/* EACH INPUT RECORD.*/
```

Figure 57. Sorting varying-length records using input and output handling routines
PUT SKIP EDIT(I,STRING) (F(2),X(3),A);
CALL PLSRETC(12);
ENDPUT: RETURN(STRING);
END;

PUTOUT:PROC(STRING);
/*OUTPUT HANDLING ROUTINE OUTPUT SORTED RECORDS*/
DCL STRING CHAR (*);
/*NOTE THAT FOR VARYING RECORDS THE STRING
PARAMETER FOR THE OUTPUT-HANDLING ROUTINE
SHOULD BE DECLARED ADJUSTABLE BUT CANNOT BE
DECLARED VARYING*/
PUT SKIP EDIT(STRING)(A); /*PRINT THE SORTED DATA*/
CALL PLSRETC(4);
END; /*ENDS PUTOUT*/
END;

/*
//GO.SYSIN DD *
003329HOOKER S.W. RIVERDALE, SATCHWELL LANE, BACONSFIELD
002886BOOKER R.R. ROTORUA, LINKEDGE LANE, TOBLEY
003077ROOKER & SON, LITTLETON NURSERIES, SHOLTSPAR
059334HOOK E.H. 109 ELMTREE ROAD, GANNET PARK, NORTHAMPTON
073872HOME TAVERN, WESTLEIGH
000931FOREST, IVER, BUCKS
*/
//GO.SYSPRINT DD SYSOUT=A
//GO.SORTOUT DD SYSOUT=A
//GO.SYSPUT DD SYSOUT=A
//GO.SORTWK01 DD UNIT=SYSDA,SPACE=(CYL,1)
//GO.SORTWK02 DD UNIT=SYSDA,SPACE=(CYL,1)
/*/ 

Sorting varying-length records using input and output handling routines (continued)
Chapter 15. ILC with C

This chapter describes some aspects of InterLanguage Communication (ILC) between PL/I and C. The examples illustrate how to use many of the data types common to both languages and should enable you to write PL/I code that either calls or is called by C.

Equivalent data types

The table Table 32 lists the common C and PL/I data type equivalents.

<table>
<thead>
<tr>
<th>C type</th>
<th>Matching PL/I type</th>
</tr>
</thead>
<tbody>
<tr>
<td>char[...]</td>
<td>char(...) varyingz</td>
</tr>
<tr>
<td>wchar[...]</td>
<td>wchar(...) varyingz</td>
</tr>
<tr>
<td>signed char</td>
<td>fixed bin(7)</td>
</tr>
<tr>
<td>unsigned char</td>
<td>unsigned fixed bin(8)</td>
</tr>
<tr>
<td>short</td>
<td>fixed bin(15)</td>
</tr>
<tr>
<td>unsigned short</td>
<td>unsigned fixed bin(16)</td>
</tr>
<tr>
<td>int</td>
<td>fixed bin(31)</td>
</tr>
<tr>
<td>unsigned int</td>
<td>unsigned fixed bin(32)</td>
</tr>
<tr>
<td>long long</td>
<td>fixed bin(63)</td>
</tr>
<tr>
<td>unsigned long long</td>
<td>unsigned fixed bin(64)</td>
</tr>
<tr>
<td>float</td>
<td>float bin(21)</td>
</tr>
<tr>
<td>double</td>
<td>float bin(53)</td>
</tr>
<tr>
<td>long double</td>
<td>float bin(p) (p &gt;= 54)</td>
</tr>
<tr>
<td>enum</td>
<td>ordinal</td>
</tr>
<tr>
<td>typedef</td>
<td>define alias</td>
</tr>
<tr>
<td>struct</td>
<td>define struct</td>
</tr>
<tr>
<td>union</td>
<td>define union</td>
</tr>
<tr>
<td>struct *</td>
<td>handle</td>
</tr>
</tbody>
</table>

Simple type equivalence

So, for example, the following illustrates the translation of the simple typedef for `time_t` from the C header file `time.h`:

```c
typedef long time_t;
define alias time_t fixed bin(31);
```

Figure 58. Simple type equivalence
**Struct type equivalence**

The following example illustrates the translation of the simple struct for `tm` from the C header file `time.h`:

```pli
struct tm {
    int tm_sec;
    int tm_min;
    int tm_hour;
    int tm_mday;
    int tm_mon;
    int tm_year;
    int tm_wday;
    int tm_yday;
    int tm_isdst;
};
```

define structure

```
1 tm
    ,2 tm_sec fixed bin(31)
    ,2 tm_min fixed bin(31)
    ,2 tm_hour fixed bin(31)
    ,2 tm_mday fixed bin(31)
    ,2 tm_mon fixed bin(31)
    ,2 tm_year fixed bin(31)
    ,2 tm_wday fixed bin(31)
    ,2 tm_yday fixed bin(31)
    ,2 tm_isdst fixed bin(31)
```

*Figure 59. Sample struct type equivalence*

**Enum type equivalence**

The following example illustrates the translation of the simple enum `__device_t` from the C header file `stdio.h`:
typedef enum {
    __disk = 0,
    __terminal = 1,
    __printer = 2,
    __tape = 3,
    __tdq = 5,
    __dummy = 6,
    __memory = 8,
    __hfs = 9,
    __hiperspace = 10
} __device_t;

define ordinal __device_t (
    __disk value(0),
    __terminal value(1),
    __printer value(2),
    __tape value(3),
    __tdq value(4),
    __dummy value(5),
    __memory value(8),
    __hfs value(9),
    __hiperspace value(10)
);

Figure 60. Sample enum type equivalence

File type equivalence
A C file declaration depends on the platform, but it often starts as follows:

```c
struct __file {
    unsigned char *__bufPtr;
    ... } FILE;
```

Figure 61. Start of the C declaration for its FILE type

All we want is a pointer (or token) for a file, so we can finesse this translation with:

```c
define struct 1 file;
define alias file_Handle handle file;
```

Figure 62. PL/I equivalent for a C file

Using C functions
Let's say we wanted to write a program to read a file and dump it as formatted hex - using the C functions fopen and fread.

The code for this program is straightforward:
Most of the declarations in the INCLUDE file `filedump` are obvious:

```pli
filedump:
  proc(fn) options(noexecops main);
  dcl fn char(*) var;

%include filedump;

file = fopen( fn, 'rb' );

if file = sysnull() then
  do;
    display( 'file could not be opened' );
    return;
  end;

do forever;
  unspec(buffer) = ''b;
  read_In = fread( addr(buffer), 1, stg(buffer), file );
  if read_In = 0 then
    leave;
  display( heximage(addr(buffer),16,' ') || ' ' || translate(buffer,(32)'.',unprintable) );
  if read_In < stg(buffer) then
    leave;
  end;
  call fclose( file );
end filedump;
```

Figure 63. Sample code to use fopen and fread to dump a file

Most of the declarations in the INCLUDE file `filedump` are obvious:

```pli
define struct 1 file;
define alias file_Handle handle file;

define alias size_t unsigned fixed bin(32);
define alias int signed fixed bin(31);

dcl file type(file_Handle);
dcl read_In fixed bin(31);
dcl buffer char(16);
dcl unprintable char(32) value( substr(collate(),1,32) );
```

Figure 64. Declarations for filedump program

**Matching simple parameter types**

It would be easy to mistranslate the declarations for the C functions. For instance, you could take the following declaration for the C function `fread`:
On some platforms, this would not link successfully because C names are case sensitive. In order to prevent this kind of linker problem, it is best to specify the name in mixed case using the extended form of the external attribute. So, for instance, the declare for *fread* would be better as:

```c
size_t fread( void *,
    size_t,
    size_t,
    FILE *);
```

*Figure 65. C declaration of fread*

and translate it to:

```fortran
dcl fread    ext
    entry( pointer,
        type size_t,
        type size_t,
        type file_Handle )
    returns( type size_t );
```

*Figure 66. First incorrect declaration of fread*

On some platforms, this would not link successfully because C names are case sensitive. In order to prevent this kind of linker problem, it is best to specify the name in mixed case using the extended form of the external attribute. So, for instance, the declare for *fread* would be better as:

```fortran
fread ext('fread')
    entry( pointer,
        type size_t,
        type size_t,
        type file_Handle )
    returns( type size_t );
```

*Figure 67. Second incorrect declaration of fread*

But this wouldn't run right, because while PL/I parameters are *byaddr* by default, C parameters are *byvalue* by default; so we fix that by adding the *byvalue* attribute to the parameters:

```fortran
fread ext('fread')
    entry( pointer byvalue,
        type size_t byvalue,
        type size_t byvalue,
        type file_Handle byvalue )
    returns( type size_t );
```

*Figure 68. Third incorrect declaration of fread*

But note how the return value is set in *Figure 69 on page 346* a fourth parameter (the address of the temporary _temp5) is passed to the function *fread*, which is then expected to place the return code in the integer at that address. This is the convention for how values are returned when the *byaddr* attribute applies to *returns*, and PL/I uses this convention by default.
This wouldn’t run right, because C return values are `byvalue`. So, we fix that with one more `byvalue` attribute.

```
dcl fread  ext('fread')
  entry( pointer byvalue,
         type size_t byvalue,
         type size_t byvalue,
         type file_Handle byvalue )
  returns( type size_t byvalue );
```

Figure 69. Code generated for RETURNS BYADDR

Note how the return value is set now in Figure 71: no extra address is passed, and the return value is simply returned in register 15.

```
*  read_In = fread( addr(buffer), 1, stg(buffer), file );
*
 L  r2,FILE(,r13,176)
 L  r1,fread(,r5,12)
 LA r7,BUFFER(,r13,184)
 L  r15,&EPA_&WSA(,r1,8)
 L  r0,&EPA_&WSA(,r1,12)
 ST r0,CEECAA (,r12,500)
 LA r1,MX_TEMP1(,r13,152)
 ST r8,MX_TEMP1(,r13,152)
 LA r8,1
 ST r8,MX_TEMP1(,r13,156)
 ST r7,MX_TEMP1(,r13,160)
 ST r4,MX_TEMP1(,r13,164)
 ST r2,MX_TEMP1(,r13,168)
 BALR r14,r15
 L  r0, Temp5(,r13,420)
 ST r0,READ_IN(,r13,180)
```

Figure 70. Correct declaration of fread

```
*  read_In = fread( addr(buffer), 1, stg(buffer), file );
*
 L  r2,FILE(,r13,176)
 L  r1,fread(,r5,12)
 LA r7,BUFFER(,r13,184)
 L  r15,&EPA_&WSA(,r1,8)
 L  r0,&EPA_&WSA(,r1,12)
 ST r0,CEECAA (,r12,500)
 LA r1,MX_TEMP1(,r13,152)
 ST r8,MX_TEMP1(,r13,152)
 LA r8,1
 ST r8,MX_TEMP1(,r13,156)
 ST r7,MX_TEMP1(,r13,160)
 ST r4,MX_TEMP1(,r13,164)
 ST r2,MX_TEMP1(,r13,168)
 BALR r14,r15
 LR r0,r15
 ST r0,READ_IN(,r13,180)
```

Figure 71. Code generated for RETURNS BYVALUE
Matching string parameter types

Now that we have translated \texttt{fread} correctly, we might try this translation for \texttt{fopen}:

\begin{verbatim}
dcl fopen  ext('fopen')
  entry( char(*) varyingz byvalue,
         char(*) varyingz byvalue )
  returns( byvalue type file_handle );
\end{verbatim}

\textit{Figure 72. First incorrect declaration of fopen}

But C really has no strings, only pointers, and these pointers would be passed \textit{byvalue}; so the strings should be \textit{byaddr}:

\begin{verbatim}
dcl fopen  ext('fopen')
  entry( char(*) varyingz byaddr,
         char(*) varyingz byaddr )
  returns( byvalue type file_handle );
\end{verbatim}

\textit{Figure 73. Second incorrect declaration of fopen}

But PL/I passes descriptors with strings and C doesn't understand them, so they must be suppressed. We can do this by adding \texttt{options(nodescriptor)} to the declaration:

\begin{verbatim}
dcl fopen  ext('fopen')
  entry( char(*) varyingz byaddr,
         char(*) varyingz byaddr )
  returns( byvalue type file_handle )
  options ( nodescriptor );
\end{verbatim}

\textit{Figure 74. Correct declaration of fopen}

This will work, but isn't optimal since the parameters are input-only; if the parameter is a constant, the \texttt{nonassignable} attribute will prevent a copy being made and passed. Hence, the best translation of the declaration of \texttt{fopen} is:

\begin{verbatim}
dcl fopen  ext('fopen')
  entry( char(*) varyingz nonasgn byaddr,
         char(*) varyingz nonasgn byaddr )
  returns( byvalue type file_handle )
  options ( nodescriptor );
\end{verbatim}

\textit{Figure 75. Optimal, correct declaration of fopen}

At this point, the declare for the \texttt{fclose} function presents few surprises except perhaps for the \texttt{optional} attribute in the returns specification. This attribute allows us to invoke the fclose function via a CALL statement and not have to dispose of the return code. But note that if the file were an output file, the return code on fclose should always be checked since the last buffer may be written out only when the file is closed and that write could fail for lack of space.
Now, on z/OS UNIX, we can compile and run the programs with the commands:

```pli
pli -qdisplay=std filedump.pli
filedump filedump.pli
```

This would produce the following output:

```
15408689 938584A4 94977A40 97999683 . fieldump: proc
40869550 409697A3 899695A2 40959685 (fn) options(noe
A7858396 97A24094 81899550 5E151540 xecops main)...
```

The C quicksort function `qsort` takes a compare routine. For instance, to sort an array of integers, the following function (which uses the `byvalue` attribute twice - for the reasons discussed above) could be used:

```
comp2:
   proc( key, element )
   returns( fixed bin(31) byvalue );

dcl (key, element) pointer byvalue;
dcl word based fixed bin(31);

select;
   when( key->word < element->word )
      return( -1 );
   when( key->word = element->word )
      return( 0 );
   otherwise
      return( +1 );
end;
end;
```

And the C `qsort` function could be used with this compare routine to sort an array of integers, as in the following code fragment:
But since C function pointers are not the same as PL/I ENTRY variables, the C `qsort` function must not be declared as simply:

```
dcl a(1:4) fixed bin(31) init(19,17,13,11);
put skip data( a );
call qsort( addr(a), dim(a), stg(a(1)), comp2 );
put skip data( a );
```

**Figure 80. Sample code to use C qsort function**

But since C function pointers are not the same as PL/I ENTRY variables, the C `qsort` function must not be declared as simply:

```
dcl qsort ext('qsort')
entry( pointer,
      fixed bin(31),
      fixed bin(31),
      entry returns( byvalue fixed bin(31) )
) 
options( byvalue nodescriptor );
```

**Figure 81. Incorrect declaration of qsort**

Recall that a PL/I ENTRY variable may point to a nested function (and thus requires a backchain address as well as an entry point address). But a C function pointer is limited in pointing to a non-nested function only and so, a PL/I ENTRY variable and a C function pointer do not even use the amount of storage.

However, a C function pointer is equivalent to the new PL/I type: a LIMITED ENTRY. and hence the C `qsort` function could be declared as:

```
dcl qsort ext('qsort')
entry( pointer,
      fixed bin(31),
      fixed bin(31),
      limited entry
      returns( byvalue fixed bin(31) )
) 
options( byvalue nodescriptor );
```

**Figure 82. Correct declaration of qsort**

### Linkages

On z/OS, there are two crucial facts about linkages:

- IBM C, JAVA and Enterprise PL/I use the same linkage by default.
- This linkage is not the system linkage.

For a traditional PL/I application where all parameters are `byaddr`, the differences between the code generated when a function has the default linkage and when it has the system linkage would usually not matter. But if the parameters are `byvalue` (as they usually are in C and JAVA), the differences can break your code.
In fact, there is only a small difference if the parameters are byaddr. In Figure 83, the only difference between the code generated for a function with the default linkage and for one with the system linkage is that the high-order bit is turned on for the last parameter of the system linkage call.

This difference would be transparent to most programs.

```
dcl dfta ext entry( fixed bin(31) byaddr
, fixed bin(31) byaddr );
dcl sysa ext entry( fixed bin(31) byaddr
, fixed bin(31) byaddr )
  options( linkage(system) );

* call dfta( n, m );
*  LA r0,M(,r13,172)
  LA r2,N(,r13,168)
  L r15,=V(DFTV)(,r3,126)
  LA r1,#MX_TEMP1(,r13,152)
  ST r2,#MX_TEMP1(,r13,152)
  ST r0,#MX_TEMP1(,r13,156)
  BALR r14,r15

* call sysa( n, m );
*  LA r0,M(,r13,172)
  LA r2,N(,r13,168)
  O r0,=X'80000000'
  L r15,=V(SYSV)(,r3,130)
  LA r1,#MX_TEMP1(,r13,152)
  ST r2,#MX_TEMP1(,r13,152)
  ST r0,#MX_TEMP1(,r13,156)
  BALR r14,r15
```

**Figure 83. Code when parameters are BYADDR**

But, there is a big difference if the parameters are byvalue rather than byaddr. In Figure 84 on page 351 for the function with the default linkage, register 1 points to the values of the integers passed, while for the function with the system linkage, register 1 points to the addresses of those values.

This difference would not be transparent to most programs.
Sharing output and input

The following topics contain the information about sharing output and input with a C program by specifying the STDSYS option.

For more information and limitations about the STDSYS option, see the following manuals:

- “STDSYS” on page 79 in Chapter 1, “Using compiler options and facilities,” on page 5.
- Stream I/O with unprintable characters in Enterprise PL/I for z/OS Compiler and Run-Time Migration Guide.

Sharing output

If you want to share SYSPRINT with a C program, you must compile your PL/I code with the STDSYS option.

By default, DISPLAY statements use WTO's to display their output. If you specify the DISPLAY(STD) compiler option, DISPLAY statements will use the C puts function to display their output. This can be particularly useful under z/OS UNIX.

Behavior of the standard C stream for sharing output under MVS batch, TSO batch, IMS batch, and IMS interactive is as follows:
stdout goes first to DD:SYSPRINT. If DD:SYSPRINT does not exist, stdout looks for DD:SYSTERM and then DD:SYSERR. If neither of these files exist, the library opens a sysout=* data set using DD SYSPRINT and sends the stdout stream to it.

Sharing input
To share SYSIN with a C program, you must compile the application with the STDSYS option and open SYSIN as an input stream file or avoid using the DD names that are reserved by the C Library. You can also copy SYSIN to a temporary data set in a prior job step and use that as SYSIN in your PL/I job step; it can be shared when it is not allocated to an instream file.

An input stream file can be opened only once when it is allocated to the SYSIN DD in JCL.

Behavior of the standard C stream for sharing input under MVS batch, TSO batch, IMS batch, and IMS interactive is as follows:

stdin goes to DD:SYSIN. If DD:SYSIN does not exist, all read operations from stdin fails.

Using the ATTACH statement
The ATTACH statement in a PL/I program uses the underlying pthread library to perform thread management. When the ATTACH statement is issued, the pthread library tries to allocate standard C stream files.

After the ATTACH statement is issued, any attempt to open SYSIN by the application fails when SYSIN is an instream file that is allocated to the SYSIN DD in JCL (for example, //SYSIN DD * or //SYSIN DD DATA). A SYSIN OPEN FAILED error is issued.

Redirecting C standard streams
Use the STDSYS option in all mixed PL/I and C applications and all multithreading applications.

If you compile your program with the NOSTDSYS option, it might conflict with the use of the SYSIN DD name and the SYSPRINT DD name in C.

For example, if your program has the ATTACH statement, it starts the C environment directly. Starting the C environment causes the SYSIN and SYSPRINT streams to be opened and closed independent of your PL/I application programs. In same cases, SYSIN might fail to open or the SYSPRINT data set might be overwritten.

Summary
What we have learned in this chapter:
• C is case sensitive.
• Parameters should be BYVALUE.
• Return values should be BYVALUE.
• String parameters should be BYADDR.
• Arrays and structures should also be BYADDR.
• No descriptors should be passed.
• Input-only strings should be NONASSIGNABLE.
• C function pointers map to LIMITED ENTRYs.
• The IBM C compilers and the IBM PL/I compilers use the same default linkage (and it matters).
Chapter 16. Interfacing with Java

This chapter gives a brief description of Java and the Java Native Interface (JNI) and explains why you might be interested in using it with PL/I. A simple Java-PL/I application will be described and information on compatibility between the two languages will also be discussed. Instructions on how to build and run the Java-PL/I sample applications assume the work is being done in the z/OS UNIX System Services environment of z/OS.

Before you can communicate with Java from PL/I you need to have Java installed on your z/OS system. Contact your local System Administrator for more information on how to set up your z/OS Java environment.

These sample programs have been compiled and tested with Java JRE Version 1.6.0. To determine the level of Java in your z/OS UNIX System Services environment enter this command from the command line:

```
java -version
```

The active Java version is then displayed and looks something like:

```
java version "1.6.0"
Java(TM) SE Runtime Environment (build pmz3160_26sr1-20111114 01(SR1))
IBM J9 VM (build 2.6, JRE 1.6.0 z/OS s390-31 20111113_94967 (JIT enabled, AOT enabled)
```

Java Native Interface (JNI)

Java is an object-oriented programming language invented by Sun Microsystems and provides a powerful way to make Internet documents interactive.

The Java Native Interface (JNI) is the Java interface to native programming languages and is part of the Java Development Kits. By writing programs that use the JNI, you ensure that your code is portable across many platforms.

The JNI allows Java code that runs within a Java Virtual Machine (JVM) to operate with applications and libraries written in other languages, such as PL/I. In addition, the Invocation API allows you to embed a Java Virtual Machine into your native PL/I applications.

Java is a fairly complete programming language; however, there are situations in which you want to call a program written in another programming language. You would do this from Java with a method call to a native language, known as a native method.

Some reasons to use native methods may include the following:

- The native language has a special capability that your application needs and that the standard Java class libraries lack.
- You already have many existing applications in your native language and you wish to make them accessible to a Java application.
- You want to implement a intensive series of complicated calculations in your native language and have your Java applications call these functions.
- You or your programmers have a broader skill set in your native language and you do not wish to lose this advantage.
Programming through the JNI lets you use native methods to do many different operations. A native method can:

- utilize Java objects in the same way that a Java method uses these objects.
- create Java objects, including arrays and strings, and then inspect and use these objects to perform its tasks.
- inspect and use objects created by Java application code.
- update Java objects that it created or were passed to it, and these updated objects can then be made available to the Java application.

Finally, native methods can also easily call already existing Java methods, capitalizing on the functionality already incorporated in the Java programming framework. In these ways, both the native language side and the Java side of an application can create, update, and access Java objects and then share these objects between them.

**Calling PL/I program from Java**

When a PL/I program is called from a Java program, all files opened in PL/I must be explicitly closed in PL/I before PL/I returns control to Java for the last time. Similarly, all modules that are fetched in the PL/I program must be released.

**JNI Sample Program #1 - 'Hello World'**

**Writing Java Sample Program #1**

The first sample program is yet another variation of the "Hello World!" program.

The "Hello World!" program has one Java class, callingPLI.java. The native method, written in PL/I, is contained in hiFromPLI.pli. Here is a brief overview of the steps for creating this sample program:

1. Write a Java program that defines a class containing a native method, loads the native load library, and calls the native method.
2. Compile the Java program to create a Java class.
3. Write a PL/I program that implements the native method and displays the "Hello!" text.
4. Compile and link the PL/I program.
5. Run the Java program which calls the native method in the PL/I program.

**Step 1: Writing the Java Program**

**Declare the Native Method:** All methods, whether Java methods or native methods, must be declared within a Java class. The only difference in the declaration of a Java method and a native method is the keyword native. The native keyword tells Java that the implementation of this method will be found in a native library that will be loaded during the execution of the program. The declaration of the native method looks like this:

```java
public native void callToPLI();
```

In the above statement, the void means that there is no return value expected from this native method call. The empty parentheses in the method name callToPLI( ), means that there are no parameters being passed on the call to the native method.
Load the Native Library: A step that loads the native library must be included so the native library will be loaded at execution time. The Java statement that loads the native library looks like this:

```java
static {
    System.loadLibrary("hiFromPLI");
}
```

In the above statement, the Java System method `System.loadLibrary(...)` is called to find and load the native library. The PL/I shared library, `libhiFromPLI.so`, will be created during the step that compiles and links the PL/I program.

Write the Java Main Method: The `callingPLI` class also includes a `main` method to instantiate the class and call the native method. The `main` method instantiates `callingPLI` and calls the `callToPLI()` native method.

The complete definition of the `callingPLI` class, including all the points addressed above in this section, looks like this:

```java
public class callingPLI {
    public native void callToPLI();
    static {
        System.loadLibrary("hiFromPLI");
    }
    public static void main(String[] argv) {
        callingPLI callPLI = new callingPLI();
        callPLI.callToPLI();
        System.out.println("And Hello from Java, too!");
    }
}
```

Step 2: Compiling the Java Program
Use the Java compiler to compile the `callingPLI` class into an executable form. The command would look like this:

```
javac callingPLI.java
```

Step 3: Writing the PL/I Program
The PL/I implementation of the native method looks much like any other PL/I subroutine.

Useful PL/I Compiler Options: The sample program contains a series of `*PROCESS` statements that define the important compiler options.

```pli
*PROCESS Limits( Extname( 100 ) ) Margins( 1, 100 ) ;
*PROCESS Display(Std) Dllinit Extrn(Short);
*PROCESS Rent Default( ASCII IEEE ) ;
```

Here is a brief description of them and why they are useful:

**Extname(100)**
Allows for longer, Java style, external names.

**Margins(1,100)**
Extending the margins gives you more room for Java style names and identifiers.

**Display(Std)**
Writes the "Hello World" text to stdout, not via WTOs. In the z/OS UNIX environment WTOs would not be seen by the user.

**Dllinit**
Includes the initialization code needed for creating a DLL.
EXTRNs are emitted only for those constants that are referenced. This option is necessary for Enterprise PL/I V3R3 and later.

**Default( ASCII IEEE );**
- **ASCII** specifies that CHARACTER and PICTURE data is held in ASCII - the form in which it is held by JAVA
- **IEEE** specifies that FLOAT data is held in IEEE format - the form in which it is held by JAVA

**RENT**
The RENT option insures that code is reentrant even if it writes on static variables.

**Correct Form of PL/I Procedure Name and Procedure Statement:** The PL/I procedure name must conform to the Java naming convention in order to be located by the Java Class Loader at execution time. The Java naming scheme consists of three parts. The first part identifies the routine to the Java environment, the second part is the name of the Java class that defines the native method, and the third part is the name of the native method itself.

Here is a breakdown of PL/I procedure name `java_callingPLI_callToPLI` in the sample program:

- **Java**
  - All native methods resident in dynamic libraries must begin with `Java`
- `_callingPLI`
  - The name of the Java class that declares the native method
- `_callToPLI`
  - The name of the native method itself.

**Note:** There is an important difference between coding a native method in PL/I and in C. The `javah` tool, which is shipped with the JDK, generates the form of the external references required for C programs. When you write your native methods in PL/I and follow the rules above for naming your PL/I external references, performing the `javah` step is not necessary for PL/I native methods.

In addition, the following options need to be specified in the OPTIONS option on the PROCEDURE statement:
- **FromAlien**
- **NoDescriptor**
- **ByValue**

The complete procedure statement for the sample program looks like this:

```pli
Java_callingPLI_callToPLI:
Proc( JNIEnv , MyJObject )
External( "Java_callingPLI_callToPLI" )
Options( FromAlien NoDescriptor ByValue );
```

**JNI Include File:** The two PL/I include files which contain the PL/I definitions of the Java Native interface are `ibmzjni.inc` which in turn includes `ibmzjnim.inc`. These include files are included with this statement:

```pli
#include ibmzjni;
```

The `ibmzjni` and `ibmzjnim` include files are provided in the PL/I SIBMZSAM data set.
The Complete PL/I Procedure: For completeness, here is the entire PL/I program that defines the native method:

```pli
*Process Limits( Extname(100) ) Margins(1,100) ;
*Process Display(Std) Dllinit Extn(Short);
*Process Rent Default( ASCII IEEE );
PliJava_Demo: Package Exports(*);
```

```pli
Java_callingPLI_callToPLI:
Proc( JNIEnv , MyJObject )
  External( "Java_callingPLI_callToPLI" )
  Options( FromAlien NoDescriptor ByValue );
```

```pli
%include ibmzjni;
Dcl myJObject Type jObject;
Display('Hello from Enterprise PL/I!');
```

End;

**Step 4: Compiling and Linking the PL/I Program**

**Compiling the PL/I Program:** Compile the PL/I sample program with the following command:

```
pli -c hiFromPLI.pli
```

**Linking the Shared Library:** Link the resulting PL/I object deck into a shared library with this command:

```
c89 -o libhiFromPLI.so hiFromPLI.o
```

Be sure to include the `lib` prefix on the name of the PL/I shared library or the Java class loader will not find it.

**Step 5: Running the Sample Program**

Run the Java - PL/I sample program with this command:

```
java callingPLI
```

The output of the sample program will look like this:

```
Hello from Enterprise PL/I!
And Hello from Java, too!
```

The first line written from the PL/I native method. The second line is from the calling Java class after returning from the PL/I native method call.

---

**JNI Sample Program #2 - Passing a String**

**Writing Java Sample Program #2**

This sample program passes a string back and forth between Java and PL/I. See Figure 85 on page 361 for the complete listing of the `jPassString.java` program. The Java portion has one Java class, `jPassString.java`. The native method, written in PL/I, is contained in `passString.pli`. Much of the information from the first sample program applies to this sample program as well. Only new or different aspects will be discussed for this sample program.

**Step 1: Writing the Java Program**

Declare the Native Method: The native method for this sample program looks like this:
public native void pliShowString();

**Load the Native Library:** The Java statement that loads the native library for this sample program looks like this:

```java
static {
    System.loadLibrary("passString");
}
```

**Write the Java Main Method:** The `jPassString` class also includes a `main` method to instantiate the class and call the native method. The `main` method instantiates `jPassString` and calls the `pliShowString()` native method.

This sample program prompts the user for a string and reads that value in from the command line. This is done within a `try/catch` statement as shown in Figure 85 on page 361.
// Read a string, call PL/I, display new string upon return
import java.io.*;

public class jPassString{

  /* Field to hold Java string */
  String myString;

  /* Load the PL/I native library */
  static {
    System.loadLibrary("passString");
  }

  /* Declare the PL/I native method */
  public native void pliShowString();

  /* Main Java class */
  public static void main(String[] arg) {

    System.out.println(" ");

    /* Instantiate Java class and initialize string */
    jPassString myPassString = new jPassString();
    myPassString.myString = " ";

    /* Prompt user for a string */
    try {
      BufferedReader in = new BufferedReader(
        new InputStreamReader(System.in));
      /* Process until 'quit' received */
      while (!myPassString.myString.equalsIgnoreCase("quit")) {
        System.out.println(
          "From Java: Enter a string or 'quit' to quit.");
        System.out.print("Java Prompt > ");
        /* Get string from command line */
        myPassString.myString = in.readLine();
        if (!myPassString.myString.equalsIgnoreCase("quit")) {
          /* Call PL/I native method */
          myPassString.pliShowString();
          /* Return from PL/I and display new string */
          System.out.println(" ");
          System.out.println(
            "From Java: String set by PL/I is: "
            + myPassString.myString );
        }
      }
    } catch (IOException e) {
    }
  }
}

Figure 85. Java Sample Program #2 - Passing a String

Step 2: Compiling the Java Program
The command to compile the Java code would look like this:
javac jPassString.java
Step 3: Writing the PL/I Program
All of the information about writing the PL/I "Hello World" sample program applies to this program as well.

Correct Form of PL/I Procedure Name and Procedure Statement: The PL/I procedure name for this program would be _Java_jPassString_pliShowString_.

The complete procedure statement for the sample program looks like this:
```
Java_jPassString_pliShowString:
  Proc( JNIEnv , myjobject )
  external( "Java_jPassString_pliShowString" )
  Options( FromAlien NoDescriptor ByValue );
```

JNI Include File: The two PL/I include files which contain the PL/I definitions of the Java Native interface are _ibmzjni.inc_ which in turn includes _ibmzjnim.inc_. These include files are included with this statement:
```
%include ibmzjni;
```

The _ibmzjni_ and _ibmzjnim_ include files are provided in the PL/I SIBMZSAM data set.

The Complete PL/I Procedure: The complete PL/I program is shown in Figure 86 on page 363. This sample PL/I program makes several calls through the JNI.

Upon entry, a reference to the calling Java Object, _myObject_ is passed into the PL/I procedure. The PL/I program will use this reference to get information from the calling object. The first piece of information is the Class of the calling object which is retrieved using the _GetObjectClass_ JNI function. This Class value is then used by the _GetFieldID_ JNI function to get the identity of the Java string field in the Java object that we are interested in. This Java field is further identified by providing the name of the field, _myString_, and the JNI field descriptor, _Ljava/lang/String;,_ which identifies the field as a Java String field. The value of the Java string field is then retrieved using the _GetObjectField_ JNI function. Before PL/I can use the Java string value, it must be unpacked into a form that PL/I can understand. The _GetStringUTFChars_ JNI function is used to convert the Java string into a PL/I varyingz string which is then displayed by the PL/I program.

After displaying the retrieved Java string, the PL/I program prompts the user for a PL/I string to be used to update the string field in the calling Java object. The PL/I string value is converted to a Java string using the _NewString_ JNI function. This new Java string is then used to update the string field in the calling Java object using the _SetObjectField_ JNI function.

When the PL/I program ends control is returned to Java, where the newly updated Java string is displayed by the Java program.
**Step 4: Compiling and Linking the PL/I Program**

**Compiling the PL/I Program:** Compile the PL/I sample program with the following command:

```
pli -c passString.pli
```
Linking the Shared Library:  Link the resulting PL/I object deck into a shared library with this command:

```
c89 -o libpassString.so passString.o
```

Be sure to include the `lib` prefix on the name or the PL/I shared library or the Java class loader will not find it.

**Step 5: Running the Sample Program**

Run the Java - PL/I sample program with this command:

```
java jPassString
```

The output of the sample program, complete with the prompts for user input from both Java and PL/I, will look like this:

```
>java jPassString

From Java: Enter a string or 'quit' to quit.
Java Prompt > A string entered in Java

From PLI: String retrieved from Java is: A string entered in Java
From PLI: Enter a string to be returned to Java:
A string entered in PL/I

From Java: String set by PL/I is: A string entered in PL/I
From Java: Enter a string or 'quit' to quit.
Java Prompt > quit
```

---

**JNI Sample Program #3 - Passing an Integer**

**Writing Java Sample Program #3**

This sample program passes an integer back and forth between Java and PL/I. See Figure 87 on page 365 for the complete listing of the `jPassInt.java` program. The Java portion has one Java class, `jPassInt.java`. The native method, written in PL/I, is contained in `passInt.pli`. Much of the information from the first sample program applies to this sample program as well. Only new or different aspects will be discussed for this sample program.

**Step 1: Writing the Java Program**

Declare the Native Method:  The native method for this sample program looks like this:

```
public native void pliShowInt();
```

Load the Native Library:  The Java statement that loads the native library for this sample program looks like this:

```
static {
   System.loadLibrary("passInt");
}
```

Write the Java Main Method:  The `jPassInt` class also includes a `main` method to instantiate the class and call the native method. The `main` method instantiates `jPassInt` and calls the `pliShowInt()` native method.

This sample program prompts the user for an integer and reads that value in from the command line. This is done within a `try/catch` statement as shown in Figure 87 on page 365.
// Read an integer, call PL/I, display new integer upon return
import java.io.*;
import java.lang.*;

public class jPassInt{

/* Fields to hold Java string and int */
int myInt;
String myString;

/* Load the PL/I native library */
static {
    System.loadLibrary("passInt");
}

/* Declare the PL/I native method */
public native void pliShowInt();

/* Main Java class */
public static void main(String[] arg) {
    System.out.println(" ");
    jPassInt pInt = new jPassInt();
pInt.myInt = 1024;
pInt.myString=" ";

/* Prompt user for an integer */
try {
    BufferedReader in = new BufferedReader(
        new InputStreamReader(System.in));
    while (!pInt.myString.equalsIgnoreCase("quit")) {
        System.out.println("From Java: Enter an Integer or 'quit' to quit.");
        System.out.print("Java Prompt > ");
        pInt.myString = in.readLine();
        if (!pInt.myString.equalsIgnoreCase("quit")) {
            pInt.myInt = Integer.parseInt( pInt.myString );
            pInt.pliShowInt();
            System.out.println(" ");
            System.out.println("From Java: Integer set by PL/I is: " + pInt.myInt );
        }
    }
    catch (IOException e) {
        }
}
}

Figure 87. Java Sample Program #3 - Passing an Integer
**Step 2: Compiling the Java Program**
The command to compile the Java code would look like this:
```
javac jPassInt.java
```

**Step 3: Writing the PL/I Program**
All of the information about writing the PL/I "Hello World" sample program applies to this program as well.

**Correct Form of PL/I Procedure Name and Procedure Statement:** The PL/I procedure name for this program would be `Java_jPassInt_pliShowInt`.

The complete procedure statement for the sample program looks like this:
```
Java_passNum_pliShowInt:
  Proc( JNIEnv , myjobject )
  external( "Java_jPassInt_pliShowInt" )
  Options( FromAlien NoDescriptor ByValValue );
```

**JNI Include File:** The two PL/I include files which contain the PL/I definitions of the Java Native interface are `ibmzjni.inc` which in turn includes `ibmzjnim.inc`. These include files are included with this statement:
```
#include ibmzjni;
```

The `ibmzjni` and `ibmzjnim` include files are provided in the PL/I SIBMZSAM data set.

**The Complete PL/I Procedure:** The complete PL/I program is shown in Figure 88 on page 367. This sample PL/I program makes several calls through the JNI.

Upon entry, a reference to the calling Java object, `myObject`, is passed into the PL/I procedure. The PL/I program will use this reference to get information from the calling object. The first piece of information is the Class of the calling object which is retrieved using the `GetObjectClass` JNI function. This Class value is then used by the `GetFieldID` JNI function to get the identity of the Java integer field in the Java object that we are interested in. This Java field is further identified by providing the name of the field, `myInt`, and the JNI field descriptor, `I`, which identifies the field as an integer field. The value of the Java integer field is then retrieved using the `GetIntField` JNI function which is then displayed by the PL/I program.

After displaying the retrieved Java integer, the PL/I program prompts the user for a PL/I integer to be used to update the integer field in the calling Java object. The PL/I integer value is then used to update the integer field in the calling Java object using the `SetIntField` JNI function.

When the PL/I program ends, control is returned to Java, where the newly updated Java integer is displayed by the Java program.
Step 4: Compiling and Linking the PL/I Program

Compiling the PL/I Program:  Compile the PL/I sample program with the following command:

```
pli -c passInt.pli
```

Linking the Shared Library:  Link the resulting PL/I object deck into a shared library with this command:

```
c89 -o libpassInt.so passInt.o
```

Be sure to include the *lib* prefix on the name or the PL/I shared library or the Java class loader will not find it.

---

Figure 88. PL/I Sample Program #3 - Passing an Integer

```pli
*Process Limits( Extname( 100 ) ) Margins( 1, 100 ) ;
*Process Display(Std) Dlimit Extn(Short);
*Process Rent Default( ASCII IEEE );
plijava_demo: package exports(*);

Java_passNum_pliShowInt:
Proc( JNIEnv , myjobject )
  external("Java_jPassInt_pliShowInt")
  Options( FromAlien NoDescriptor ByValue );

%include ibmzjni;
Dcl myClazz Type jClass;
Dcl myFID Type jFieldID;
Dcl myJInt Type jint;
dcl rtnJInt Type jint;
Dcl myJObject Type jobject;
Dcl pliReply Char(132) Varz;
Dcl nullPtr Pointer;

Display(' ');
/* Get information about the calling Class      */
myClazz = GetObjectClass(JNIEnv, myJObject);
/* Get Field ID for int field from Java        */
myFID = GetFieldID(JNIEnv, myClazz, "myInt", "I");
/* Get Integer value from Java                */
myJInt = GetIntField(JNIEnv, myJObject, myFID);
display('From PLI: Integer retrieved from Java is: ' || trim(myJInt ));
display('From PLI: Enter an integer to be returned to Java:
       reply(pliReply);

rtnJInt = pliReply;
/* Set Integer value in Java from PL/I       */
nullPtr = SetIntField(JNIEnv, myJObject, myFID, rtnJInt);

End;
end;
```
Step 5: Running the Sample Program
Run the Java - PL/I sample program with this command:

```
java jPassInt
```

The output of the sample program, complete with the prompts for user input from both Java and PL/I, will look like this:

```
>java jPassInt
From Java: Enter an Integer or 'quit' to quit.
Java Prompt > 12345
From PL/I: Integer retrieved from Java is: 12345
From PL/I: Enter an integer to be returned to Java:
54321
From Java: Integer set by PL/I is: 54321
From Java: Enter an Integer or 'quit' to quit.
Java Prompt > quit
>```

JNI Sample Program #4 - Java Invocation API

Writing Java Sample Program #4
This sample program is a little different from the previous samples. In this sample PL/I invokes Java first, through the Java Invocation API, creating an embedded Java Virtual Machine (JVM). PL/I then calls a Java method passing it a string which the Java method then displays.

The PL/I sample program is named `createJVM.pli` and the Java method it calls is contained in `javaPart.java`.

Step 1: Writing the Java Program
Since this sample does not use a PL/I native method there is no need to declare one. Instead, the Java portion for this sample is just a simple Java method.

Write the Java Main Method: The `javaPart` class contains only one statement. This statement prints out a short 'Hello World...' from Java then appends the string that was passed to it from the PL/I program. The entire class is shown in Figure 89.

```java
// Receive a string from PL/I then display it after saying "Hello"
public class javaPart {
    public static void main(String[] args) {
        System.out.println("From Java - Hello World... "+ args[0]);
    }
}
```

Figure 89. Java Sample Program #4 - Receiving and printing a String

Step 2: Compiling the Java Program
The command to compile the Java code looks like this:

```
javac javaPart.java
```

Step 3: Writing the PL/I Program
Most of the information about writing the PL/I "Hello World" sample program applies to this program as well. However, since in this sample PL/I is calling Java, there are some additional points to consider.
Correct Form of PL/I Procedure Name and Procedure Statement: Since in this sample the PL/I program is calling Java, the PL/I program is MAIN. There is no need to be concerned about the external name of this PL/I program because it is not referenced.

The complete procedure statement for the sample program looks like this:
```
createJVM: Proc Options(Main);
```

JNI Include File: The two PL/I include files which contain the PL/I definitions of the Java Native interface are `ibmzjni.inc` which in turn includes `ibmzjnim.inc`. Even though in this sample PL/I is calling Java, these include files are still necessary. These include files are included with this statement:
```
%include ibmzjni;
```

The `ibmzjni` and `ibmzjnim` include files are provided in the PL/I SIBMZSAM data set.

Linking the PL/I program with the Java library: Since this PL/I sample program calls Java we need to prepare it to link to the Java library. The Java libraries are linked with XPLINK and the PL/I modules are not. PL/I can still link to and call XPLINK libraries but must use the PLIXOPT variable to specify the `XPLINK=ON` run-time option. This is what the declare of the PLIXOPT variable would look like:
```
Dcl PLIXOPT Char(40) Varying Ext Static Init('XPLINK(ON)'e);
```

For a description of using PLIXOPT, see z/OS Language Environment Programming Guide.

Using the Java Invocation API: This PL/I sample program calls the Java Invocation API `JNI_CreateJavaVM` to create its own embedded JVM. This API requires certain structures to be set up and initialized correctly as shown in Figure 90 on page 371. First, `JNI_GetDefaultJavaVMInitArgs` is called to get the default initialization options. Next, these default options are modified with the addition of the `java.class.path` information. Finally, `JNI_CreateJavaVM` is called to create the embedded JVM.

The Complete PL/I Program: The complete PL/I program is shown in Figure 90 on page 371. This sample PL/I program makes several calls through the JNI.

In this sample the reference to the Java object, a newly created JVM in this case, is not passed in but is instead returned from the call to the `JNI_CreateJavaVM` API. The PL/I program will use this reference to get information from the JVM. The first piece of information is the Class containing the Java method we wish to call. This Class is found by using the `FindClass` JNI function. The Class value is then used by the `GetStaticMethodID` JNI function to get the identity of the Java method that will be called.

Before we call this Java method we need to convert the PL/I string into a format that Java understands. The PL/I program holds the string in ASCII format. Java strings are stored in UTF format. In addition, Java strings are not really strings as PL/I programmers understand them but are themselves a Java class and can only be modified through methods. We create a Java string by using the `NewStringUTF` JNI function. This function returns a Java object called `myJString` that contains the PL/I string converted to UTF. Next we create a Java object array by calling the `NewObjectArray` JNI function passing it the reference to the `myJString` object. This function returns a reference to a Java object array containing the string we want the Java method to display.
Now the Java method can be called with the `CallStaticVoidMethod` JNI function and will then display the string passed to it. After displaying the string, the PL/I program destroys the embedded JVM with the `DestroyJavaVM` JNI function and the PL/I program completes.

The complete source of the PL/I program is in [Figure 90 on page 371](#).
%include ibmzjni;

Dcl myJObjArray Type jobjectArray;
Dcl myClass Type jclass;
Dcl myMethodID Type jmethodID;
Dcl myJString Type jstring;
Dcl myRC Fixed Bin(31) Init(0);
Dcl myPLIStr Char(50) Varz
    Init('... a PLI string in a Java Virtual Machine!');
Dcl OptStr1 char(1024) varz;
Dcl OptStr2 char(1024) varz;
Dcl myNull Pointer;
Dcl VM_Args Like JavaVMInitArgs;
Dcl myOptions Like JavaVMOption;
Dcl PLIXOPT Char(40) Varying Ext Static Init( 'XPLINK(ON)'e );

Display('From PL/I - Beginning execution of createJVM...');
VM_Args.version = JNI_VERSION_1_6;
myRC = JNI_GetDefaultJavaVMInitArgs( addr(VM_Args) );
OptStr1 = "-Djava.class.path=.:";
OptStr2 = "-Djava.compiler=NONE";

myOptions(1).theOptions = addr(OptStr1);
myOptions(2).theOptions = addr(OptStr2);
VM_Args.nOptions = 2;
VM_Args.JavaVMOption = addr(myOptions);

/* Create the Java VM */
myrc = JNI_CreateJavaVM(
    addr(jvm_ptr),
    addr(JNIEnv),
    addr(VM_Args) );

/* Get the Java Class for the javaPart class */
myClass = FindClass(JNIEnv, "javaPart");
/* Get static method ID */
myMethodID = GetStaticMethodID(JNIEnv, myClass,"main",
    "([Ljava/lang/String;)V" );
/* Create a Java String Object from the PL/I string. */
myJString = NewStringUTF(JNIEnv, myPLIStr);
myJObjArray = NewObjectArray(JNIEnv, 1,
    FindClass(JNIEnv,"java/lang/String"), myJString);
Display('From PL/I - Calling Java method in new JVM from PL/I...');
Display(' ');
myNull = CallStaticVoidMethod(JNIEnv, myClass,
    myMethodID, myJObjArray );
/* destroy the Java VM */
Display(' '); Display('From PL/I - Destroying the new JVM from PL/I...');
myRC = DestroyJavaVM( JavaVM );
end;

Figure 90. PL/I Sample Program #4 - Calling the Java Invocation API
Step 4: Compiling and Linking the PL/I Program

Compiling the PL/I Program:  Compile the PL/I sample program with the following command:

```
pli -c createJVM.pli
```

Linking the PL/I program to the Java Library:  Link the resulting PL/I object deck with the Java library with this command:

```
c89 -o createJVM createJVM.o $JAVA_HOME/bin/classic/libjvm.x
```

Notice the reference to the $JAVA_HOME environment variable. This variable should point to the directory that your Java 1.4 product is installed in. For example, to set up this variable in your environment you would use the following command:

```
export JAVA_HOME="/usr/lpp/java/J6.0"
```

In this case the Java 1.4 product is assumed to be installed in /usr/lpp/java/J6.0.

Step 5: Running the Sample Program

Run the PL/I - Java sample program with this command:

```
createJVM
```

The output of the sample program will look like this:

```from PL/I - Beginning execution of createJVM...
from PL/I - Calling Java method in new JVM from PL/I...
from Java - Hello World... ... a PLI string in a Java Virtual Machine!
from PL/I - Destroying the new JVM from PL/I...
```

Attaching programs to an existing Java VM

You can use the Java Invocation API to attach your program to an existing Java VM that was not created by your program.

When your PL/I application is running within an IMS JMP (Java messaging processing) region, IMS already created a Java VM. You can take the following steps to attach your program to the Java VM. Then you can call any functions you need through the Java native interface.

1. Locate the Java VM instance to attach your program to by using the JNI_GetCreatedJavaVMS API. In this IMS environment, only one Java VM is created, so you ask for only one Java VM pointer to be returned. You can code the call as follows:

   ```
   rc = JNI_GetCreatedJavaVMS( jvm_ptr, 1, nVMs );
   ```

   `jvm_ptr`  Is a pointer to a Java VM and is declared in the IBMZJNI include file. Java returns the address of the Java VM in this variable when the call returns successfully.

   `nVMs`  Is a fixed bin(31) variable that Java updates with the number of Java VM addresses when it returns. `nVMs` is 1 if the call is successful.

2. Acquire JNI environment pointer for the Java VM. You can use the JGetEnv function in the JNIInvokeInterface_structure that is declared in the IBMZJNI include file:

   ```
   rc = JGetEnv( jvm_ptr, JNIEnv, JNI_VERSION_1_6 );
   ```

   `jvm_ptr`  Is the pointer to the Java VM instance.
JNIEnv

Is a pointer that Java sets to the JNI environment pointer for the Java VM instance.

JNI_VERSION_1_6

Is a constant holding the value of an interface version number which is also declared in the IBMZJNI include file.

Now that you have the JNI environment pointer for the Java VM, you can make calls to any of the JNI functions through the Java native interface.

---

### Determining equivalent Java and PL/I data types

When you communicate with Java from PL/I you will need to match the data types between the two programming languages. This table shows Java primitive types and their PL/I equivalents:

<table>
<thead>
<tr>
<th>Java Type</th>
<th>PL/I Type</th>
<th>Size in Bits</th>
</tr>
</thead>
<tbody>
<tr>
<td>Boolean</td>
<td>jboolean</td>
<td>8, unsigned</td>
</tr>
<tr>
<td>byte</td>
<td>jbyte</td>
<td>8</td>
</tr>
<tr>
<td>char</td>
<td>jchar</td>
<td>16, unsigned</td>
</tr>
<tr>
<td>short</td>
<td>jshort</td>
<td>16</td>
</tr>
<tr>
<td>int</td>
<td>jint</td>
<td>32</td>
</tr>
<tr>
<td>long</td>
<td>jlong</td>
<td>64</td>
</tr>
<tr>
<td>float</td>
<td>jfloat</td>
<td>21</td>
</tr>
<tr>
<td>double</td>
<td>jdouble</td>
<td>53</td>
</tr>
<tr>
<td>void</td>
<td>jvoid</td>
<td>n/a</td>
</tr>
</tbody>
</table>

---

Chapter 16. Interfacing with Java 373
Part 5. Specialized programming tasks

Chapter 17. Using the PLISAXA and PLISAXB XML parsers

Overview .......................................................... 377
The PLISAXA built-in subroutine ......................... 378
The PLISAXB built-in subroutine ......................... 378
The SAX event structure .................................... 379
  start_of_document ......................................... 379
  version_information ...................................... 380
  encoding_declaration .................................... 380
  standalone_declaration .................................. 380
  document_type_declaration ............................. 380
  end_of_document ......................................... 380
  start_of_element ........................................ 380
  attribute_name .......................................... 380
  attribute_characters ................................. 381
  attribute_predefined_reference ...................... 381
  attribute_character_reference ....................... 381
  end_of_element ......................................... 381
  start_of_CDATA_section ............................... 382
  end_of_CDATA_section ................................ 382
  content_characters .................................... 382
  content_predefined_reference ....................... 382
  content_character_reference ......................... 382
  processing_instruction ............................... 382
  comment .................................................. 382
  unknown_attribute_reference ......................... 382
  unknown_content_reference ............................ 382
  start_of_prefix_mapping .............................. 382
  end_of_prefix_mapping ................................ 382
  exception ................................................ 382
Parameters to the event functions .................... 382
  namespace_declare .................................... 382
  end_of_input ............................................ 382
  unresolved_reference .................................. 382
  exception ................................................ 382
  Parameters to the event functions .................. 382
  Differences in the events .............................. 382
Coded character sets for XML documents .................. 382
  Supported EBCDIC code pages ......................... 383
  Supported ASCII code pages ......................... 384
  Specifying the code page ............................. 384
Using a number ............................................. 385
Using an alias .............................................. 385
Exceptions .................................................. 385
  Parsing XML documents with validation .................. 385
  Creating an OSR ......................................... 386
Example with a simple document ..................... 386
  Example of using the PLISAXC built-in subroutine .... 389
  Example of using the PLISAXD built-in subroutine .... 429

Chapter 18. Using the PLISAXC and PLISAXD XML parsers

Overview .......................................................... 407
The PLISAXC built-in subroutine ......................... 408
The PLISAXD built-in subroutine ......................... 408
The SAX event structure .................................... 409
  start_of_document ......................................... 409
  version_information ...................................... 410
  encoding_declaration .................................... 410
  standalone_declaration .................................. 410
  document_type_declaration ............................. 410
  end_of_document ......................................... 410
  start_of_element ........................................ 410
  attribute_name .......................................... 410
  attribute_characters ................................. 410
  attribute_predefined_reference ...................... 410
  attribute_character_reference ....................... 410
  end_of_element ......................................... 410
  start_of_CDATA_section ............................... 411
  end_of_CDATA_section ................................ 411
  content_characters .................................... 411
  content_predefined_reference ....................... 411
  content_character_reference ......................... 411
  processing_instruction ............................... 411
  comment .................................................. 411
  namespace_declare .................................... 412
  end_of_input ............................................ 412
  unresolved_reference .................................. 412
  exception ................................................ 412
  Parameters to the event functions .................. 412
  Differences in the events .............................. 412
Coded character sets for XML documents .................. 412
  Supported code pages .................................. 413
  Specifying the code page ............................. 413
Using a number ............................................. 414
Using an alias .............................................. 414
Exceptions .................................................. 414
  Parsing XML documents with validation .................. 414
  Creating an OSR ......................................... 415
Example with a simple document ..................... 415
  Example of using the PLISAXC built-in subroutine .... 418
  Example of using the PLISAXD built-in subroutine .... 429

Chapter 19. Using PLIDUMP

PLIDUMP usage notes .......................................... 441
Locating variables in the PLIDUMP output .................. 441
  Locating AUTOMATIC variables ....................... 441
  Locating STATIC variables ............................ 442
  Locating CONTROLLED variables ..................... 442
  Under NORENT WRITABLE .............................. 443
  Under NORENT NOWRITABLE(FWS) ...................... 443
  Under NORENT NOWRITABLE(PRV) ....................... 444
Saved Compilation Data .................................... 444
  Copyright ............................................... 444
  Timestamp ............................................... 445
  Saved options string .................................. 445

Chapter 20. Interrupts and attention processing

Using ATTENTION ON-units ................................ 455
Interaction with a debugging tool ...................... 455

Chapter 21. Using the Checkpoint/Restart facility

Requesting a checkpoint record ....................... 457
  Defining the checkpoint data set .................... 457
Requesting a restart ...................................... 458
  Automatic restart after a system failure ............ 458
  Automatic restart within a program ................. 458
Getting a deferred restart .............................. 459
  Modifying checkpoint/restart activity ............... 459

© Copyright IBM Corp. 1999, 2015
Chapter 17. Using the PLISAXA and PLISAXB XML parsers

The PLISAXx \( (x = A \text{ or } B) \) built-in subroutines provide basic XML parsing capability which allows programs to consume inbound XML documents, check them for well-formedness, and react to their contents.

These subroutines do not provide XML generation, which must instead be accomplished by PL/I program logic or by using the XMLCHAR built-in function.

PLISAXA and PLISAXB have no special environmental requirements. They execute in all the principal runtime environments, including CICS, IMS, MQ Series, z/OS batch and TSO.

PLISAXA and PLISAXB do have some important limits: they have no support for XML name spaces, no support for Unicode UTF-8 documents, and they require that the entire XML document be passed to them (either in a buffer or a file) before they do any parsing of it. The PLISAXC and PLISAXD built-in subroutines do not have these limits; the next chapter discusses PLISAXC and PLISAXD in detail.

Overview

There are two major types of interfaces for XML parsing: event-based and tree-based.

For an event-based API, the parser reports events to the application through callbacks. Such events include the start of the document, the beginning of an element, and so on. The application provides handlers to deal with the events reported by the parser. The Simple API for XML or SAX is an example of an industry-standard event-based API.

For a tree-based API (such as the Document Object Model or DOM), the parser translates the XML into an internal tree-based representation. Interfaces are provided to navigate the tree.

IBM PL/I provides a SAX-like event-based interface for parsing XML documents. The parser invokes an application-supplied handler for parser events, passing references to the corresponding document fragments.

The parser has the following characteristics:

- It provides high-performance, but non-standard interfaces.
- It supports XML files encoded in either Unicode UTF-16 or any of several single-byte code pages listed in the section "Coded character sets for XML documents" on page 383.
- The parser is non-validating, but does partially check well-formedness.

XML documents have two levels of conformance: well-formedness and validity, both of which are defined in the XML standard, which you can find at http://www.w3c.org/XML/. Recapitulating these definitions, an XML document is well-formed if it complies with the basic XML grammar, and with a few specific rules, such as the requirement that the names on start and end element tags must
match. A well-formed XML document is also valid if it has an associated document type declaration (DTD) and if it complies with the constraints expressed in the DTD.

The XML parser is non-validating, but does partially check for well-formedness errors, and generates exception events if it discovers any.

For each parser event, you must provide a PL/I function that accepts the appropriate parameters and returns the appropriate return value - as in the example code below. Note in particular that the return value must be returned BYVALUE. Also, these functions must all use the OPTLINK linkage. You can use the DEFAULT(LINKAGE(OPTLINK)) option to specify this linkage, or you can specify it on the individual PROCEDUREs and ENTRYs via the OPTIONS(LINKAGE(OPTLINK)) attribute.

The PLISAXA built-in subroutine

The PLISAXA built-in subroutine allows you to invoke the XML parser for an XML document residing in a buffer in your program.

\[
\text{PLISAXA}(e, p, x, n, c)
\]

- \(e\): An event structure
- \(p\): A pointer value or "token" that the parser will pass back to the event functions
- \(x\): The address of the buffer containing the input XML
- \(n\): The number of bytes of data in that buffer
- \(c\): A numeric expression specifying the purported codepage of that XML

Note that if the XML is contained in a CHARACTER VARYING or a WIDECHAR VARYING string, then the ADDRDATA built-in function should be used to obtain the address of the first data byte.

Also note that if the XML is contained in a WIDECHAR string, the value for the number of bytes is twice the value returned by the LENGTH built-in function.

The PLISAXB built-in subroutine

The PLISAXB built-in subroutine allows you to invoke the XML parser for an XML document residing in a file.

\[
\text{PLISAXB}(e, p, x, c)
\]

- \(e\): An event structure
- \(p\): A pointer value or "token" that the parser will pass back to the event functions
- \(x\): A character string expression specifying the input file
- \(c\): A numeric expression specifying the purported codepage of that XML
Under batch, the character string specifying the input file should have the form 'file://dd:ddname', where ddname is the name of the DD statement specifying the file.

Under z/OS UNIX, the character string specifying the input file should have the form 'file://filename', where filename is the name of a z/OS UNIX file.

Under both batch and z/OS UNIX, the character string specifying the input file should have no leading or trailing blanks.

The input XML file must be less than 2G in size. Moreover, since the parser will read the entire file into memory, the REGION for your program must be large enough for the parser to obtain a piece of storage that can contain all of the document.

### The SAX event structure

The event structure is a structure consisting of 24 LIMITED ENTRY variables which point to functions that the parser will invoke for various "events".

All these ENTRYs must use the OPTLINK linkage.

The library routines that support the XML parsers recognize if an event in any of the PLISAX event structures is set to null (via the NULLENTRY built-in function or via use of UNSPEC) and do not call it in that case. This allows you to limit your XML parsing code to only the events in which you are interested and to improve the performance of the overall parse at the same time.

The descriptions below of each event refer to the example of an XML document in Figure 91. In these descriptions, the term "XML text" refers to the string based on the pointer and length passed to the event.

```xml
xmlDocument =
'<?xml version="1.0" standalone="yes"?>'
|-- '<!--This document is just an example-->'
|-- '<sandwich>'
|-- '<bread type="baker\'s best"/>
|-- '<?spread please use real mayonnaise ?>'
|-- '<meat>Ham & turkey</meat>'
|-- '<filling>Cheese, lettuce, tomato, etc.</filling>'
|-- '<![CDATA[We should add a <relish> element in future!]]>'
|-- '</sandwich>'
|-- 'junk';
```

**Figure 91. Sample XML document**

In the order of their appearance in this structure, the parser may recognize the following events:

**start_of_document**

This event occurs once, at the beginning of parsing the document. The parser passes the address and length of the entire document, including any line-control characters, such as LF (Line Feed) or NL (New Line). For the above example, the document is 305 characters in length.
version_information
This event occurs within the optional XML declaration for the version information. The parser passes the address and length of the text containing the version value, "1.0" in the example above.

encoding_declaration
This event occurs within the XML declaration for the optional encoding declaration. The parser passes the address and length of the text containing the encoding value.

standalone_declaration
This event occurs within the XML declaration for the optional standalone declaration. The parser passes the address and length of the text containing the standalone value, "yes" in the example above.

document_type_declaration
This event occurs when the parser finds a document type declaration. Document type declarations begin with the character sequence "<!DOCTYPE" and end with a ">" character, with some fairly complicated grammar rules describing the content in between. The parser passes the address and length of the text containing the entire declaration, including the opening and closing character sequences, and is the only event where XML text includes the delimiters. The example above does not have a document type declaration.

derm_of_document
This event occurs once, when document parsing has completed.

start_of_element
This event occurs once for each element start tag or empty element tag. The parser passes the address and length of the text containing the element name. For the first start_of_element event during parsing of the example, this would be the string "sandwich".

attribute_name
This event occurs for each attribute in an element start tag or empty element tag, after recognizing a valid name. The parser passes the address and length of the text containing the attribute name. The only attribute name in the example is "type".

attribute_characters
This event occurs for each fragment of an attribute value. The parser passes the address and length of the text containing the fragment. An attribute value normally consists of a single string only, even if it is split across lines:

```
<element attribute="This attribute value is split across two lines"/>
```

The attribute value might consist of multiple pieces, however. For instance, the value of the "type" attribute in the "sandwich" example at the beginning of the section consists of three fragments: the string "baker", the single character "'", and the string "'s best". The parser passes these fragments as three separate events. It
passes each string, "baker" and "s best" in the example, as attribute_characters events, and the single character "" as an attribute_predefined_reference event, described next.

**attribute_predefined_reference**

This event occurs in attribute values for the five pre-defined entity references ",&", ",", ",">", ",<" and "". The parser passes a CHAR(1) or WIDECHAR(1) value that contains one of ",&", ",", ",">", ",<" or ",", respectively.

**attribute_character_reference**

This event occurs in attribute values for numeric character references (Unicode code points or "scalar values") of the form ",&#dd;" or ",&#xhh;", where "d" and "h" represent decimal and hexadecimal digits, respectively. The parser passes a FIXED BIN(31) value that contains the corresponding integer value.

**end_of_element**

This event occurs once for each element end tag or empty element tag when the parser recognizes the closing angle bracket of the tag. The parser passes the address and length of the text containing the element name.

**start_of_CDATA_section**

This event occurs at the start of a CDATA section. CDATA sections begin with the string "<![CDATA[" and end with the string "]]>", and are used to "escape" blocks of text containing characters that would otherwise be recognized as XML markup. The parser passes the address and length of the text containing the opening characters "<![CDATA[". The parser passes the content of a CDATA section between these delimiters as a single content-characters event. In the preceding example, the content-characters event is passed the text "We should add a <relish> element in future!".

**end_of_CDATA_section**

This event occurs when the parser recognizes the end of a CDATA section. The parser passes the address and length of the text containing the closing character sequence, "]]>".

**content_characters**

This event represents the "heart" of an XML document: the character data between element start and end tags. The parser passes the address and length of the text containing the this data, which usually consists of a single string only, even if it is split across lines:

```xml
<element1>This character content is split across two lines</element1>
```

If the content of an element includes any references or other elements, the complete content may comprise several segments. For instance, the content of the "meat" element in the example consists of the string "Ham ", the character "&" and the string ", turkey". Notice the trailing and leading spaces, respectively, in these two string fragments. The parser passes these three content fragments as separate events. It passes the string content fragments, "Ham " and ", turkey", as content_characters events, and the single "&" character as a content_predefined_reference event. The parser also uses the content_characters event to pass the text of CDATA sections to the application.
**content_predefined_reference**

This event occurs in element content for the five pre-defined entity references "&", "'", ">", "<" and "". The parser passes a CHAR(1) or WIDECHAR(1) value that contains one of "&", "'", ">", "" or "", respectively.

**content_character_reference**

This event occurs in element content for numeric character references (Unicode code points or 'scalar values') of the form "&#dd;" or "&#xhh;", where "d" and "h" represent decimal and hexadecimal digits, respectively. The parser passes a FIXED BIN(31) value that contains the corresponding integer value.

**processing_instruction**

Processing instructions (PIs) allow XML documents to contain special instructions for applications. This event occurs when the parser recognizes the name following the PI opening character sequence, "<?". The event also covers the data following the processing instruction (PI) target, up to but not including the PI closing character sequence, ">">. Trailing, but not leading white space characters in the data are included. The parser passes the address and length of the text containing the target, "spread" in the example, and the address and length of the text containing the data, "please use real mayonnaise" in the example.

**comment**

This event occurs for any comments in the XML document. The parser passes the address and length of the text between the opening and closing comment delimiters, "<!--" and "-->", respectively. In the example, the text of the only comment is "This document is just an example".

**unknown_attribute_reference**

This event occurs within attribute values for entity references other than the five pre-defined entity references, listed for the event attribute_predefined_character. The parser passes the address and length of the text containing the entity name.

**unknown_content_reference**

This event occurs within element content for entity references other than the five pre-defined entity references listed for the content_predefined_character event. The parser passes the address and length of the text containing the entity name.

**start_of_prefix_mapping**

This event is currently not generated.

**end_of_prefix_mapping**

This event is currently not generated.

**exception**

The parser generates this event when it detects an error in processing the XML document.

**Parameters to the event functions**

All of these functions must return a BYVALUE FIXED BIN(31) value that is a return code to the parser. For the parser to continue normally, this value should be zero.
All of these functions will be passed as the first argument a BYVALUE POINTER that is the token value passed originally as the second argument to the built-in function.

With the following exceptions, all of the functions will also be passed a BYVALUE POINTER and a BYVALUE FIXED BIN(31) that supply the address and length of the text element for the event. The following functions and events are different:

**end_of_document**
- No argument other than the user token is passed.

**attribute_predefined_reference**
- In addition to the user token, one additional argument is passed: a BYVALUE CHAR(1) or, for a UTF-16 document, a BYVALUE WIDECHAR(1) that holds the value of the predefined character.

**content_predefined_reference**
- In addition to the user token, one additional argument is passed: a BYVALUE CHAR(1) or, for a UTF-16 document, a BYVALUE WIDECHAR(1) that holds the value of the predefined character.

**attribute_character_reference**
- In addition to the user token, one additional argument is passed: a BYVALUE FIXED BIN(31) that holds the value of the numeric reference.

**content_character_reference**
- In addition to the user token, one additional argument is passed: a BYVALUE FIXED BIN(31) that holds the value of the numeric reference.

**processing_instruction**
- In addition to the user token, four additional arguments are passed:
  1. a BYVALUE POINTER that is the address of the target text
  2. a BYVALUE FIXED BIN(31) that is the length of the target text
  3. a BYVALUE POINTER that is the address of the data text
  4. a BYVALUE FIXED BIN(31) that is the length of the data text

**exception**
- In addition to the user token, three additional arguments are passed:
  1. a BYVALUE POINTER that is the address of the offending text
  2. a BYVALUE FIXED BIN(31) that is the byte offset of the offending text within the document
  3. a BYVALUE FIXED BIN(31) that is the value of the exception code

---

**Coded character sets for XML documents**

The PLISAX built-in subroutine supports only XML documents in WIDECHAR encoded using Unicode UTF-16, or in CHARACTER encoded using one of the explicitly supported single-byte character sets listed below. The parser uses up to three sources of information about the encoding of your XML document, and signals an exception XML event if it discovers any conflicts between these sources:

1. The parser determines the basic encoding of a document by inspecting its initial characters.
2. If step 1 succeeds, the parser then looks for any encoding declaration.
3. Finally, it refers to the codepage value on the PLISAX built-in subroutine call. If this parameter was omitted, it defaults to the value provided by the CODEPAGE compiler option value that you specified explicitly or by default.
If the XML document begins with an XML declaration that includes an encoding declaration specifying one of the supported code pages listed below, the parser honors the encoding declaration if it does not conflict with either the basic document encoding or the encoding information from the PLISAX built-in subroutine. If the XML document does not have an XML declaration at all, or if the XML declaration omits the encoding declaration, the parser uses the encoding information from the PLISAX built-in subroutine to process the document, as long as it does not conflict with the basic document encoding.

**Supported EBCDIC code pages**

In the following table, the first number is for the Euro Country Extended Code Page (ECECP), and the second is for Country Extended Code Page (CECP).

<table>
<thead>
<tr>
<th>CCSID</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>01047</td>
<td>Latin 1 / Open Systems</td>
</tr>
<tr>
<td>01140, 00037</td>
<td>USA, Canada, etc.</td>
</tr>
<tr>
<td>01141, 00273</td>
<td>Austria, Germany</td>
</tr>
<tr>
<td>01142, 00277</td>
<td>Denmark, Norway</td>
</tr>
<tr>
<td>01143, 00278</td>
<td>Finland, Sweden</td>
</tr>
<tr>
<td>01144, 00280</td>
<td>Italy</td>
</tr>
<tr>
<td>01145, 00284</td>
<td>Spain, Latin America (Spanish)</td>
</tr>
<tr>
<td>01146, 00285</td>
<td>UK</td>
</tr>
<tr>
<td>01147, 00297</td>
<td>France</td>
</tr>
<tr>
<td>01148, 00500</td>
<td>International</td>
</tr>
<tr>
<td>01149, 00871</td>
<td>Iceland</td>
</tr>
</tbody>
</table>

**Supported ASCII code pages**

<table>
<thead>
<tr>
<th>CCSID</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>00813</td>
<td>ISO 8859-7 Greek / Latin</td>
</tr>
<tr>
<td>00819</td>
<td>ISO 8859-1 Latin 1 / Open Systems</td>
</tr>
<tr>
<td>00920</td>
<td>ISO 8859-9 Latin 5 (ECMA-128, Turkey TS-5881)</td>
</tr>
</tbody>
</table>

**Specifying the code page**

If your document does not include an encoding declaration in the XML declaration, or does not have an XML declaration at all, the parser uses the encoding information provided by the PLISAX built-in subroutine call in conjunction with the basic encoding of the document.

You can also specify the encoding information for the document in the XML declaration, with which most XML documents begin. An example of an XML declaration that includes an encoding declaration is:

```xml
<?xml version="1.0" encoding="ibm-1140"?>
```

If your XML document includes an encoding declaration, ensure that it is consistent with the encoding information provided by the PLISAX built-in subroutine and with the basic encoding of the document. If there is any conflict
between the encoding declaration, the encoding information provided by the PLISAX built-in subroutine and the basic encoding of the document, the parser signals an exception XML event.

Specify the encoding declaration as follows:

**Using a number**

You can specify the CCSID number (with or without any number of leading zeroes), prefixed by any of the following (in any mixture of upper or lower case):

<table>
<thead>
<tr>
<th>IBM_</th>
<th>CP</th>
<th>CCSID_</th>
</tr>
</thead>
<tbody>
<tr>
<td>IBM-</td>
<td>CP_</td>
<td>CCSID-</td>
</tr>
</tbody>
</table>

**Using an alias**

You can use any of the following supported aliases (in any mixture of lower and upper case):

<table>
<thead>
<tr>
<th>Code page</th>
<th>Supported aliases</th>
</tr>
</thead>
<tbody>
<tr>
<td>037</td>
<td>EBCDIC-CP-US, EBCDIC-CP-CA, EBCDIC-CP-WT, EBCDIC-CP-NL</td>
</tr>
<tr>
<td>500</td>
<td>EBCDIC-CP-BE, EBCDIC-CP-CH</td>
</tr>
<tr>
<td>813</td>
<td>ISO-8859-7, ISO_8859-7</td>
</tr>
<tr>
<td>819</td>
<td>ISO-8859-1, ISO_8859-1</td>
</tr>
<tr>
<td>920</td>
<td>ISO-8859-9, ISO_8859-9</td>
</tr>
<tr>
<td>1200</td>
<td>UTF-16</td>
</tr>
</tbody>
</table>

**Exceptions**

For most exceptions, the XML text contains the part of the document that was parsed up to and including the point where the exception was detected. For encoding conflict exceptions, which are signaled before parsing begins, the length of the XML text is either zero or the XML text contains just the encoding declaration value from the document. The example above contains one item that causes an exception event, the superfluous “junk” following the "sandwich" element end tag.

There are two kinds of exceptions:

1. Exceptions that allow you to continue parsing optionally. Continuable exceptions have exception codes in the range 1 through 99, 100,001 through 165,535, or 200,001 to 265,535. The exception event in the example above has an exception number of 1 and thus is continuable.

2. Fatal exceptions, which don't allow continuation. Fatal exceptions have exception codes greater than 99 (but less than 100,000).

Returning from the exception event function with a non-zero return code normally causes the parser to stop processing the document, and return control to the program that invoked the PLISAXA or PLISAXB built-in subroutine.

For continuable exceptions, returning from the exception event function with a zero return code requests the parser to continue processing the document, although
further exceptions might subsequently occur. See “Continuable exception codes” on page 398 for details of the actions that the parser takes when you request continuation.

A special case applies to exceptions with exception numbers in the ranges 100,001 through 165,535 and 200,001 through 265,535. These ranges of exception codes indicate that the document’s CCSID (determined by examining the beginning of the document, including any encoding declaration) is not identical to the CCSID value provided (explicitly or implicitly) by the PLISAXA or PLISAXB built-in subroutine, even if both CCSIDs are for the same basic encoding, EBCDIC or ASCII.

For these exceptions, the exception code passed to the exception event contains the document’s CCSID, plus 100,000 for EBCDIC CCSIDs, or 200,000 for ASCII CCSIDs. For instance, if the exception code contains 101,140, the document’s CCSID is 01140. The CCSID value provided by the PLISAXA or PLISAXB built-in subroutine is either set explicitly as the last argument on the call or implicitly when the last argument is omitted and the value of the CODEPAGE compiler option is used.

Depending on the value of the return code after returning from the exception event function for these CCSID conflict exceptions, the parser takes one of three actions:

1. If the return code is zero, the parser proceeds using the CCSID provided by the built-in subroutine.
2. If the return code contains the document’s CCSID (that is, the original exception code value minus 100,000 or 200,000), the parser proceeds using the document’s CCSID. This is the only case where the parser continues after a non-zero value is returned from one of the parsing events.
3. Otherwise, the parser stops processing the document, and returns control to the PLISAXA or PLISAXB built-in subroutine which will raise the ERROR condition.

Example

The following example illustrates the use of the PLISAXA built-in subroutine and uses the example XML document cited above:
saxtest: package exports(saxtest);

define alias event
    limited entry( pointer, pointer, fixed bin(31) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_end_of_document
    limited entry( pointer )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_predefined_ref
    limited entry( pointer, char(1) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) nodescriptor );

define alias event_character_ref
    limited entry( pointer, fixed bin(31) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_pi
    limited entry( pointer, pointer, fixed bin(31),
                  pointer, fixed bin(31) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_exception
    limited entry( pointer, pointer, fixed bin(31),
                   fixed bin(31) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

---

Figure 92. PLISAXA coding example - type declarations
saxtest: proc options( main );

dcl
  1 eventHandler static

    ,2 e01 type event
    init( start_of_document )
    ,2 e02 type event
    init( version_information )
    ,2 e03 type event
    init( encoding_declaration )
    ,2 e04 type event
    init( standalone_declaration )
    ,2 e05 type event
    init( document_type_declaration )
    ,2 e06 type event_end_of_document
    init( end_of_document )
    ,2 e07 type event
    init( start_of_element )
    ,2 e08 type event
    init( attribute_name )
    ,2 e09 type event
    init( attribute_characters )
    ,2 e10 type event_predefined_ref
    init( attribute_predefined_reference )
    ,2 e11 type event_character_ref
    init( attribute_character_reference )
    ,2 e12 type event
    init( end_of_element )
    ,2 e13 type event
    init( start_of_CDATA )
    ,2 e14 type event
    init( end_of_CDATA )
    ,2 e15 type event
    init( content_characters )
    ,2 e16 type event_predefined_ref
    init( content_predefined_reference )
    ,2 e17 type event_character_ref
    init( content_character_reference )
    ,2 e18 type event_pi
    init( processing_instruction )
    ,2 e19 type event
    init( comment )
    ,2 e20 type event
    init( unknown_attribute_reference )
    ,2 e21 type event
    init( unknown_content_reference )
    ,2 e22 type event
    init( start_of_prefix_mapping )
    ,2 e23 type event
    init( end_of_prefix_mapping )
    ,2 e24 type event_exception
    init( exception )
    ;

Figure 93. PLISAXA coding example - event structure
dcl token char(8);

dcl xmlDocument char(4000) var;

xmlDocument = 
  '<?xml version="1.0" standalone="yes"?>' 
  '<!--This document is just an example-->' 
  '<sandwich>' 
  '<bread type="baker\’s best" />' 
  '<?spread please use real mayonnaise ?>' 
  '<meat>Ham & turkey</meat>' 
  '<filling>Cheese, lettuce, tomato, etc.</filling>' 
  '<![CDATA[We should add a <relish> element in future!]]>' 
  '</sandwich>' 
  'junk';

call plisaxa( eventHandler, 
  addr(token), 
  addrdata(xmlDocument), 
  length(xmlDocument) );

end;

Figure 94. PLISAXA coding example - main routine
dcl chars char(32000) based;

start_of_document:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

dcl userToken    pointer;
dcl xmlToken     pointer;
dcl tokenLength  fixed bin(31);

  put skip list( lowercase( procname() )
                || ' length=' || tokenLength );

  return(0);
end;

version_information:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

dcl userToken    pointer;
dcl xmlToken     pointer;
dcl tokenLength  fixed bin(31);

  put skip list( lowercase( procname() )
                || ' <' || substr(xmltoken->chars,1,tokenlength ) || '>' );

  return(0);
end;

encoding_declaration:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

dcl userToken    pointer;
dcl xmlToken     pointer;
dcl tokenLength  fixed bin(31);

  put skip list( lowercase( procname() )
                || ' <' || substr(xmltoken->chars,1,tokenlength ) || '>' );

  return(0);
end;

Figure 95. PLISAXA coding example - event routines
standalone_declaration:
    proc( userToken, xmlToken, TokenLength )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

    put skip list( lowercase( procname() )
        || ' <' || substr(xmltoken->chars,1,tokenlength ) || ' >' );

    return(0);
end;

document_type_declaration:
    proc( userToken, xmlToken, TokenLength )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

    put skip list( lowercase( procname() )
        || ' <' || substr(xmltoken->chars,1,tokenlength ) || ' >' );

    return(0);
end;

document_type_declaration:
    proc( userToken, xmlToken, TokenLength )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

    put skip list( lowercase( procname() )
        || ' <' || substr(xmltoken->chars,1,tokenlength ) || ' >' );

    return(0);
end;

end_of_document:
    proc( userToken )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

dcl userToken pointer;

    put skip list( lowercase( procname() ) );

    return(0);
end;

PLISAXA coding example - event routines (continued)
start_of_element:
   proc( userToken, xmlToken, TokenLength )
   returns( byvalue fixed bin(31) )
   options( byvalue linkage(optlink) );

   dcl userToken pointer;
   dcl xmlToken pointer;
   dcl tokenLength fixed bin(31);

   put skip list( lowercase( procname() )
                  || ' ' || substr(xmltoken->chars,1,tokenlength ) || '>' );

   return(0);
end;

attribute_name:
   proc( userToken, xmlToken, TokenLength )
   returns( byvalue fixed bin(31) )
   options( byvalue linkage(optlink) );

   dcl userToken pointer;
   dcl xmlToken pointer;
   dcl tokenLength fixed bin(31);

   put skip list( lowercase( procname() )
                  || ' ' || substr(xmltoken->chars,1,tokenlength ) || '>' );

   return(0);
end;

attribute_characters:
   proc( userToken, xmlToken, TokenLength )
   returns( byvalue fixed bin(31) )
   options( byvalue linkage(optlink) );

   dcl userToken pointer;
   dcl xmlToken pointer;
   dcl tokenLength fixed bin(31);

   put skip list( lowercase( procname() )
                  || ' ' || substr(xmltoken->chars,1,tokenlength ) || '>' );

   return(0);
end;

**PLISAXA coding example - event routines (continued)**
attribute_predefined_reference:
proc( userToken, reference )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) nodescriptor );

dcl userToken pointer;
dcl reference char(1);

put skip list( lowercase( procname() )
|| ' ' || hex(reference ) );

return(0);
end;

attribute_character_reference:
proc( userToken, reference )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl reference fixed bin(31);

put skip list( lowercase( procname() )
|| '<' || hex(reference ) );

return(0);
end;

d_end_of_element:
proc( userToken, xmlToken, TokenLength )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

put skip list( lowercase( procname() )
|| '<' || substr(xmltoken->chars,1,tokenlength ) || '>' );

return(0);
end;

PLISAXA coding example - event routines (continued)
start_ofCDATA:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

  dcl userToken pointer;
  dcl xmlToken pointer;
  dcl tokenLength fixed bin(31);

  put skip list( lowercase( procname() )
    || ' <' || substr(xmlToken->chars,1,tokenLength ) || '>' );

  return(0);
end;

end_ofCDATA:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

  dcl userToken pointer;
  dcl xmlToken pointer;
  dcl tokenLength fixed bin(31);

  put skip list( lowercase( procname() )
    || ' <' || substr(xmlToken->chars,1,tokenLength ) || '>' );

  return(0);
end;

content_characters:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

  dcl userToken pointer;
  dcl xmlToken pointer;
  dcl tokenLength fixed bin(31);

  put skip list( lowercase( procname() )
    || ' <' || substr(xmlToken->chars,1,tokenLength ) || '>' );

  return(0);
end;

PLISAXA coding example - event routines (continued)
content_predefined_reference:
proc( userToken, reference )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) nodescriptor );

dcl userToken pointer;
dcl reference char(1);

put skip list( lowercase( procname() )
|| ' ' || hex(reference ) );
return(0);
end;

content_character_reference:
proc( userToken, reference )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl reference fixed bin(31);

put skip list( lowercase( procname() )
|| '<' || hex(reference ) );
return(0);
end;

processing_instruction:
proc( userToken, piTarget, piTargetLength,
piData, piDataLength )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl piTarget pointer;
dcl piTargetLength fixed bin(31);
dcl piData pointer;
dcl piDataLength fixed bin(31);

put skip list( lowercase( procname() )
|| '<' || subst(piTarget->chars,1,piTargetLength ) || '>' );
return(0);
end;

PLISAXA coding example - event routines (continued)
comment:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

  dcl userToken      pointer;
  dcl xmlToken      pointer;
  dcl tokenLength   fixed bin(31);

  put skip list( lowercase( procname() )
                 || ' <' || substr(xmltoken->chars,1,tokenlength ) || '>' );

  return(0);
end;

unknown_attribute_reference:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

  dcl userToken      pointer;
  dcl xmlToken      pointer;
  dcl tokenLength   fixed bin(31);

  put skip list( lowercase( procname() )
                 || ' <' || substr(xmltoken->chars,1,tokenlength ) || '>' );

  return(0);
end;

unknown_content_reference:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

  dcl userToken      pointer;
  dcl xmlToken      pointer;
  dcl tokenLength   fixed bin(31);

  put skip list( lowercase( procname() )
                 || ' <' || substr(xmltoken->chars,1,tokenlength ) || '>' );

  return(0);
end;

PLISAXA coding example - event routines (continued)
The preceding program would produce the following output:
Continuable exception codes

For each value of the exception code parameter passed to the exception event (listed under the heading "Number"), the following table describes the exception, and the actions that the parser takes when you request it to continue after the exception. In these descriptions, the term "XML text" refers to the string based on the pointer and length passed to the event.

**Table 34. Continuable Exceptions**

<table>
<thead>
<tr>
<th>Number</th>
<th>Description</th>
<th>Parser Action on Continuation</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>The parser found an invalid character while scanning white space outside element content.</td>
<td>The parser generates a content_characters event with XML text containing the (single) invalid character. Parsing continues at the character after the invalid character.</td>
</tr>
<tr>
<td>2</td>
<td>The parser found an invalid start of a processing instruction, element, comment or document type declaration outside element content.</td>
<td>The parser generates a content_characters event with the XML text containing the 2- or 3-character invalid initial character sequence. Parsing continues at the character after the invalid sequence.</td>
</tr>
</tbody>
</table>

Figure 96. PLISAXA coding example - program output
<table>
<thead>
<tr>
<th>Number</th>
<th>Description</th>
<th>Parser Action on Continuation</th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>The parser found a duplicate attribute name.</td>
<td>The parser generates an attribute_name event with the XML text containing the duplicate attribute name.</td>
</tr>
<tr>
<td>4</td>
<td>The parser found the markup character &quot;&lt;&quot; in an attribute value.</td>
<td>Prior to generating the exception event, the parser generates an attribute_characters event for any part of the attribute value prior to the &quot;&lt;&quot; character. After the exception event, the parser generates an attribute_characters event with XML text containing &quot;&lt;&quot;. Parsing then continues at the character after the &quot;&lt;&quot;.</td>
</tr>
<tr>
<td>5</td>
<td>The start and end tag names of an element did not match.</td>
<td>The parser generates an end_of_element event with XML text containing the mismatched end name.</td>
</tr>
<tr>
<td>6</td>
<td>The parser found an invalid character in element content.</td>
<td>The parser includes the invalid character in XML text for the subsequent content_characters event.</td>
</tr>
<tr>
<td>7</td>
<td>The parser found an invalid start of an element, comment, processing instruction or CDATA section in element content.</td>
<td>Prior to generating the exception event, the parser generates a content_characters event for any part of the content prior to the &quot;&lt;&quot; markup character. After the exception event, the parser generates a content_characters event with XML text containing 2 characters: the &quot;&lt;&quot; followed by the invalid character. Parsing continues at the character after the invalid character.</td>
</tr>
<tr>
<td>8</td>
<td>The parser found in element content the CDATA closing character sequence &quot;]]&gt;&quot; without the matching opening character sequence “&lt;![CDATA[“.</td>
<td>Prior to generating the exception event, the parser generates a content_characters event for any part of the content prior to the &quot;]]&gt;&quot; character sequence. After the exception event, the parser generates a content_characters event with XML text containing the 3-character sequence &quot;]]&gt;&quot;. Parsing continues at the character after this sequence.</td>
</tr>
<tr>
<td>9</td>
<td>The parser found an invalid character in a comment.</td>
<td>The parser includes the invalid character in XML text for the subsequent comment event.</td>
</tr>
<tr>
<td>10</td>
<td>The parser found in a comment the character sequence &quot;--&quot; not followed by &quot;&gt;&quot;.</td>
<td>The parser assumes that the &quot;--&quot; character sequence terminates the comment, and generates a comment event. Parsing continues at the character after the &quot;--&quot; sequence.</td>
</tr>
<tr>
<td>11</td>
<td>The parser found an invalid character in a processing instruction data segment.</td>
<td>The parser includes the invalid character in XML text for the subsequent processing_instruction event.</td>
</tr>
<tr>
<td>12</td>
<td>A processing instruction target name was &quot;xml&quot; in lower-case, upper-case or mixed-case.</td>
<td>The parser generates a processing_instruction event with XML text containing &quot;xml&quot; in the original case.</td>
</tr>
<tr>
<td>13</td>
<td>The parser found an invalid digit in a hexadecimal character reference (of the form &amp;#dddd).</td>
<td>The parser generates an attribute_characters or content_characters event with XML text containing the invalid digit. Parsing of the reference continues at the character after this invalid digit.</td>
</tr>
<tr>
<td>Number</td>
<td>Description</td>
<td>Parser Action on Continuation</td>
</tr>
<tr>
<td>--------</td>
<td>-------------</td>
<td>------------------------------</td>
</tr>
<tr>
<td>14</td>
<td>The parser found an invalid digit in a decimal character reference (of the form &amp;#dddr).</td>
<td>The parser generates an attribute_characters or content_characters event with XML text containing the invalid digit. Parsing of the reference continues at the character after this invalid digit.</td>
</tr>
<tr>
<td>15</td>
<td>The encoding declaration value in the XML declaration did not begin with lower- or upper-case A through Z</td>
<td>The parser generates the encoding event with XML text containing the encoding declaration value as it was specified.</td>
</tr>
<tr>
<td>16</td>
<td>A character reference did not refer to a legal XML character.</td>
<td>The parser generates an attribute_character_reference or content_character_reference event with XML-NTEXT containing the single Unicode character specified by the character reference.</td>
</tr>
<tr>
<td>17</td>
<td>The parser found an invalid character in an entity reference name.</td>
<td>The parser includes the invalid character in the XML text for the subsequent unknown_attribute_reference or unknown_content_reference event.</td>
</tr>
<tr>
<td>18</td>
<td>The parser found an invalid character in an attribute value.</td>
<td>The parser includes the invalid character in XML text for the subsequent attribute_characters event.</td>
</tr>
<tr>
<td>50</td>
<td>The document was encoded in EBCDIC, and the CODEPAGE compiler option specified a supported EBCDIC code page, but the document encoding declaration did not specify a recognizable encoding.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
<tr>
<td>51</td>
<td>The document was encoded in EBCDIC, and the document encoding declaration specified a supported EBCDIC encoding, but the parser does not support the code page specified by the CODEPAGE compiler option.</td>
<td>The parser uses the encoding specified by the document encoding declaration.</td>
</tr>
<tr>
<td>52</td>
<td>The document was encoded in EBCDIC, and the CODEPAGE compiler option specified a supported EBCDIC code page, but the document encoding declaration specified an ASCII encoding.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
<tr>
<td>53</td>
<td>The document was encoded in EBCDIC, and the CODEPAGE compiler option specified a supported EBCDIC code page, but the document encoding declaration specified a supported Unicode encoding.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
<tr>
<td>54</td>
<td>The document was encoded in EBCDIC, and the CODEPAGE compiler option specified a supported EBCDIC code page, but the document encoding declaration specified a Unicode encoding that the parser does not support.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
<tr>
<td>55</td>
<td>The document was encoded in EBCDIC, and the CODEPAGE compiler option specified a supported EBCDIC code page, but the document encoding declaration specified an encoding that the parser does not support.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
</tbody>
</table>
### Table 34. Continuable Exceptions (continued)

<table>
<thead>
<tr>
<th>Number</th>
<th>Description</th>
<th>Parser Action on Continuation</th>
</tr>
</thead>
<tbody>
<tr>
<td>56</td>
<td>The document was encoded in ASCII, and the CODEPAGE compiler option specified a supported ASCII code page, but the document encoding declaration did not specify a recognizable encoding.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
<tr>
<td>57</td>
<td>The document was encoded in ASCII, and the document encoding declaration specified a supported ASCII encoding, but the parser does not support the code page specified by the CODEPAGE compiler option.</td>
<td>The parser uses the encoding specified by the document encoding declaration.</td>
</tr>
<tr>
<td>58</td>
<td>The document was encoded in ASCII, and the CODEPAGE compiler option specified a supported ASCII code page, but the document encoding declaration specified a supported EBCDIC encoding.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
<tr>
<td>59</td>
<td>The document was encoded in ASCII, and the CODEPAGE compiler option specified a supported ASCII code page, but the document encoding declaration specified a supported Unicode encoding.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
<tr>
<td>60</td>
<td>The document was encoded in ASCII, and the CODEPAGE compiler option specified a supported ASCII code page, but the document encoding declaration specified a Unicode encoding that the parser does not support.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
<tr>
<td>61</td>
<td>The document was encoded in ASCII, and the CODEPAGE compiler option specified a supported ASCII code page, but the document encoding declaration specified an encoding that the parser does not support.</td>
<td>The parser uses the encoding specified by the CODEPAGE compiler option.</td>
</tr>
<tr>
<td>100,001 through 165,535</td>
<td>The document was encoded in EBCDIC, and the encodings specified by the CODEPAGE compiler option and the document encoding declaration are both supported EBCDIC code pages, but are not the same. The exception code contains the CCSID for the encoding declaration plus 100,000.</td>
<td>If you return zero from the exception event, the parser uses the encoding specified by the CODEPAGE compiler option. If you return the CCSID from the document encoding declaration (by subtracting 100,000 from the exception code), the parser uses this encoding.</td>
</tr>
<tr>
<td>200,001 through 265,535</td>
<td>The document was encoded in ASCII, and the encodings specified by the CODEPAGE compiler option and the document encoding declaration are both supported ASCII code pages, but are not the same. The exception code contains the CCSID for the encoding declaration plus 200,000.</td>
<td>If you return zero from the exception event, the parser uses the encoding specified by the CODEPAGE compiler option. If you return the CCSID from the document encoding declaration (by subtracting 200,000 from the exception code), the parser uses this encoding.</td>
</tr>
</tbody>
</table>

### Terminating exception codes

#### Table 35. Terminating Exceptions

<table>
<thead>
<tr>
<th>Number</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>100</td>
<td>The parser reached the end of the document while scanning the start of the XML declaration.</td>
</tr>
<tr>
<td>Number</td>
<td>Description</td>
</tr>
<tr>
<td>--------</td>
<td>-------------</td>
</tr>
<tr>
<td>101</td>
<td>The parser reached the end of the document while looking for the end of the XML declaration.</td>
</tr>
<tr>
<td>102</td>
<td>The parser reached the end of the document while looking for the root element.</td>
</tr>
<tr>
<td>103</td>
<td>The parser reached the end of the document while looking for the version information in the XML declaration.</td>
</tr>
<tr>
<td>104</td>
<td>The parser reached the end of the document while looking for the version information value in the XML declaration.</td>
</tr>
<tr>
<td>106</td>
<td>The parser reached the end of the document while looking for the encoding declaration value in the XML declaration.</td>
</tr>
<tr>
<td>108</td>
<td>The parser reached the end of the document while looking for the standalone declaration value in the XML declaration.</td>
</tr>
<tr>
<td>109</td>
<td>The parser reached the end of the document while scanning an attribute name.</td>
</tr>
<tr>
<td>110</td>
<td>The parser reached the end of the document while scanning an attribute value.</td>
</tr>
<tr>
<td>111</td>
<td>The parser reached the end of the document while scanning a character reference or entity reference in an attribute value.</td>
</tr>
<tr>
<td>112</td>
<td>The parser reached the end of the document while scanning an empty element tag.</td>
</tr>
<tr>
<td>113</td>
<td>The parser reached the end of the document while scanning the root element name.</td>
</tr>
<tr>
<td>114</td>
<td>The parser reached the end of the document while scanning an element name.</td>
</tr>
<tr>
<td>115</td>
<td>The parser reached the end of the document while scanning character data in element content.</td>
</tr>
<tr>
<td>116</td>
<td>The parser reached the end of the document while scanning a processing instruction in element content.</td>
</tr>
<tr>
<td>117</td>
<td>The parser reached the end of the document while scanning a comment or CDATA section in element content.</td>
</tr>
<tr>
<td>118</td>
<td>The parser reached the end of the document while scanning a comment in element content.</td>
</tr>
<tr>
<td>119</td>
<td>The parser reached the end of the document while scanning a CDATA section in element content.</td>
</tr>
<tr>
<td>120</td>
<td>The parser reached the end of the document while scanning a character reference or entity reference in element content.</td>
</tr>
<tr>
<td>121</td>
<td>The parser reached the end of the document while scanning after the close of the root element.</td>
</tr>
<tr>
<td>122</td>
<td>The parser found a possible invalid start of a document type declaration.</td>
</tr>
<tr>
<td>123</td>
<td>The parser found a second document type declaration.</td>
</tr>
<tr>
<td>124</td>
<td>The first character of the root element name was not a letter, '_' or ':'.</td>
</tr>
<tr>
<td>125</td>
<td>The first character of the first attribute name of an element was not a letter, '_' or ':'.</td>
</tr>
<tr>
<td>126</td>
<td>The parser found an invalid character either in or following an element name.</td>
</tr>
<tr>
<td>127</td>
<td>The parser found a character other than '=' following an attribute name.</td>
</tr>
<tr>
<td>128</td>
<td>The parser found an invalid attribute value delimiter.</td>
</tr>
<tr>
<td>130</td>
<td>The first character of an attribute name was not a letter, '_' or ':'.</td>
</tr>
<tr>
<td>131</td>
<td>The parser found an invalid character either in or following an attribute name.</td>
</tr>
<tr>
<td>132</td>
<td>An empty element tag was not terminated by a '&gt;' following the '/'.</td>
</tr>
<tr>
<td>133</td>
<td>The first character of an element end tag name was not a letter, '_' or ':'.</td>
</tr>
<tr>
<td>134</td>
<td>An element end tag name was not terminated by a '&gt;'.</td>
</tr>
<tr>
<td>Number</td>
<td>Description</td>
</tr>
<tr>
<td>--------</td>
<td>-------------</td>
</tr>
<tr>
<td>135</td>
<td>The first character of an element name was not a letter, <code>_</code> or <code>:</code>.</td>
</tr>
<tr>
<td>136</td>
<td>The parser found an invalid start of a comment or CDATA section in element content.</td>
</tr>
<tr>
<td>137</td>
<td>The parser found an invalid start of a comment.</td>
</tr>
<tr>
<td>138</td>
<td>The first character of a processing instruction target name was not a letter, <code>_</code> or <code>:</code>.</td>
</tr>
<tr>
<td>139</td>
<td>The parser found an invalid character in or following a processing instruction target name.</td>
</tr>
<tr>
<td>140</td>
<td>A processing instruction was not terminated by the closing character sequence <code>?&gt;</code>.</td>
</tr>
<tr>
<td>141</td>
<td>The parser found an invalid character following <code>&amp;</code> in a character reference or entity reference.</td>
</tr>
<tr>
<td>142</td>
<td>The version information was not present in the XML declaration.</td>
</tr>
<tr>
<td>143</td>
<td><code>version</code> in the XML declaration was not followed by a <code>=</code>.</td>
</tr>
<tr>
<td>144</td>
<td>The version declaration value in the XML declaration is either missing or improperly delimited.</td>
</tr>
<tr>
<td>145</td>
<td>The version information value in the XML declaration specified a bad character, or the start and end delimiters did not match.</td>
</tr>
<tr>
<td>146</td>
<td>The parser found an invalid character following the version information value closing delimiter in the XML declaration.</td>
</tr>
<tr>
<td>147</td>
<td>The parser found an invalid attribute instead of the optional encoding declaration in the XML declaration.</td>
</tr>
<tr>
<td>148</td>
<td><code>encoding</code> in the XML declaration was not followed by a <code>=</code>.</td>
</tr>
<tr>
<td>149</td>
<td>The encoding declaration value in the XML declaration is either missing or improperly delimited.</td>
</tr>
<tr>
<td>150</td>
<td>The encoding declaration value in the XML declaration specified a bad character, or the start and end delimiters did not match.</td>
</tr>
<tr>
<td>151</td>
<td>The parser found an invalid character following the encoding declaration value closing delimiter in the XML declaration.</td>
</tr>
<tr>
<td>152</td>
<td>The parser found an invalid attribute instead of the optional standalone declaration in the XML declaration.</td>
</tr>
<tr>
<td>153</td>
<td><code>standalone</code> in the XML declaration was not followed by a <code>=</code>.</td>
</tr>
<tr>
<td>154</td>
<td>The standalone declaration value in the XML declaration is either missing or improperly delimited.</td>
</tr>
<tr>
<td>155</td>
<td>The standalone declaration value was neither <code>yes</code> nor <code>no</code> only.</td>
</tr>
<tr>
<td>156</td>
<td>The standalone declaration value in the XML declaration specified a bad character, or the start and end delimiters did not match.</td>
</tr>
<tr>
<td>157</td>
<td>The parser found an invalid character following the standalone declaration value closing delimiter in the XML declaration.</td>
</tr>
<tr>
<td>158</td>
<td>The XML declaration was not terminated by the proper character sequence <code>?&gt;</code>, or contained an invalid attribute.</td>
</tr>
<tr>
<td>159</td>
<td>The parser found the start of a document type declaration after the end of the root element.</td>
</tr>
<tr>
<td>160</td>
<td>The parser found the start of an element after the end of the root element.</td>
</tr>
<tr>
<td>300</td>
<td>The document was encoded in EBCDIC, but the CODEPAGE compiler option specified a supported ASCII code page.</td>
</tr>
<tr>
<td>301</td>
<td>The document was encoded in EBCDIC, but the CODEPAGE compiler option specified Unicode.</td>
</tr>
<tr>
<td>Number</td>
<td>Description</td>
</tr>
<tr>
<td>--------</td>
<td>-------------</td>
</tr>
<tr>
<td>302</td>
<td>The document was encoded in EBCDIC, but the CODEPAGE compiler option specified an unsupported code page.</td>
</tr>
<tr>
<td>303</td>
<td>The document was encoded in EBCDIC, but the CODEPAGE compiler option is unsupported and the document encoding declaration was either empty or contained an unsupported alphabetic encoding alias.</td>
</tr>
<tr>
<td>304</td>
<td>The document was encoded in EBCDIC, but the CODEPAGE compiler option is unsupported and the document did not contain an encoding declaration.</td>
</tr>
<tr>
<td>305</td>
<td>The document was encoded in EBCDIC, but the CODEPAGE compiler option is unsupported and the document encoding declaration did not specify a supported EBCDIC encoding.</td>
</tr>
<tr>
<td>306</td>
<td>The document was encoded in ASCII, but the CODEPAGE compiler option specified a supported EBCDIC code page.</td>
</tr>
<tr>
<td>307</td>
<td>The document was encoded in ASCII, but the CODEPAGE compiler option specified an unsupported code page.</td>
</tr>
<tr>
<td>308</td>
<td>The document was encoded in ASCII, but the CODEPAGE compiler option did not specify a supported EBCDIC code page, ASCII or Unicode.</td>
</tr>
<tr>
<td>309</td>
<td>The CODEPAGE compiler option specified a supported ASCII code page, but the document was encoded in Unicode.</td>
</tr>
<tr>
<td>310</td>
<td>The CODEPAGE compiler option specified a supported EBCDIC code page, but the document was encoded in Unicode.</td>
</tr>
<tr>
<td>311</td>
<td>The CODEPAGE compiler option specified an unsupported code page, but the document was encoded in Unicode.</td>
</tr>
<tr>
<td>312</td>
<td>The document was encoded in ASCII, but both the encodings provided externally and within the document encoding declaration are unsupported.</td>
</tr>
<tr>
<td>313</td>
<td>The document was encoded in ASCII, but the CODEPAGE compiler option is unsupported and the document did not contain an encoding declaration.</td>
</tr>
<tr>
<td>314</td>
<td>The document was encoded in ASCII, but the CODEPAGE compiler option is unsupported and the document encoding declaration did not specify a supported ASCII encoding.</td>
</tr>
<tr>
<td>315</td>
<td>The document was encoded in UTF-16 Little Endian, which the parser does not support on this platform.</td>
</tr>
<tr>
<td>316</td>
<td>The document was encoded in UCS4, which the parser does not support.</td>
</tr>
<tr>
<td>317</td>
<td>The parser cannot determine the document encoding. The document may be damaged.</td>
</tr>
<tr>
<td>318</td>
<td>The document was encoded in UTF-8, which the parser does not support.</td>
</tr>
<tr>
<td>319</td>
<td>The document was encoded in UTF-16 Big Endian, which the parser does not support on this platform.</td>
</tr>
<tr>
<td>501, 502, 503</td>
<td>An internal error occurred in PLISAX(A\B). Contact IBM support.</td>
</tr>
<tr>
<td>500</td>
<td>Memory allocation failed for the PLISAXA internal data structures. Increase the amount of storage available to the application program.</td>
</tr>
<tr>
<td>520</td>
<td>Memory allocation failed for the PLISAXB internal data structures. Increase the amount of storage available to the application program.</td>
</tr>
<tr>
<td>521</td>
<td>An internal error occurred in PLISAX(A\B). Contact IBM support.</td>
</tr>
<tr>
<td>523</td>
<td>PLISAXB encountered a file I/O error.</td>
</tr>
<tr>
<td>524</td>
<td>Memory allocation failed in PLISAXB while attempting to cache the XML document from the file system. Increase the amount of storage available to the application program.</td>
</tr>
<tr>
<td>525</td>
<td>An unsupported URI scheme was specified to PLISAXB.</td>
</tr>
</tbody>
</table>
Table 35. Terminating Exceptions (continued)

<table>
<thead>
<tr>
<th>Number</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>526</td>
<td>The XML document provided to PLISAXB was less than the minimum of 4 characters or was too large.</td>
</tr>
<tr>
<td>527, 560</td>
<td>An internal error occurred in PLISAX(A</td>
</tr>
<tr>
<td>561</td>
<td>No event handler was specified to either PLISAX(A</td>
</tr>
<tr>
<td>562, 563, 580, 581</td>
<td>An internal error occurred in PLISAX(A</td>
</tr>
<tr>
<td>600 to 99,999</td>
<td>Internal error. Report the error to your service representative.</td>
</tr>
</tbody>
</table>
Chapter 18. Using the PLISAXC and PLISAXD XML parsers

The PLISAXC and PLISAXD built-in subroutines provide basic XML parsing capability, which allows programs to consume inbound XML documents, check them for well-formedness, and react to their contents.

The XML parser used by PLISAXC is non-validating, but does partially check for well-formedness errors, and generates exception events if it discovers any.

The PLISAXD built-in subroutine provides XML parsing with validation capability. It determines whether an inbound XML documentation conforms to a set of rules specified in an inbound XML schema.

The PLISAXC and PLISAXD subroutines do not provide XML generation, which must instead be accomplished by PL/I program logic or by using the XMLCHAR built-in function.

PLISAXC and PLISAXD have no special environmental requirements except that it is not supported in AMODE 24. It executes in all the principal runtime environments, including CICS, IMS, MQ Series, z/OS batch, and TSO.

Because the PLISAXC and PLISAXD built-in subroutines and the PLISAXA and PLISAXB built-in subroutines do have much similarity, some of the discussion below repeats material from the previous chapter.

Overview

There are two major types of interfaces for XML parsing: event-based and tree-based.

For an event-based API, the parser reports events to the application through callbacks. Such events include the start of the document, the beginning of an element, and so on. The application provides handlers to deal with the events reported by the parser. The Simple API for XML or SAX is an example of an industry-standard event-based API.

For a tree-based API (such as the Document Object Model or DOM), the parser translates the XML into an internal tree-based representation. Interfaces are provided to navigate the tree.

IBM PL/I compiler provides, by using PLISAXC or PLISAXD, a SAX-like event-based interface for parsing XML documents. The parser invokes an application-supplied handler for parser events, passing references to the corresponding document fragments.

The parser has the following characteristics:

- It provides high-performance, but non-standard interfaces.
- It supports XML files encoded in either Unicode UTF-16, UTF-8 or any of several single-byte code pages listed in the section “Coded character sets for XML documents” on page 415.
XML documents have two levels of conformance: well-formedness and validity, both of which are defined in the XML standard, which you can find at http://www.w3c.org/XML/. Recapitulating these definitions, an XML document is well-formed if it complies with the basic XML grammar, and with a few specific rules, such as the requirement that the names on start and end element tags must match. A well-formed XML document is also valid if it has an associated document type declaration (DTD) and if it complies with the constraints expressed in the DTD.

For each parser event, you must provide a PL/I function that accepts the appropriate parameters and returns the appropriate return value - as in the example code below. Note in particular that for these functions

- the return value must be returned BYVALUE.
- the linkage used must be the OPTLINK linkage. You can use the DEFAULT(LINKAGE(OPTLINK)) option to specify this linkage, or you can specify it on the individual PROCEDUREs and ENTRYs by using the OPTIONS(LINKAGE(OPTLINK)) attribute.

The PLISAXC built-in subroutine

The PLISAXC built-in subroutine allows you to invoke the XML parser for an XML document residing in one or more buffers in your program.

```
PLISAXC(e, p, x, n, c)
```

- `e` An event structure
- `p` A pointer value or "token" that the parser will pass back to the event functions
- `x` The address of the initial buffer containing the input XML
- `n` The number of bytes of data in that buffer
- `c` A numeric expression specifying the codepage of that XML

Note that if the XML is contained in a CHARACTER VARYING or a WIDECHAR VARYING string, then the ADDRDATA built-in function should be used to obtain the address of the first data byte.

Also note that if the XML is contained in a WIDECHAR string, the value for the number of bytes is twice the value returned by the LENGTH built-in function.

The PLISAXD built-in subroutine

The PLISAXD built-in subroutine allows you to invoke the XML parser with the validation capability. Both the XML document and the Optimized Schema Representation (OSR) file are in one or more buffers in your program.

```
PLISAXD(e, p, x, n, o, c)
```

- `e` An event structure
- `p` A pointer value or "token" that the parser passes back to the event functions
- `x` The address of the initial buffer containing the input XML
n  The number of bytes of data in that buffer
o  The address of the buffer containing the input OSR
c  A numeric expression specifying the codepage of the XML document

If the XML is contained in a CHARACTER VARYING or a WIDECHAR VARYING string, the ADDRDATA built-in function should be used to obtain the address of the first data byte.

If the XML is contained in a WIDECHAR string, the value for the number of bytes is twice the value returned by the LENGTH built-in function.

Note: An OSR is the preprocessed version of a schema. For more information about OSR, see *XML System Services User’s Guide and Reference*.

### The SAX event structure

The event structure is a structure consisting of 19 LIMITED ENTRY variables, which point to functions that the parser invokes for various "events".

All of these ENTRYs must use the OPTLINK linkage.

The library routines that support the XML parsers recognize if an event in any of the PLISAX event structures is set to null (via the NULLENTRY built-in function or via use of UNSPEC) and do not call it in that case. This allows you to limit your XML parsing code to only the events in which you are interested and to improve the performance of the overall parse at the same time.

All of these ENTRYs have a first (and sometimes the only) parameter: the user token passed by the program to PLISAXC and PLISAXD.

The descriptions below of these 19 events refer to the example of an XML document in Figure 97. In these descriptions, the term "XML text" refers to the string based on the pointer and length passed to the event.

```sql
xmlDocument =
  '<!DOCTYPE xml SYSTEM &mySchema;> 
  <sandwich>
    <bread type="baker&s best"/>
    <spread please use real mayonnaise/>
    <meat>Ham & turkey</meat>
    <filling>Cheese, lettuce, tomato, etc.</filling>
  </sandwich> 
</xmlDocument>
```

*Figure 97. Sample XML document*

Depending on the contents of the XML documents, the parser might recognize the following events:

**start_of_document**

This event occurs once, at the beginning of parsing the document. The parser passes no parameters to this event (except the user token).
**version_information**

This event occurs within the optional XML declaration for the version information. The parser passes the address and length of the text containing the version value, "1.0" in the previous example.

**encoding_declaration**

This event occurs within the XML declaration for the optional encoding declaration. The parser passes the address and length of the text containing the encoding value.

**standalone_declaration**

This event occurs within the XML declaration for the optional standalone declaration. The parser passes the address and length of the text containing the standalone value, "yes" in the previous example.

**document_type_declaration**

This event occurs when the parser finds a document type declaration. Document type declarations begin with the character sequence "<!DOCTYPE" and end with a ">" character, with some fairly complicated grammar rules describing the content in between. The parser passes the address and length of the text containing the entire declaration, including the opening and closing character sequences, and is the only event where XML text includes the delimiters. The example above does not have a document type declaration.

**end_of_document**

This event occurs once, when document parsing has completed. The parser passes no parameters to this event (except the user token).

**start_of_element**

This event occurs once for each element start tag or empty element tag. The parser passes the address and length of the text containing the element name as well as any applicable namespace information. For the first start_of_element event during parsing of the example, this would be the string "sandwich".

**attribute_name**

This event occurs for each attribute in an element start tag or empty element tag, after recognizing a valid name. The parser passes the address and length of the text containing the attribute name as well as any applicable namespace information. The only attribute name in the example is "type".

**attribute_characters**

This event occurs for each attribute value. The parser passes the address and length of the text containing the fragment. An attribute value normally consists of a single string only, even if it is split across lines:

```
<element attribute="This attribute value is split across two lines"/>
```

The parser also passes a flag byte, which indicates whether the next event provides additional characters that form part of the content. This can be true when there is a lot of data between the start and end tags.
**end_of_element**

This event occurs once for each element end tag or empty element tag when the parser recognizes the closing angle bracket of the tag. The parser passes the address and length of the text containing the element name as well as any applicable namespace information.

**start_of_CDATA_section**

This event occurs at the start of a CDATA section. CDATA sections begin with the string “<![CDATA[” and end with the string “]]>”, and are used to "escape" blocks of text containing characters that would otherwise be recognized as XML markup. The parser passes no parameters to this event (except the user token). After this event, the parser passes the content of the CDATA section between these delimiters as one or more content-characters events. In the preceding example, the content-characters event is passed the text "We should add a <relish> element in future!".

**end_of_CDATA_section**

This event occurs when the parser recognizes the end of a CDATA section. The parser passes no parameters to this event (except the user token).

**content_characters**

This event represents the "heart" of an XML document: the character data between element start and end tags. The parser passes the address and length of the text containing the data, which usually consists of a single string only, even if it is split across lines:

```xml
<element1>This character content is
split across two lines</element1>
```

The parser also passes a flag byte, which indicates if the next event provides additional characters that form part of the content. This can be true when there is a lot of data between the start and end tags.

The parser also uses the content_characters event to pass the text of CDATA sections to the application.

**processing_instruction**

Processing instructions (PIs) allow XML documents to contain special instructions for applications. This event occurs when the parser recognizes the name following the PI opening character sequence "<?". The event also covers the data following the processing instruction (PI) target, up to but not including the PI closing character sequence "?>". Trailing, but not leading white space characters in the data are included. The parser passes the address and length of the text containing the target, "spread" in the example, and the address and length of the text containing the data, "please use real mayonnaise" in the example.

The parser also passes a flag byte, which indicates whether the next event provides additional characters that form part of the content. This can be true when there is a lot of data between the start and end tags.

**comment**

This event occurs for any comments in the XML document. The parser passes the address and length of the text between the opening and closing comment.
delimiters, "<!-" and "-->", respectively. In the example, the text of the only comment is "This document is just an example".

The parser also passes a flag byte, which indicates whether the next event provides additional characters that form part of the content. This can be true when there is a lot of data between the start and end tags.

**namespace_declare**

This event occurs for any namespace declarations in the XML document. The parser passes the address and length of the namespace prefix (if any) as well as the address and length of the namespace uri. If there is no namespace prefix, the passed length will be zero and the value of the address should not be used. There is no corresponding event in the PLIXSAXA and PLISAXB built-in subroutines.

**end_of_input**

This event occurs whenever the parser reaches the end of the current input buffer. The parser passes (along with the BYVALUE user token) two BYADDR parameters: the address and length of the next buffer for it to process. Note that this and the content character events are the only events that have any BYADDR parameters, but this is the only event that has parameters that the called event should change. There is no corresponding event in the PLIXSAXA and PLISAXB built-in subroutines, and it is this event that allows PLISAXC and PLISAXD to parse an XML document of arbitrary size.

**unresolved_reference**

This event occurs for any unresolved references in the XML document. The parser passes the address and length of the unresolved reference.

**exception**

The parser generates this event when it detects an error in processing the XML document.

**Parameters to the event functions**

All of these functions must return a BYVALUE FIXED BIN(31) value that is a return code to the parser. If any value other than zero is returned, the parser will terminate.

All of these functions are passed as the first argument a BYVALUE POINTER. This pointer is the token value that is passed originally as the second argument to the built-in function.

With the following exceptions, all of the functions will also be passed a BYVALUE POINTER and a BYVALUE FIXED BIN(31) that supply the address and length of the text element for the event. The following functions and events are different:

**start_of_document**

No argument other than the user token is passed.

**end_of_document**

No argument other than the user token is passed.

**start_ofCDATA**

No argument other than the user token is passed.
end_of_CDATA
No argument other than the user token is passed.

start_of_element
In addition to the usual three parameters, four additional arguments are passed:
1. a BYVALUE POINTER that is the address of the namespace prefix
2. a BYVALUE FIXED BIN(31) that is the length of the namespace prefix
3. a BYVALUE POINTER that is the address of the namespace uri
4. a BYVALUE FIXED BIN(31) that is the length of the namespace uri

end_of_element
In addition to the usual three parameters, four additional arguments are passed:
1. a BYVALUE POINTER that is the address of the namespace prefix
2. a BYVALUE FIXED BIN(31) that is the length of the namespace prefix
3. a BYVALUE POINTER that is the address of the namespace uri
4. a BYVALUE FIXED BIN(31) that is the length of the namespace uri

attribute_name
In addition to the usual three parameters, four additional arguments are passed:
1. a BYVALUE POINTER that is the address of the namespace prefix
2. a BYVALUE FIXED BIN(31) that is the length of the namespace prefix
3. a BYVALUE POINTER that is the address of the namespace uri
4. a BYVALUE FIXED BIN(31) that is the length of the namespace uri

attribute_characters
In addition to the usual three parameters, one additional argument is passed:
• a BYVALUE ALIGNED BIT(8) flag byte that indicates
  – whether more content characters will be presented in the next event - this is true if the first bit is on, that is, this is true if this field anded with '80'BX is non-null
  – whether there are no characters that need to be escaped if converted back to XML - this is true if the second bit is on, that is, this is true if this field anded with '40'BX is non-null

Note that this entry must also be declared with OPTIONS(NODESCRIPTOR).

namespace_declare
In addition to the user token, four additional arguments passed:
1. a BYVALUE POINTER that is the address of the namespace prefix
2. a BYVALUE FIXED BIN(31) that is the length of the namespace prefix
3. a BYVALUE POINTER that is the address of the namespace uri
4. a BYVALUE FIXED BIN(31) that is the length of the namespace uri

content_characters
In addition to the usual three parameters, one additional argument is passed:
• a BYVALUE ALIGNED BIT(8) flag byte that indicates
  – whether more content characters will be presented in the next event - this is true if the first bit is on, that is, this is true if this field anded with '80'BX is non-null
whether there are no characters that need to be escaped if converted back to XML - this is true if the second bit is on, that is, this is true if this field anded with '40'BX is non-null

Note that this entry must also be declared with OPTIONS(NODESCRIPTOR).

eof_input
In addition to the user token, two additional arguments are passed:
1. a BYADDR POINTER that is the address of the next input buffer
2. a BYADDR FIXED BIN(31) that is the length of the next input buffer

processing_instruction
In addition to the user token, five additional arguments are passed:
1. a BYVALUE POINTER that is the address of the target text
2. a BYVALUE FIXED BIN(31) that is the length of the target text
3. a BYVALUE POINTER that is the address of the data text
4. a BYVALUE FIXED BIN(31) that is the length of the data text
5. a BYVALUE ALIGNED BIT(8) flag byte that indicates
   - whether more content characters will be presented in the next event - this is true if the first bit is on, that is, this is true if this field anded with '80'BX is non-null
   - whether there are no characters that need to be escaped if converted back to XML - this is true if the second bit is on, that is, this is true if this field anded with '40'BX is non-null

Note that this entry must also be declared with OPTIONS(NODESCRIPTOR).

comment
In addition to the usual three parameters, one additional argument is passed:
• a BYVALUE ALIGNED BIT(8) flag byte that indicates
  - whether more content characters will be presented in the next event - this is true if the first bit is on, that is, this is true if this field anded with '80'BX is non-null
  - whether there are no characters that need to be escaped if converted back to XML - this is true if the second bit is on, that is, this is true if this field anded with '40'BX is non-null

Note that this entry must also be declared with OPTIONS(NODESCRIPTOR).

exception
In addition to the user token, three additional arguments are passed:
1. a BYVALUE FIXED BIN(31) that is the byte offset of the offending text within the document
2. a BYVALUE FIXED BIN(31) that is the return code for the exception
3. a BYVALUE FIXED BIN(31) that is the reason code for the exception

Differences in the events
The following events are part of the PLISAXA and PLISAXB event structure, but are not in PLISAXC and PLISAXD:
• attribute_predefined_reference
• attribute_character_reference
• content_predefined_reference
• content_character_reference
The following events are not part of the PLISAXA and PLISAXB event structure, but are in PLISAXC and PLISAXD:

- namespace_declare
- unresolved_reference
- end_of_input

Some of the events that are common to PLISAXA and PLISAXB, PLISAXC and PLISAXD are passed different parameters (apart from the omnipresent user token):

- **start_of_document**
  - is passed no parameters
- **start_of_element**
  - is passed namespace data as well
- **end_of_element**
  - is passed namespace data as well
- **attribute_name**
  - is passed namespace data as well
- **attribute_characters**
  - is passed a flag byte as well
- **start_of_cdata**
  - is passed no parameters
- **end_of_cdata**
  - is passed no parameters
- **content_characters**
  - is passed a flag byte as well
- **processing_instruction**
  - is passed a flag byte as well
- **comment**
  - is passed a flag byte as well
- **exception**
  - is passed a return and reason code instead of an error id

### Coded character sets for XML documents

The PLISAXC and PLISAXD built-in subroutines support only XML documents in WIDECHAR encoded using Unicode UTF-16 or in CHARACTER encoded using either UTF-8 or one of the explicitly supported single-byte character sets listed below. The parser uses up to three sources of information about the encoding of your XML document, and signals an exception XML event if it discovers any conflicts between these sources:

1. The parser determines the basic encoding of a document by inspecting its initial characters.
2. If step 1 succeeds, the parser then looks for any encoding declaration.
3. Finally, it refers to the codepage value on the PLISAXC or PLISAXD built-in subroutine call. If this parameter was omitted, it defaults to the value provided by the CODEPAGE compiler option value that you specified explicitly or by default.

If the XML document begins with an XML declaration that includes an encoding declaration specifying one of the supported code pages listed below, the parser honors the encoding declaration if it does not conflict with either the basic document encoding or the encoding information from the PLISAXC or PLISAXD built-in subroutines. If the XML document does not have an XML declaration at all, or if the XML declaration omits the encoding declaration, the parser uses the encoding information from the PLISAXC or PLISAXD built-in subroutine to process the document, as long as it does not conflict with the basic document encoding.

**Supported code pages**

In the following table, the first number is for the Euro Country Extended Code Page (ECECP), and the second is for Country Extended Code Page (CECP).

<table>
<thead>
<tr>
<th>CCSID</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>01208</td>
<td>Unicode UTF-8</td>
</tr>
<tr>
<td>01047</td>
<td>Latin 1 / Open Systems</td>
</tr>
<tr>
<td>01140, 00037</td>
<td>USA, Canada, etc.</td>
</tr>
<tr>
<td>01141, 00273</td>
<td>Austria, Germany</td>
</tr>
<tr>
<td>01142, 00277</td>
<td>Denmark, Norway</td>
</tr>
<tr>
<td>01143, 00278</td>
<td>Finland, Sweden</td>
</tr>
<tr>
<td>01144, 00280</td>
<td>Italy</td>
</tr>
<tr>
<td>01145, 00284</td>
<td>Spain, Latin America (Spanish)</td>
</tr>
<tr>
<td>01146, 00285</td>
<td>UK</td>
</tr>
<tr>
<td>01147, 00297</td>
<td>France</td>
</tr>
<tr>
<td>01148, 00500</td>
<td>International</td>
</tr>
<tr>
<td>01149, 00871</td>
<td>Iceland</td>
</tr>
</tbody>
</table>

**Specifying the code page**

If your document does not include an encoding declaration in the XML declaration, or does not have an XML declaration at all, the parser uses the encoding information provided by the PLISAXC or PLISAXD built-in subroutine call in conjunction with the basic encoding of the document.

You can also specify the encoding information for the document in the XML declaration, with which most XML documents begin. An example of an XML declaration that includes an encoding declaration is:

```xml
<?xml version="1.0" encoding="ibm-1140"?>
```

If your XML document includes an encoding declaration, ensure that it is consistent with the encoding information provided by the PLISAXC or PLISAXD built-in subroutine and with the basic encoding of the document. If there is any conflict between the encoding declaration, the encoding information provided by the PLISAXC or PLISAXD built-in subroutine and the basic encoding of the document, the parser signals an exception XML event.
Specify the encoding declaration as follows:

**Using a number**
You can specify the CCSID number (with or without any number of leading zeroes), prefixed by any of the following (in any mixture of upper or lower case):

- IBM_
- IBM-
- CP
- CP_
- CP-
- CCSID_
- CCSID-

**Using an alias**
You can use any of the following supported aliases (in any mixture of lower and upper case):

<table>
<thead>
<tr>
<th>Code page</th>
<th>Supported aliases</th>
</tr>
</thead>
<tbody>
<tr>
<td>037</td>
<td>EBCDIC-CP-US, EBCDIC-CP-CA, EBCDIC-CP-WT, EBCDIC-CP-NL</td>
</tr>
<tr>
<td>500</td>
<td>EBCDIC-CP-BE, EBCDIC-CP-CH</td>
</tr>
<tr>
<td>813</td>
<td>ISO-8859-7, ISO_8859-7</td>
</tr>
<tr>
<td>819</td>
<td>ISO-8859-1, ISO_8859-1</td>
</tr>
<tr>
<td>920</td>
<td>ISO-8859-9, ISO_8859-9</td>
</tr>
<tr>
<td>1200</td>
<td>UTF-16</td>
</tr>
</tbody>
</table>

**Exceptions**
If an exception event occurs, the reason and return codes passed to it are those from the XML System Services parser, and the documentation provided with that parser explains what these return and reason codes mean.

**Parsing XML documents with validation**
The PLISAXD built-in subroutine not only parses XML documents in the same manner as PLISAXC, but also determines whether an inbound XML document conforms to a set of rules specified in an inbound XML schema.

The inbound schema used for XML validation when using the PLISAXD built-in subroutine must be in a preprocessed format known as an Optimized Schema Representation (OSR).

The following information contains the description of the XML schema and the way to build an OSR for an XML schema. For an example of using PLISAXD built-in subroutine to parse XML documents with validation, see [“Example of using the PLISAXD built-in subroutine”](#) on page 429.

**XML schema**
An XML schema is a mechanism, defined by the W3C, for describing and constraining the structure and content of XML documents. Through its support for datatypes and namespaces, an XML schema has the potential to provide the standard structure for XML elements and attributes. Therefore, an XML schema, which is itself expressed in XML, can effectively define a class of XML documents of a given type, for example, stock item.
The following sample XML document describes an item for stock keeping purposes:

```xml
<stockItem itemNumber="453-SR">
  <itemName>Stainless steel rope thimbles</itemName>
  <quantityOnHand>23</quantityOnHand>
</stockItem>
```

The stock keeping example document above is both well formed and valid according to the following schema called stock.xsd. (The numbers that precede each line are not part of the schema, but are used in the explanation after the schema.)

```xml
1. <xsd:schema xmlns:xsd="http://www.w3.org/2001/XMLSchema">
2.  
3.  <xsd:element name="stockItem" type="stockItemType"/>
4.  
5.  <xsd:complexType name="stockItemType">
6.   
7.   <xsd:sequence>
8.   
9.   <xsd:element name="itemName" type="xsd:string" minOccurs="0"/>
10.   
11.   <xsd:element name="quantityOnHand">
12.     
13.     <xsd:simpleType>
14.       
15.       <xsd:restriction base="xsd:nonNegativeInteger">
16.         
17.         <xsd:maxExclusive value="100"/>
18.         
19.       </xsd:restriction>
20.     </xsd:element>
21.   </xsd:sequence>
22.   
23.   <xsd:attribute name="itemNumber" type="SKU" use="required"/>
24.  </xsd:complexType>
25. </xsd:schema>
```

The schema declares (line 3) that the root element is stockItem, which has a mandatory itemNumber attribute (line 16) of type SKU, and includes a sequence (lines 6 - 15) of other elements:

- An optional itemName element of type string (line 7)
- A required quantityOnHand element that has a constrained range of 1 - 99 based on the type nonNegativeInteger (lines 8 - 14)

Type declarations can be inline and unnamed, as in lines 9 - 13, which include the maxExclusive facet to specify the legal values for the quantityOnHand element.

For the itemNumber attribute, by contrast, the named type SKU is declared separately in lines 19 - 23, which include a pattern facet that uses regular expression syntax to specify that the legal values for that type consist of (in order): 3 digits, a hyphen-minus, then two uppercase letters.

**Creating an OSR**

To generate a schema in the OSR format from a text-form schema, use the z/OS UNIX command xsdosr, which invokes the OSR generator provided by z/OS System Service.
For example, to convert the text-form schema in the stock.xsd file to a schema in preprocessed format in the stock.osr file, you can use the following z/OS UNIX command:

```
xsdosrg -v -o /u/HLQ/xml/stock.osr /u/HLQ/xml/stock.xsd
```

where /u/HLQ/xml/ is the directory in which the stock.osr and stock.xsd files are located.

If you want to copy the generated OSR file into a PDS, use the z/OS UNIX cp command. To specify an MVS data set name, precede the name with double slashes (/). For example, to copy the HFS file stock.osr into a fixed block record format PDS called HLQ.XML.OSR with record length 80, you can use the following command:

```
cp -p /u/HLQ/xml/stock.osr "//HLQ.XML.OSR(STOCK)"
```

To omit the fully qualified name of the PDS, you can eliminate the single quotation mark in the cp statement such as:

```
cp -p /u/HLQ/xml/stock.osr "//XML.OSR(STOCK)"
```

For more information, see the related reference in XML System Services User’s Guide and Reference.

---

**Example with a simple document**

The following two examples illustrate the use of the PLISAXC and PLISAXD built-in subroutines.

---

**Example of using the PLISAXC built-in subroutine**

The following example illustrates the use of the PLISAXC built-in subroutine and uses the example XML document in the preceding example. This example is essentially the same as that used in the previous chapter: it does not use namespaces, and all the input is passed when PLISAXC is first invoked (and as a result, the end_of_input event should not be invoked).
saxtest: package exports(saxtest);

define alias event
    limited entry( pointer, pointer, fixed bin(31) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_with_flag
    limited entry( pointer, pointer, fixed bin(31),
          bit(8) aligned )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_with_namespace
    limited entry( pointer, pointer, fixed bin(31),
                  pointer, fixed bin(31),
                  pointer, fixed bin(31) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_without_data
    limited entry( pointer )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_pi
    limited entry( pointer, pointer, fixed bin(31),
                  pointer, fixed bin(31) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_namespace_dcl
    limited entry( pointer, pointer, fixed bin(31),
                  pointer, fixed bin(31) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_exception
    limited entry( pointer, fixed bin(31),
                  fixed bin(31),
                  fixed bin(31) )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

define alias event_end_of_input
    limited entry( pointer,
                  pointer byaddr,
                  fixed bin(31) byaddr )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

Figure 98. PLISAXC coding example - type declarations
saxtest: proc options( main );

dcl
  1 eventHandler static
    2 e01 type event_without_data
        init( start_of_document )
    2 e02 type event
        init( version_information )
    2 e03 type event
        init( encoding_declaration )
    2 e04 type event
        init( standalone_declaration )
    2 e05 type event
        init( document_type_declaration )
    2 e06 type event_without_data
        init( end_of_document )
    2 e07 type event_with_namespace
        init( start_of_element )
    2 e08 type event_with_namespace
        init( attribute_name )
    2 e09 type event_with_flag
        init( attribute_characters )
    2 e10 type event_with_namespace
        init( end_of_element )
    2 e11 type event_without_data
        init( start_of_CDATA )
    2 e12 type event_without_data
        init( end_of_CDATA )
    2 e13 type event_with_flag
        init( content_characters )
    2 e14 type event_pi
        init( processing_instruction )
    2 e15 type event_with_flag
        init( comment )
    2 e16 type event_namespace_dcl
        init( namespace_declare )
    2 e17 type event_end_of_input
        init( end_of_input )
    2 e18 type event
        init( unresolved_reference )
    2 e19 type event_exception
        init( exception )
;
Figure 99. PLISAXC coding example - event structure
dcl token char(8);

dcl xmlDocument char(4000) var;

xmlDocument =
  '<?xml version="1.0" standalone="yes"?>'
  '<!--This document is just an example-->'
  '<sandwich>'
  '<bread type="baker's best"/>'
  '<?spread please use real mayonnaise ?>'
  '<meat>Ham & turkey</meat>'
  '<filling>Cheese, lettuce, tomato, etc.</filling>'
  '<![CDATA[We should add a <relish> element in future!]]>'
  '</sandwich>'

';

call plisaxc( eventHandler,
                addr(token),
                addrdata(xmlDocument),
                length(xmlDocument) );

end;

Figure 100. PLISAXC coding example - main routine
dcl chars char(32000) based;

start_of_document:
proc( userToken )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

put skip list( lowercase( procname() )
|| ' length=' || tokenLength );

return(0);
end;

version_information:
proc( userToken, xmlToken, TokenLength )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

put skip list( lowercase( procname() )
|| ' ' || substr(xmlToken->chars,1,tokenLength ) || '>' );

return(0);
end;

encoding_declaration:
proc( userToken, xmlToken, TokenLength )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

put skip list( lowercase( procname() )
|| ' ' || substr(xmlToken->chars,1,tokenLength ) || '>' );

return(0);
end;

standalone_declaration:
proc( userToken, xmlToken, TokenLength )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

put skip list( lowercase( procname() )
|| ' ' || substr(xmlToken->chars,1,tokenLength ) || '>' );

return(0);
end;

Figure 101. PLISAXC coding example - event routines
document_type_declaration:
    proc( userToken, xmlToken, TokenLength )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );
    dcl userToken pointer;
    dcl xmlToken pointer;
    dcl tokenLength fixed bin(31);
    put skip list( lowercase( procname() )
        || ' ' || substr(xmlToken->chars,1,tokenLength ) || '>' );
    return(0);
end;

namespace_declare:
    proc( userToken, nsPrefix, nsPrefixLength, nsUri, nsUriLength )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );
    dcl userToken pointer;
    dcl nsPrefix pointer;
    dcl nsPrefixLength fixed bin(31);
    dcl nsUri pointer;
    dcl nsUriLength fixed bin(31);
    put skip list( lowercase( procname() ) );
    put skip list( 'prefix = ' || ' ' || substr(nsPrefix->chars,1,nsPrefixlength ) || '>' );
    put skip list( 'Uri = ' || ' ' || substr(nsUri->chars,1,nsUriLength ) || '>' );
    return(0);
end;

end_of_document:
    proc( userToken )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );
    dcl userToken pointer;
    put skip list( lowercase( procname() ) );
    return(0);
end;

PLISAXC coding example - event routines (continued)
start_of_element:
  proc( userToken, xmlToken, TokenLength, 
       nsPrefix, nsPrefixLength, 
       nsUri, nsUriLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

dcl userToken pointer;
  dcl xmlToken pointer;
  dcl tokenLength fixed bin(31);
  dcl nsPrefix pointer;
  dcl nsPrefixLength fixed bin(31);
  dcl nsUri pointer;
  dcl nsUriLength fixed bin(31);

  put skip list( lowercase( procname() ) 
                 || ' ' || substr(xmltoken->chars,1,tokenlength ) || '>' );

  return(0);
end;

attribute_name:
  proc( userToken, xmlToken, TokenLength, 
       nsPrefix, nsPrefixLength, 
       nsUri, nsUriLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

dcl userToken pointer;
  dcl xmlToken pointer;
  dcl tokenLength fixed bin(31);
  dcl nsPrefix pointer;
  dcl nsPrefixLength fixed bin(31);
  dcl nsUri pointer;
  dcl nsUriLength fixed bin(31);

  put skip list( lowercase( procname() ) 
                 || ' ' || substr(xmltoken->chars,1,tokenlength ) || '>' );

  return(0);
end;

del end_of_element:
  proc( userToken, xmlToken, TokenLength, 
       nsPrefix, nsPrefixLength, 
       nsUri, nsUriLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

dcl userToken pointer;
  dcl xmlToken pointer;
  dcl tokenLength fixed bin(31);
  dcl nsPrefix pointer;
  dcl nsPrefixLength fixed bin(31);
  dcl nsUri pointer;
  dcl nsUriLength fixed bin(31);

  put skip list( lowercase( procname() ) 
                 || ' ' || substr(xmltoken->chars,1,tokenlength ) || '>' );

  return(0);
end;

PLISAXC coding example - event routines (continued)
content_characters:
  proc(userToken, xmlToken, TokenLength, flags )
  returns( byvalue fixed bin(31) )
  options( nodelector, byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);
dcl flags bit(8) aligned;

  put skip list( lowercase( procname() )
  || ' ' || substr(xmltoken->chars,1,tokenlength ) || ' ' );

  if flags = 'b' then;
  else
    put skip list( '!!flags = ' || flags );

  return(0);
end;

attribute_characters:
  proc(userToken, xmlToken, TokenLength, flags )
  returns( byvalue fixed bin(31) )
  options( nodelector, byvalue linkage(optlink) );

  dcl userToken pointer;
  dcl xmlToken pointer;
  dcl tokenLength fixed bin(31);
  dcl flags bit(8);

  put skip list( lowercase( procname() )
  || ' ' || substr(xmltoken->chars,1,tokenlength ) || ' ' );

  if flags = 'b' then;
  else
    put skip list( '!!flags = ' || flags );

  return(0);
end;

start_of_CDATA:
  proc(userToken )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

dcl userToken pointer;

  put skip list( lowercase( procname() ) );

  return(0);
end;

del_end_of_CDATA:
  proc(userToken )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

dcl userToken pointer;

  put skip list( lowercase( procname() ) );

  return(0);
end;

PLISAXC coding example - event routines (continued)
processing_instruction:
proc( userToken,
     piTarget, piTargetLength,
     piData, piDataLength,
     flags )
returns( byvalue fixed bin(31) )
options( nodescriptor, byvalue linkage(optlink) );

dcl userToken pointer;
dcl piTarget pointer;
dcl piTargetLength fixed bin(31);
dcl piData pointer;
dcl piDataLength fixed bin(31);
dcl flags bit(8);

put skip list( lowercase( procname() )
  || ' <' || substr(piTarget->chars,1,piTargetLength) || '>' );

if flags = 'b' then;
else
  put skip list( '!!flags = ' || flags );

return(0);
end;

comment:
proc( userToken, xmlToken, TokenLength, flags )
returns( byvalue fixed bin(31) )
options( nodescriptor, byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);
dcl flags bit(8);

put skip list( lowercase( procname() )
  || ' <' || substr(xmlToken->chars,1,tokenLength) || '>' );

if flags = 'b' then;
else
  put skip list( '!!flags = ' || flags );

return(0);
end;

unresolved_reference:
proc( userToken, xmlToken, TokenLength )
returns( byvalue fixed bin(31) )
options( nodescriptor, byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

put skip list( lowercase( procname() )
  || ' <' || substr(xmlToken->chars,1,tokenLength) || '>' );

return(0);
end;

PLISAXC coding example - event routines (continued)
exception:
    proc( userToken, currentOffset, return_code, reason_code )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

dcl userToken    pointer;
dcl currentOffset fixed bin(31);
dcl return_code  fixed bin(31);
dcl reason_code  fixed bin(31);

    put skip list( lowercase( procname() )
        || ' return_code ='' || hex(return_code)
        || ', reason_code ='' || hex(reason_code)
        || ', offset ='' || currentOffset );

    return(0);
end;
end_of_input:
    proc( userToken, addr_xml, length_xml )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

dcl userToken    pointer;
dcl addr_xml byaddr pointer;
dcl length_xml byaddr fixed bin(31);

    return(1);
end;
end;

PLISAXC coding example - event routines (continued)

The preceding program would produce the following output:
Example of using the PLISAXD built-in subroutine

The following example illustrates the use of the PLISAXD built-in subroutine and uses the example XML document and the XML schema cited in the previous examples.

This example includes the validation of 8 different XML files against the same stock.osr schema. In the following output, you can see which XML documents in the saxdtest program fail validation against the schema.

The PLISAXD built-in subroutine requires the XML schema file to be read into a buffer. The OSR file in the following example is in a PDS. The initial size of the OSR buffer is set to 4096. If you have a larger OSR file, you can increase the initial size of the OSR buffer accordingly.
If the inbound schema file were in an HFS file instead, you could use the following code to read the OSR file into the buffer:

dcl osrin file input stream environment(u);
dcl fileddint builtin;
dcl fileread builtin;

/* Read the HFS OSR file into buffer*/

open file(osrin);
osr_length = fileddint(osrin,'filesize');
osr_ptr = allocate(osr_length);
rc = fileread(osrin,osr_ptr,osr_length);

To run a program by using an OSR in a PDS, you can specify the following DD statement in the JCL:

//OSRIN DD DSN=HLQ.XML.OSR(STOCK),DISP=SHR

If the associated ddname OSRIN is an HFS file, then use the following JCL statement instead:

//OSRIN DD PATH="/u/HLQ/xml/stock.osr"

---

saxdtest: package exports(saxdtest);

/*****************************/
/* saxdtest: Test PL/I XML validation support */
/* expected output: */
/* */
/* SAXDTEST: PL/I XML Validation sample */
/* SAXDTEST: Document Successfully parsed */
/* SAXDTEST: Document Successfully parsed */
/* Invalid: missing attribute itemNumber. */
/* exception return_code =00000018, reason_code =8613 */
/* Invalid: unexpected attribute warehouse. */
/* exception return_code =00000018, reason_code =8612 */
/* Invalid: illegal attribute value 123-Ab. */
/* exception return_code =00000018, reason_code =8809 */
/* invalid: missing element quantityOnHand. */
/* exception return_code =00000018, reason_code =8611 */
/* invalid: unexpected element comment. */
/* exception return_code =00000018, reason_code =8607 */
/* invalid: out-of-range element value 100 */
/* exception return_code =00000018, reason_code =8803 */
/* */
/*****************************/

define alias event
limited entry( pointer, pointer, fixed bin(31) )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

define alias event_with_flag
limited entry( pointer, pointer, fixed bin(31),
bit(8) aligned )
returns( byvalue fixed bin(31) )
options( nodelist byvalue linkage(optlink) );

Figure 103. PLISAXD coding example - event routines
define alias event_with_namespace
  limited entry( pointer, pointer, fixed bin(31),
  pointer, fixed bin(31),
  pointer, fixed bin(31) )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

define alias event_without_data
  limited entry( pointer )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

define alias event_pi
  limited entry( pointer, pointer, fixed bin(31),
  pointer, fixed bin(31) )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

define alias event_namespace_dcl
  limited entry( pointer, pointer, fixed bin(31),
  pointer, fixed bin(31) )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

define alias event_exception
  limited entry( pointer, fixed bin(31),
  fixed bin(31),
  fixed bin(31) )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

define alias event_end_of_input
  limited entry( pointer,
  pointer byaddr,
  fixed bin(31) byaddr )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

PLISAXD coding example – event routines (continued)
saxdtest: proc options(main);

dcl
  1 eventHandler static
  ,2 e01 type event_without_data
    init( start_of_document )
  ,2 e02 type event
    init( version_information )
  ,2 e03 type event
    init( encoding_declaration )
  ,2 e04 type event
    init( standalone_declaration )
  ,2 e05 type event
    init( document_type_declaration )
  ,2 e06 type event_without_data
    init( end_of_document )
  ,2 e07 type event_with_namespace
    init( start_of_element )
  ,2 e08 type event_with_namespace
    init( attribute_name )
  ,2 e09 type event_with_flag
    init( attribute_characters )
  ,2 e10 type event_with_namespace
    init( end_of_element )
  ,2 e11 type event_without_data
    init( start_of_CDATA )
  ,2 e12 type event_without_data
    init( end_of_CDATA )
  ,2 e13 type event_with_flag
    init( content_characters )
  ,2 e14 type pi
    init( processing_instruction )
  ,2 e15 type event_with_flag
    init( comment )
  ,2 e16 type event_namespace_dcl
    init( namespace_declare )
  ,2 e17 type event_end_of_input
    init( end_of_input )
  ,2 e18 type event
    init( unresolved_reference )
  ,2 e19 type event_exception
    init( exception )
;

PLISAXD coding example - event routines (continued)
dcl token char(45);

dcl rc fixed bin(31);
dcl i fixed bin(31);
dcl xml_document(8) char(300) var;
dcl xml_valid_msg(8) char(45) var;
dcl osr_ptr pointer;
dcl record char(80);
dcl osr_index fixed bin;
dcl osr_buf_tail fixed bin;
dcl temp_osr pointer;
dcl buf_size fixed bin(31) init(4096);
dcl rec_size fixed bin(31);
dcl osr_length fixed bin(31);
dcl buf_length fixed bin(31);
dcl eof fixed bin(31) init(0);
dcl osrin file input;

on endfile(osrin) begin;
oef = 1;
end;

/* read the entire PDS osr file into the buffer */
put skip list ('SAXDTEST: PL/I XML Validation sample ');

osr_length = buf_size;
osr_index = 0;
osr_buf_tail = 0;
rec_size = length(record);
osr_ptr = allocate(osr_length);

do while (eof = 0);
read file(osrin) into(record);

osr_buf_tail += rec_size;
if osr_buf_tail > buf_size then do;
   buf_length = osr_length;
   osr_length += buf_size;
temp_osr = allocate(osr_length);

call plimove(temp_osr, osr_ptr, buf_length);
call plifree(osr_ptr);
   osr_ptr = temp_osr;
   osr_buf_tail = rec_size;

   call plimove(osr_ptr+osr_index, addr(record), rec_size);
   osr_index += rec_size;
end;
else do;
   call plimove(osr_ptr+osr_index, addr(record), rec_size);
   osr_index += rec_size;
end;
eof = 1;
end;

PLISAXD coding example - event routines (continued)
/* Valid XMLFILE */
xml_document(1) = '<stockItem itemNumber="453-SR">
  '<itemName>Stainless steel rope thimbles</itemName>
  '<quantityOnHand>23</quantityOnHand>
  '</stockItem>;
xml_valid_msg(1) = 'Valid XMLFILE ';
/* Valid: the ITEMNAME element can be omitted */
xml_document(2) = '<stockItem itemNumber="453-SR">
  '<quantityOnHand>23</quantityOnHand>
  '</stockItem>;
xml_valid_msg(2) = 'Valid: the ITEMNAME element can be omitted.';
/* Invalid: missing attribute itemNumber */
xml_document(3) = '<stockItem>
  '<itemName>Stainless steel rope thimbles</itemName>
  '<quantityOnHand>23</quantityOnHand>
  '</stockItem>;
xml_valid_msg(3) = 'Invalid: missing attribute itemNumber.';
/* Invalid: unexpected attribute warehouse */
xml_document(4) = '<stockItem itemNumber="453-SR" warehouse="NY">
  '<itemName>Stainless steel rope thimbles</itemName>
  '<quantityOnHand>23</quantityOnHand>
  '</stockItem>;
xml_valid_msg(4) = 'Invalid: unexpected attribute warehouse.';
/* Invalid: illegal attribute value 123-Ab */
xml_document(5) = '<stockItem itemNumber="123-Ab">
  '<itemName>Stainless steel rope thimbles</itemName>
  '<quantityOnHand>23</quantityOnHand>
  '</stockItem>;
xml_valid_msg(5) = 'Invalid: illegal attribute value 123-Ab.';
/* Invalid: missing element quantityOnHand */
xml_document(6) = '<stockItem itemNumber="074-UN">
  '<itemName>Stainless steel rope thimbles</itemName>
  '</stockItem>;
xml_valid_msg(6) = 'Invalid: missing element quantityOnHand.';
/* Invalid: unexpected element comment */
xml_document(7) = '<stockItem itemNumber="453-SR">
  '<itemName>Stainless steel rope thimbles</itemName>
  '<quantityOnHand>1</quantityOnHand>
  '<comment>Nylon bristles</comment>
  '</stockItem>;
xml_valid_msg(7) = 'Invalid: unexpected element comment.';
/* Invalid: out-of-range element value 100 */
xml_document(8) = '<stockItem itemNumber="123-AB">
  '<itemName>Paintbrush</itemName>
  '<quantityOnHand>100</quantityOnHand>
  '</stockItem>;
xml_valid_msg(8) = 'Invalid: out-of-range element value 100';

PLISAXD coding example - event routines (continued)
do i = 1 to hbound(xml_document);
  token = xml_valid_msg(i);
  call plisaxd( eventHandler,
    addr(token),
    addrdata(xml_document(i)),
    length(xml_document(i)),
    osr_ptr,
    37 );
end;

close file(osrin);
call plifree(osr_ptr);
end;

start_of_document:
  proc( userToken )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );
    dcl userToken    pointer;
    return(0);
end;

version_information:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );
    dcl userToken    pointer;
    dcl xmlToken    pointer;
    dcl tokenLength fixed bin(31);
    return(0);
end;

encoding_declaration:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );
    dcl userToken    pointer;
    dcl xmlToken    pointer;
    dcl tokenLength fixed bin(31);
    return(0);
end;

standalone_declaration:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );
    dcl userToken    pointer;
    dcl xmlToken    pointer;
    dcl tokenLength fixed bin(31);
    return(0);
end;

PLISAXD coding example - event routines (continued)
document_type_declaration:
  proc( userToken, xmlToken, TokenLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

dcl userToken pointer;
 dcl xmlToken pointer;
 dcl tokenLength fixed bin(31);

  return(0);
 end;

end_of_document:
  proc( userToken )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

  dcl userToken pointer;

  put skip list('SAXDTEST: Document Successfully parsed ');

  return(0);
 end;

start_of_element:
  proc( userToken, xmlToken, TokenLength,
        nsPrefix, nsPrefixLength,
        nsUri, nsUriLength )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );

  dcl userToken pointer;
  dcl xmlToken pointer;
  dcl tokenLength fixed bin(31);
  dcl nsPrefix pointer;
  dcl nsPrefixLength fixed bin(31);
  dcl nsUri pointer;
  dcl nsUriLength fixed bin(31);

  return(0);
 end;

PLISAXD coding example - event routines (continued)
attribute_name:
    proc( userToken, xmlToken, tokenLength, nsPrefix, nsPrefixLength, nsUri, nsUriLength )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);
dcl nsPrefix pointer;
dcl nsPrefixLength fixed bin(31);
dcl nsUri pointer;
dcl nsUriLength fixed bin(31);

return(0);
end;

attribute_characters:
    proc( userToken, xmlToken, TokenLength, flags )
    returns( byvalue fixed bin(31) )
    options( nodescriptor, byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);
dcl flags bit(8);

put skip list( lowercase( procname() )
    || '<' || substr(xmltoken->chars,1,tokenlength) || '>' );

if flags = 'b' then;
else
    put skip list( '!!flags=' || flags );

return(0);
end;

end_of_element:
    proc( userToken, xmlToken, TokenLength, nsPrefix, nsPrefixLength, nsUri, nsUriLength )
    returns( byvalue fixed bin(31) )
    options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);
dcl nsPrefix pointer;
dcl nsPrefixLength fixed bin(31);
dcl nsUri pointer;
dcl nsUriLength fixed bin(31);

return(0);
end;

PLISAXD coding example - event routines (continued)
start_of_CDATA:
  proc( userToken )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );
  dcl userToken pointer;

  return(0);
end;

end_of_CDATA:
  proc( userToken )
  returns( byvalue fixed bin(31) )
  options( byvalue linkage(optlink) );
  dcl userToken pointer;

  return(0);
end;

content_characters:
  proc( userToken, xmlToken, TokenLength, flags )
  returns( byvalue fixed bin(31) )
  options( nodescriptor, byvalue linkage(optlink) );
  dcl userToken pointer;
  dcl xmlToken pointer;
  dcl tokenLength fixed bin(31);
  dcl flags bit(8) aligned;

  if flags = 'b' then;
    else
      return(0);
  end;

processing_instruction:
  proc( userToken, piTarget, piTargetLength, piData, piDataLength, flags )
  returns( byvalue fixed bin(31) )
  options( nodescriptor, byvalue linkage(optlink) );
  dcl userToken pointer;
  dcl piTarget pointer;
  dcl piTargetLength fixed bin(31);
  dcl piData pointer;
  dcl piDataLength fixed bin(31);
  dcl flags bit(8);

  put skip list( lowercase( procmname() )
    || ' ' || substr(piTarget->chars,1,piTargetLength ) || '>' );

  if flags = 'b' then;
    else
      put skip list( '!!flags = ' || flags );

      return(0);
  end;

PLISAXD coding example - event routines (continued)
comment:
proc( userToken, xmlToken, TokenLength, flags )
returns( byvalue fixed bin(31) )
options( nodelist, byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);
dcl flags bit(8);

put skip list( lowercase( procname() )
|| ' '|| substr(xmlToken->chars,1,TokenLength) || ' ');

if flags = ''b then;
else
  put skip list( ''!!flags='' || flags ); return(0);
end;

namespace_declare:
proc( userToken, nsPrefix, nsPrefixLength,
nsUri, nsUriLength )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl nsPrefix pointer;
dcl nsPrefixLength fixed bin(31);
dcl nsUri pointer;
dcl nsUriLength fixed bin(31);

return(0);
end;

unresolved_reference:
proc( userToken, xmlToken, TokenLength )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl xmlToken pointer;
dcl tokenLength fixed bin(31);

return(0);
end;

exception:
proc( userToken, currentOffset, return_code, reason_code )
returns( byvalue fixed bin(31) )
options( byvalue linkage(optlink) );

dcl userToken pointer;
dcl currentOffset fixed bin(31);
dcl return_code fixed bin(31);
dcl reason_code fixed bin(31);
dcl validmsg char(45) based;

put skip list( userToken -> validmsg);
put skip list( lowercase( procname() )
|| ' return_code =' || hex(return_code)
|| ', reason_code =' || substr(hex(reason_code),5,4));

return(0);
end;
PLISAXD coding example - event routines (continued)

The following output shows the result of the sample program. For those documents that are not valid, the PLISAXD built-in subroutine invokes the XML exception event with the return code and reason code listed. For a detailed description of each return code and reason code, see XML System Services User's Guide and Reference.

SAXDTEST: PL/I XML Validation sample
SAXDTEST: Document Successfully parsed
SAXDTEST: Document Successfully parsed
Invalid: missing attribute itemNumber.
exception return_code =00000018, reason_code =8613
Invalid: unexpected attribute warehouse.
exception return_code =00000018, reason_code =8612
Invalid: illegal attribute value 123-Ab.
exception return_code =00000018, reason_code =8809
exception return_code =00000018, reason_code =8611
Invalid: unexpected element comment.
exception return_code =00000018, reason_code =8607
Invalid: out-of-range element value 100
exception return_code =00000018, reason_code =8803

Figure 104. Output from PLISAXD sample
Chapter 19. Using PLIDUMP

This section provides information about dump options and the syntax used to call PLIDUMP, and describes PL/I-specific information included in the dump that can help you debug your routine.

**Note:** PLIDUMP conforms to National Language Support standards.

Figure 105 shows an example of a PL/I routine calling PLIDUMP to produce a z/OS Language Environment dump. In this example, the main routine PLIDMP calls PLIDMPA, which then calls PLIDMPB. The call to PLIDUMP is made in routine PLIDMPB.

The syntax and options for PLIDUMP are shown below.

```plaintext
/SM590000/SM590000
PLIDUMP(character-string-expression 1,character-string-expression 2)

character-string-expression 1
is a dump options character string consisting of one or more of the following:

A Requests information relevant to all tasks in a multitasking program.
B BLOCKS (PL/I hexadecimal dump).
C Continue. The routine continues after the dump.
```

© Copyright IBM Corp. 1999, 2015
Exit from current task of a multitasking program. Program continues to run after requested dump is completed.

FILES.

FILES.

This includes all Language Environment storage, and hence all the BASED and CONTROLLED storage acquired via ALLOCATE statements.

Note: A ddname of CEESNAP should be specified with the H option to produce a SNAP dump of a PL/I routine, but if this is omitted, Language Environment will issue a message but still produce a dump with much very useful information.

BLOCKS (when running under CICS). The Transaction Work Area is included.

NOSTORAGE.

NOFILES.

NOBLOCKS.

NOBLOCKS (when running under CICS).

NOTRACEBACK.

Only information relevant to the current task in a multitasking program.

Stop. The enclave is terminated with a dump.

TRACEBACK.

T, F, and C are the default options.

character-string-expression 2

is a user-identified character string up to 80 characters long that is printed as the dump header.

PLIDUMP usage notes

If you use PLIDUMP, the following considerations apply:

- If a routine calls PLIDUMP a number of times, use a unique user-identifier for each PLIDUMP invocation. This simplifies identifying the beginning of each dump.
- A DD statement with the ddname PLIDUMP, PL1DUMP, or CEEDUMP can be used to define the data set for the dump.
- The data set defined by the PLIDUMP, PL1DUMP, or CEEDUMP DD statement should specify a logical record length (LRECL) of at least 133 to prevent dump records from wrapping. If SYSOUT is used as the target in any one of these DDs, you must specify MSGFILE(SYSOUT,FBA,133,0) or MSGFILE(SYSOUT,VBA,137,0) to ensure that the lines are not wrapped.
- When you specify the H option in a call to PLIDUMP, the PL/I library issues an OS SNAP macro to obtain a dump of virtual storage. The first invocation of PLIDUMP results in a SNAP identifier of 0. For each successive invocation, the ID is increased by one to a maximum of 256, after which the ID is reset to 0.
- Support for SNAP dumps using PLIDUMP is only provided under z/OS. SNAP dumps are not produced in a CICS environment.
– If the SNAP is not successful, the CEE3DMP DUMP file displays the message:
  Snap was unsuccessful
– If the SNAP is successful, CEE3DMP displays the message:
  Snap was successful; snap ID = nnn
  where nnn corresponds to the SNAP identifier described above. An
  unsuccessful SNAP does not increment the identifier.

- If you want the program unit name, program unit address and program unit
  offset to be listed correctly in the dump traceback table, you need to make sure
  that your PL/I program unit is compiled with a compile-time option other than
  TEST(NONE,NOSYM). For example, you can specify the option as
  TEST(NOSYM,NOHOOK,BLOCK).

If you want to ensure portability across system platforms, use PLIDUMP to
generate a dump of your PL/I routine.

Locating variables in the PLIDUMP output

To find variables in the PLIDUMP output, you should compile your program with
the MAP option. The MAP option will cause the compiler to add to the listing a
table showing the offset within AUTOMATIC and STATIC storage of all level-1
variables that are AUTOMATIC or STATIC.

To find a variable that is an element in a structure, it is also useful if you compile
your program with the AGGREGATE option. This option will cause the compiler
to add to the listing a table showing the offsets of all the elements of all the
structures in your program.

Locating AUTOMATIC variables

To find an AUTOMATIC variable in the dump, you should find its offset within
automatic using the output from the MAP option (and if necessary the
AGGREGATE option). If PLIDUMP has been invoked with the B option, the dump
output will contain a hex dump of the "Dynamic save area" (or DSA) for each
block. This is the automatic storage for that block.

For example, consider the following simple program:

```plaintext
Compiler Source
Line.File
2.0    test: proc options(main);
3.0
4.0    dcl a fixed bin(31);
5.0    dcl b fixed bin(31);
6.0
7.0    on error
8.0    begin;
9.0    call plidump('TFBC');
10.0   end;
11.0
12.0   a = 0;
13.0   b = 29;
14.0   b = 17 / a;
```

The result of the compiler MAP option for this program looks like this, except that
there is actually one more column on the right and the columns are actually
spaced much further apart:
***** STORAGE OFFSET LISTING *****
IDENTIFIER DEFINITION ATTRIBUTES

A  1-0:4  Class = automatic, Location = 160 : 0xA0(r13),
B  1-0:5  Class = automatic, Location = 164 : 0xA4(r13),

So, A is located at hex A0 off of register 13 and B is located at hex A4 off of register 13, where register 13 points to the DSA.

Since in this program PLIDUMP is called with the B option, it will include a hexadecimal dump of automatic storage for each block in the current calling chain. This will look like (again with the right columns cutoff):

Dynamic save area (TEST): 0AD963C8
+000000 0AD963C8 10000000 0A96188 00000000 00000000
+000020 0AD963E8 00000000 00000000 00000000 00000000
+000040 0AD96408 00000000 00000000 00000000 00000000
+000060 0AD96428 00000000 00000000 00000000 00000000
+000080 0AD96448 = +00009F 0AD96467 same as above
+0000A0 0AD96468 00000000 0000001D 00100000 00000000
+0000C0 0AD96488 0B300000 0A700930 0AD963C8 00000000
+0000E0 0AD964A8 00000000 00000000 00000000 00000000
+000100 0AD964C8 0A47810 0A70E6D0 0AD96540 0AD960F0
+000120 0AD964E8 00000001 0A70F4F8 0AD96318 0AD96318
+000140 0AD96508 00000000 00000000 00000000 00000000

Since A is at hex offset A0 and B is at hex offset A4 in AUTOMATIC, the dump shows that A and B have the (expected) hex values of 00000000 and 0000001D respectively.

Please note that under the compiler options OPT(2) and OPT(3), some variables, particularly FIXED BIN and POINTER scalar variables, may never be allocated storage and thus could not be found in the dump output.

### Locating STATIC variables

If you compiled your code with the RENT option, static variables are located in the WSA (the Writeable Static Area) for the current load module. The offset of a variable within the WSA can be found from the output of the MAP option, and the WSA is held in the Language Environment control block called the CAA. The value of the WSA is also listed in the Language Environment dump.

However, if you compiled your code with the NORENT option, EXTERNAL STATIC is found as usual (using the linker listing and the output of the compiler’s MAP option). INTERNAL STATIC will be dumped as part of the Language Environment dump (if PLIDUMP was called with the B option).

Please also note that unlike the older PL/I compilers, the address of static is not dedicated to any one register.

For example, consider the program above with the variables changed to STATIC:

**Compiler Source**

<table>
<thead>
<tr>
<th>Line</th>
<th>File</th>
</tr>
</thead>
<tbody>
<tr>
<td>2.0</td>
<td>test: proc options(main);</td>
</tr>
<tr>
<td>3.0</td>
<td></td>
</tr>
<tr>
<td>4.0</td>
<td>dcl a fixed bin(31) static;</td>
</tr>
<tr>
<td>5.0</td>
<td>dcl b fixed bin(31) static;</td>
</tr>
<tr>
<td>6.0</td>
<td></td>
</tr>
<tr>
<td>7.0</td>
<td>on error</td>
</tr>
<tr>
<td>8.0</td>
<td>begin;</td>
</tr>
<tr>
<td>9.0</td>
<td>call plidump('TFBC');</td>
</tr>
</tbody>
</table>
When compiled with the NORENT option, the result of the compiler MAP option for this program looks like this, except that there is actually one more column on the right and the columns are actually spaced much further apart:

```
***** STORAGE OFFSET LISTING *****
IDENTIFIER DEFINITION ATTRIBUTES
A 1-0:4  Class = static, Location = 0 : 0x0 + CSECT ***TEST2
B 1-0:5  Class = static, Location = 4 : 0x4 + CSECT ***TEST2
```

So, A is located at hex offset 00 into the static CSECT for the compilation unit TEST while B is located at hex offset 04.

Since in this program PLIDUMP is called with the B option, it will include a hexadecimal dump of static storage for each compilation in the current calling chain. This will look like (again with the right columns cutoff):

```
Static for procedure TEST  Timestamp: 2004.08.12
+000000 0FC00AA0 00000000 0000001D 0FC00DC8 0FC00AC0
+000020 0FC00AC0 0FC00AA8 00444042 00A3AE01 0FC009C8
+000040 0FC00AE0 6E3BFFE0 00000000 00000000 00000000
+000060 0FC00B00 00000000 00000000 00000000 00000000
+000000 0AD963C8 10000000 0AD96188 00000000 00000000
```

So, A at hex offset 00 has the (expected) hex value 00000000, and B at hex offset 04 has the (also expected) hex value 0000001D or the decimal value 29.

### Locating CONTROLLED variables

CONTROLLED variables are essentially LIFO stacks, and each CONTROLLED variable has an "anchor" that points to the top of that stack. The key to locating a CONTROLLED variable is to locate this anchor, and its location depends on the compiler options as described below.

In the rest of this discussion of CONTROLLED variables, the program source will be the same program as above, but with the storage class changed to CONTROLLED:

```
Compiler Source
Line.File
2.0  test: proc options(main);
3.0
4.0  dcl a fixed bin(31) controlled;
5.0  dcl b fixed bin(31) controlled;
6.0
7.0  on error
8.0     begin;
9.0     call plidump('TFBHC');
10.0    end;
11.0
12.0  allocate a, b;
13.0  a = 0;
14.0  b = 29;
15.0  b = 17 / a;
```
Under NORENT WRITABLE

The result of the compiler MAP option looks like this, except that again there is actually one more column on the right and the columns are actually spaced much further apart:

```
***** STORAGE OFFSET LISTING *****
IDENTIFIER DEFINITION ATTRIBUTES
A  1-0:4  Class = static, Location = 8 : 0x8 + CSECT ***TEST2
B  1-0:5  Class = static, Location = 12 : 0xC + CSECT ***TEST2
```

Note that these lines describe the location of the anchors for A and B (not the location of A and B themselves). So, A’s anchor is located at hex 08 into the static CSECT for the compilation unit TEST while B’s anchor is located at hex 0C.

If PLIDUMP is called with the B option, it will include a hexadecimal dump of static storage for each compilation in the current calling chain. This will look like (again with the right columns cutoff):

```
Static for procedure TEST  Timestamp: ...
+000000 0FC00A88 0FC00B00 0FC00A88 102B8A30 102B8A50
+000020 0FC00A88 10444042 00A3AE01 0FC00980
+000040 0FC00AC8 6E3BFFE0 00000000 00000000 00000000 00000000
```

So, A’s anchor is at 102B8A30 and B’s is at 102B8A50. But since these are CONTROLLED variables, their storage was obtained via ALLOCATE statements, and hence these addresses point into heap storage. But if PLIDUMP is called with the H option, it will include a hexadecimal dump of heap storage. This will look like (again with the right columns cutoff):

```
Enclave Storage:
Initial (User) Heap
+000000 102B7018 C8C1D5C3 0FC0F990 0FC0F990 00000000 ...
+001A00 102B8A18 102B7018 00000020 0FC00A90 00000014
00000000 00000000 00000000 00000000
+001A20 102B8A38 102B7018 00000020 0FC00A94 00000014
00000000 00000000 0000001D 00000000
```

Since A’s anchor was at 102B8A30, A has the hex value 00000000, and since B’s anchor was at 102B8A50, B has the (expected) hex value 0000001D.

Under NORENT NOWRITABLE(FWS)

The result of the compiler MAP option under these options looks like this, except that again there is actually one more column on the right and the columns are actually spaced much further apart:

```
***** STORAGE OFFSET LISTING *****
IDENTIFIER DEFINITION ATTRIBUTES
A  1-0:4  Class = automatic, Location = 236 : 0xEC(r13)
B  1-0:5  Class = automatic, Location = 240 : 0xF0(r13)
```

Note: Under these options, there is an extra level of indirection in locating CONTROLLED variables, and hence the lines above describe the locations of the addresses of the anchors for A and B. So, the address of A’s anchor is located at hex EC into the automatic for the block TEST while B’s is located at hex F0.

Since PLIDUMP is called with the B option, it will include a hexadecimal dump of automatic storage for each block in the current calling chain. This will look like (again with the right columns cutoff):
Dynamic save area (TEST): 102973CB
+000000 102973CB 10000000 10297188 00000000 8FC007DA
....
+0000E0 102974A8 0FC00998 00000000 00000000 102B8A40
102B8A28 10297030 102977D0 8FD3D7E

So, the address of A's anchor is 102B8A40 and B's is 102B8A28.

Since PLIDUMP was also called with the H option, it will include a hexadecimal dump of heap storage. This will look like (again with the right columns cutoff):

Enclave Storage:
Initial (User) Heap
+000000 102B7018 C8C1D5C3 0FC0F990 0FC0F990 00000000
....
+001A00 102B8A18 102B7018 00000018 00000000 0FC00A78
102B8A80 00000000 102B7018 0000018
+001A20 102B8A38 102B8A20 0FC00A74 102B8A60 00000000
102B7018 00000020 102B8A40 0000014
+001A40 102B8A58 00000000 00000000 00000000 00000000
102B7018 00000020 102B8A28 0000014
+001A60 102B8A78 00000000 00000000 00000000 00000000
00000000 00000000 00000000 00000000

Since B's anchor address was at 102B8A28, B's anchor is at 102B8A80, and B has, as expected, the hex value 000001D or decimal 29.

Under NORENT NOWRITABLE(PRV)
The MAP listing when compiled with these options would look like this:

<table>
<thead>
<tr>
<th>IDENTIFIER</th>
<th>DEFINITION</th>
<th>ATTRIBUTES</th>
</tr>
</thead>
<tbody>
<tr>
<td>***TEST3</td>
<td>1-0:4</td>
<td>Class = ext def, Location = CSECT ***TEST3</td>
</tr>
<tr>
<td>***TEST4</td>
<td>1-0:5</td>
<td>Class = ext def, Location = CSECT ***TEST4</td>
</tr>
<tr>
<td>_PRV_OFFSETS</td>
<td>1-0:1</td>
<td>Class = static, Location = B : 0x8 + CSECT ***TEST2</td>
</tr>
</tbody>
</table>

The key here is the last line in this output: _PRV_OFFSETS is a static table that holds the offset into the PRV table for each CONTROLLED variable. This static table is generated only if the MAP option is specified.

To interpret this table, the compiler will also produce, immediately after the block names table, another, usually small, listing, which for our program would look like:

<table>
<thead>
<tr>
<th>Number</th>
<th>Offset</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>8</td>
<td>A</td>
</tr>
<tr>
<td>1</td>
<td>9</td>
<td>B</td>
</tr>
</tbody>
</table>

This table lists the hex offset within the runtime _PRV_OFFSETS table for each of the named CONTROLLED variables. The block number (in the first column) can be used to distinguish variables with the same name but declared in different blocks.

Since the _PRV_OFFSETS table is in static storage (at hex offset 8) and since PLIDUMP was called with the B option, it will appear in the dump output, which would look like this:
Static for procedure TEST   Timestamp: ...
+000000  10908EC8  02020240  00000005  6DD79EC5  6DD6C6C6
          00000000  00000004  000000A0  00100000
+000020  10908EE8  6E3BFFE0  00000000  00000000  00000000
          00000000  90010000  00000000  00000000

So, the offset of A in the PRV table is 0, and the offset of B in the PRV table is 4. Note also the eyecatcher "_PRV_OFF" that occupies the first 8 bytes of the _PRV_OFFSETS table.

The PRV table is always located at offset 4 within the CAA which, since PLIDUMP was called with the H option, will be in the dump output and which looks like:

Control Blocks Associated with the Thread:
CAA: 0A7107D0
+000000  0A7107D0  00000800  0ADB7DE0  0AD97018  0ADB7018
          00000000  00000000  00000000  00000000

So, the address of the PRV table is 0ADB7DE0, and it will also be in the dump output amongst the HEAP storage:

Enclave Storage:
Initial (User) Heap
+000000  102B7018  C8C1D5C3  0FC0F990  0FC0F990  00000000
          ... 
+0000C0  0ADB7D08  00000000  00000000  0ADB8A38  0ADB8A58
          0ADB7018  00000020  00000014  0A7107D4  00000000

So, the PRV table contains 0ADB8A38 0ADB8A58 etc, and since, as derived from the _PRV_OFFSETS table, A’s offset into the PRV table is 0 and B’s offset is 4, these are also the addresses of A and B respectively.

These addresses will also appear in the HEAP storage in the dump:

Enclave Storage:
Initial (User) Heap
+000000  102B7018  C8C1D5C3  0FC0F990  0FC0F990  00000000
          ... 
+001A00  0ADB8A18  00000000  00000000  0ADB7018  00000020
          00000000  00000014  0A7107D4  00000000
+001A20  0ADB8A38  00000000  00000000  0ADB7018  00000020
          00000004  00000014  0A7107D4  00000000
+001A40  0ADB8A58  0000001D  00000000  00000000  00000000
          00000000  00000000  00000000  00000000

So, since A’s address is 0ADB8A38, its hex value is, as expected, 00000000, and since B’s address is 0ADB8A58, its hex value is, also as expected, 0000001D.

Saved Compilation Data

During a compilation, the compiler saves various information about the compilation in the load module. This information can be very useful in debugging and in future migration efforts. This chapter describes the information saved.

Copyright

If you specify the COPYRIGHT compiler option, the compiler will save the COPYRIGHT string as a CHARACTER VARYING string placed immediately before the timestamp data in the object. This string will be followed by as many blanks as necessary so that the string plus these blanks will occupy a multiple of 4 bytes.
**Timestamp**

The compiler saves in every load module a "timestamp", which is a 20-byte character string of the form "YYYYMMDDHHMISSVNRNML" which records the date and time of the compilation as well as the version of the compiler that produced this string. The elements of the string have the following meanings:

- **YYYY**
  - the year of the compilation
- **MM**
  - the month of the compilation
- **DD**
  - the day of the compilation
- **HH**
  - the hour of the compilation
- **MI**
  - the minute of the compilation
- **SS**
  - the second of the compilation
- **VN**
  - the version number of the compiler
- **RN**
  - the release number of the compiler
- **ML**
  - the maintenance level of the compiler

The timestamp can be located from the PPA2: at offset 12 in the PPA2 is a four-byte integer giving the offset (possibly negative) to the timestamp from the address of the PPA2.

The PPA2, in turn, can be located from the PPA1: at offset 4 in the PPA1 is a four-byte integer giving the offset (possibly negative) to the PPA2 from the entry point address corresponding to that PPA1.

The PPA1, in turn, can be located from the entry point address for a block: at offset 12 from the entry point address is a four-byte integer giving the offset (possibly negative) to the PPA1 from the entry point address.

**Saved options string**

The compiler also stores in the load module a 32-byte string that records the compiler options used in building the load module.

A PL/I declaration for the saved options string is given in Figure 106 on page 450. For most of the fields in the structure, the meaning of the field is obvious given its name, but a few of the fields need some explanation:

- **sos_words** holds the number of the bytes in the structure divided by 4.
- **sos_version** is the version number for this structure. It is not a compiler version number.
- The size of the structure and what fields have been set depends on the version number.
Declare for the saved options string (continued)

Figure 106. Declare for the saved options string
Declare for the saved options string (continued)
Declare for the saved options string (continued)
Declare for the saved options string (continued)

The possible values for the sos_cmpat field are given by these declares:

```
dcl sos_cmpat_le bit(4) value('0000'b);
dcl sos_cmpat_v1 bit(4) value('0001'b);
dcl sos_cmpat_v2 bit(4) value('0010'b);
dcl sos_cmpat_v3 bit(4) value('0011'b);
```

The possible values for the sos_system field are given by these declares:

```
dcl sos_system_mvs bit(4) value('0001'b);
dcl sos_system_tso bit(4) value('0010'b);
dcl sos_system_cics bit(4) value('0011'b);
dcl sos_system_ims bit(4) value('0100'b);
dcl sos_system_os bit(4) value('0101'b);
```

The possible values for the sos_test_hooks field are given by these declares:

```
dcl sos_test_hooks_none bit(4) value('0000'b);
dcl sos_test_hooks_block bit(4) value('0001'b);
dcl sos_test_hooks_stmt bit(4) value('0011'b);
dcl sos_test_hooks_path bit(4) value('0101'b);
dcl sos_test_hooks_all bit(4) value('0111'b);
```
The possible values for the sos_linkage field are given by these declares:

```
dcl sos_linkage_optlink fixed bin(8) unsigned value(1);
dcl sos_linkage_system fixed bin(8) unsigned value(2);
```

The possible values for the sos_bifprec field are given by these declares:

```
dcl sos_bifprec_15 bit(2) value('01'b);
dcl sos_bifprec_31 bit(2) value('10'b);
```

The possible values for the sos_floatinmath field are given by these declares:

```
dcl sos_floatinmath_asis bit(2) value('00'b);
dcl sos_floatinmath_long bit(2) value('10'b);
dcl sos_floatinmath_extndd bit(2) value('11'b);
```

The saved options string is located after the timestamp in one of two ways:

1. if the service option has been specified, the string specified in the service option follows immediately after the timestamp as a character varying string. Then the saved options string follows after the service string as a second character varying string.
2. if the service option has not been specified, the saved options string follows immediately after the timestamp as a character varying string.

The length of the varying string that holds the saved options string may be longer than the size of the saved options string itself.

The presence (or absence) of the service string is indicated in the PPA2 by the flag byte at decimal offset 20 in the PPA2: if the result of anding this byte with '20'bx is not zero, then the service string is present.

In some earlier releases of the PL/I compiler, the compiler did not place a saved options string in the load module. The presence (or absence) of the saved options string is indicated in the PPA2 by the flag byte at decimal offset 20 in the PPA2: if the result of anding this byte with '02'bx is not zero, then the saved options string is present.
Chapter 20. Interrupts and attention processing

To enable a PL/I program to recognize attention interrupts, two operations must be possible:

- You must be able to create an interrupt. This is done in different ways depending upon both the terminal you use and the operating system.
- Your program must be prepared to respond to the interrupt. You can write an ON ATTENTION statement in your program so that the program receives control when the ATTENTION condition is raised.

**Note:** If the program has an ATTENTION ON-unit that you want invoked, you must compile the program with either of the following:
- The INTERRUPT option (supported only in TSO)
- A TEST option other than NOTEST or TEST(NONE,NOSYM).
  Compiling this way causes INTERRUPT(ON) to be in effect, unless you explicitly specify INTERRUPT(OFF) in PLIXOPT.

You can find the procedure used to create an interrupt in the IBM instruction manual for the operating system and terminal that you are using.

There is a difference between the interrupt (the operating system recognized your request) and the raising of the ATTENTION condition.

An interrupt is your request that the operating system notify the running program. If a PL/I program was compiled with the INTERRUPT compile-time option, instructions are included that test an internal interrupt switch at discrete points in the program. The internal interrupt switch can be set if any program in the load module was compiled with the INTERRUPT compile-time option.

The internal switch is set when the operating system recognizes that an interrupt request was made. The execution of the special testing instructions (polling) raises the ATTENTION condition. If a debugging tool hook (or a CALL PLITEST) is encountered before the polling occurs, the debugging tool can be given control before the ATTENTION condition processing starts.

Polling ensures that the ATTENTION condition is raised between PL/I statements, rather than within the statements.

Figure 107 on page 456 shows a skeleton program, an ATTENTION ON-unit, and several situations where polling instructions will be generated. In the program polling will occur at:

- LABEL1
- Each iteration of the DO
- The ELSE PUT SKIP ... statement
- Block END statements
Using ATTENTION ON-units

You can use processing within the ATTENTION ON-unit to terminate potentially endless looping in a program.

Control is given to an ATTENTION ON-unit when polling instructions recognize that an interrupt has occurred. Normal return from the ON-unit is to the statement following the polling code.

Interaction with a debugging tool

If the program has the TEST(ALL) or TEST(ERROR) run-time option in effect, an interrupt causes the debugging tool to receive control the next time a hook is encountered. This might be before the program’s polling code recognizes that the interrupt occurred.

Later, when the ATTENTION condition is raised, the debugging tool receives control again for condition processing.

Figure 107. Using an ATTENTION ON-unit
Chapter 21. Using the Checkpoint/Restart facility

This chapter describes the PL/I Checkpoint/Restart feature which provides a convenient method of taking checkpoints during the execution of a long-running program in a batch environment.

At points specified in the program, information about the current status of the program is written as a record on a data set. If the program terminates due to a system failure, you can use this information to restart the program close to the point where the failure occurred, avoiding the need to rerun the program completely.

This restart can be either automatic or deferred. An automatic restart is one that takes place immediately (provided the operator authorizes it when requested by a system message). A deferred restart is one that is performed later as a new job.

You can request an automatic restart from within your program without a system failure having occurred.

PL/I Checkpoint/Restart uses the Advanced Checkpoint/Restart Facility of the operating system. This facility is described in the books listed in "Bibliography" on page 533.

To use checkpoint/restart you must do the following:

- Request, at suitable points in your program, that a checkpoint record is written. This is done with the built-in subroutine PLICKPT.
- Provide a data set on which the checkpoint record can be written.
- Also, to ensure the desired restart activity, you might need to specify the RD parameter in the EXEC or JOB statement (see z/OS JCL Reference).

Note: You should be aware of the restrictions affecting data sets used by your program. These are detailed in the "Bibliography" on page 533.

Requesting a checkpoint record

Each time you want a checkpoint record to be written, you must invoke, from your PL/I program, the built-in subroutine PLICKPT.

```plaintext
CALL PLICKPT((ddname),check-id,org,code)
```

The four arguments are all optional. If you do not use an argument, you need not specify it unless you specify another argument that follows it in the given order. In this case, you must specify the unused argument as a null string (''). The following paragraphs describe the arguments.

**ddname**

is a character string constant or variable specifying the name of the DD
statement defining the data set that is to be used for checkpoint records. If you omit this argument, the system will use the default ddname SYSCHK.

**check-id**

is a character string constant or variable specifying the name that you want to assign to the checkpoint record so that you can identify it later. If you omit this argument, the system will supply a unique identification and print it at the operator’s console.

**org**

is a character string constant or variable with the attributes CHARACTER(2) whose value indicates, in operating system terms, the organization of the checkpoint data set. PS indicates sequential (that is, CONSECUTIVE) organization; PO represents partitioned organization. If you omit this argument, PS is assumed.

**code**

is a variable with the attributes FIXED BINARY (31), which can receive a return code from PLICKPT. The return code has the following values:

- 0  A checkpoint has been successfully taken.
- 4  A restart has been successfully made.
- 8  A checkpoint has not been taken. The PLICKPT statement should be checked.
- 12 A checkpoint has not been taken. Check for a missing DD statement, a hardware error, or insufficient space in the data set. A checkpoint will fail if taken while a DISPLAY statement with the REPLY option is still incomplete.
- 16 A checkpoint has been taken, but ENQ macro calls are outstanding and will not be restored on restart. This situation will not normally arise for a PL/I program.

**Defining the checkpoint data set**

You must include a DD statement in the job control procedure to define the data set in which the checkpoint records are to be placed. This data set can have either CONSECUTIVE or partitioned organization. You can use any valid ddname. If you use the ddname SYSCHK, you do not need to specify the ddname when invoking PLICKPT.

You must specify a data set name only if you want to keep the data set for a deferred restart. The I/O device can be any direct access device.

To obtain only the last checkpoint record, then specify status as NEW (or OLD if the data set already exists). This will cause each checkpoint record to overwrite the previous one.

To retain more than one checkpoint record, specify status as MOD. This will cause each checkpoint record to be added after the previous one.

If the checkpoint data set is a library, “check-id” is used as the member-name. Thus a checkpoint will delete any previously taken checkpoint with the same name.

For direct access storage, you should allocate enough primary space to store as many checkpoint records as you will retain. You can specify an incremental space
allocation, but it will not be used. A checkpoint record is approximately 5000 bytes longer than the area of main storage allocated to the step.

No DCB information is required, but you can include any of the following, where applicable:

OPTCD=W, OPTCD=C, RECFM=UT

These subparameters are described in the *z/OS JCL User’s Guide*.

---

### Requesting a restart

A restart can be automatic or deferred. You can make automatic restarts after a system failure or from within the program itself. The system operator must authorize all automatic restarts when requested by the system.

#### Automatic restart after a system failure

If a system failure occurs after a checkpoint has been taken, the automatic restart will occur at the last checkpoint if you have specified RD=R (or omitted the RD parameter) in the EXEC or JOB statement.

If a system failure occurs before any checkpoint has been taken, an automatic restart, from the beginning of the job step, can still occur if you have specified RD=R in the EXEC or JOB statement.

After a system failure occurs, you can still force automatic restart from the beginning of the job step by specifying RD=RNC in the EXEC or JOB statement. By specifying RD=RNC, you are requesting an automatic step restart without checkpoint processing if another system failure occurs.

#### Automatic restart within a program

You can request a restart at any point in your program. The rules for the restart are the same as for a restart after a system failure. To request the restart, you must execute the statement:

```
CALL PLIREST;
```

To effect the restart, the compiler terminates the program abnormally, with a system completion code of 4092. Therefore, to use this facility, the system completion code 4092 must not have been deleted from the table of eligible codes at system generation.

#### Getting a deferred restart

To ensure that automatic restart activity is canceled, but that the checkpoints are still available for a deferred restart, specify RD=NR in the EXEC or JOB statement when the program is first executed.

```
RESTART=(stepname,check-id)
```

If you subsequently require a deferred restart, you must submit the program as a new job, with the RESTART parameter in the JOB statement. Use the RESTART parameter to specify the job step at which the restart is to be made and, if you want to restart at a checkpoint, the name of the checkpoint record.
For a restart from a checkpoint, you must also provide a DD statement that defines
the data set containing the checkpoint record. The DD statement must be named
SYSCHK. The DD statement must occur immediately before the EXEC statement
for the job step.

**Modifying checkpoint/restart activity**

You can cancel automatic restart activity from any checkpoints taken in your
program by executing the statement:

```
CALL PLICANC;
```

However, if you specified RD=R or RD=RNC in the JOB or EXEC statement,
automatic restart can still take place from the beginning of the job step.

Also, any checkpoints already taken are still available for a deferred restart.

You can cancel any automatic restart and the taking of checkpoints, even if they
were requested in your program, by specifying RD=NC in the JOB or EXEC
statement.
Chapter 22. Using user exits

PL/I provides a number of user exits that allow you to customize the PL/I product to suit your needs. The PL/I products supply default exits and the associated source files.

If you want the exits to perform functions that are different from those supplied by the default exits, we recommend that you modify the supplied source files as appropriate.

At times, it is useful to be able to tailor the compiler to meet the needs of your organization. For example, you might want to suppress certain messages or alter the severity of others. You might want to perform a specific function with each compilation, such as logging statistical information about the compilation into a file. A compiler user exit handles this type of function.

With PL/I, you can write your own user exit or use the exit provided with the product, either 'as is' or modified, depending on what you want to do with it. The user exit source code provided with the product can be seen in Figure 18 on page 172.

The purpose of this chapter is to describe:

- Procedures that the compiler user exit supports
- How to activate the compiler user exit
- IBMUEXIT, the IBM-supplied compiler user exit
- Requirements for writing your own compiler user exit

Procedures performed by the compiler user exit

The compiler user exit performs three specific procedures:

- Initialization
- Interception and filtering of compiler messages
- Termination

As illustrated in Figure 108 on page 462, the compiler passes control to the initialization procedure, the message filter procedure, and the termination procedure. Each of these three procedures, in turn, passes control back to the compiler when the requested procedure is completed.
Each of the three procedures is passed two different control blocks:

- A global control block that contains information about the compilation. This is passed as the first parameter. For specific information on the global control block, see “Structure of global control blocks” on page 463.

- A function-specific control block that is passed as the second parameter. The content of this control block depends upon which procedure has been invoked. For detailed information, see “Writing the initialization procedure” on page 464, “Writing the message filtering procedure” on page 465, and “Writing the termination procedure” on page 467.

Activating the compiler user exit

In order to activate the compiler user exit, you must specify the EXIT compile-time option. For more information on the EXIT option, see “EXIT” on page 33.

The EXIT compile-time option allows you to specify a user-option-string which specifies the DDname for the user exit input file. If you do not specify a string, SYSUEXIT is used as the DDname for the user exit input file.

The user-option-string is passed to the user exit functions in the global control block which is discussed in “Structure of global control blocks” on page 463. See the field “Uex_UIB_User_char_str” in the section “Structure of global control blocks” on page 463 for additional information.

The IBM-supplied compiler exit, IBMUEXIT

IBM supplies you with the sample compiler user exit, IBMUEXIT, which filters messages for you. It monitors messages and, based on the message number that you specify, suppresses the message or changes the severity of the message.

Customizing the compiler user exit

As was mentioned earlier, you can write your own compiler user exit or simply use the one shipped with the compiler. In either case, the name of the fetchable file for the compiler user exit must be IBMUEXIT.
This section describes how to:
- Modify the user exit input file for customized message filtering
- Create your own compiler user exit

**Modifying SYSUEXIT**

Rather than spending the time to write a completely new compiler user exit, you can simplify modify the user exit input file.

Edit the file to indicate which message numbers you want to suppress, and which message number severity levels you want to change. A sample file is shown in Figure 109.

<table>
<thead>
<tr>
<th>Fac Id</th>
<th>Msg No</th>
<th>Severity</th>
<th>Suppress</th>
<th>Comment</th>
</tr>
</thead>
<tbody>
<tr>
<td>'IBM'</td>
<td>1042</td>
<td>-1</td>
<td>1</td>
<td>String spans multiple lines</td>
</tr>
<tr>
<td>'IBM'</td>
<td>1044</td>
<td>-1</td>
<td>1</td>
<td>FIXED BIN 7 mapped to 1 byte</td>
</tr>
<tr>
<td>'IBM'</td>
<td>1047</td>
<td>8</td>
<td>0</td>
<td>Order inhibits optimization</td>
</tr>
<tr>
<td>'IBM'</td>
<td>1052</td>
<td>-1</td>
<td>1</td>
<td>Nodelist with = extent arg</td>
</tr>
<tr>
<td>'IBM'</td>
<td>1059</td>
<td>0</td>
<td>0</td>
<td>Select without OTHERWISE</td>
</tr>
<tr>
<td>'IBM'</td>
<td>1169</td>
<td>0</td>
<td>1</td>
<td>Precision of result determined</td>
</tr>
</tbody>
</table>

*Figure 109. Example of an user exit input file*

The first two lines are header lines and are ignored by IBMUEXIT. The remaining lines contain input separated by a variable number of blanks.

Each column of the file is relevant to the compiler user exit:
- The first column should contain the letters 'IBM' in single quotation marks for all compiler messages to which you want the exit to apply.
- The second column contains the four digit message number.
- The third column shows the new message severity. Severity -1 indicates that the severity should be left as the default value.
- The fourth column indicates whether or not the message is to be suppressed. A '1' indicates the message is to be suppressed, and a '0' indicates that it should be printed.
- The comment field, found in the last column, is for your information, and is ignored by IBMUEXIT.

**Writing your own compiler exit**

To write your own user exit, you can use IBMUEXIT (see the source in Figure 18 on page 172) as a model. As you write the exit, make sure it covers the areas of initialization, message filtering, and termination.

As noted in that section, the compiler user exit must be compiled with the RENT option and linked as a DLL.

**Structure of global control blocks**

The global control block is passed to each of the three user exit procedures (initialization, filtering, and termination) whenever they are invoked. The following code and accompanying explanations describe the contents of each field in the global control block.

```plaintext
dcl
  1 Uex UIB native based(null()),
  2 Uex UIB Length fixed bin(31),
```

Chapter 22. Using user exits 463
Data Entry Fields

- **Uex_UIB_Length**: Contains the length of the control block in bytes. The value is storage (Uex_UIB).
- **Uex_UIB.Exit_token**: Used by the user exit procedure. For example, the initialization may set it to a data structure which is used by both the message filter, and the termination procedures.
- **Uex_UIB_User_char_str**: Points to an optional character string, if you specify it. For example, in pli filename (EXIT ('string'))...fn can be a character string up to thirty-one characters in length.
- **Uex_UIB_char_len**: Contains the length of the string pointed to by the User_char_str. The compiler sets this value.
- **Uex_UIB_Filename_str**: Contains the name of the source file that you are compiling, and includes the drive and subdirectories as well as the filename. The compiler sets this value.
- **Uex_UIB_Filename_len**: Contains the length of the name of the source file pointed to by the Filename_str. The compiler sets this value.
- **Uex_UIB_return_code**: Contains the return code from the user exit procedure. The user sets this value.
- **Uex_UIB_reason_code**: Contains the procedure reason code. The user sets this value.
- **Uex_UIB_Exit_Routs**: Contains the exit entries set up by the initialization procedure.
- **Uex_UIB_Termination**: Contains the entry that is to be called by the compiler at termination time. The user sets this value.
- **Uex_UIB_Message_Filter**: Contains the entry that is to be called by the compiler whenever a message needs to be generated. The user sets this value.

**Writing the initialization procedure**

Your initialization procedure should perform any initialization required by the exit, such as opening files and allocating storage. The initialization procedure-specific control block is coded as follows:

```pli
Dcl 1 Uex_ISA native based( null() ),
2 Uex_ISA_Length_fixed bin(31); /* storage(Uex_ISA) */
```
The global control block syntax for the initialization procedure is discussed in the section “Structure of global control blocks” on page 463.

Upon completion of the initialization procedure, you should set the return/reason codes to the following:

0/0  
Continue compilation

4/n  
Reserved for future use

8/n  
Reserved for future use

12/n  
Reserved for future use

16/n  
Abort compilation

Writing the message filtering procedure

The message filtering procedure permits you to either suppress messages or alter the severity of messages. You can increase the severity of any of the messages but you can decrease the severity only of ERROR (severity code 8) or WARNING (severity code 4) messages.

The procedure-specific control block contains information about the messages. It is used to pass information back to the compiler indicating how a particular message should be handled.

The following is an example of a procedure-specific message filter control block:

Dcl 1 Uex_MFX native based( null() ),
  2 Uex_MFX_Len fixed bin(31),
  2 Uex_MFX_Facility_Id char(3), /* of component writing message */
  2 * char(1),
  2 Uex_MFX_Message_No fixed bin(31),
  2 Uex_MFX_Severity fixed bin(15),
  2 Uex_MFX_New_Severity fixed bin(15), /* set by exit proc */
  2 Uex_MFX_Inserts fixed bin(15),
  2 Uex_MFX_Inserts_Data( 6 refer(Uex_MFX_Inserts) ),
  3 Uex_MFX_Insert_Type fixed bin(7),
  3 Uex_MFX_Insert_Type_Data union unaligned,
  4 * char(8),
  4 Uex_MFX_Insert_Bin8 fixed bin(63),
  4 Uex_MFX_Insert_Bin fixed bin(31),
  4 Uex_MFX_Insert_Str,
    5 Uex_MFX_Insert_Str_Len fixed bin(15),
    5 Uex_MFX_Insert_Str_Addr pointer,
  4 Uex_MFX_Insert_Series,
    5 Uex_MFX_Insert_Series_Sep char(1),
    5 Uex_MFX_Insert_Series_Addr pointer;

Data Entry Fields

• **Uex_MFX_Len**: Contains the length of the control block in bytes. The value is storage (Uex_MFX).

• **Uex_MFX_Facility_Id**: Contains the ID of the facility; for the compiler, the ID is IBM. The compiler sets this value.
- **Uex_MFX_Message_no**: Contains the message number that the compiler is going to generate. The compiler sets this value.
- **Uex_MFX_Severity**: Contains the severity level of the message; it can be from one to fifteen characters in length. The compiler sets this value.
- **Uex_MFX_New_Severity**: Contains the new severity level of the message; it can be from one to fifteen characters in length. The user sets this value.
- **Uex_MFX_Inserts**: Contains the number of inserts for the message; it can range from zero to six. The compiler sets this value.
- **Uex_MFX_Inserts_Data**: Contains fields to describe each of the inserts. The compiler sets these values.
- **Uex_MFX_Ins_Type**: Contains the type of the insert. The possible insert types are:
  - **Uex_Ins_Type_Xb31**: Is used for a FIXED BIN(31) and has the value 1.
  - **Uex_Ins_Type_Char**: Is used for a CHAR string and has the value 2.
  - **Uex_Ins_Type_Series**: Is used for a series of CHAR strings and has the value 3.
  - **Uex_Ins_Type_Xb63**: Is used for a FIXED BIN(63) and has the value 4.
  The compiler sets this value.
- **Uex_MFX_Ins_Bin**: Contains the integer value for an insert that has integer type. The compiler sets this value.
- **Uex_MFX_Ins_Str_Len**: Contains the length (in bytes) for an insert that has character type. The compiler sets this value.
- **Uex_MFX_Ins_Str_Addr**: Contains the address of the character string for an insert that has character type. The compiler sets this value.
- **Uex_MFX_Ins_Series_Sep**: Contains the character that should be inserted between each element for an insert that has series type. Typically, this is a blank, period or comma. The compiler sets this value.
- **Uex_MFX_Ins_Series_Addr**: Contains the address of the series of varying character strings for an insert that has series type. The address points to a FIXED BIN(31) field holding the number of strings to concatenate followed by the addresses of those strings. The compiler sets this value.

Upon completion of the message filtering procedure, set the return/reason codes to one of the following:

0/0  
Continue compilation, output message

0/1  
Continue compilation, do not output message

4/n  
Reserved for future use

8/n  
Reserved for future use

16/n  
Abort compilation
Writing the termination procedure

You should use the termination procedure to perform any cleanup required, such as closing files. You might also want to write out final statistical reports based on information collected during the error message filter procedures and the initialization procedures.

The termination procedure-specific control block is coded as follows:

```
Dcl 1 Uex_ISA native based,
    2 Uex_ISA_Length_fixed bin(31); /* storage(Uex_ISA) */
```

The global control block syntax for the termination procedure is discussed in "Structure of global control blocks" on page 463. Upon completion of the termination procedure, set the return/reason codes to one of the following:

- 0/0  
  Continue compilation
- 4/n  
  Reserved for future use
- 8/n  
  Reserved for future use
- 12/n  
  Reserved for future use
- 16/n  
  Abort compilation
Example of suppressing SQL messages

*Figure 110* shows how to modify the user exit to examine the message inserts and suppress two SQL informational messages and one SQL warning message.

```pli
*Process dft(nodescriptor connected);
*Process or('!') not('!');
*Process limits(exname(31)) rent;

/*************************************************************************/
/* */
/* NAME - IBMUEXIT.PLI */
/* */
/* DESCRIPTION */
/* User-exit sample program. */
/* */
/* Licensed Materials - Property of IBM */
/* 5639-A83, 5639-A24 (C) Copyright IBM Corp. 1992,2011. */
/* All Rights Reserved. */
/* US Government Users Restricted Rights--Use, duplication or */
/* disclosure restricted by GSA ADP Schedule Contract with */
/* IBM Corp. */
/* */
/* DISCLAIMER OF WARRANTIES */
/* The following enclosed code is sample code created by IBM */
/* Copyright. This sample code is not part of any standard */
/* IBM product and is provided to you solely for the purpose of */
/* assisting you in the development of your applications. The */
/* code is provided "AS IS", without warranty of any kind. */
/* IBM shall not be liable for any damages arising out of your */
/* use of the sample code, even if IBM has been advised of the */
/* possibility of such damages. */
/* */
/*************************************************************************/
```

*Figure 110. Suppressing SQL messages*
During initialization, IBMUEXIT is called. It reads information about the messages being screened from a text file and stores the information in a linked list. IBMUEXIT also sets up the entry points for the message filter service and termination service.

For each message generated by the compiler, the compiler calls the message filter registered by IBMUEXIT. The filter looks the message up in the linked list previously created to see if it is one for which some action should be taken.

The termination service is called at the end of the compile but does nothing. It could be enhanced to generates reports or do other cleanup work.

pack: package exports(*);

Dcl
  1 Uex_UIB native Based( null() ),
     Uex_UIB_Length fixed bin(31),
     Uex_UIB_Exit_token pointer, /* for user exit's use */
     Uex_UIB_User_char_str pointer, /* to exit option str */
     Uex_UIB_User_char_len fixed bin(31),
     Uex_UIB_Filename_str pointer, /* to source filename */
     Uex_UIB_Filename_len fixed bin(31),
     Uex_UIB_return_code fixed bin(31), /* set by exit procs */
     Uex_UIB_reason_code fixed bin(31), /* set by exit procs */
  2 Uex_UIB_Exit_Routs, /* exit entries set at initialization */
     limited entry ( Uex_UIB_Termination,
                    Uex_UIB_Message_Filter, /* call for each msg */
                    *,*,*,* )
     limited entry ( *, /* to Uex_UIB */
                    *, /* to a request area */
     );

Suppressing SQL messages (continued)
Dcl 1 Uex_MFX based( null() ),
  2 Uex_MFX_Length fixed bin(31),
  2 Uex_MFX_Facility_Id char(3),
  2 Uex_MFX_Version fixed bin(7),
  2 Uex_MFX_Message_no fixed bin(31),
  2 Uex_MFX_Severity fixed bin(15),
  2 Uex_MFX_New_Severity fixed bin(15),
  2 Uex_MFX_Inserts fixed bin(15),
  3 Uex_MFX_Ins_Type fixed bin(7),
  3 Uex_MFX_Ins_Type_Data union unaligned,
  4 * char(8),
  4 Uex_MFX_Ins_Bin fixed bin(31),
  4 Uex_MFX_Ins_Str,
  5 Uex_MFX_Ins_Str_Len fixed bin(15),
  5 Uex_MFX_Ins_Str_Addr pointer,
  4 Uex_MFX_Ins_Series,
  5 Uex_MFX_Ins_Series_Sep char(1),
  5 Uex_MFX_Ins_Series_Addr pointer;

dcl uex_Ins_Type_Xb31 fixed bin(15) value(1);

dcl uex_Ins_Type_Char fixed bin(15) value(2);

dcl uex_Ins_Type_Series fixed bin(15) value(3);

Dcl 1 Uex_TSA native based( null() ),
  2 Uex_TSA_Length fixed bin(31);

Dcl 1 Uex_Severity_Normal fixed bin(15) value(0);

dcl uex_Severity_Warning fixed bin(15) value(4);

dcl uex_Severity_Error fixed bin(15) value(8);

dcl uex_Severity_Severe fixed bin(15) value(12);

dcl uex_Severity_Unrecoverable fixed bin(15) value(16);

Dcl 1 Uex_Return_Normal fixed bin(15) value(0);

dcl uex_Return_Warning fixed bin(15) value(4);

dcl uex_Return_Error fixed bin(15) value(8);

dcl uex_Return_Severe fixed bin(15) value(12);

dcl uex_Return_Unrecoverable fixed bin(15) value(16);

Suppressing SQL messages (continued)
/** reason codes */

dcl uex_Reason_Output fixed bin(15) value(0);
dcl uex_Reason_Suppress fixed bin(15) value(1);

dcl header pointer;

dcl
1 message_item native based,
2 message_Info,
3 facid char(3),
3 msgno fixed bin(31),
3 newsev fixed bin(15),
3 reason fixed bin(31),
3 variable char(31) var,
2 link pointer;

ibmuexit: proc ( ue, ia ) options( fetchable );

dcl 1 ue like uex_Uib byaddr;
dcl 1 ia like uex_Isa byaddr;

dcl sysuexit file stream input env(recsize(80));
dcl next pointer;
dcl based_Chars char(8) based;
dcl title_Str char(8) var;
dcl eof bit(1);

on error
begin;
on error system;
call plidump('TFBHS');
end;

on undefinedfile(sysuexit)
begin;
put edit ('** User exit unable to open exit file ')
(A) skip;
put skip;
signal error;
end;

if ue.uex_Uib_User_Char_Len = 0 then
do;
open file(sysuexit);
end;
else

if ue.uex_Uib_User_Char_Len = 0 then
do;
open file(sysuexit);
end;
else

title_Str = substr( ue.uex_Uib_User_Char_Str->based_Chars,
1, ue.uex_Uib_User_Char_Len );
open file(sysuexit) title(title_Str);
end;

Suppressing SQL messages (continued)
ue.Uex_UIB_Message_Filter = message_filter;

header = sysnull();
allocate message_item set(next);

next->link = header;
header = next;

allocate message_item set(next);
get file(sysuexit) skip;
get file(sysuexit) list(next->message_info);
dowhile( eof = '0'b );

next->link = header;
header = next;

allocate message_item set(next);
get file(sysuexit) skip;
get file(sysuexit) list(next->message_info);
end;

Suppressing SQL messages (continued)
free next->message_Item;
close file(sysuexit);

message_Filter: proc ( ue, mf );

dcl 1 ue like uex_Uib byaddr;
dcl 1 mf like uex_Mfx byaddr;

dcl next pointer;
dcl jx fixed bin(31);
dcl insert char(256) var;
dcl based_Chars char(256) based;

on error
begin;
  on error system;
  call plidump('TFBHS');
end;

//***************************************************************/
/*
/* by default, leave the reason code etc unchanged */
/*
/*****************************************************************************/

ue.uex_Uib_Reason_Code = uex_Reason_Output;
ue.uex_Uib_Return_Code = 0;

mf.uex_Mfx_New_Severity = mf.uex_Mfx_Severity;

//***************************************************************/
/*
/* save the first insert if it has character type */
/*
/*****************************************************************************/

insert = '*';
if mf.Uex_MFX_Length < stg(mf) then;
else
  if mf.Uex_MFX_Inserts = 0 then;
  else
    do jx = 1 to mf.Uex_MFX_Inserts;
      select( mf.Uex_MFX_Ins_Type(jx) );
      when( uex_Ins_Type_Char )
      do;
        if jx = 1 then
          insert = substr( mf.Uex_MFX_Ins_Str_Addr(jx)->based_Chars,
                          1,mf.Uex_MFX_Ins_Str_Len(jx));
          end;
        otherwise;
        end;
    end;
  end;
end;

Suppressing SQL messages (continued)
search list:
do next = header repeat( next->link ) while( next != sysnull() );

if next->msgno = mf.uex_Mfx_Message_No
 & next->facid = mf.Uex_Mfx_Facility_Id then
   do;
   if next->variable = '*' then
     leave search_list;
   if next->variable
     = substr(insert,1,length(next->variable)) then
     leave search_list;
   end;
end;

if next = sysnull() then;
else
  do;
    ue.uex_Uib_Reason_Code = next->reason;
    if next->newsev < 0 then;
      else
        mf.uex_Mfx_New_Severity = next->newsev;
      end;
    end;
exitterm: proc ( ue, ta );

dcl 1 ue like uex_Uib byaddr;
dcl 1 ta like uex_Tsa byaddr;

ue.uex_Uib_return_Code = 0;
ue.uex_Uib_reason_Code = 0;
end;
end;

Suppressing SQL messages (continued)
Chapter 23. PL/I descriptors

This chapter describes PL/I parameter passing conventions between PL/I routines at run time. For additional information about Language Environment run-time environment considerations, other than descriptors, see z/OS Language Environment Programming Guide. This includes run-time environment conventions and assembler macros supporting these conventions.

Passing an argument

When a string, an array, or a structure is passed as an argument, the compiler passes a descriptor for that argument unless the called routine is declared with OPTIONS(NODESCRIPTOR). There are two methods for passing such descriptors:

- By descriptor list
- By descriptor locator

The following key features should be noted about each of these two methods:

- **When arguments are passed with a descriptor list**
  - The number of arguments passed is one greater than the number of arguments specified if any of the arguments needs a descriptor.
  - An argument passed with a descriptor can be received as a pointer passed by value (BYVALUE).
  - The compiler uses this method when the DEFAULT(DESCLIST) compiler option is in effect.

- **When arguments are passed by descriptor locator**
  - The number of arguments passed always matches the number of arguments specified.
  - An argument passed with a descriptor can be received as a pointer passed by address (BYADDR).
  - The compiler uses this method when the DEFAULT(DESCLOCATOR) compiler option is in effect.

Argument passing by descriptor list

When arguments and their descriptors are passed with a descriptor list, an extra argument is passed whenever at least one argument needs a descriptor. This extra argument is a pointer to a list of pointers. The number of entries in this list equals the number of arguments passed. For arguments that don’t require a descriptor, the corresponding pointer in the descriptor list is set to SYSNULL. For arguments that do require a descriptor, the corresponding pointer in the descriptor list is set to the address of that argument’s descriptor.

So, for example, suppose the routine sample is declared as

```plaintext
declare sample entry( fixed bin(31), varying char(*) )
oneptions( byaddr descriptor );
```

Then, if sample is called as in the following statement:

```plaintext
call sample( 1, 'test' );
```

The following three arguments are passed to the routine:

- Address of a fixed bin(31) temporary with the value 1
• Address of a varying char(4) temporary with the value test
• Address of a descriptor list consisting of the following:
  – SYSNULL()
  – Address of the descriptor for a varying char(4) string

**Argument passing by descriptor-locator**

When arguments and their descriptors are passed by descriptor-locator, whenever an argument requires a descriptor, the address of a locator/descriptor for it is passed instead.

Except for strings, the locator/descriptor is a pair of pointers. The first pointer is the address of the data; the second pointer is the address of the descriptor. For strings, under CMPAT(LE), the locator/descriptor is still such a pair of pointers. But under the other CMPAT options, the locator/descriptor consists of the address of the string and then the string descriptor itself.

So, for example, suppose the routine `sample` is declared again as

```pli
declare sample entry( fixed bin(31), varying char(*) )
  options( byaddr descriptor );
```

Then, if `sample` is called as in the following statement

```pli
call sample( 1, 'test' );
```

The following two arguments are passed to the routine:

• The address of a fixed bin(31) temporary with the value 1
• The address of a locator-descriptor consisting of the following:
  – The address of a varying char(4) temporary with the value test
  – Under CMPAT(LE), the address of the CMPAT(LE) descriptor for a varying char(4) string
  – Under CMPAT(V*), the CMPAT(V*) descriptor for a varying char(4) string

### CMPAT(V*) descriptors

Unlike LE descriptors, the CMPAT(V*) descriptors are not self-describing. However, the string descriptors are the same for all CMPAT(V*) options, and they also share the same codepage encoding as the LE string descriptors.

### String descriptors

In a string descriptor, the first 2 bytes specify the maximum length for the string. This maximum length is always held in native format.

The third byte contains various flags (to indicate, for example, if the string length in a VARYING string is held in littleendian or bigendian format or if the data in a WIDECHAR string is held in littleendian or bigendian format).

In a string descriptor for a nonvarying bit string, the fourth byte gives the bit offset.

In a string descriptor for a CHARACTER string, the fourth byte encodes the compiler CODEPAGE option.

The declare for a string descriptor under CMPAT(V1) or CMPAT(V2) is:
The declare for a string descriptor under CMPAT(V3) is:

```
declare
  1 dso_longstr based( null() ),
  2 dso_longstr_info,
  3 * fixed bin(8) unsigned,
  3 dso_datatype fixed bin(8) unsigned,
  3 * union,
  4 dso_longstr_bitofs fixed bin(8) unsigned,
  4 dso_longstr_codepage type ccs_Codepage_Enum,
  3 dso_longstr_info2,
  4 dso_longstr_has_nonnative_len bit(1), /* for varying */
  4 dso_longstr_is_ebcdic bit(1), /* for char */
  4 dso_longstr_has_nonnative_data bit(1), /* for wchar */
  4 * bit(1),
  2 dso_longstr_length fixed bin(31),
  2 dso_longstr_end char(0);
```

The possible values for the codepage encoding are defined via:

```
define ordinal
ccs_Codepage_Enum
  ( ccs_Codepage_01047 value(1)
    ,ccs_Codepage_01140
    ,ccs_Codepage_01141
    ,ccs_Codepage_01142
    ,ccs_Codepage_01143
    ,ccs_Codepage_01144
    ,ccs_Codepage_01145
    ,ccs_Codepage_01146
    ,ccs_Codepage_01147
    ,ccs_Codepage_01148
    ,ccs_Codepage_01149
    ,ccs_Codepage_00819
    ,ccs_Codepage_00813
    ,ccs_Codepage_00920
    ,ccs_Codepage_00037
    ,ccs_Codepage_00273
    ,ccs_Codepage_00277
    ,ccs_Codepage_00278
    ,ccs_Codepage_00280
    ,ccs_Codepage_00284
    ,ccs_Codepage_00285
    ,ccs_Codepage_00297
    ,ccs_Codepage_00500
    ...
    ...
  )
```
Array descriptors

In the following declares, the upper bound for the arrays is declared as 15, but it should be understood that the actual upper bound will always match the number of dimensions in the array it describes.

The declare for a CMPAT(V1) array descriptor is:

```
declare
1 dso_v1 based( null() ),
  2 dso_v1_rvo fixed bin(31), /* relative virtual origin */
  2 dso_v1_data(1:15),
  3 dso_v1_stride fixed bin(31), /* multiplier */
  3 dso_v1_hbound fixed bin(15), /* hbound */
  3 dso_v1_lbound fixed bin(15); /* lbound */
```

The declare for a CMPAT(V2) array descriptor is:

```
declare
1 dso_v2 based( null() ),
  2 dso_v2_rvo fixed bin(31), /* relative virtual origin */
  2 dso_v2_data(1:15),
  3 dso_v2_stride fixed bin(31), /* multiplier */
  3 dso_v2_hbound fixed bin(15), /* hbound */
  3 dso_v2_lbound fixed bin(15); /* lbound */
```

The declare for a CMPAT(V3) array descriptor is:

```
declare
1 dso_v3 based( null() ),
  2 dso_v3_rvo fixed bin(63), /* relative virtual origin */
  2 dso_v3_data(1:15),
  3 dso_v3_stride fixed bin(63), /* multiplier */
  3 dso_v3_hbound fixed bin(63), /* hbound */
  3 dso_v3_lbound fixed bin(63); /* lbound */
```

CMPAT(LE) descriptors

Every LE descriptor starts with a 4-byte field. The first byte specifies the descriptor type (scalar, array, structure or union). The remaining three bytes are zero unless they are set by the particular descriptor type.

The declare for a descriptor header is:

```
declare
1 dsc_Header based( sysnull() ),
  2 dsc_Type fixed bin(8) unsigned,
  2 dsc_Datatype fixed bin(8) unsigned,
  2 * fixed bin(8) unsigned,
  2 * fixed bin(8) unsigned;
```

The possible values for the dsc_Type field are:

```
declare
dsc_Type_Unset fixed bin(8) value(0),
  dsc_Type_Element fixed bin(8) value(2),
  dsc_Type_Array fixed bin(8) value(3),
  dsc_Type_Structure fixed bin(8) value(4),
  dsc_Type_Union fixed bin(8) value(4);
```
String descriptors

In a string descriptor, the second byte of the header indicates the string type (bit, character or graphic as well as nonvarying, varying or varyingz).

In a string descriptor for a nonvarying bit string, the third byte of the header gives the bit offset.

In a string descriptor for a CHARACTER string, the third byte of the header encodes the compiler CODEPAGE option.

In a string descriptor for a varying string, the fourth byte has a bit indicating if the string length is held in nonnative format.

In a string descriptor for a character string, the fourth byte also has a bit indicating if the string data is in EBCDIC.

The declare for a string descriptor is:

\[
\text{declare}
\]
\[
\begin{align*}
1 & \text{dsc\_String based( sysnull() ),} \\
2 & \text{dsc\_String\_Header,} \\
3 & \text{dsc\_String\_Type fixed bin(8) unsigned,} \\
4 & \text{dsc\_String\_Codepage ordinal ccs\_Codepage\_Enum,} \\
5 & \text{dsc\_String\_BitOfs fixed bin(8) unsigned,} \\
6 & \text{dsc\_String\_Has\_Nonnative\_Len bit(1),} \\
7 & \text{dsc\_String\_Is\_Ebcdic bit(1),} \\
8 & \text{dsc\_String\_Has\_Nonnative\_Data bit(1),} \\
9 & \text{dsc\_String\_Length fixed bin(31); /* max length of string */}
\end{align*}
\]

The possible values for the dsc\_String\_Type field are:

\[
\text{declare}
\]
\[
\begin{align*}
\text{dsc\_String\_Type\_Unset fixed bin(8) value(0),} \\
\text{dsc\_String\_Type\_Char\_Nonvarying fixed bin(8) value(2),} \\
\text{dsc\_String\_Type\_Char\_Varyingz fixed bin(8) value(3),} \\
\text{dsc\_String\_Type\_Char\_Varying2 fixed bin(8) value(4),} \\
\text{dsc\_String\_Type\_Bit\_Nonvarying fixed bin(8) value(6),} \\
\text{dsc\_String\_Type\_Bit\_Varying2 fixed bin(8) value(7),} \\
\text{dsc\_String\_Type\_Graphic\_Nonvarying fixed bin(8) value(9),} \\
\text{dsc\_String\_Type\_Graphic\_Varyingz fixed bin(8) value(10),} \\
\text{dsc\_String\_Type\_Graphic\_Varying2 fixed bin(8) value(11),} \\
\text{dsc\_String\_Type\_Widechar\_Nonvarying fixed bin(8) value(13),} \\
\text{dsc\_String\_Type\_Widechar\_Varyingz fixed bin(8) value(14),} \\
\text{dsc\_String\_Type\_Widechar\_Varying2 fixed bin(8) value(15);}
\end{align*}
\]

Array descriptors

The declare for an array descriptor is:

\[
\text{declare}
\]
\[
\begin{align*}
1 & \text{dsc\_Array based( sysnull() ),} \\
2 & \text{dsc\_Array\_Header like dsc\_Header,} \\
3 & \text{dsc\_Array\_EltLen fixed bin(31), /* Length of array element */} \\
4 & \text{dsc\_Array\_Rank fixed bin(31), /* Count of dimensions */} \\
5 & \text{dsc\_Array\_RVO fixed bin(31), /* Relative virtual origin */} \\
6 & \text{dsc\_Array\_Data( 1: 1 refer(dsc\_Array\_Rank) ),} \\
7 & \text{dsc\_Array\_LBound fixed bin(31), /* LBound */} \\
8 & \text{dsc\_Array\_Extent fixed bin(31), /* HBound - LBound + 1 */} \\
9 & \text{dsc\_Array\_Stride fixed bin(31); /* Multiplier */}
\end{align*}
\]
Part 6. Appendixes
Appendix. SYSADATA message information

When you specify the MSG suboption of the XINFO compile-time option, the compiler generates a SYSADATA file that contains:

- Counter records
- Literal records
- File records
- Message records

You should note that the records in the file are not necessarily produced in the order listed above, for example, literal and file records might be interleaved. If you are writing code that reads a SYSADATA file, you should not rely on the order of the records in the file except for the following exceptions:

- Counter records are the first records in the file.
- Each literal record precedes any reference to the literal it defines.
- Each file record precedes any reference to the file it describes.

Understanding the SYSADATA file

The SYSADATA file is a sequential binary file. Under z/OS batch, the compiler writes the SYSADATA records to the file specified by the SYSADATA DD statement, and that file must not be a member of a PDS. On all other systems, the compiler writes to a file with the extension "adt".

Each record in the file contains a header. This 8-byte header has fields that are the same for all records in the file:

**Compiler**
- A number representing the compiler that produced the data. For PL/I, the number is 40.

**Edition number**
- The edition number of the compiler that produced the data. For this product, it is the number 2.

**SYSADATA level**
- A number representing the level of SYSADATA that this file format represents. For this product, it is the number 4.

The header also has some fields that vary from record to record:

- Record type
- Whether or not the record is continued onto the next record

Possible record types are encoded as an ordinal value as shown in Figure 111 on page 484.
The declare for the header part of a record is shown in Figure 112.

```
Define ordinal xin_Rect
(Xin_Rect_Msg value(50), /* Message record */
 Xin_Rect_Fil value(57), /* File record */
 Xin_Rect_Sum value(61), /* Summary record */
 Xin_Rect_Rep value(62), /* Replace record */
 Xin_Rect_Src value(63), /* Source record */
 Xin_Rect_Tok value(64), /* Token record */
 Xin_Rect_Sym value(66), /* Symbol record */
 Xin_Rect_Lit value(67), /* Literal record */
 Xin_Rect_Syn value(69), /* Syntax record */
 Xin_Rect_Ord_Type value(80), /* ordinal type record */
 Xin_Rect_Ord_Elem value(81), /* ordinal element record */
 Xin_Rect_Ctr value(82) ) /* counter record */
prec(15);
```

Figure 111. Record types encoded as an ordinal value

The declare for the header part of a record is shown in Figure 112.

```
Dcl
1 Xin_Hdr Based( null() ), /* Header portion */
   /* */
2 Xin_Hdr_Prod /* Language code */
   fixed bin(8) unsigned, /* */
   /* */
2 Xin_Hdr_Rect /* Record type */
   unal ordinal xin_Rect, /* */
   /* */
2 Xin_Hdr_Level /* SYSADATA level */
   fixed bin(8) unsigned, /* */
   /* */
2 * union, /* */
3 Xin_Hdr_Flags bit(8), /* flags */
3 /* */
4 * bit(6), /* Reserved */
4 Xin_Hdr_LittleEndian /* ints are little endian */
   bit(1), /* */
4 Xin_Hdr_Cont bit(1), /* Record continued in next rec */
   /* */
2 Xin_Hdr_Edition /* compiler "edition" */
   fixed bin(8) unsigned, /* */
   /* */
2 Xin_Hdr_Fill bit(32), /* reserved */
   /* */
2 Xin_Hdr_Data_Len /* length of data part */
   fixed bin(16) unsigned, /* */
   /* */
2 Xin_Hdr_End char(0); /* */
```

Figure 112. Declare for the header part of a record

**Summary Record**

The record with type Xin_Rect_Sum, the summary record, will be the first record in the file and its declare is shown in Figure 113 on page 485.
Counter records

Each counter record specifies, for a subsequent record type, how many records of that type the file contains and how many bytes those records occupy.

Literal records

Each literal record assigns a number, called a literal index, that is used by later records to refer to the characters named in this particular record.
File records

Each file record assigns a number, called a file index, that is used by later records to refer to the file described by this record. The described file may be the primary PL/I source file or an INCLUDEd file. Each file record specifies a literal index for the fully qualified name of the file.

For an INCLUDEd file, each file record also contains the file index and source line number from whence the INCLUDE request came. (For primary source files, these fields are zero.)

```
Dcl
1 xin_Lit Based( null() ), /* literal record */
   /*
   2 xin_LitHdr /* standard header */
   /* like xin_Hdr, */
   /* */
   /*
   2 xin_Lit_Inx /* adata index for literal */
   fixed bin(31) unsigned, /* */
   /* */
   /*
   2 xin_Lit_Len /* length of literal */
   fixed bin(31) unsigned, /* */
   /* */
   /*
   2 xin_Lit_Val char(2000); /* literal value */
Figure 115. Declare for a literal record
```

```
Dcl
1 xin_Fil Based( null() ), /* file record */
   /*
   2 xin_FilHdr /* standard header */
   /* like xin_Hdr, */
   /* */
   /*
   2 xin_Fil_File_Id /* file id from whence it */
   fixed bin(31) unsigned, /* was INCLUDEd */
   /* */
   /*
   2 xin_Fil_Line_No /* line no within that file */
   fixed bin(31) unsigned, /* */
   /* */
   /*
   2 xin_Fil_Id /* id assigned to this file */
   fixed bin(31) unsigned, /* */
   /* */
   /*
   2 xin_Fil_Name /* literal index of the */
   fixed bin(31) unsigned; /* fully qualified file name */
Figure 116. Declare for a file record
```

Message records

Each message record describes a message issued during the compilation. Message records are not generated for suppressed messages.

Each message record contains:

- The file index and source line number for the file and line to which the message is attributed. If the message pertains to the compilation as a whole, these fields are zero.
• The identifier (e.g. IBM1502) and severity associated with the message.
• The text of the message.

The declare for a message record is as follows:

```
Dcl
  1 xin_Msg  Based( null() ), /* message record */
  2 xin_Msg_Hdr /* standard header */
    like xin_Hdr, /* */
  2 xin_Msg_File_Id /* file id */
    fixed bin(31) unsigned, /* */
  2 xin_Msg_Line_No /* line no within file */
    fixed bin(31) unsigned, /* */
  2 xin_Msg_Id /* identifier (i.e. IBM1502) */
    char(16), /* */
  2 xin_Msg_Severity /* severity (0, 4, 8, 12 or 16) */
    fixed bin(15) signed, /* */
  2 xin_Msg_Length /* length of message */
    fixed bin(16) unsigned, /* */
  2 Xin_Msg_Text /* actual message */
    char(100 refer(xin_Msg_Length));
```

Figure 117. Declare for a message record

Understanding SYSADATA symbol information

When you specify the SYM suboption of the XINFO compile-time option, the compiler generates a SYSADATA file that contains the following information in addition to the records generated for the MSG suboption:

• Ordinal type records
• Ordinal element records
• Symbol records

Symbol records are not generated for built-in functions, generic variables or variables with non-constant extents.

Ordinal type records

Each ordinal type record assigns a number, called an ordinal type index, that is used by later records to refer to an ordinal type described by this record. The name of the type is indicated by a literal index. Each ordinal type record contains the file index and source line number for the file and line in which the ordinal type was declared.

Each ordinal type record contains:
• A count of the number of values defined by that type
• The precision associated with the type
• Bits indicating if it is signed or unsigned
Ordinal element records

Each ordinal type record is immediately followed by a series of records (as many as specified by the ordinal type count) that describes the values named by that ordinal.

Each ordinal element record assigns a number, called an ordinal element index, that is used by later records to refer to an ordinal element described by this record. The name of the element is indicated by a literal index. Each ordinal element record contains the file index and source line number for the file and line in which the ordinal element was declared.

Additionally, each ordinal element record contains

- The ordinal type index of the ordinal type to which it belongs
- The value for that element
Symbol records

Each symbol record assigns a number that is called a symbol index. The index is used by later records to refer to the symbol described by this record. For example, the index can be used as the name of a user variable or constant. The name of the identifier is indicated by a literal index. Each symbol record contains the file index and source line number for the file and line in which the symbol was declared.

If the identifier is part of a structure or union, the symbol record contains a symbol index for each of the following:
- The first sibling, if any
- The parent, if any
- The first child, if any

Consider the following structure:

```plaintext
dcl
  1 a,
    3 b fixed bin,
    3 c fixed bin,
    3 d,
    5 e fixed bin,
    5 f fixed bin;
```

The symbol indices assigned to the elements of the preceding structure would be as follows:

```
declare /* */
  1 xin_Ord_Elem based( null() ), /* */
  2 xin_Ord_Elem_Hdr /* standard header */
    like xin_Hdr,
    /* */
  2 xin_Ord_Elem_File_Id /* file id */
    fixed bin(31) unsigned,
    /* */
  2 xin_Ord_Elem_Line_No /* line no within file */
    fixed bin(31) unsigned,
    /* */
  2 xin_Ord_Elem_Id /* identifying number */
    fixed bin(31),
    /* */
  2 xin_Ord_Elem_Type_Id /* id of ordinal type */
    fixed bin(31),
    /* */
  2 xin_Ord_Elem_Value /* ordinal value */
    fixed bin(31),
    /* */
  2 xin_Ord_Elem_Name /* ordinal name */
    fixed bin(31); /* */
```

Figure 119. Declare for an ordinal element record
Each symbol record also contains a series of bit(1) fields that indicate if various attributes apply to this variable.

Each symbol record also contains the following elements:

**User-given structure level**
This is a user-given structure level for the identifier. For the element c of the structure above, the value is 3. For non-structure members, the value is set to 1.

**Logical structure level**
The logical structure level for the identifier. For the element c of the structure above, the value is 2. For non-structure members, the value is set to 1.

**Dimensions**
The number of dimensions declared for the variable not counting any inherited dimensions.

The number of dimensions for the variable including all inherited dimensions.

**Offset**
The offset into the outermost parent structure.

**Elemental size**
Elemental size is in bytes unless the variable is bit aligned, in which case it is in bits. In either case, this does not factor any in dimensions.

**Size**
Size in bytes with its dimensions factored in.

**Alignment**
Identified by the following:

- 0 for bit-aligned
- 7 for byte-aligned
- 15 for halfword-aligned
- 31 for fullword-aligned
- 63 for quadword-aligned

A union within the record is dedicated to describing information that is dependent on the variable’s storage class:

**Static variables**
If the variable is declared as external with a separate external name (dcl x ext('y')), the literal index of that name is specified.

**Based variables**
If the variable is declared as based on another mapped variable that is not an element of an array, the symbol index of that variable is specified.
Defined variables

If the variable is declared as defined on another mapped variable that is not an element of an array, the symbol index of that variable is specified here. If its position attribute is constant, it is also specified.

The variable's data type is specified by the ordinal shown in Figure 121.

```c
define
    ordinal
    xin_Data_Kind
    (    xin_Data_Kind_Unset
         ,xin_Data_Kind_Character
         ,xin_Data_Kind_Bit
         ,xin_Data_Kind_Graphic
         ,xin_Data_Kind_Fixed
         ,xin_Data_Kind_Float
         ,xin_Data_Kind_Picture
         ,xin_Data_Kind_Pointer
         ,xin_Data_Kind_Offset
         ,xin_Data_Kind_Entry
         ,xin_Data_Kind_File
         ,xin_Data_Kind_Label
         ,xin_Data_Kind_Format
         ,xin_Data_Kind_Area
         ,xin_Data_Kind_Task
         ,xin_Data_Kind_Event
         ,xin_Data_Kind_Econdition
         ,xin_Data_Kind_Structure
         ,xin_Data_Kind_Union
         ,xin_Data_Kind_Descriptor
         ,xin_Data_Kind_Ordinal
         ,xin_Data_Kind_Handle
         ,xin_Data_Kind_Type
         ,xin_Data_Kind_Builtin
         ,xin_Data_Kind_Generic
         ,xin_Data_Kind_Widechar
     ) prec(8) unsigned;
```

Figure 121. Data type of a variable

A union within the record is dedicated to describing information that is dependent on the variable's data type. Most of this information is self-explanatory (for example, the precision for an arithmetic type) except perhaps for the following:

Picture variables

The literal index of the picture specification is specified.

Entry variables

If the variable has the returns attribute, the symbol index of the returns description is specified.

Ordinal variables

The ordinal type index is specified.

Typed variables and handles

The symbol index of the underlying type is specified.

String and area variables

The type and value of the extent is specified in addition to the symbol index of the returns description. The type of the extent is encoded by the values:
If the element has any dimensions, the type and values for its lower and upper bounds are specified at the very end of the record. These fields are not present if the element has no dimensions. Figure 122 shows the declare for a symbol record.

```pli
declare
  ( xin_Extent_Constant value(01),
    xin_Extent_Star value(02),
    xin_Extent_Nonconstant value(04),
    xin_Extent_Refer value(08),
    xin_Extent_In_Error value(16)
  )
fixed bin;
```

Figure 122. Declare for a symbol record
Declare for a symbol record (continued)
Declare for a symbol record (continued)
Declare for a symbol record (continued)
Declare for a symbol record (continued)
Declare for a symbol record (continued)
Declare for a symbol record (continued)
The definition of the ordinal xin_Bif_Kind can be found in Figure 123.

```c
define
ordinal
xin_Bif_Kind
  ("xin_Bif_Unknown
   ,xin_bif_abs
   ,xin_bif_acos
   ,xin_bif_add
   ,xin_bif_addr
   ,xin_bif_all
   ,xin_bif_allocation
   ,xin_bif_allocn
   ,xin_bif_any
   ,xin_bif_asin
   ,xin_bif_atan
   ,xin_bif_atand
   ,xin_bif_atanh
   ,xin_bif_bin
   ,xin_bif_binvalue
   ,xin_bif_binary
   ,xin_bif_binaryvalue
   ,xin_bif_bit
   ,xin_bif_bool
   ,xin_bif ceil
   ,xin_bif_char
   ,xin_bif_completion
   ,xin_bif_complex
   ,xin_bif_conjg
   ,xin_bif_copy
   ,xin_bif cos
   ,xin_bif cosd
   ,xin_bif_cosh
   ,xin_bif_count
   ,xin_bif_cpln
   ,xin_bif_cplx
   ,xin_bif_cstg
   ,xin_bif_currentstorage
   ,xin_bif_datafield
   ,xin_bif_date
   ,xin_bif_datetime
   ,xin_bif_dec
   ,xin_bif_decimal
   ,xin_bif_dim
   ,xin_bif_divide
   ,xin_bif_empty
   ,xin_bif_entryaddr
   ,xin_bif_erf
   ,xin_bif_erfc
   ,xin_bif_exp
   ,xin_bif_fixed
   ,xin_bif_float
   ,xin_bif_floor
   ,xin_bif_graphic
   ,xin_bif_hbound

Figure 123. Declare for xin_Bif_Kind
```
Declare for xin_Bif_Kind (continued)
Declare for xin_Bif_Kind (continued)
Declare for xin_Bif_Kind (continued)
Declare for xin_Bif_Kind (continued)
Declare for xin_Bif_Kind (continued)
Declare for xin_Bif_Kind (continued)

Note that the attributes flags reflect the attributes after the compiler has applied all
defaults. So, for example, every numeric variable (including numeric PICTURE
variables) has either the REAL or COMPLEX attribute flag set.

Understanding SYSADATA syntax information

When you specify the SYN suboption of the XINFO compile-time option, the
compiler generates a SYSADATA file that contains the following information in
addition to the records generated for the MSG and SYM suboptions:

- Source records
- Token records
- Syntax records

Source records

Each source record assigns a number, called a source id, that is used by later
records to refer to the source line described by this record. The line may be from
the primary PL/1 source file or an INCLUDEd file, as indicated by the source file
id and linenumber fields in the record. The rest of the record holds the actual data
in the source line.
Token records

Each token record assigns a number, called a token index, that is used by later records to refer to a token recognized by the PL/I compiler. The record also identifies the type of the token plus the column and line on which it started and ended.

```
Dcl
1 Xin_Src    Based( null() ), /* source record */
  /* */
2 Xin_Src_Hdr /* standard header */
  /* */
  like Xin_Hdr, /* */
  /* */
2 Xin_Src_File_Id /* file id */
  fixed bin(32) unsigned, /* */
  /* */
2 Xin_Src_Line_No /* line no within file */
  fixed bin(32) unsigned, /* */
  /* */
2 Xin_Src_Id /* id for this source record */
  fixed bin(32) unsigned, /* */
  /* */
2 Xin_Src_Length /* length of text */
  fixed bin(16) unsigned, /* */
  /* */
2 Xin_Src_Text /* actual text */
  char(137 refer(xin_Src_Length));
```

**Figure 124. Declare for a source record**

```
Dcl
1 Xin_Tok    Based( null() ), /* token record */
  /* */
2 Xin_Tok_Hdr /* standard header */
  /* */
  like Xin_Hdr, /* */
  /* */
2 Xin_Tok_Inx /* adata index for token */
  fixed bin(32) unsigned, /* */
  /* */
2 Xin_Tok_Begin_Line /* starting line no within file */
  fixed bin(32) unsigned, /* */
  /* */
2 Xin_Tok_End_Line_Offset /* offset of end line from first */
  fixed bin(16) unsigned, /* */
  /* */
2 Xin_Tok_KIND_Value /* token kind */
  ordinal xin_Tok_KIND, /* */
  /* */
2 Xin_Tok_Rsvd /* reserved */
  fixed bin(8) unsigned, /* */
  /* */
2 Xin_Tok_Begin_Col /* starting column */
  fixed bin(16) unsigned, /* */
  /* */
2 Xin_Tok_End_Col /* ending column */
  fixed bin(16) unsigned; /* */
```

**Figure 125. Declare for a token record**
The ordinal xin_Tok_Kind identifies the type of the token record.

```c
Define
   ordinal
      xin_Tok_Kind
      ( xin_Tok_Kind_Unset
        ,xin_Tok_Kind_Lexeme
        ,xin_Tok_Kind_Comment
        ,xin_Tok_Kind_Literal
        ,xin_Tok_Kind_Identifier
        ,xin_Tok_Kind_Keyword
      ) prec(8) unsigned;
```

Figure 126. Declare for the token record kind

**Syntax records**

Each syntax record assigns a number, called a node id, that is used by later records to refer to other syntax records.

The first syntax record will have kind xin_Syn.Kind.Package, and if the compilation unit has any procedures, the child node of this record will point to the first of these procedures. The parent, sibling and child nodes will then provide a map with the appropriate relationships of all the procedures and begin bocks in the compilation unit.

Consider the following simple program:

```plaintext
a: proc;
call b;
call c;
b: proc;
end b;
c: proc;
call d;
d: proc;
end d;
end c;
end a;
```

The node indices assigned to the blocks of the preceding program would be as follows:

<table>
<thead>
<tr>
<th>symbol</th>
<th>index</th>
<th>sibling</th>
<th>parent</th>
<th>child</th>
</tr>
</thead>
<tbody>
<tr>
<td>-</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>2</td>
</tr>
<tr>
<td>a</td>
<td>2</td>
<td>0</td>
<td>1</td>
<td>3</td>
</tr>
<tr>
<td>b</td>
<td>3</td>
<td>4</td>
<td>2</td>
<td>0</td>
</tr>
<tr>
<td>c</td>
<td>4</td>
<td>0</td>
<td>2</td>
<td>5</td>
</tr>
<tr>
<td>d</td>
<td>5</td>
<td>0</td>
<td>4</td>
<td>0</td>
</tr>
</tbody>
</table>

Figure 127. Node indices assigned to the blocks in a program
Dcl
  1 Xin_Syn Based( null() ), /* syntax record */
    / * */
  2 Xin_Syn_Hdr /* standard header */
    like Xin_Hdr,
    / * */
  2 Xin_Syn_Node_Id /* node id */
    fixed bin(32) unsigned, /* */
    / * */
  2 Xin_Syn_Node_Kind /* node type */
    ordinal xin_syn_kind, /* */
    / * */
  2 Xin_Syn_Node_Exp_Kind /* node sub type */
    ordinal xin_exp_kind, /* */
    / * */
  2 * /* reserved */
    fixed bin(16) unsigned, /* */
    / * */
  2 Xin_Syn_Parent_Node_Id /* node id of parent */
    fixed bin(32) unsigned, /* */
    / * */
  2 Xin_Syn_Sibling_Node_Id /* node id of sibling */
    fixed bin(32) unsigned, /* */
    / * */
  2 Xin_Syn_Child_Node_Id /* node id of child */
    fixed bin(32) unsigned, /* */
    / * */
  2 xin_Syn_First_Tok /* id of first spanned token */
    fixed bin(32) unsigned, /* */
    / * */
  2 xin_Syn_Last_Tok /* id of last spanned token */
    fixed bin(32) unsigned, /* */

Figure 128. Declare for a syntax record
Declare for a syntax record (continued)
The ordinal xin_Syn.Kind identifies the type of the syntax record.
Define
ordinal
xin_Syn.Kind
( xin_Syn.Kind_Unset
, xin_Syn.Kind_Lexeme
, xin_Syn.Kind_Asterisk
, xin_Syn.Kind_Int
, xin_Syn.Kind_Name
, xin_Syn.Kind_Expression
, xin_Syn.Kind_Parenthesized_Expr
, xin_Syn.Kind_Argument_List
, xin_Syn.Kind_Keyword
, xin_Syn.Kind_ProcStmt
, xin_Syn.Kind_BeginStmt
, xin_Syn.Kind_Stmt
, xin_Syn.Kind_Substmt
, xin_Syn.Kind_Label
, xin_Syn.Kind_InvokeBegin
, xin_Syn.Kind_Assignment
, xin_Syn.Kind_Assignment_Byname
, xin_Syn.Kind_DoFragment
, xin_Syn.Kind_KeyedList
, xin_Syn.Kind_IterationFactor
, xin_Syn.Kind_IfClause
, xin_Syn.Kind_ElseClause
, xin_Syn.Kind_DoStmt
, xin_Syn.Kind_SelectStmt
, xin_Syn.Kind_WhenStmt
, xin_Syn.Kind_OtherwiseStmt
, xin_Syn.Kind_Procedure
, xin_Syn.Kind_Package
, xin_Syn.Kind_BeginBlock
, xin_Syn.Kind_Picture
, xin_Syn.Kind_RawRfrnc
, xin_Syn.Kind_GenericDesc
) prec(8) unsigned;

Figure 129. Declare for the syntax record kind

Consider the following simple program:

```
a: proc(x);
dcl x char(8);
x = substr(datetime(),1,8);
end;
```

The node indices assigned to the blocks of the preceding program would be as follows:
The procedure record contains the identifier (in the block_sym field) for the symbol record for ENTRY A. This symbol record contains, in turn, the node identifier (in the first_stmt_id field) for the first statement in that procedure.

Note that for the statement records
- the sibling node identifier points to the next statement record, if any
- the child node identifier points to the first element of that statement record

The records for the PROCEDURE statement consists of 4 records:
- a label record
- a keyword record (for the PROCEDURE keyword)
- an expression record (for the parameter X) with expression kind of unsub_rfrnc and a sym_id for the symbol X
- a lexeme record (for the semicolon)

The records for the assignment statement consists of 2 records:
- an assignment record which has 2 children:
  - an expression record (for the target X) with expression kind of unsub_rfrnc and a sym_id for the symbol X
  - an expression record (for the source) with expression kind of builtin_rfrnc and a sym_id for the symbol SUBSTR, and this record has itself 3 children:
    - an expression record (for the first argument) with expression kind of builtin_rfrnc and a sym_id for the symbol DATETIME
    - an expression record (for the second argument) with expression kind of number and a literal_id for the value 1
    - an expression record (for the third argument) with expression kind of number and a literal_id for the value 8
  - a lexeme record (for the semicolon)

The records for the END statement consists of 2 records:

---

<table>
<thead>
<tr>
<th>node_kind</th>
<th>index</th>
<th>sibling</th>
<th>parent</th>
<th>child</th>
</tr>
</thead>
<tbody>
<tr>
<td>package</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>2</td>
</tr>
<tr>
<td>procedure</td>
<td>2</td>
<td>0</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>expression</td>
<td>3</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>stmt</td>
<td>4</td>
<td>5</td>
<td>2</td>
<td>6</td>
</tr>
<tr>
<td>stmt</td>
<td>5</td>
<td>10</td>
<td>2</td>
<td>11</td>
</tr>
<tr>
<td>label</td>
<td>6</td>
<td>7</td>
<td>4</td>
<td>0</td>
</tr>
<tr>
<td>keyword</td>
<td>7</td>
<td>8</td>
<td>4</td>
<td>0</td>
</tr>
<tr>
<td>expression</td>
<td>8</td>
<td>9</td>
<td>4</td>
<td>0</td>
</tr>
<tr>
<td>lexeme</td>
<td>9</td>
<td>0</td>
<td>4</td>
<td>0</td>
</tr>
<tr>
<td>stmt</td>
<td>10</td>
<td>0</td>
<td>2</td>
<td>18</td>
</tr>
<tr>
<td>assignment</td>
<td>11</td>
<td>12</td>
<td>5</td>
<td>13</td>
</tr>
<tr>
<td>lexeme</td>
<td>12</td>
<td>0</td>
<td>5</td>
<td>0</td>
</tr>
<tr>
<td>expression</td>
<td>13</td>
<td>14</td>
<td>11</td>
<td>0</td>
</tr>
<tr>
<td>expression</td>
<td>14</td>
<td>0</td>
<td>11</td>
<td>15</td>
</tr>
<tr>
<td>expression</td>
<td>15</td>
<td>16</td>
<td>14</td>
<td>0</td>
</tr>
<tr>
<td>expression</td>
<td>16</td>
<td>17</td>
<td>14</td>
<td>0</td>
</tr>
<tr>
<td>expression</td>
<td>17</td>
<td>0</td>
<td>14</td>
<td>0</td>
</tr>
<tr>
<td>keyword</td>
<td>18</td>
<td>19</td>
<td>10</td>
<td>0</td>
</tr>
<tr>
<td>lexeme</td>
<td>19</td>
<td>0</td>
<td>10</td>
<td>0</td>
</tr>
</tbody>
</table>

Figure 130. Node indices assigned to the syntax records in a program
• a keyword record (for the END keyword)
• a lexeme record (for the semicolon)

The ordinal `xin_Exp_Kind` identifies the type of an expression for a syntax record that describes an expression. Some of these records will have non-zero child nodes, for example:
• an infix op, such as a minus for a subtraction, will have a child node that describes its left-hand operand (and the sibling node of that operand will describe the right-hand operator)
• a prefix op, such as a minus for a negation, will have a child node that describes its operand

```
Define
ordinal
xin_Exp_Kind
  ('xin_Exp_Kind_Unset
   ,xin_Exp_Kind_Bit_String
   ,xin_Exp_Kind_Char_String
   ,xin_Exp_Kind_Graphic_String
   ,xin_Exp_Kind_Number
   ,xin_Exp_Kind_Infix_Op
   ,xin_Exp_Kind_Prefix_Op
   ,xin_Exp_Kind_Builtin_Rfrnc
   ,xin_Exp_Kind_Entry_Rfrnc
   ,xin_Exp_Kind_Qualified_Rfrnc
   ,xin_Exp_Kind_Unsub_Rfrnc
   ,xin_Exp_Kind_Subscripted_Rfrnc
   ,xin_Exp_Kind_Type_Func
   ,xin_Exp_Kind_Widechar_String
  ) prec(8) unsigned;
```

Figure 131. Declare for the expression kind

The ordinal `xin_Number_Kind` identifies the type of a number for a syntax record that describes a number.

```
Define
ordinal
xin_Number_Kind
  ('xin_Number_Kind_Unset
   ,xin_Number_Kind_Real_Fixed_Bin
   ,xin_Number_Kind_Real_Fixed_Dec
   ,xin_Number_Kind_Real_Float_Bin
   ,xin_Number_Kind_Real_Float_Dec
   ,xin_Number_Kind_Cplx_Fixed_Bin
   ,xin_Number_Kind_Cplx_Fixed_Dec
   ,xin_Number_Kind_Cplx_Float_Bin
   ,xin_Number_Kind_Cplx_Float_Dec
  ) prec(8) unsigned;
```

Figure 132. Declare for the number kind

The ordinal `xin_Lex_Kind` identifies the type of a lexeme for a syntax record that describes a lexical unit. In these ordinal names,
• "vrule" means "vertical rule" which is used, for instance, as the "or" symbol
"dbl" means "double", so that dbl_Vrule is a doubled vertical rule which is used, for instance, as the "concatenate" symbol

```c
Define ordinal xin_Lex_Kind
    ( xin_Lex_Undefined,
      xin_Lex_Period,
      xin_Lex_Colon,
      xin_Lex_Semicolon,
      xin_Lex_Lparen,
      xin_Lex_Rparen,
      xin_Lex_Comma,
      xin_Lex_Equals,
      xin_Lex_Gt,
      xin_Lex_Ge,
      xin_Lex_Lt,
      xin_Lex_Le,
      xin_Lex_Ne,
      xin_Lex_Lctr,
      xin_Lex_Star,
      xin_Lex_Dbl_Colon,
      xin_Lex_Not,
      xin_Lex_Vrule,
      xin_Lex_Dbl_Vrule,
      xin_Lex_And,
      xin_Lex_Dbl_Slash,
      xin_Lex_Equals_Gt,
      xin_Lex_Equals_Le,
      xin_Lex_Equals_Lctr,
      xin_Lex_Equals_Minus,
      xin_Lex_Equals_Plus,
      xin_Lex_Equals_Star,
      xin_Lex_Equals_Slash,
      xin_Lex_Equals_Slash,
      xin_Lex_Equals_Vrule,
      xin_Lex_Equals_And,
      xin_Lex_Equals_Dbl_Star,
      xin_Lex_Equals_Dbl_Vrule,
      xin_Lex_Equals_Dbl_Slash
    ) unsigned prec(16);
```

Figure 133. Declare for the lexeme kind

The ordinal xin_Voc_Kind identifies the keyword for a syntax record that describes an item from the compiler's "vocabulary".
Define
ordinal
xin_Voc_Kind
  { xin_Voc_Undefined
  , xin_Voc_a
  , xin_Voc_abnormal
  , xin_Voc_act
  , xin_Voc_activate
  , xin_Voc_adata
  , xin_Voc_addbuff
  , xin_Voc_aggregate
  , xin_Voc_aix
  , xin_Voc_alias
  , xin_Voc_amber
  , xin_Voc_aggregated
  , xin_Voc_alien
  , xin_Voc_aligned
  , xin_Voc_all
  , xin_Voc_alloc
  , xin_Voc_allocate
  , xin_Voc_annonce
  , xin_Voc_ans
  , xin_Voc_any
  , xin_Voc_anycond
  , xin_Voc_anycondition
  , xin_Voc_area
  , xin_Voc_as
  , xin_Voc_ascii
  , xin_Voc_asgn
  , xin_Voc_asm
  , xin_Voc_asmtdli
  , xin_Voc_assembler
  , xin_Voc_assignable
  , xin_Voc_attach
  , xin_Voc_attention
  , xin_Voc_attn
  , xin_Voc_attribute
  , xin_Voc_attributes
  , xin_Voc_auto
  , xin_Voc_automatic
  , xin_Voc_b
  , xin_Voc_backwards
  , xin_Voc_based
  , xin_Voc_begin
  , xin_Voc_beta
  , xin_Voc_bigendian
  , xin_Voc_bin
  , xin_Voc_binary
  , xin_Voc_bind
  , xin_Voc_bit
  , xin_Voc_bkwd
  , xin_Voc_blksize
  , xin_Voc_block
  , xin_Voc_buf

Figure 134. Declare for the voc kind
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Declare for the voc kind (continued)
Notices

This information was developed for products and services offered in the U.S.A. IBM may not offer the products, services, or features discussed in this document in other countries. Consult your local IBM representative for information on the products and services currently available in your area. Any reference to an IBM product, program, or service is not intended to state or imply that only that IBM product, program, or service may be used. Any functionally equivalent product, program, or service that does not infringe any IBM intellectual property right may be used instead. However, it is the user's responsibility to evaluate and verify the operation of any non-IBM product, program, or service.

IBM may have patents or pending patent applications covering subject matter described in this document. The furnishing of this document does not give you any license to these patents. You can send license inquiries, in writing, to:

IBM Corporation
J74/G4
555 Bailey Avenue
San Jose, CA 95141-1099
U.S.A.

For license inquiries regarding double-byte (DBCS) information, contact the IBM Intellectual Property Department in your country or send inquiries, in writing, to:

Intellectual Property Licensing
Legal and Intellectual Property Law
IBM Japan, Ltd.
3-2-12, Roppongi, Minato-ku, Tokyo 106-8711

The following paragraph does not apply to the United Kingdom or any other country where such provisions are inconsistent with local law:

INTERNATIONAL BUSINESS MACHINES CORPORATION PROVIDES
THIS PUBLICATION "AS IS" WITHOUT WARRANTY OF ANY KIND,
EITHER EXPRESS OR IMPLIED, INCLUDING, BUT NOT LIMITED
TO, THE IMPLIED WARRANTIES OF NON-INFRINGEMENT,
MERCHANTABILITY OR FITNESS FOR A PARTICULAR PURPOSE.

Some states do not allow disclaimer of express or implied warranties in certain transactions, therefore, this statement may not apply to you.

This information could include technical inaccuracies or typographical errors.

Changes are periodically made to the information herein; these changes will be incorporated in new editions of the publication. IBM may make improvements and/or changes in the product(s) and/or the program(s) described in this publication at any time without notice.

Any references in this publication to non-IBM Web sites are provided for convenience only and do not in any manner serve as an endorsement of those Web sites. The materials at those Web sites are not part of the materials for this IBM product and use of those Web sites is at your own risk.
Trademarks

IBM, the IBM logo, and ibm.com are trademarks or registered trademarks of International Business Machines Corp., registered in many jurisdictions worldwide. Other product and service names might be trademarks of IBM or other companies. A current list of IBM trademarks is available on the Web at "Copyright and trademark information" at www.ibm.com/legal/copytrade.shtml.

Intel is a registered trademark of Intel Corporation in the United States and other countries.

Java and all Java-based trademarks and logos are trademarks or registered trademarks of Oracle and/or its affiliates.

Microsoft, Windows, and Windows NT are trademarks of Microsoft Corporation in the United States and other countries.

Pentium is a registered trademark of Intel Corporation in the United States and other countries.

Unicode is a trademark of the Unicode Consortium.

UNIX is a registered trademark of The Open Group in the United States and other countries.

Other company, product or service names may be the trademarks or service marks of others.

If you are viewing this information in softcopy, the photographs and color illustrations may not appear.
Bibliography

PL/I publications

**Enterprise PL/I for z/OS**
- Programming Guide, GI11-9145
- Language Reference, SC14-7285
- Messages and Codes, GC14-7286
- Compiler and Run-Time Migration Guide, GC14-7284

**PL/I for MVS & VM**
- Installation and Customization under MVS, SC26-3119
- Language Reference, SC26-3114
- Compile-Time Messages and Codes, SC26-3229
- Diagnosis Guide, SC26-3149
- Migration Guide, SC26-3118
- Programming Guide, SC26-3113
- Reference Summary, SX26-3821

**PL/I for AIX**
- Programming Guide, SC14-7319
- Language Reference, SC14-7320
- Messages and Codes, GC14-7321
- Installation Guide, GC14-7322

Related publications

**DB2 for z/OS**
- Administration Guide, SC19-2968
- Application Programming and SQL Guide, SC19-2969
- Command Reference, SC19-2972
- Messages, GC19-2979
- Codes, GC19-2971
- SQL Reference, SC19-2983
See also the Information Center: publib.boulder.ibm.com/infocenter/dzichelp/v2r2/topic/com.ibm.db2z10.doc/src/alltoc/db2z_10_prodhome.htm

**LOBs with DB2 for z/OS: Stronger and Faster**, SG24-7270

**DFSORT™**
- Application Programming Guide, SC33-4035
- Installation and Customization, SC33-4034

**IMS/ESA®**
- Application Programming: Database Manager, SC26-8015
- Application Programming: Database Manager Summary, SC26-8037
- Application Programming: Design Guide, SC26-8016
TXSeries for Multiplatforms

Encina Administration Guide Volume 2: Server Administration, SC09-4474
Encina SFS Programming Guide, SC09-4483
See also the Information Center: publib.boulder.ibm.com/infocenter/txformp/v7r1/index.jsp

z/Architecture

Principles of Operation, SA22-7832

z/OS Language Environment

Concepts Guide, SA22-7567
Debugging Guide, GA22-7560
Run-Time Messages, SA22-7566
Customization, SA22-7564
Programming Guide, SA22-7561
Programming Reference, SA22-7562
Run-Time Application Migration Guide, GA22-7565
Writing Interlanguage Communication Applications, SA22-7563

z/OS MVS

JCL Reference, SA22-7597
JCL User’s Guide, SA22-7598
System Commands, SA22-7627

z/OS TSO/E

Command Reference, SA22-7782
User’s Guide, SA22-7794

z/OS UNIX System Services

z/OS UNIX System Services Command Reference, SA22-7802
z/OS UNIX System Services Programming: Assembler Callable Services Reference, SA22-7803
z/OS UNIX System Services User’s Guide, SA22-7801

Unicode® and character representation

z/OS Support for Unicode: Using Conversion Services, SC33-7050
Glossary

This glossary defines terms for all platforms and releases of PL/I. It might contain terms that this manual does not use. If you do not find the terms for which you are looking, see the index in this manual or IBM Dictionary of Computing, SC20-1699.

A

access  To reference or retrieve data.

action specification  In an ON statement, the ON-unit or the single keyword SYSTEM, either of which specifies the action to be taken whenever the appropriate condition is raised.

activate (a block)  To initiate the execution of a block. A procedure block is activated when it is invoked. A begin-block is activated when it is encountered in the normal flow of control, including a branch. A package cannot be activated.

activate (a preprocessor variable or preprocessor entry point)  To make a macro facility identifier eligible for replacement in subsequent source code. The %ACTIVATE statement activates preprocessor variables or preprocessor entry points.

active  The state of a block after activation and before termination. The state in which a preprocessor variable or preprocessor entry name is said to be when its value can replace the corresponding identifier in source program text. The state in which an event variable is said to be during the time it is associated with an asynchronous operation. The state in which a task variable is said to be when its associated task is attached. The state in which a task is said to be before it has been terminated.

actual origin (AO)  The location of the first item in the array or structure.

additive attribute  A file description attribute for which there are no defaults, and which, if required, must be stated explicitly or implied by another explicitly stated attribute. Contrast with alternative attribute.

adjustable extent  The bound (of an array), the length (of a string), or the size (of an area) that might be different for different generations of the associated variable. Adjustable extents are specified as expressions or asterisks (or by REFER options for based variables), which are evaluated separately for each generation. They cannot be used for static variables.

aggregate  See data aggregate.

aggregate expression  An array, structure, or union expression.

aggregate type  For any item of data, the specification whether it is structure, union, or array.

allocated variable  A variable with which main storage is associated and not freed.

allocation  The reservation of main storage for a variable. A generation of an allocated variable. The association of a PL/I file with a system data set, device, or file.

alignment  The storing of data items in relation to certain machine-dependent boundaries (for example, a fullword or halfword boundary).

alphabetic character  Any of the characters A through Z of the English alphabet and the alphabetic extenders #, $, and @ (which can have a different graphic representation in different countries).

alphabetic character  An alphabetic character or a digit.

alternative attribute  A file description attribute that is chosen from a group of attributes. If none is specified, a default is assumed. Contrast with additive attribute.
ambiguous reference
A reference that is not sufficiently qualified to identify one and only one name known at the point of reference.

area
A portion of storage within which based variables can be allocated.

argument
An expression in an argument list as part of an invocation of a subroutine or function.

argument list
A parenthesized list of zero or more arguments, separated by commas, following an entry name constant, an entry name variable, a generic name, or a built-in function name. The list becomes the parameter list of the entry point.

arithmetic comparison
A comparison of numeric values. See also bit comparison, character comparison.

arithmetic constant
A fixed-point constant or a floating-point constant. Although most arithmetic constants can be signed, the sign is not part of the constant.

arithmetic conversion
The transformation of a value from one arithmetic representation to another.

arithmetic data
Data that has the characteristics of base, scale, mode, and precision. Coded arithmetic data and pictured numeric character data are included.

arithmetic operators
Either of the prefix operators + and -, or any of the following infix operators: + - * / **

array
A named, ordered collection of one or more data elements with identical attributes, grouped into one or more dimensions.

array expression
An expression whose evaluation yields an array of values.

array of structures
An ordered collection of identical structures specified by giving the dimension attribute to a structure name.

array variable
A variable that represents an aggregate of data items that must have identical attributes. Contrast with structure variable.

ASCII

assignment
The process of giving a value to a variable.

asynchronous operation
The overlap of an input/output operation with the execution of statements. The concurrent execution of procedures using multiple flows of control for different tasks.

attachment of a task
The invocation of a procedure and the establishment of a separate flow of control to execute the invoked procedure (and procedures it invokes) asynchronously, with execution of the invoking procedure.

attention
An occurrence, external to a task, that could cause a task to be interrupted.

attribute
A descriptive property associated with a name to describe a characteristic represented. A descriptive property used to describe a characteristic of the result of evaluation of an expression.

automatic storage allocation
The allocation of storage for automatic variables.

automatic variable
A variable whose storage is allocated automatically at the activation of a block and released automatically at the termination of that block.

B

base
The number system in which an arithmetic value is represented.

base element
A member of a structure or a union that is itself not another structure or union.

base item
The automatic, controlled, or static variable or the parameter upon which a defined variable is defined.
based reference
A reference that has the based storage class.

based storage allocation
The allocation of storage for based variables.

based variable
A variable whose storage address is provided by a locator. Multiple generations of the same variable are accessible. It does not identify a fixed location in storage.

begin-block
A collection of statements delimited by BEGIN and END statements, forming a name scope. A begin-block is activated either by the raising of a condition (if the begin-block is the action specification for an ON-unit) or through the normal flow of control, including any branch resulting from a GOTO statement.

binary
A number system whose only numerals are 0 and 1.

binary digit
See bit.

binary fixed-point value
An integer consisting of binary digits and having an optional binary point and optional sign. Contrast with decimal fixed-point value.

binary floating-point value
An approximation of a real number in the form of a significand, which can be considered as a binary fraction, and an exponent, which can be considered as an integer exponent to the base of 2. Contrast with decimal floating-point value.

bit
A 0 or a 1. The smallest amount of space of computer storage.

bit comparison
A left-to-right, bit-by-bit comparison of binary digits. See also arithmetic comparison, character comparison.

bit string constant
A series of binary digits enclosed in and followed immediately by the suffix B. Contrast with character constant. A series of hexadecimal digits enclosed in single quotes and followed by the suffix B4.

bit string
A string composed of zero or more bits.

bit string operators
The logical operators not and exclusive-or (¬), and (¬), and or (¬).

bit value
A value that represents a bit type.

block
A sequence of statements, processed as a unit, that specifies the scope of names and the allocation of storage for names declared within it. A block can be a package, procedure, or a begin-block.

bounds
The upper and lower limits of an array dimension.

break character
The underscore symbol (_) . It can be used to improve the readability of identifiers. For instance, a variable could be called OLD_INVENTORY_TOTAL instead of OLDINVENTORYTOTAL.

built-in function
A predefined function supplied by the language, such as SQRT (square root).

built-in function reference
A built-in function name, which has an optional argument list.

built-in name
The entry name of a built-in subroutine.

built-in subroutine
Subroutine that has an entry name that is defined at compile-time and is invoked by a CALL statement.

buffer
Intermediate storage, used in input/output operations, into which a record is read during input and from which a record is written during output.

C
call
To invoke a subroutine by using the CALL statement or CALL option.

character comparison
A left-to-right, character-by-character comparison according to the collating sequence. See also arithmetic comparison, bit comparison.
character string constant
A sequence of characters enclosed in single quotes; for example, 'Shakespeare's 'Hamlet'.

character set
A defined collection of characters. See also ASCII and EBCDIC.

character string picture data
Picture data that has only a character value. This type of picture data must have at least one A or X picture specification character. Contrast with numeric picture data.

closing (of a file)
The dissociation of a file from a data set or device.

coded arithmetic data
Data items that represent numeric values and are characterized by their base (decimal or binary), scale (fixed-point or floating-point), and precision (the number of digits each can have). This data is stored in a form that is acceptable, without conversion, for arithmetic calculations.

combined nesting depth
The deepest level of nesting, determined by counting the levels of PROCEDURE/BEGIN/ON, DO, SELECT, and IF..THEN..ELSE nestings in the program.

comment
A string of zero or more characters used for documentation that are delimited by /* and */.

commercial character
• CR (credit) picture specification character
• DB (debit) picture specification character

comparison operator
An operator that can be used in an arithmetic, string locator, or logical relation to indicate the comparison to be done between the terms in the relation. The comparison operators are:
  = (equal to)
  > (greater than)
  < (less than)
  >= (greater than or equal to)
  <= (less than or equal to)
  /= (not equal to)
  > (not greater than)
  < (not less than)

compile time
In general, the time during which a source program is translated into an object module. In PL/I, it is the time during which a source program can be altered, if desired, and then translated into an object program.

compiler options
Keywords that are specified to control certain aspects of a compilation, such as: the nature of the object module generated, the types of printed output produced, and so forth.

complex data
Arithmetic data, each item of which consists of a real part and an imaginary part.

composite operator
An operator that consists of more than one special character, such as <=, *+, and */.

compound statement
A statement that contains other statements. In PL/I, IF, ON, OTHERWISE, and WHEN are the only compound statements. See statement body.

concatenation
The operation that joins two strings in the order specified, forming one string whose length is equal to the sum of the lengths of the two original strings. It is specified by the operator ||.

condition
An exceptional situation, either an error (such as an overflow), or an expected situation (such as the end of an input file). When a condition is raised (detected), the action established for it is processed. See also established action and implicit action.

condition name
Name of a PL/I-defined or programmer-defined condition.

condition prefix
A parenthesized list of one or more condition names prefixed to a statement.
It specifies whether the named conditions are to be enabled or disabled.

**connected aggregate**
An array or structure whose elements occupy contiguous storage without any intervening data items. Contrast with *nonconnected aggregate*.

**connected reference**
A reference to connected storage. It must be apparent, prior to execution of the program, that the storage is connected.

**connected storage**
Main storage of an uninterrupted linear sequence of items that can be referred to by a single name.

**constant**
An arithmetic or string data item that does not have a name and whose value cannot change. An identifier declared with the VALUE attribute. An identifier declared with the FILE or the ENTRY attribute but without the VARIABLE attribute.

**constant reference**
A value reference which has a constant as its object

**contained block, declaration, or source text**
All blocks, procedures, statements, declarations, or source text inside a begin, procedure, or a package block. The entire package, procedure, and the BEGIN statement and its corresponding END statements are not contained in the block.

**containing block**
The package, procedure, or begin-block that contains the declaration, statement, procedure, or other source text in question.

**contextual declaration**
The appearance of an identifier that has not been explicitly declared in a DECLARE statement, but whose context of use allows the association of specific attributes with the identifier.

**control character**
A character in a character set whose occurrence in a particular context specifies a control function. One example is the end-of-file (EOF) marker.

**control format item**
A specification used in edit-directed transmission to specify positioning of a data item within the stream or printed page.

**control variable**
A variable that is used to control the iterative execution of a DO statement.

**controlled parameter**
A parameter for which the CONTROLLED attribute is specified in a DECLARE statement. It can be associated only with arguments that have the CONTROLLED attribute.

**controlled storage allocation**
The allocation of storage for controlled variables.

**controlled variable**
A variable whose allocation and release are controlled by the ALLOCATE and FREE statements, with access to the current generation only.

**control sections**
Grouped machine instructions in an object module.

**conversion**
The transformation of a value from one representation to another to conform to a given set of attributes. For example, converting a character string to an arithmetic value such as FIXED BINARY (15,0).

**cross section of an array**
The elements represented by the extent of at least one dimension of an array. An asterisk in the place of a subscript in an array reference indicates the entire extent of that dimension.

**current generation**
The generation of an automatic or controlled variable that is currently available by referring to the name of the variable.

**D**

**data**
Representation of information or of value in a form suitable for processing.

**data aggregate**
A data item that is a collection of other data items.
data attribute
A keyword that specifies the type of data that the data item represents, such as FIXED BINARY.

data-directed transmission
The type of stream-oriented transmission in which data is transmitted. It resembles an assignment statement and is of the form name = constant.

data item
A single named unit of data.

data list
In stream-oriented transmission, a parenthesized list of the data items used in GET and PUT statements. Contrast with format list.

data set
A collection of data external to the program that can be accessed by reference to a single file name. A device that can be referenced.

data specification
The portion of a stream-oriented transmission statement that specifies the mode of transmission (DATA, LIST, or EDIT) and includes the data list(s) and, for edit-directed mode, the format list(s).

data stream
Data being transferred from or to a data set by stream-oriented transmission, as a continuous stream of data elements in character form.

data transmission
The transfer of data from a data set to the program or vice versa.

data type
A set of data attributes.

DBCS
In the character set, each character is represented by two consecutive bytes.

deactivated
The state in which an identifier is said to be when its value cannot replace a preprocessor identifier in source program text. Contrast with active.

debugging
Process of removing bugs from a program.

decimal
The number system whose numerals are 0 through 9.
descriptor
A control block that holds information about a variable, such as area size, array bounds, or string length.
digit
One of the characters 0 through 9.
dimension attribute
An attribute that specifies the number of dimensions of an array and indicates the bounds of each dimension.
disabled
The state of a condition in which no interrupt occurs and no established action will take place.
do-group
A sequence of statements delimited by a DO statement and ended by its corresponding END statement, used for control purposes. Contrast with block.
do-loop
See iterative do-group.
dummy argument
Temporary storage that is created automatically to hold the value of an argument that cannot be passed by reference.
dump
Printout of all or part of the storage used by a program as well as other program information, such as a trace of an error's origin.

E
EBCDIC
(Extended Binary-Coded Decimal Interchange Code). A coded character set consisting of 8-bit coded characters.
edit-directed transmission
The type of stream-oriented transmission in which data appears as a continuous stream of characters and for which a format list is required to specify the editing desired for the associated data list.
element
A single item of data as opposed to a collection of data items such as an array; a scalar item.
element expression
An expression whose evaluation yields an element value.
the environment of the activation that is associated with the entry constant.

**environment (of an activation)**
Information associated with and used in the invoked block regarding data declared in containing blocks.

**environment (of a label constant)**
Identity of the particular activation of a block to which a reference to a statement-label constant applies. This information is determined at the time a statement-label constant is passed as an argument or is assigned to a statement-label variable, and it is passed or assigned along with the constant.

**established action**
The action taken when a condition is raised. See also implicit action and ON-statement action.

**epilogue**
Those processes that occur automatically at the termination of a block or task.

**evaluation**
The reduction of an expression to a single value, an array of values, or a structured set of values.

**event**
An activity in a program whose status and completion can be determined from an associated event variable.

**event variable**
A variable with the EVENT attribute that can be associated with an event. Its value indicates whether the action has been completed and the status of the completion.

**explicit declaration**
The appearance of an identifier (a name) in a DECLARE statement, as a label prefix, or in a parameter list. Contrast with implicit declaration.

**exponent characters**
The following picture specification characters:

1. K and E, which are used in floating-point picture specifications to indicate the beginning of the exponent field.
2. F, the scaling factor character, specified with an integer constant that indicates the number of decimal positions the decimal point is to be moved from its assumed position to the right (if the constant is positive) or to the left (if the constant is negative).

**expression**
A notation, within a program, that represents a value, an array of values, or a structured set of values. A constant or a reference appearing alone, or a combination of constants and/or references with operators.

**extended alphabet**
The uppercase and lowercase alphabetic characters A through Z, $, @ and #, or those specified in the NAMES compiler option.

**extent**
The range indicated by the bounds of an array dimension, by the length of a string, or by the size of an area. The size of the target area if this area were to be assigned to a target area.

**external name**
A name (with the EXTERNAL attribute) whose scope is not necessarily confined only to one block and its contained blocks.

**external procedure**
A procedure that is not contained in any other procedure. A level-2 procedure contained in a package that is also exported.

**external symbol**
Name that can be referred to in a control section other than the one in which it is defined.

**External Symbol Dictionary (ESD)**
Table containing all the external symbols that appear in the object module.

**extralingual character**
Characters (such as $, @, and #) that are not classified as alphanumeric or special. This group includes characters that are determined with the NAMES compiler option.

**F**

**factoring**
The application of one or more attributes to a parenthesized list of names in a DECLARE statement, eliminating the repetition of identical attributes for multiple names.
field (in the data stream)
That portion of the data stream whose width, in number of characters, is defined by a single data or spacing format item.

field (of a picture specification)
Any character-string picture specification or that portion (or all) of a numeric character picture specification that describes a fixed-point number.

file
A named representation, within a program, of a data set or data sets. A file is associated with the data set(s) for each opening.

file constant
A name declared with the FILE attribute but not the VARIABLE attribute.

file description attributes
Keywords that describe the individual characteristics of each file constant. See also alternative attribute and additive attribute.

file expression
An expression whose evaluation yields a value of the type file.

file name
A name declared for a file.

file variable
A variable to which file constants can be assigned. It has the attributes FILE and VARIABLE and cannot have any of the file description attributes.

fixed-point constant
See arithmetic constant.

fix-up
A solution, performed by the compiler after detecting an error during compilation, that allows the compiled program to run.

floating-point constant
See arithmetic constant.

flow of control
Sequence of execution.

format
A specification used in edit-directed data transmission to describe the representation of a data item in the stream (data format item) or the specific positioning of a data item within the stream (control format item).

format constant
The label prefix on a FORMAT statement.

format data
A variable with the FORMAT attribute.

format label
The label prefix on a FORMAT statement.

format list
In stream-oriented transmission, a list specifying the format of the data item on the external medium. Contrast with data list.

fully qualified name
A name that includes all the names in the hierarchical sequence above the member to which the name refers, as well as the name of the member itself.

function (procedure)
A procedure that has a RETURNS option in the PROCEDURE statement. A name declared with the RETURNS attribute. It is invoked by the appearance of one of its entry names in a function reference and it returns a scalar value to the point of reference. Contrast with subroutine.

function reference
An entry constant or an entry variable, either of which must represent a function, followed by a possibly empty argument list. Contrast with subroutine call.

G

generation (of a variable)
The allocation of a static variable, a particular allocation of a controlled or automatic variable, or the storage indicated by a particular locator qualification of a based variable or by a defined variable or parameter.

generic descriptor
A descriptor used in a GENERIC attribute.

generic key
A character string that identifies a class of keys. All keys that begin with the string are members of that class. For example, the recorded keys 'ABCD', 'ABCE', and 'ABDF', are all members of the classes identified by the generic keys 'A' and 'AB', and the first two are also members of the class 'ABC'; and the three recorded
keys can be considered to be unique members of the classes 'ABCD', 'ABCE', 'ABDF', respectively.

generic name
The name of a family of entry names. A reference to the generic name is replaced by the entry name whose parameter descriptors match the attributes of the arguments in the argument list at the point of invocation.

group
A collection of statements contained within larger program units. A group is either a do-group or a select-group and it can be used wherever a single statement can appear, except as an on-unit.

H
hex
See hexadecimal digit.

hexadecimal
Pertaining to a numbering system with a base of sixteen; valid numbers use the digits 0 through 9 and the characters A through F, where A represents 10 and F represents 15.

hexadecimal digit
One of the digits 0 through 9 and A through F. A through F represent the decimal values 10 through 15, respectively.

I
identifier
A string of characters, not contained in a comment or constant, and preceded and followed by a delimiter. The first character of the identifier must be one of the 26 alphabetic characters and extralingual characters, if any. The other characters, if any, can additionally include extended alphabetic, digit, or the break character.

IEEE
Institute of Electrical and Electronics Engineers.

implicit
The action taken in the absence of an explicit specification.

implicit action
The action taken when an enabled condition is raised and no ON-unit is currently established for the condition. Contrast with ON-statement action.

implicit declaration
A name not explicitly declared in a DECLARE statement or contextually declared.

implicit opening
The opening of a file as the result of an input or output statement other than the OPEN statement.

infix operator
An operator that appears between two operands.

inherited dimensions
For a structure, union, or element, those dimensions that are derived from the containing structures. If the name is an element that is not an array, the dimensions consist entirely of its inherited dimensions. If the name is an element that is an array, its dimensions consist of its inherited dimensions plus its explicitly declared dimensions. A structure with one or more inherited dimensions is called a nonconnected aggregate. Contrast with connected aggregate.

input/output
The transfer of data between auxiliary medium and main storage.

insertion point character
A picture specification character that is, on assignment of the associated data to a character string, inserted in the indicated position. When used in a P-format item for input, the insertion character is used for checking purposes.

integer
An optionally signed sequence of digits or a sequence of bits without a decimal or binary point. An optionally signed whole number, commonly described as FIXED BINARY (p,0) or FIXED DECIMAL (p,0).

integral boundary
A byte multiple address of any 8-bit unit on which data can be aligned. It usually is a halfword, fullword, or doubleword (2-, 4-, or 8-byte multiple respectively) boundary.

interleaved array
An array that refers to nonconnected storage.
**interleaved subscripts**
Subscripts that exist in levels other than the lowest level of a subscripted qualified reference.

**internal block**
A block that is contained in another block.

**internal name**
A name that is known only within the block in which it is declared, and possibly within any contained blocks.

**internal procedure**
A procedure that is contained in another block. Contrast with *external procedure*.

**interrupt**
The redirection of the program’s flow of control as the result of raising a condition or attention.

**invocation**
The activation of a procedure.

**invoke**
To activate a procedure.

**invoked procedure**
A procedure that has been activated.

**invoking block**
A block that activates a procedure.

**iteration factor**
In an INITIAL attribute specification, an expression that specifies the number of consecutive elements of an array that are to be initialized with the given value. In a format list, an expression that specifies the number of times a given format item or list of format items is to be used in succession.

**iterative do-group**
A do-group whose DO statement specifies a control variable and/or a WHILE or UNTIL option.

**K**

**key**
Data that identifies a record within a direct access data set. See *source key* and *recorded key*.

**keyword**
An identifier that has a specific meaning in PL/I when used in a defined context.

**keyword statement**
A simple statement that begins with a keyword, indicating the function of the statement.

**known (applied to a name)**
Recognized with its declared meaning. A name is known throughout its scope.

**L**

**label**
A name prefixed to a statement. A name on a PROCEDURE statement is called an entry constant; a name on a FORMAT statement is called a format constant; a name on other kinds of statements is called a label constant. A data item that has the LABEL attribute.

**label constant**
A name written as the label prefix of a statement (other than PROCEDURE, ENTRY, FORMAT, or PACKAGE) so that, during execution, program control can be transferred to that statement through a reference to its label prefix.

**label data**
A label constant or the value of a label variable.

**label prefix**
A label prefixed to a statement.

**label variable**
A variable declared with the LABEL attribute. Its value is a label constant in the program.

**leading zeroes**
Zeros that have no significance in an arithmetic value. All zeros to the left of the first nonzero in a number.

**level number**
A number that precedes a name in a DECLARE statement and specifies its relative position in the hierarchy of structure names.

**level-one variable**
A major structure or union name. Any unsubscripted variable not contained within a structure or union.

**lexically**
Relating to the left-to-right order of units.

**library**
An MVS partitioned data set or a CMS
MACLIB that can be used to store other data sets called members.

**list-directed**
The type of stream-oriented transmission in which data in the stream appears as constants separated by blanks or commas and for which formatting is provided automatically.

**locator**
A control block that holds the address of a variable or its descriptor.

**locator/descriptor**
A locator followed by a descriptor. The locator holds the address of the variable, not the address of the descriptor.

**locator qualification**
In a reference to a based variable, either a locator variable or function reference connected by an arrow to the left of a based variable to specify the generation of the based variable to which the reference refers. It might be an implicit reference.

**locator value**
A value that identifies or can be used to identify the storage address.

**locator variable**
A variable whose value identifies the location in main storage of a variable or a buffer. It has the POINTER or OFFSET attribute.

**locked record**
A record in an EXCLUSIVE DIRECT UPDATE file that has been made available to one task only and cannot be accessed by other tasks until the task using it relinquishes it.

**logical level (of a structure or union member)**
The depth indicated by a level number when all level numbers are in direct sequence (when the increment between successive level numbers is one).

**logical operators**
The bit-string operators not and exclusive-or (¬), and (∧), and or (∨).

**loop**
A sequence of instructions that is executed iteratively.

**lower bound**
The lower limit of an array dimension.

**main procedure**
An external procedure whose PROCEDURE statement has the OPTIONS (MAIN) attribute. This procedure is invoked automatically as the first step in the execution of a program.

**major structure**
A structure whose name is declared with level number 1.

**member**
A structure, union, or element name in a structure or union. Data sets in a library.

**minor structure**
A structure that is contained within another structure or union. The name of a minor structure is declared with a level number greater than one and greater than its parent structure or union.

**mode (of arithmetic data)**
An attribute of arithmetic data. It is either real or complex.

**multiple declaration**
Two or more declarations of the same identifier internal to the same block without different qualifications. Two or more external declarations of the same identifier.

**multiprocessing**
The use of a computing system with two or more processing units to execute two or more programs simultaneously.

**multiprogramming**
The use of a computing system to execute more than one program concurrently, using a single processing unit.

**multitasking**
A facility that allows a program to execute more than one PL/I procedure simultaneously.

**N**

**name**
Any identifier that the user gives to a variable or to a constant. An identifier appearing in a context where it is not a keyword. Sometimes called a user-defined name.

**nesting**
The occurrence of:
- A block within another block
- A group within another group
• An IF statement in a THEN clause or in an ELSE clause
• A function reference as an argument of a function reference
• A remote format item in the format list of a FORMAT statement
• A parameter descriptor list in another parameter descriptor list
• An attribute specification within a parenthesized name list for which one or more attributes are being factored

nonconnected storage
Storage occupied by nonconnected data items. For example, interleaved arrays and structures with inherited dimensions are in nonconnected storage.

null locator value
A special locator value that cannot identify any location in internal storage. It gives a positive indication that a locator variable does not currently identify any generation of data.

null statement
A statement that contains only the semicolon symbol (;). It indicates that no action is to be taken.

null string
A character, graphic, or bit string with a length of zero.

numeric-character data
See decimal picture data.

numeric picture data
Picture data that has an arithmetic value as well as a character value. This type of picture data cannot contain the characters 'A' or 'X.'

ON-condition
An occurrence, within a PL/I program, that could cause a program interrupt. It can be the detection of an unexpected error or of an occurrence that is expected, but at an unpredictable time.

ON-statement action
The action explicitly established for a condition that is executed when the condition is raised. When the ON-statement is encountered in the flow of control for the program, it executes, establishing the action for the condition. The action executes when the condition is raised if the ON-unit is still established or a RESIGNAL statement reestablishes it. Contrast with implicit action.

ON-unit
The specified action to be executed when the appropriate condition is raised.

opening (of a file)
The association of a file with a data set.

operand
The value of an identifier, constant, or an expression to which an operator is applied, possibly in conjunction with another operand.

operational expression
An expression that consists of one or more operators.

operator
A symbol specifying an operation to be performed.

option
A specification in a statement that can be used to influence the execution or interpretation of the statement.

P
package constant
The label prefix on a PACKAGE statement.

packed decimal
The internal representation of a fixed-point decimal data item.

padding
One or more characters, graphics, or bits concatenated to the right of a string to extend the string to a required length. One or more bytes or bits inserted in a structure or union so that the following element within the structure or union is aligned on the appropriate integral boundary.
parameter
A name in the parameter list following
the PROCEDURE statement, specifying an
argument that will be passed when the
procedure is invoked.

parameter descriptor
The set of attributes specified for a
parameter in an ENTRY attribute
specification.

parameter descriptor list
The list of all parameter descriptors in an
ENTRY attribute specification.

parameter list
A parenthesized list of one or more
parameters, separated by commas and
following either the keyword
PROCEDURE in a procedure statement or
the keyword ENTRY in an ENTRY
statement. The list corresponds to a list of
arguments passed at invocation.

partially qualified name
A qualified name that is incomplete. It
includes one or more, but not all, of the
names in the hierarchical sequence above
the structure or union member to which
the name refers, as well as the name of
the member itself.

picture data
Numeric data, character data, or a mix of
both types, represented in character form.

picture specification
A data item that is described using the
picture characters in a declaration with
the PICTURE attribute or in a P-format
item.

picture specification character
Any of the characters that can be used in
a picture specification.

PL/I character set
A set of characters that has been defined
to represent program elements in PL/I.

PL/I prompter
Command processor program for the PLI
command that checks the operands and
allocates the data sets required by the
compiler.

point of invocation
The point in the invoking block at which
the reference to the invoked procedure
appears.

pointer
A type of variable that identifies a
location in storage.

pointer value
A value that identifies the pointer type.

pointer variable
A locator variable with the POINTER
attribute that contains a pointer value.

precision
The number of digits or bits contained in
a fixed-point data item, or the minimum
number of significant digits (excluding
the exponent) maintained for a
floating-point data item.

prefix
A label or a parenthesized list of one or
more condition names included at the
beginning of a statement.

prefix operator
An operator that precedes an operand
and applies only to that operand. The
prefix operators are plus (+), minus (-),
and not (¬).

preprocessor
A program that examines the source
program before the compilation takes
place.

preprocessor statement
A special statement appearing in the
source program that specifies the actions
to be performed by the preprocessor. It is
executed as it is encountered by the
preprocessor.

primary entry point
The entry point identified by any of the
names in the label list of the
PROCEDURE statement.

priority
A value associated with a task, that
specifies the precedence of the task
relative to other tasks.

problem data
Coded arithmetic, bit, character, graphic,
and picture data.

problem-state program
A program that operates in the problem
state of the operating system. It does not
contain input/output instructions or other
privileged instructions.

procedure
A collection of statements, delimited by
PROCEDURE and END statements. A procedure is a program or a part of a program, delimits the scope of names, and is activated by a reference to the procedure or one of its entry names. See also external procedure and internal procedure.

**procedure reference**
An entry constant or variable. It can be followed by an argument list. It can appear in a CALL statement or the CALL option, or as a function reference.

**program**
A set of one or more external procedures or packages. One of the external procedures must have the OPTIONS(MAIN) specification in its procedure statement.

**program control data**
Area, locator, label, format, entry, and file data that is used to control the processing of a PL/I program.

**prologue**
The processes that occur automatically on block activation.

**pseudovariable**
Any of the built-in function names that can be used to specify a target variable. It is usually on the left-hand side of an assignment statement.

**qualified name**
A hierarchical sequence of names of structure or union members, connected by periods, used to identify a name within a structure. Any of the names can be subscripted.

**recorded key**
A character string identifying a record in a direct access data set where the character string itself is also recorded as part of the data.

**record-oriented data transmission**
The transmission of data in the form of separate records. Contrast with stream data transmission.

**recursive procedure**
A procedure that can be called from within itself or from within another active procedure.

**reentrant procedure**
A procedure that can be activated by multiple tasks, threads, or processes simultaneously without causing any interference between these tasks, threads, and processes.

**REFER expression**
The expression preceding the keyword REFER, which is used as the bound, length, or size when the based variable containing a REFER option is allocated, either by an ALLOCATE or LOCATE statement.

**REFER object**
The variable in a REFER option that holds or will hold the current bound, length, or size for the member. The REFER object must be a member of the same structure or union. It must not be locator-qualified or subscripted, and it must precede the member with the REFER option.

**reference**
The appearance of a name, except in a context that causes explicit declaration.

**relative virtual origin (RVO)**
The actual origin of an array minus the virtual origin of an array.

**remote format item**
The letter R followed by the label (enclosed in parentheses) of a FORMAT statement. The format statement is used by edit-directed data transmission statements to control the format of data being transmitted.

**repetition factor**
A parenthesized unsigned integer constant that specifies:
1. The number of times the string constant that follows is to be repeated.
2. The number of times the picture character that follows is to be repeated.

repetitive specification
An element of a data list that specifies controlled iteration to transmit one or more data items, generally used in conjunction with arrays.

restricted expression
An expression that can be evaluated by the compiler during compilation, resulting in a constant. Operands of such an expression are constants, named constants, and restricted expressions.

returned value
The value returned by a function procedure.

RETURNS descriptor
A descriptor used in a RETURNS attribute, and in the RETURNS option of the PROCEDURE and ENTRY statements.

S
scalar variable
A variable that is not a structure, union, or array.

scale
A system of mathematical notation whose representation of an arithmetic value is either fixed-point or floating-point.

scale factor
A specification of the number of fractional digits in a fixed-point number.

scaling factor
See scale factor.

scope (of a condition prefix)
The portion of a program throughout which a particular condition prefix applies.

scope (of a declaration or name)
The portion of a program throughout which a particular name is known.

secondary entry point
An entry point identified by any of the names in the label list of an entry statement.

select-group
A sequence of statements delimited by SELECT and END statements.

selection clause
A WHEN or OTHERWISE clause of a select-group.

self-defining data
An aggregate that contains data items whose bounds, lengths, and sizes are determined at program execution time and are stored in a member of the aggregate.

separator
See delimiter.

shift
Change of data in storage to the left or to the right of original position.

shift-in
Symbol used to signal the compiler at the end of a double-byte string.

shift-out
Symbol used to signal the compiler at the beginning of a double-byte string.

sign and currency symbol characters
The picture specification characters. S, +, -, and $ (or other national currency symbols enclosed in < and >).

simple parameter
A parameter for which no storage class attribute is specified. It can represent an argument of any storage class, but only the current generation of a controlled argument.

simple statement
A statement other than IF, ON, WHEN, and OTHERWISE.

source
Data item to be converted for problem data.

source key
A key referred to in a record-oriented transmission statement that identifies a particular record within a direct access data set.

source program
A program that serves as input to the source program processors and the compiler.

source variable
A variable whose value participates in some other operation, but is not modified by the operation. Contrast with target variable.
spill file
   Data set named SYSUT1 that is used as a temporary workfile.

standard default
   The alternative attribute or option assumed when none has been specified and there is no applicable DEFAULT statement.

standard file
   A file assumed by PL/I in the absence of a FILE or STRING option in a GET or PUT statement. SYSIN is the standard input file and SYSPRINT is the standard output file.

standard system action
   Action specified by the language to be taken for an enabled condition in the absence of an ON-unit for that condition.

statement
   A PL/I statement, composed of keywords, delimiters, identifiers, operators, and constants, and terminated by a semicolon (;). Optionally, it can have a condition prefix list and a list of labels. See also keyword statement and null statement.

statement body
   A statement body can be either a simple or a compound statement.

statement label
   See label constant.

static storage allocation
   The allocation of storage for static variables.

static variable
   A variable that is allocated before execution of the program begins and that remains allocated for the duration of execution.

stream-oriented data transmission
   The transmission of data in which the data is treated as though it were a continuous stream of individual data values in character form. Contrast with record-oriented data transmission.

string
   A contiguous sequence of characters, graphics, or bits that is treated as a single data item.

string variable
   A variable declared with the BIT, CHARACTER, or GRAPHIC attribute, whose values can be either bit, character, or graphic strings.

structure
   A collection of data items that need not have identical attributes. Contrast with array.

structure expression
   An expression whose evaluation yields a structure set of values.

structure of arrays
   A structure that has the dimension attribute.

structure member
   See member.

structuring
   The hierarchy of a structure, in terms of the number of members, the order in which they appear, their attributes, and their logical level.

subroutine
   A procedure that has no RETURNS option in the PROCEDURE statement. Contrast with function.

subroutine call
   An entry reference that must represent a subroutine, followed by an optional argument list that appears in a CALL statement. Contrast with function reference.

subscript
   An element expression that specifies a position within a dimension of an array. If the subscript is an asterisk, it specifies all of the elements of the dimension.

subscript list
   A parenthesized list of one or more subscripts, one for each dimension of the array, which together uniquely identify either a single element or cross section of the array.

subtask
   A task that is attached by the given task or any of the tasks in a direct line from the given task to the last attached task.

synchronous
   A single flow of control for serial execution of a program.
target  Attributes to which a data item (source) is converted.

target reference  A reference that designates a receiving variable (or a portion of a receiving variable).

target variable  A variable to which a value is assigned.

task  The execution of one or more procedures by a single flow of control.

task name  An identifier used to refer to a task variable.

task variable  A variable with the TASK attribute whose value gives the relative priority of a task.

termination (of a block)  Cessation of execution of a block, and the return of control to the activating block by means of a RETURN or END statement, or the transfer of control to the activating block or to some other active block by means of a GO TO statement.

termination (of a task)  Cessation of the flow of control for a task.

truncation  The removal of one or more digits, characters, graphics, or bits from one end of an item of data when a string length or precision of a target variable has been exceeded.

type  The set of data attributes and storage attributes that apply to a generation, a value, or an item of data.

U

undefined  Indicates something that a user must not do. Use of a undefined feature is likely to produce different results on different implementations of a PL/I product. In that case, the application program is in error.

union  A collection of data elements that overlay each other, occupying the same storage. The members can be structures, unions, elementary variables, or arrays. They need not have identical attributes.

union of arrays  A union that has the DIMENSION attribute.

upper bound  The upper limit of an array dimension.

V

value reference  A reference used to obtain the value of an item of data.

variable  A named entity used to refer to data and to which values can be assigned. Its attributes remain constant, but it can refer to different values at different times.

variable reference  A reference that designates all or part of a variable.

virtual origin (VO)  The location where the element of the array whose subscripts are all zero are held. If such an element does not appear in the array, the virtual origin is where it would be held.

Z

zero-suppression characters  The picture specification characters Z and *, which are used to suppress zeros in the corresponding digit positions and replace them with blanks or asterisks respectively.
Index

Special characters

/ (forward slash) 192
*PROCESS, specifying options in 93
% INCLUDE statement 94, 160
% control statement 94
%source statement library 160
%NOPRINT 94
%NOPRINT statement 94
%PAGE 94
%PAGE statement 94
%PRINT 94
%PRINT statement 94
%PROCESS, specifying options in 93
%POP statement 94
%PRINT statement 94
%PRINT statement 94
%PROCESS, specifying options in 93
%SKIP 94
%SKIP statement 94
%SKIP statement 94

A
access
ESDS 278
REGIONAL(1) data set 259
relative-record data set 297
access method services
regional data set 261
REGIONAL(1) data set
direct access 259
sequential access 259
accessibility
assistive technologies xxxv
keyboard navigation xxxv
of Enterprise PL/I for z/OS xxxv
of this information xxxvi
ACCT EXEC statement parameter 152
aggregate
length table 98
AGGREGATE compiler option 9
ALIGNED compiler suboption 24
ALL option
hooks location suboption 83
ALLOCATE statement 98
alternate ddname under z/OS UNIX, in
TITLE option 192
AMP parameter 265
ANS
compiler suboption 24
APPEND option under z/OS UNIX 194
ARCH compiler option 9, 303
argument
sort program 326
argument passing
by descriptor list 475
by descriptor-locator 476
array descriptor
array descriptor 478, 479
ASCII
compiler suboption
description 24
assembler routines
FETCHing 179
ASSIGNABLE compiler suboption 24
assistance technologies xxxv
ATTENTION ON-units 456
attention processing
attention interrupt, effect of 44
ATTENTION ON-units 456
debugging tool 456
main description 455
attribute table 98
ATTRIBUTES option 11
automatic
padding 167
prompting
overriding 166
using 165
restart
after system failure 459
checkpoint/restart facility 457
within a program 459
auxiliary storage for sort 326
avoiding calls to library routines 315

B
BACKREG compiler option 11
batch compile
OS/390 155, 157
BIPFREC compiler option 11
BIN1ARG compiler suboption 25
BKWD option 207, 272
BLANK compiler option 12
BLKOFF compiler option 13
BLKSIZE
consecutive data sets 247
ENVIRONMENT 207
comparison with DCB
subparameter 209
for record I/O 210
option of ENVIRONMENT
for stream I/O 224
subparameter 204
block
and record 200
size
consecutive data sets 247
maximum 211
object module 159
PRINT files 232
record length 211
regional data sets 262
specifying 200
BUFFERS option
for stream I/O 224
BUFND option 272
BUFFN option 273
BUFSIZE option under z/OS UNIX 195
BUFSP option 273
BYADDR
description 306
effect on performance 307
using with DEFAULT option 25
BYVALUE
description 306
effect on performance 307
using with DEFAULT option 25
C
C routines
FETCHing 179
capacity record
REGIONAL(1) 257
carriage control character 50, 231
carriage return-line feed (CR - LF) 198
caserules compiler option 13
cataloged procedure
compile and bind 141
compile only 140
compile, bind, and run 143
compile, input data for 143, 146
compile, prelink and link-edit 145
compile, prelink, link-edit, and
run 146
compile, prelink, load and run 148
description of 139
invoking 150
listing 150
modifying
DD statement 152
EXEC statement 152
multiple invocations 150
under OS/390
IBM-supplied 139
to invoke 150
to modify 151
ceestart compiler option 13
caracter string attribute table 98
caracters
 carriage control 50, 231
 print control 50, 231
 CHECK compiler option 14
checkpoint data
for sort 330
checkpoint data, defining, PLICKPT
built-in suboption 458
checkpoint/restart
deferred restart 459
PLICANC statement 460
checkpoint/restart facility
CALL PLIREST statement 459
checkpoint data set 458
description of 457
modify activity 460
PLICKPT built-in subroutine 457
request checkpoint record 457

© Copyright IBM Corp. 1999, 2015
553
consecutive data sets controlling input from the terminal capital and lowercase letters 239 condition format 239 COPY option of GET statement 239 defining QSAM files 238 end-of-file 239 format of data 237 stream and record files 238 controlling output to the terminal conditions 241 format of PRINT files 241 output from the PUT EDIT command 241 defining and using 223 input from the terminal 236 output to the terminal 241 record-oriented data transmission accessing and updating a data set 248 creating a data set 247 defining files 244 specifying ENVIRONMENT options 244 statements and options allowed 243 record-oriented I/O 243 stream-oriented data transmission accessing a data set 230 creating a data set 226 defining files 223, 224 specifying ENVIRONMENT options 224 using PRINT files 231 using SYIN and SYSPRINT files 235 control area 266 CONTROL option EXEC statement 161 interval 266 control blocks function-specific 462 global control 463 control characters carriage 50, 231 print 50, 231 COPY option 239 COPYRIGHT compiler option 17 COPYRIGHT option 448 counter records, SYSADATA 485 cross-reference table compiler listing 98 using XREF option 97 CSECT compiler option 18 CSECTCUT compiler option 18 CTLASA and CTL360 options ENVIRONMENT option for consecutive data sets 245 SCALARVARYING 214 CURRENCY compiler option 19 customizing user exit modifying SYSEXIT 463 customizing (continued) user exit (continued) structure of global control blocks 463 writing your own compiler exit 463 CYLOFL subparameter DCB parameter 204 data conversion under z/OS UNIX 191 creating under z/OS UNIX 193 sort program 330 PLISRT(x) command 335 sorting description of 319 types equivalent Java and PL/I 373 equivalent SQL and PL/I 126 data definition (DD) information under z/OS UNIX 191 data set associating PL/I files with closing a file 206 opening a file 205 specifying characteristics in the ENVIRONMENT attribute 207 associating several data sets with one file 189 blocks and records 200 checkpoint 458 conditional subparameter characteristics 204 consecutive stream-oriented data 223 data set control block (DSCB) 203 ddnames 158 defining for dump DD statement 442 logical record length 442 defining relative-record direct 202 dissociating from a file 206 system-determined block size 206 dissociating from PL/I file 189 establishing characteristics 199 indexed sequential 202 information interchange codes input in cataloged procedures label modification 205 input in cataloged procedures label modification 205 labels 203, 217 libraries extracting information 221 SPACE parameter 218 types of 217 updating 219 using 217 organization conditional subparameters 203 data definition (DD) statement 203 types of 202 partitioned 217 record format defaults 209 data set (continued) record formats fixed-length 200 undefined-length 202 variable-length 201 records 200 regional 253 REGIONAL_Level 256 accessing and updating 258 creating 257 sequential 202 sort program checkpoint data set 330 input data set 329 output data set 330 sort work data set 329 sorting 329 SORTWK 326 source statement library 160 SPACE parameter 158 stream files 223 temporary 159 to establish characteristics 199 types of comparison 215 organization 202 used by PL/I record I/O 214 unlabeled 203 using 185 VSAM blocking 266 data set type 269 defining 275 defining files 271 dummy data set 270 file attribute 270 indexed data set 279 keys 268 mass sequential insert 285 organization 266 running a program 265 specifying ENVIRONMENT options 271 VSAM option 274 VSAM performance options 274 data set under OS/390 associating one data set with several files 188 concatenating 189 HFS 189 data set under z/OS UNIX associating a PL/I file with a data set how PL/I finds data sets 193 using environment variables under 191 using the TITLE option of the OPEN statement 192 using unassociated files 193 DD_DDNAME environment variable 191 default identification 191 establishing a path 193 establishing characteristics DD_DDNAME environment variable 194 extending on output 194
data set under z/OS UNIX (continued)
maximum number of regions 197
number of regions 197
recreating output 194
data sets
allocating files 185
associating data sets with files 186
defining data sets under OS/390 187
data-directed I/O 310
coding for performance 309
DBCS compiler option 19
DBCS identifier compilation 39
DCB subparameter 206, 209
equivalent ENVIRONMENT options 209
main discussion of 204
overriding in cataloged procedure 153
regional data set 262
DD (data definition) information under z/OS UNIX 191
DD compiler option 19
DD information under z/OS UNIX
TITLE statement 192
DD statement 203
%INCLUDE 94
add to cataloged procedure 152
cataloged procedure, modifying 152
checkpoint/restart 457
create a library 218
input data set in cataloged procedure 139
modifying cataloged procedure 151
OS/390 batch compile 158
regional data set 261
standard data set 158
input (SYSIN) 159
output (SYSLIN, SYSPUNCH) 159
DD Statement
modify cataloged procedure 152
DD_DDNAME environment variables
alternate ddname under z/OS UNIX 192
APPEND 194
DELAY 195
DELIMIT 196
LRECL 196
LRMSKIP 196
PROMPT 196
PUTPAGE 196
REACCOUNT 197
RECSIZE 197
SAMELINE 197
SKIP0 198
specifying characteristics under z/OS UNIX 194
TYPE 198
ddname
%INCLUDE 94
standard data sets 158
DDSQL compiler option 20
deblocking of records 200
declaration of files under OS/390 185, 186
declaring host variables, SQL preprocessor 124
DEFAULT compiler option
description and syntax 22
suboptions
ALIGNED 24
ASCII or EBCDIC 24
ASSIGNABLE or NONASSIGNABLE 24
BINARYARG or NOBINARYARG 25
BYADDR or BYVALUE 25
CONNECTED or NONCONNECTED 26
DESCRIPT or DESLOCATOR 26
DESCRIPTOR or NODESCRIPTOR 26
DUMMY 26
E 27
EVENDEC or NOEVENDEC 27
HEXADEC 27
IBM or ANSI 24
INITFILL or NOINITFILL 27
INLINE or NOINLINE 28
LINKAGE 28
LOWERINC | UPPERINC 28
NATIVE or NONATIVE 28
NATIVEADDR or NONATIVEADDR 29
NULLSTRADDR or NONULLSTRADDR 29
NULLSTRPTR 29
NULLSYS or NULL370 29
ORDER or REORDER 30
ORDINAL(MIN | MAX) 30
OVERLAP | NOOVERLAP 30
PSEUDODUMMY or NOPSEUDODUMMY 30
RECURSIVE or NONRECURSIVE 30
RETCODE 30
RETURNS 31
SHORT 31
deferred restart 459
define data set
associating several data sets with one file 189
associating several files with one data set 188
closing a file 206
concatenating several data sets 189
ENVIRONMENT attribute 207
ESDS 277
opening a file 205
system-determined block size 206
specifying characteristics 207
define file
associating several files with one data set 188
closing a file 206
concatenating several data sets 189
ENVIRONMENT attribute 207
ESDS 277
opening a file 205
system-determined block size 206
regional data set 255
ENV options 255
keys 256
specifying characteristics 207
define file (continued)
VSAM data set 271
define file under OS/390
associating several data sets with one file 189
DEFINED
vs UNION 313
DELAY option under z/OS UNIX
description 195
DELEMIT option under z/OS UNIX
description 196
DEPRECATES compiler option 31
DEPRECATENEXT compiler option 32
DESCLIST compiler suboption 26
DESCLOCATOR compiler suboption 26
descriptor 475
descriptor area, SQL 121
DESCRIPTOR compiler option
effect on performance 307
DESCRIPTOR compiler suboption
description 26
descriptor list, argument passing 475
descriptor-locator, argument passing 476
DFSORT 319
direct data sets 202
DIRECT file
indexed ESDS with VSAM
accessing data set 283
updating data set 285
RRDS
access data set 297
DISP parameter
consecutive data sets 249
for consecutive data sets 247
to delete a data set 217
DISPLAY compiler option 33
DLL
DYNAM=DLL linker option 170
linking considerations and side-decks 163
RENT compiler option and fetching 169
DLLINIT compiler option 33
DSA
saved in PLIDUMP built-in subroutine for each block 443
DSCB (data set control block) 203, 219
DSNAME parameter
for consecutive data sets 247, 249
DSORG subparameter 204
DUMMY compiler suboption 26
dummy records
REGIONAL(1) data set 257
VSAM 270
dump
calling PLIDUMP 441
defining data set for DD statement 442
logical record length 442
identifying beginning of 442
PLIDUMP built-in subroutine 441
producing z/OS Language Environment dump 441
SNAP 442
DYNALLOC sort option 324
Index 557
hook
location suboptions 83
host
structures 130
variables, using in SQL statements 123, 124

I
I compiler message 102
IBM compiler suboption 24
IBMUEXIT compiler exit 462
IBMZC cataloged procedure 140
IBMZCB cataloged procedure 141
IBMZCPG cataloged procedure 148
IBMZCPL cataloged procedure 145
IBMZCPLG cataloged procedure 146
identifiers
not referenced 11
source program 11
IEC225I 204, 261
IGNORE option 40
ILC
linkage considerations 349
with C 341
data types 341
Enum data type 342
File type 343
Functions returning ENTRYs 348
parameter matching 344
Redirecting C standard streams 352
sharing input 352
sharing output 351
string parameter type matching 347
structure data type 342
Using the ATTACH statement 352
improving application performance 303
INCAFTER compiler option 40
INCDIR compiler option 40
INCLUDE compiler option 41
include preprocessor syntax 109
INCLUDE statement compiler 94
INDEXAREA option 207
indexed data sets
indexed sequential data set 202
indexed ESDS (entry-sequenced data set) DIRECT file 283
loading 281
SEQUENTIAL file 283
indicator variables, SQL 130
information interchange codes 200
INITFILL compiler suboption 27
initial volume label 203
initialization procedure of compiler user exit 464
INLINE compiler suboption 28
input
data for PLISRTA 335
data for sort 330
defining data sets for stream files 223, 224
redirecting 216
input (continued)
routines for sort program 330
SEQUENTIAL 247
skeletal code for sort 334
sort data set 330
input/output
compiler
data sets 159
data for compile and bind 141
data for compile, prelink, and link-edit 145
in cataloged procedures 140
OS/390, punctuating long lines 166
skeletal code for sort 332
sort data set 329
INSOURCE option 43
Inter Language Communication linkage considerations 349
Redirecting C standard streams 352
Using the ATTACH statement 352
with C 341
data types 341
Enum data type 342
File type 343
Functions returning ENTRYs 348
parameter matching 344
sharing input 352
sharing output 351
string parameter type matching 347
structure data type 342
Using C functions 343
interrupts
attention interrupt 44
interblock gap (IBG) 200
interchange codes 200
INTERNAL attribute 98
INTERRUPT compiler option 44
interrupts
attention interrupts under interactive system 44
ATTENTION ON-units 456
debugging tool 456
main description 455
invoking
cataloged procedure 150
link-editing multitasking programs 152
multiple invocations 150

J
Java 356, 357, 359, 361, 362, 363, 364, 366, 367, 368, 372
JAVA 355
Java code, compiling 357, 361, 366, 368
Java code, writing 356, 359, 364, 368
JCL (job control language)
improving efficiency 139
using during compilation 157
jni
JNI sample program 356, 359, 364, 368

K
key indexed VSAM data set 268
key-sequenced data sets
accessing with a DIRECT file 283
accessing with a SEQUENTIAL file 283
loading 281
statements and options for 279
keyboard navigation xxxv
KEYLEN subparameter 204
KEYLENGTH option 207, 214
KEYLOC option usage 207
keys
alternate index
nonunique 286
unique 286
REGIONAL(1) data set 256
dummy records 257
VSAM indexed data set 268
relative byte address 269
relative record number 269
KEYTO option under VSAM 276
KSDS (key-sequenced data set)
define and load 281
unique key alternate index path 287
updating 283
VSAM DIRECT file 283
loading 281
SEQUENTIAL file 283
KSDS with VSAM
indexed ESDS with VSAM access data set 283
accessing data set 283
updating data set 285

L
label
for data sets 203
LANG_LVL compiler option 45
Language Environment library xiv
LEAVE and REREAD options
ENVIRONMENT option
for consecutive data sets 246
length of record
specifying under z/OS UNIX 197
library
compiled object modules 219
creating a data set library 218
creating a member 220
creating and updating a library member 219
creating, examples of 219
directory 218
extracting information from a library directory 221
general description of 202
how to use 217
information required to create 218
placing a load module 220
source statement 160
source statement library 155
library (continued)

SPACE parameter 218
structure 221
system procedure
(SYSI.PROCLIB) 217
types of 217
updating a library member 221

using 217

LIMCT subparameter 204, 262
LIMITS compiler option 46
line
length 232
numbers in messages 38
line feed (LF)
definition 199
LINE option 225, 232
LINECOUNT compiler option 47
LINEDIIR compiler option 47
LINESIZE option
for tab control table 234
OPEN statement 225

link-editing
description of 163
LINKAGE compiler suboption
effect on performance 308
syntax 28
linkage considerations with ILC 349
LIST compiler option 47
listing
cataloged procedure 150
compiler
aggregate length table 98
ATTRIBUTE and cross-reference
table 97
ddbname list 5
file reference table 102
heading information 96
messages 102
options 96
preprocessor input 97
return codes 103
source program 78
statement offset address 99
storage offset listing 101
OS/390 batch compile 155, 159
source program 78
statement offset address 99
storage offset listing 101
SYSPRINT 159
LISTVIEW compiler option 48
literal records, SYSADATA 485
loader program, using 148
logical not 56
logical or 59
loops
control variables 311
LOWERINC compiler suboption 28
LRECL option under z/OS UNIX 196
LRECL subparameter 200, 204
LRMSKIP option under z/OS UNIX 196

MACRO option 49
macro preprocessor
macro definition 110
main storage for sort 325
MAP compiler option 49
MARGINI compiler option 49
MARGINS compiler option 50
mass sequential insert 285
MAXBRANCH compiler option 51
MAXGEN compiler option 51
MAXMEM compiler option 52
MAXMSC compiler option 52, 53
MAXSTMT compiler option 53
MAXTEMP compiler option 53
MDECK compiler option
description 54
message
compiler list 102
printed format 235
runtime message line numbers 38
message records, SYSADATA 486
messages
filter function 465
modifying in compiler user exit 463
mixed string constant compilation 39
MODE subparameter
usage 204
module
create and store object module 57
MSCSUMMARY compiler option 54
multiple
invocations
cataloged procedure 150
Multitasking
options in PLIDUMP 441
NAME compiler option 55
named constants
defining 313
versus static variables 313
NAMES compiler option 55
COMPILER suboption
description 28
NATIVEADDR compiler suboption 29
NATLANG compiler option 55
negative value
block-size 211
record length 210
NEST option 56
NODESCRIPTOR compiler suboption 25
NODESCRIPTOR compiler suboption 26
NOEQUALS sort option 324
NOEVENDEC compiler suboption 27
NONCONNECT compiler suboption 28
NONCONNECT compiler suboption 29
NONCONNECT compiler suboption 29
NONNULLSTRADDR compiler
suboption 29
NONNULLSTRADDR compiler
suboption 29
NONOVERLAP compiler suboption 30
effect on performance 309
NOPSEUDO DUMMY compiler
suboption 30
NOSYNTAX compiler option 80
NOT compiler option 56
note statement 102
NTM subparameter
usage 204
NULL370 compiler suboption 29
NULLSTRADDR compiler suboption 29
NULLSTRPTR compiler suboption 29
NULLSYS compiler suboption 29
NUMBER compiler option 56

O

object
module
create and store 57
record size 159
OBJECT compiler option
definition 57
offset
tab count 234
table 99
OFFSET compiler option 57
ONSNAPE compiler option 57
OPEN statement
subroutines of PL/I library 205
TITLE option 204
Operating system
data definition (DD) information
under z/OS UNIX 191

OPTCD subparameter 203, 204
optimal coding
coding style 309
compiler options 303
OPTIMIZE compiler option 303
OPTIMIZE option 58
options
for compiling 96
for creating regional data set 253
saved options string in
PLIDUMP 449
specifying comments within 92
specifying strings within 92
to specify for compilation 160
OPTIONS option 59
options under z/OS UNIX

DD_DDNAME environment variables
APPEND 194
DELAY 195
DELIMIT 196
LRECL 196
LRMSKIP 196
PROMPT 196
PUTPAGE 196
RECCOUNT 197
RECSIZE 197
SAMELINE 197
SKIP 198
TYPE 198
PL/I ENVIRONMENT attribute
BUFSIZE 195
options under z/OS UNIX (continued) using
	DD information 192
TITLE 192
OR compiler option 59
ORDER compiler suboption description 30
effect on performance 308
ORDINAL compiler suboption 30
ordinal element records, SYSADATA 487
ordinal type records, SYSADATA 488
ORGANIZATION option 214
usage 207
OS/390
batch compilation
DD statement 158
EXEC statement 157, 161
listing (SYSPRINT) 159
source statement library (SYSLIB) 160
specifying options 160
temporary workfile (SYSUT1) 159
general compilation 155
output data for PLISRTA 335
data for sort 330
defining data sets for stream files 223, 224
limit preprocessor output 54
redirecting routines for sort program 330
SEQUENTIAL 247
skeletal code for sort 334
sort data set 330
SYSIN 159
SYSOUT 159
OVERLAP compiler suboption 30

P

PACKAGES versus nested
PROCEDURES 311
PAGE option 225
PAGELENGTH, for tab control table 234
PAGESIZE, for tab control table 234
parameter passing
argument passing 475
CMPAT(LE) descriptors 478
CMPAT(V) descriptors 476
PARM parameter
for cataloged procedure 152
specify options 161
passing an argument 475
PASSWORD option 273
performance
VSAM options 274
performance improvement (continued)
coding for performance (continued)
named constants versus static variables 313
PACKAGES versus nested
PROCEDURES 311
preloading calls to library routines 315
REDUCIBLE functions 312
string assignments 310
selecting compiler options
ARCH 303
DEFAULT 306
GONUMBER 303
OPTIMIZE 303
PREFIX 305
REDUCE 304
RULES 304
PL/I compiler
user exit procedures 462
files associating with a data set under z/OS UNIX 191
PL/I code, compiling 359, 363, 367, 372
PL/I code, linking 359, 363, 367, 372
PL/I code, writing 357, 362, 366, 368
PL/I dynamic allocation
access HFS files under z/OS 189
allocate data sets 185
associating data sets with files under z/OS UNIX 191
define files
QSAM files 238
REGIONAL(1) data sets 255
stream files 224
VSAM file 265
PL/I MAIN Routines
FETCHing 178
PLICANC statement, and checkpoint/request 460
PLICPT built-in subroutine 457
PLIDUMP built-in subroutine calling to produce a z/OS Language Environment dump 441
DSA saved for each block 443
H option 442
output locating variables in 443
program unit name in dump traceback table 443
saved load module timestamp 449
saved options string 449
syntax of 441
user-identifier 442
variables locating AUTOMATIC variables in 443
locating CONTROLLED variables in 445
locating in PLIDUMP output 443
locating STATIC variables in 444
PLIREST statement 459
PLIRETC built-in subroutine return codes for sort 329
PLISAXA 377, 378
PLISAXB 377, 378
PLISAXC 407, 408
PLISAXD 407, 408
PLISRTA interface 335
PLISRTB interface 336
PLISRTC interface 337
PLISRTD interface 338
PLITABS external structure control section 235
declaration 164
PLIXOPT variable 164
PP compiler option 60
PPCICS compiler option 61
PPINCLUDE compiler option 61
PPLIST compiler option 62
PPMACRO compiler option 62
IPSQGL compiler option 63
PPTRACE compiler option 63
PRECTYPE compiler option 63
PREFIX compiler option 64, 305
using default suboptions 305
preloading library routines 315
preprocessing
%INCLUDE statement 94
input 97
limit output to 80 bytes 54
source program 49
with MACRO 49
preprocessor options
CCSID0 118
CODEPAGE 118
DEPRECREATE 118
EMPTYDBRM 119
INCONLY 119
WARNDECP 119
preprocessors
available with PL/I 109
CICS options 137
include 109
macro preprocessor 110
SQL options 117
SQL preprocessor 114
print
PRINT file
format 241
line length 232
stream I/O 231
print control character 50, 231
PRINT file
formatting conventions 164
punctuating output 164
record I/O 250
procedure
cataloged, using under OS/390 139
compile and bind (IBMZCB) 141
compile and link-edit (IBMZCPL) 145
compile only (IBMZC) 140
compile, bind, and run (IBMZCBG) 143
compile, prelink, link-edit, and run (IBMZCPLG) 146
record, prelink, and run (IBMZCPO) 148
PROCEED compiler option 64
PROCESS statement description 93
override option defaults 160
program unit name in dump traceback table 443
PROMPT option under z/OS UNIX 196
prompting automatic, overriding 166
automatic, using 165
PRTSP subparameter usage 204
PSEUDODUMMY compiler suboption 30
punctuation automatic prompting overriding 166
using 165
OS/390 automatic padding for GET EDIT 167
continuation character 166
entering ENDFILE at terminal 167
GET DATA statement 166
GET LIST statement 166
long input lines 166
SKIP option 167
output from PRINT files 164
PUT EDIT command 241
PUTPAGE option under z/OS UNIX 196
Q
QUOTE compiler option 65
R
REAL attribute 98
RECCOUNT option under z/OS UNIX 197
RECFM subparameter 204
in organization of data set 204
usage 204
record checkpoint 457
data set 458
deblocking 200
maximum size for compiler input 159
sort program 324
record format fixed-length records 200
options 225
stream 1/O 230
to specify 244
types 200
undefined-length records 202
variable-length records 201
record 1/O data set access 248
consecutive data sets 249
create 247
types of 214
data transmission 243
ENVIRONMENT option 244
format 209
record format 244
record length regional data sets 253
specify 200
value of 210
RECORD statement 324
recorded key regional data set 256
records length under z/OS UNIX 197
RECSIZE option consecutive data set 225
defaults 225
definition 210
description under z/OS UNIX 197
for stream 1/O 224, 225
RECURSIVE compiler suboption 30
REDUCE compiler option 66
effect on performance 304
reduce storage requirement 58
REDUCIBLE functions 312
region REGION parameter 152
size, EXEC statement 157
REGION size, minimum required 139
regional data sets DD statement
accessing 263
creating 261
defining files for regional data set 255
specifying ENVIRONMENT options 255
using keys 256
operating system requirement 261
REGIONAL(1) data set accessing and updating 258
creating 257
using 256
REGIONAL option of ENVIRONMENT 256
REGIONAL(1) data sets defining files for regional data set 255
regions under z/OS UNIX 197
relative byte address (RBA) 269
relative record number 269
relative-record data sets accessing with a DIRECT file 297
accessing with a SEQUENTIAL file 296
loading 294
statements and options for 292
RENT compiler option 66
REORDER compiler suboption description 30
effect on performance 308
REEXP compiler option 68
RESPECT compiler option 68
restarting requesting 459
RESTART parameter 459
to request automatic after system failure 459
automatic within a program 459
deferred restart 459
to cancel 459
to modify 460
RETCODE compiler suboption 30
return code
checkpoint/restart routine 457
PLIRETC 329
return codes in compiler listing 103
RETURNS compiler suboption 31, 309
REUSE option 207, 273
RRDS (relative record data set) define 295
load statements and options 292
load with VSAM 294
updating 297
VSAM DIRECT file 297
loading 294
SEQUENTIAL file 296
RTCHECK compiler option 68
RULES compiler option 69
effect on performance 304
run time message line numbers 38
run-time
OS/390 considerations automatic prompting 165
formatting conventions 164
GET EDIT statement 167
GET LIST and GET DATA statements 166
punctuating long lines 166
SKIP option 167
using PLIXOPT 164
S
S compiler message 102
SAMELINE option under z/OS UNIX 197
sample program, running 359, 364, 368, 372
saved options string in PLIDUMP 449
SAX parser 377, 407
SCALARVARYING option 213
SEMANTIC compiler option 77
sequential access REGIONAL(1) data set 259
sequential data set 202
SEQUENTIAL file ESDS with VSAM defining and loading 277
updating 278
indexed ESDS with VSAM access data set 283
RRDS, access data set 296
serial number volume label 203
SERVICE compiler option 78
shift code compilation 39
SHORT compiler suboption 31
SKIP option 274
in stream 1/O 225
under OS/390 167
SKIP0 option under z/OS UNIX 198
SKIPREC sort option 324
sorting assessing results 328
calling sort 326
CHKPT option 324
Index 561
sorting (continued)
choosing type of sort 320
CKPT option 324
data 319
data input and output 330
description of 319
DYNALLOC option 324
E15 input handling routine 331
EQUALS option 324
FILSZ option 324
maximum record length 325
PLISRT 319
PLISRTA(x) command 335, 339
preparation 319
RECORD statement 331
RETURN statement 331
SKIPREC option 324
SORTCKPT 330
SORTCNTL 330
SORTIN 329
sorting field 323
SORTLIB 329
SORTOUT 330
SORTWK 326, 329
storage
auxiliary 326
main 325
writing input/output routines 330
source
key
in REGIONAL(1) data sets 256
listing
location 50
program
compiler list 97
data set 159
identifiers 11
included in compiler list 43
list 78
preprocessor 49
shifting outside text 49
SOURCE compiler option 78
source records, SYSADATA 505
source statement library 160
SPACE parameter
library 218
standard data sets 158
specifying compile-time options
using flags under 156
SPILL compiler option 78
spill file 159
SQL preprocessor
communications area 120
descriptor area 121
EXEC SQL statements 114
options 117
using host structures 130
using host variables 123
using IBMUEXIT with 135
using indicator variables 130
SQL preprocessor options 118, 119
SQLCA 120
SQLDA 121
STACK subparameter
usage 204
standard data set 158
standard files (SYSPRINT and SYSIN) 216
statement
nesting level 97
offset addresses 99
statements 94
STATIC compiler option 78
STDSYS compiler option 79
step abend 203
STMT compiler option 79
STMT suboption of test 83
storage
blocking print files 232
library data sets 218
report in listing 79
sort program 325
auxiliary storage 326
main storage 325
standard data sets 158
to reduce requirement 58
STORAGE compiler option 79
stream and record files 238, 241
STREAM attribute 223
stream I/O
consecutive data sets 223
data set
access 230
create 226
record format 230
DD statement 228, 231
ENVIRONMENT options 224
file
define 223
PRINT file 231
SYSLIB
%INCLUDE 94
preprocessing 160
SYSIN 159
SYSLIB
%INCLUDE 94
preprocessing 160
SYSLIN 159
SYSOUT 329
SYSPARM compiler option 81
SYSPUNCH 159
system
failure 459
restart after failure 459
SYSTEM compiler options
SYSTEM(CICS) 81
SYSTEM(IMS) 81
SYSTEM(MVS) 81
SYSTEM(OS) 81
SYSTEM(TSO) 81
type of parameter list 81
SYSUTI compiler data set 159

T

tab control table 234
temporary workfile
SYSUTI 159
terminal
input 236
capital and lowercase letters 239
COPY option of GET
statement 239
defining QSAM files 238
terminal (continued)
input (continued)
end of file 239
format of data 237
stream and record files 238
output 241
format of PRINT file 241
output from PUT EDIT command 241
stream and record files 241
TERMINAL compiler option 82
terminating
compilation 17
termination procedure
compiler user exit 467
example of procedure-specific control block 467
syntax
global 463
specific 467
TEST compiler option
definition 83
TIME parameter 152
TIMESTAMP
saved load module timestamp in PLIDUMP 449
TITLE option
associating standard SYSPRINT file 168
description under z/OS UNIX 192
using 204
TITLE option under OS/390
specifying character string value 187
TITLE option under z/OS UNIX
using files not associated with data sets 193
token records, SYSADATA 506
traceback table
program unit name in 443
trailer label 203
TYPE option under z/OS UNIX 198
U
U compiler message 102
U option of ENVIRONMENT
for record I/O 209
for stream I/O 224
variable-length records
format 201
sort program 339
VB option of ENVIRONMENT
for record I/O 209
for stream I/O 224
VB-format records 201
VBS option of ENVIRONMENT
for stream I/O 224
VOLUME parameter
consecutive data sets 247, 249
volume serial number
direct access volumes 203
regional data sets 261
VS option of ENVIRONMENT
for stream I/O 224
VSAM (virtual storage access method)
data sets
alternate index paths 274
alternate indexes 285
blocking 266
choosing a type 269
defining 265, 275
defining files for 271
dummy data set 270
entry-sequenced 276
file attribute 270
key-sequenced and indexed
entry-sequenced 279
keys for 268
organization 266
performance options 274
relative record 292
running a program with 265
specifying ENVIRONMENT options 271
using 265
defining files 271
ENV option 271
performance option 274
VSAM (virtual storage access method) (continued)
indexed data set
load statement and options 279
mass sequential insert 285
relative-record data set 294
VTOC 203
W
W compiler message 102
WIDECHAR compiler option 87
WINDOW compiler option 87
work data sets for sort 329
WRITABLE compiler option 88
X
XINFO compiler option 89
XML
support in the SAX parser 377, 407
XML compiler option 45, 91
XREF compiler option 91
Z
z/OS UNIX
compile-time options
specifying 156
compiling under 155
DD_DDNAME environment variable 194
export command 194
setting environment variables 215
specifying compile-time options
command line 156
using flags 156
zero value 210
Readers’ Comments — We'd Like to Hear from You

Enterprise PL/I for z/OS
Programming Guide
Version 4 Release 5

Publication No. GI11-9145-04

We appreciate your comments about this publication. Please comment on specific errors or omissions, accuracy, organization, subject matter, or completeness of this book. The comments you send should pertain to only the information in this manual or product and the way in which the information is presented.

For technical questions and information about products and prices, please contact your IBM branch office, your IBM business partner, or your authorized remarketer.

When you send comments to IBM, you grant IBM a nonexclusive right to use or distribute your comments in any way it believes appropriate without incurring any obligation to you. IBM or any other organizations will only use the personal information that you supply to contact you about the issues that you state on this form.

Comments:

Thank you for your support.
Send your comments to the address on the reverse side of this form.
If you would like a response from IBM, please fill in the following information:

Name

Address

Company or Organization

Phone No.

Email address